Gottesman-Kitaev-Preskill bosonic error correcting codes: a lattice perspective

Quantum 6, 648 (2022)

Jonathan Conrad, Jens Eisert, Francesco Arzani







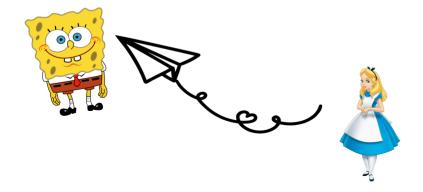




Outline

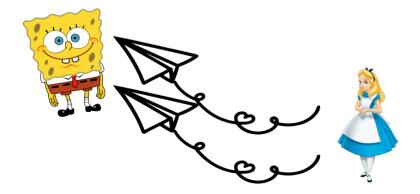
- 1) Introduction and motivation
- 2) Some definitions: Lattices in a nutshell
- 3) Results: Lattice bases and symplectically equivalent codes
- 4) Conclusions

Part 1: Introduction and motivation



Information always encoded in phys. syst.

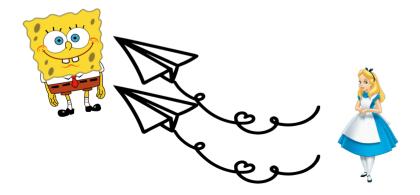
Always subject to **noise**



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Error correction: redundancy

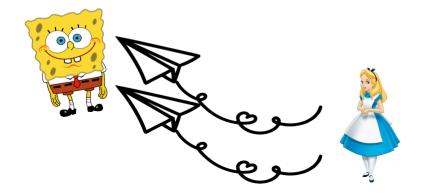


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Quantum: mostly qubits $|\alpha| |0\rangle + \beta |1\rangle$



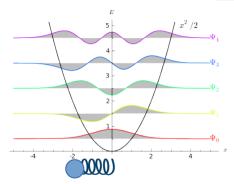
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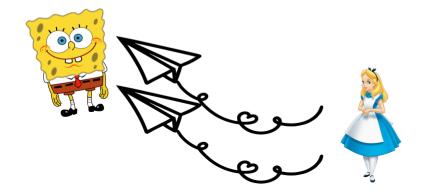
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EM field mode, LC circuit, ...



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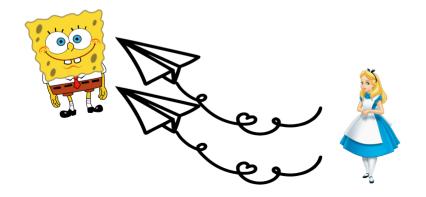
$$\frac{E}{\sqrt{x^2/2}}$$
 Ψ_1
 Ψ_2
 Ψ_2
 Ψ_1
 Ψ_2
 Ψ_3

$$\hat{\boldsymbol{x}} = (q_1, \dots, q_n, p_1, \dots, p_n)^T$$

$$[\hat{x}_j, \hat{x}_k] = iJ_{jk}$$

$$J = \left(egin{array}{cc} \mathbf{0} & \mathbb{I} \ -\mathbb{I} & \mathbf{0} \end{array}
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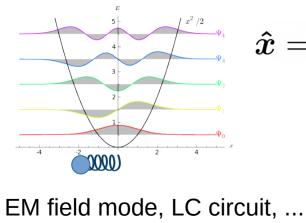
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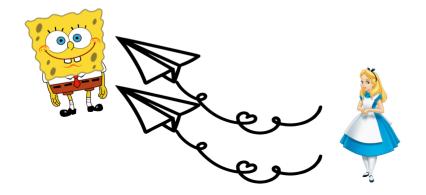
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$$|\psi\rangle = \int \mathrm{d}q\psi(q) \, |q\rangle$$

Infinite dimensional space! How to exploit?



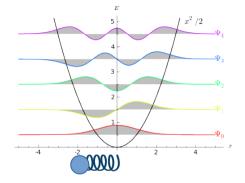
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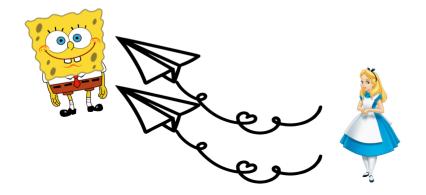
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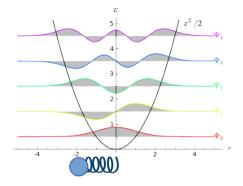
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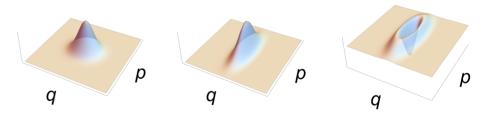
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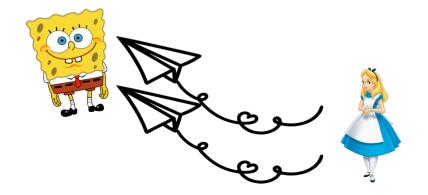
Symmetries!

In phase space: Wigner Function



Quasi-probability distribution

EM field mode, LC circuit, ...

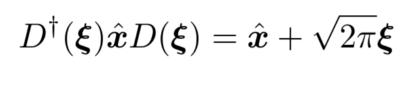


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$$D(\boldsymbol{\xi})D(\boldsymbol{\eta}) = e^{-i2\pi\boldsymbol{\xi}^T J\boldsymbol{\eta}}D(\boldsymbol{\eta})D(\boldsymbol{\xi})$$

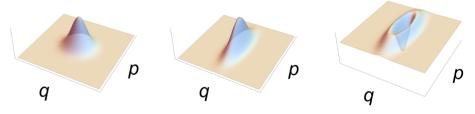
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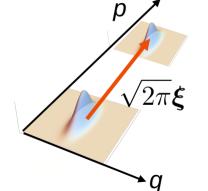
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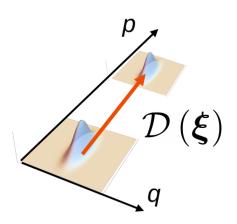
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"Logical subspace": finite photon number, odd/even photon number, or...translation symmetries!

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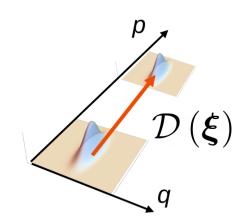
Grid codes: **stabilized** by (commuting) **displacement** operators



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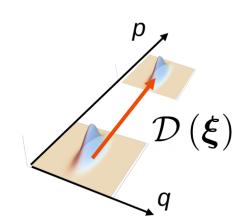


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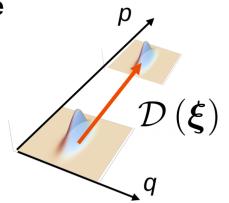
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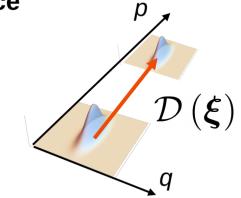
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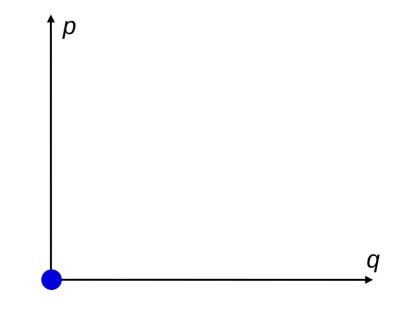
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Ex:1 qubit in 1 mode
$$\rightarrow \boldsymbol{\xi}_q, \ \boldsymbol{\xi}_p$$



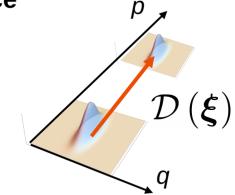
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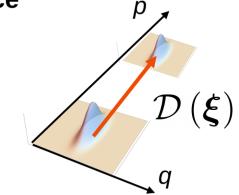
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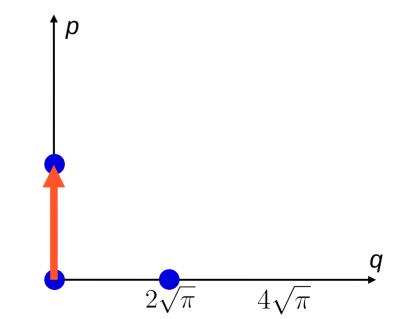
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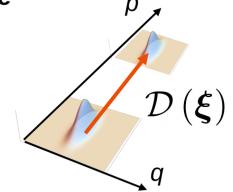
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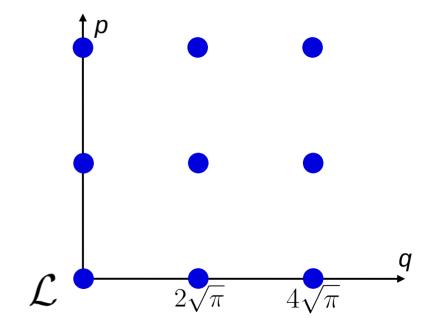
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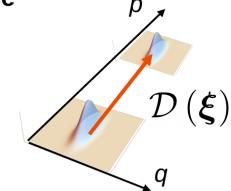
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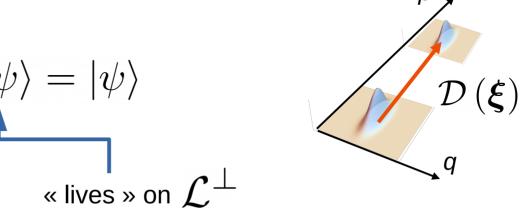
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Gottesman, Kitaev, Preskill PRA 64 (2001) GKP codes

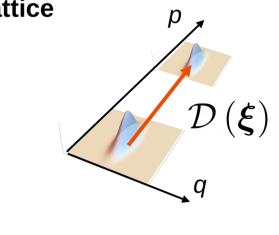
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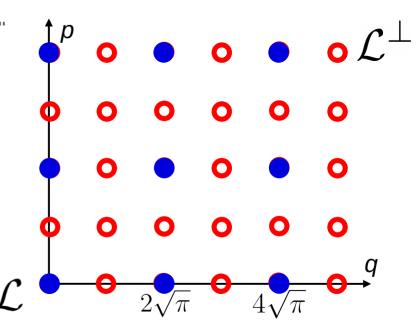
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Good protection against common noise processes
 Albert et al, PRA 97 (2018)



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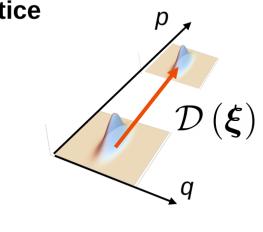
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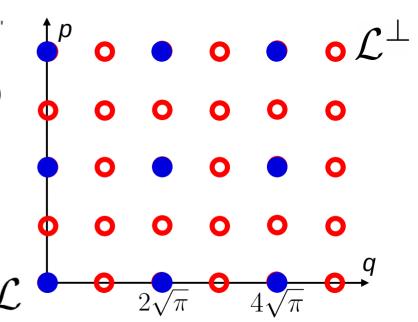
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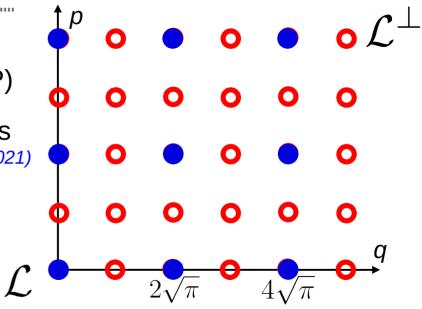
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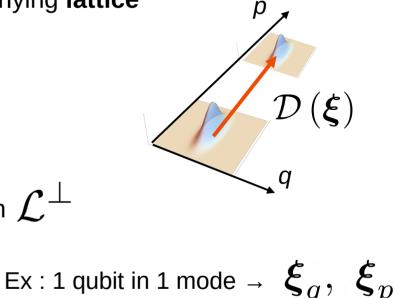
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<u>Grid codes</u>: stabilized by (commuting) displacement operators → underlying lattice

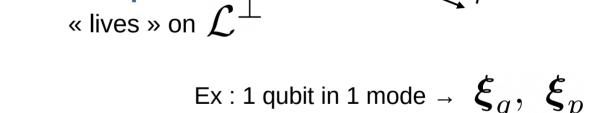
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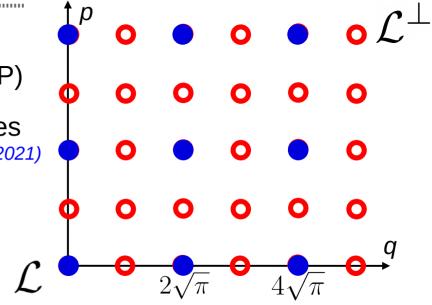
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- Can be used as effective qubits and combined with stabilizer codes Vuillot et al, PRA 99 (2019) Noh&Chamberland PRA 101 (2020) Bourassa et al, Quantum 5 (2021)
- Can protect CV systems

Noh et al, PRL 125 (2020)



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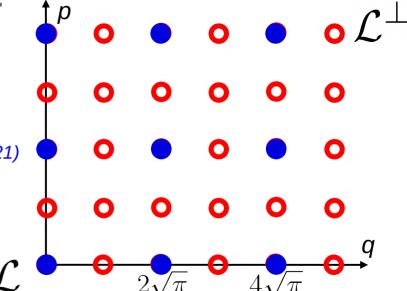
$$\mathcal{S} \cong \mathcal{L} = \{\sum_j z_j oldsymbol{\xi}_j \; : \; z_j \in \mathbb{Z} \}$$





« lives » on \mathcal{L}^{\perp}

- Good protection against common noise processes Albert et al, PRA 97 (2018)
- Logical Clifford = Gaussian operations ("easy" good for EC & QIP) Gottesman, Kitaev, Preskill PRA 64 (2001)
- Can be used as effective qubits and combined with stabilizer codes Vuillot et al, PRA 99 (2019) Noh&Chamberland PRA 101 (2020) Bourassa et al, Quantum 5 (2021)
- Can protect CV systems Noh et al. PRL 125 (2020)
- Logical states thought hard to realize, now there are experiments! Flühmann et al, Nature 566 (2019) Campagne-Ibarcg et al, Nature 584 (2020)



For exponential noise suppression: more oscillators

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Up to now: *concatenation* → regard as effective qubits, add qubit-level code

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Results

- 1. Code properties from lattice bases
- 2. Symplectic operations
- 3. Distance bounds for GKP codes
- 4. Decoding problem and Θ functions
- 5. GKP codes beyond concatenation

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Results

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Part 2 : Definitions

Start from

$$S = \langle D(\boldsymbol{\xi}_1), \dots, D(\boldsymbol{\xi}_{2n}) \rangle \quad \leftrightarrow \quad M = (\boldsymbol{\xi}_1, \dots, \boldsymbol{\xi}_{2n})^T$$

$$\Longrightarrow \left[S \simeq \mathcal{L} = \left\{ \boldsymbol{\xi} \in \mathbb{R}^{2n} | \boldsymbol{\xi}^T = \boldsymbol{a}^T M, \ \boldsymbol{a} \in \mathbb{Z}^{2n} \right\} \right]$$

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Commutation \leftrightarrow Integer symplectic product $\left[D\left(\boldsymbol{\xi}_{j}\right),D\left(\boldsymbol{\xi}_{j}\right)\right]=0\Leftrightarrow\boldsymbol{\xi}_{j}^{T}J\boldsymbol{\xi}_{k}\in\mathbb{Z}$

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$$A_{jk} = \left(MJM^T\right)_{jk} \in \mathbb{Z}$$
 Symplectically integral lattices

Symplectic Gram matrix

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 Symplectically integral lattices

Symplectic Gram matrix

Each linearly independent generator "discretizes" 1 direction in phase space

Focus on full rank lattices: $\det M \neq 0 \Rightarrow d < \infty$ Logical dimension

Logical Pauli operators:

all displacements that commute with stabilizers

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all displacements that commute with stabilizers

all vectors that have integer symplectic product with lattice vectors

$$\left\{\boldsymbol{\xi}^{\perp} \in \mathbb{R}^{2n} \mid \left(\boldsymbol{\xi}^{\perp}\right)^{T} J \boldsymbol{\xi} \in \mathbb{Z} \ \forall \boldsymbol{\xi} \in \mathcal{L}\right\}$$

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$$igsplace$$
 (symplectically) dual lattice $\mathcal{L}^\perp = \left\{m{\xi}^\perp \in \mathbb{R}^{2n} \mid \left(m{\xi}^\perp
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One can show that the logical dimension *d* is computed from:

$$d^2 = |\mathcal{L}^{\perp}/\mathcal{L}| = |\det M|/|\det M^{\perp}| = \det A = (\det M)^2$$

Part 3: Results

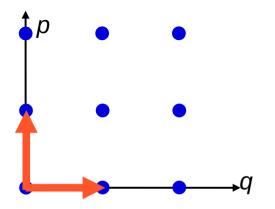
Basis (stabilizer generators) not unique given lattice (code)

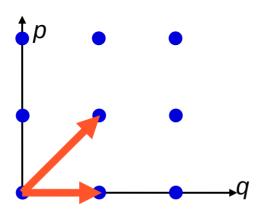
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Example: square lattice



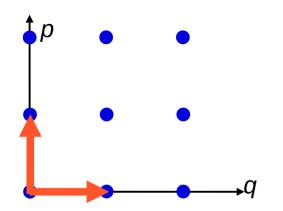


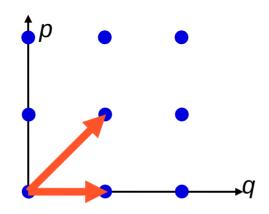
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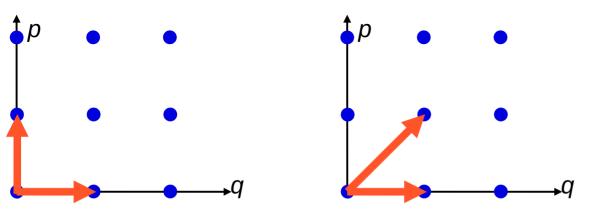
Exploit basis choice and properties to study codes

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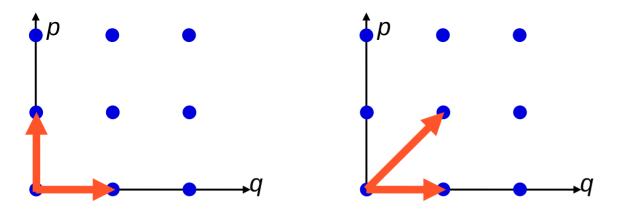
Exploit basis choice and properties to study codes

Theorem (Hadamard's bound):

Let
$$C = \max_{j} ||\boldsymbol{\xi}_{j}||, \ d = 2^{k} = \det M$$
. Then $k \leq 2n \log_{2} C$

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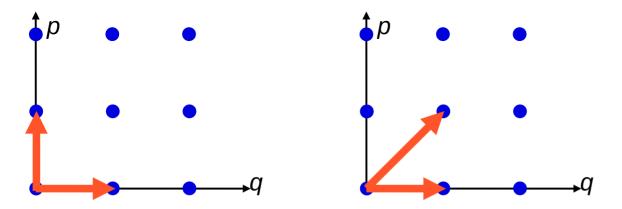
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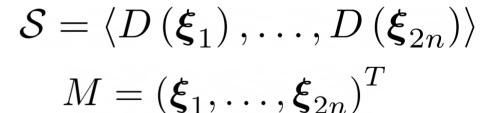
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Intuitively : $C\sim$ maximum interaction strength with ancilla to measure a stabilizer generator

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Exploit basis choice and properties to study codes

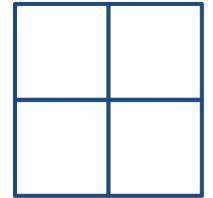
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Intuitively : $C \sim$ maximum interaction strength with ancilla to measure a stabilizer generator Encoding ratio related to "hardness" of measuring generators in the chosen basis

Concatenation:

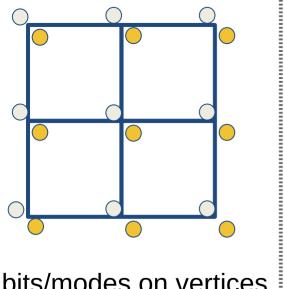
Example : L=3 GKP-surface code (n = 9 modes, logical dimension d = 2)



qubits/modes on vertices

Concatenation: 1) fix a qubit in each of *n* modes (logical dimension $d = 2^n$) (think 1 square lattice for each mode) $\rightarrow 2n$ generators

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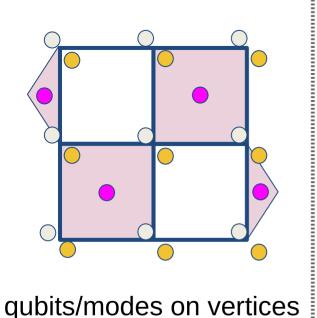


$$\circ \exp\left(i2\sqrt{\pi}\hat{q}_j\right)$$
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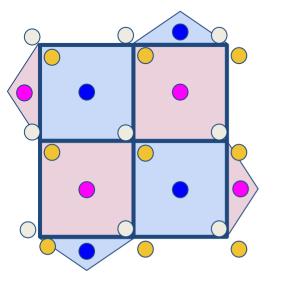
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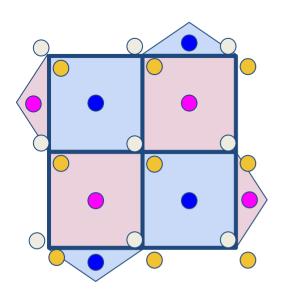
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- **Concatenation**: 1) fix a qubit in each of *n* modes (logical dimension $d = 2^n$) (think 1 square lattice for each mode) $\rightarrow 2n$ generators
 - 2) enlarge stabilizer to include additional *logical* Pauli strings → displacements in <u>dual lattice</u>

Example : L=3 GKP-surface code (n = 9 modes, logical dimension d = 2)



$$2n + n - 1 = 26$$
 stab.

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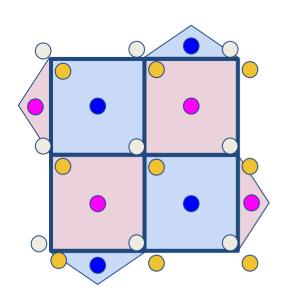
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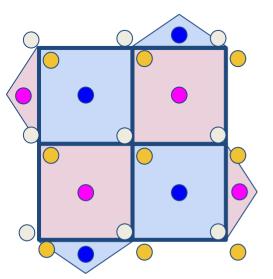
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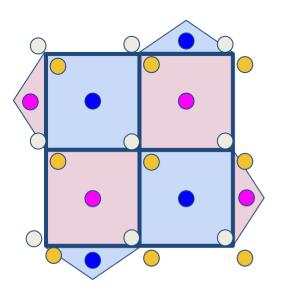
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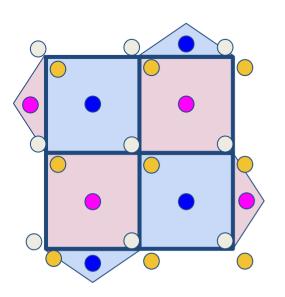
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 lattice basis reduction $M_{
m min}$ $2n$ = 18 stab. gen. can do respecting weights and geometric locality!

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Higher level code

$$M_{
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m Q} \end{array} \right)$$

lattice basis reduction

$$M_{
m min}$$

2n = 18 stab. gen.

can do respecting weights and geometric locality!



less measurements per EC cycle

$$egin{align} U_S &= \exp\left(-ioldsymbol{\hat{x}}^T H oldsymbol{\hat{x}}
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Theorem (symplectically equivalent codes):

Given
$$\mathcal{L}(M), \mathcal{L}(N), \exists S \mid M = NS^T \text{ iff } MJM^T = NJN^T \text{ (in canonical form)}$$

Multi-mode generalization of Hänggli, Heinze, König, PRA 102 (2020)

$$A_{M,N} = \left(\begin{array}{cc} 0 & D \\ -D & 0 \end{array}\right)$$

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Corollary:

Any code with
$$d = 2$$
 is s.e. to $S_{\square}^{(2)} = \left\langle e^{i2\sqrt{\pi}\hat{q}_1}, e^{-i2\sqrt{\pi}\hat{p}_1}, e^{i\sqrt{\pi}\hat{q}_j}, e^{i\sqrt{\pi}\hat{p}_j} \right\rangle, \ j > 1$

one qubit in mode 1, a fixed state on others

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- Generalizes to higher logical dimensions
- Can be used to bound distance of (any) given grid code (see paper)

Part 4: Conclusions

Presentation

- Bosonic error correction
- Lattice formalism
- Basis reduction
- Symplectic equivalence

Paper

Presentation

- Bosonic error correction
- Lattice formalism
- Basis reduction
- Symplectic equivalence

- Code distance
 - Maximum-Likelihood decoding
 - New code constructions

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Lattices for bosonic codes

Paper

Presentation

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Thank you!