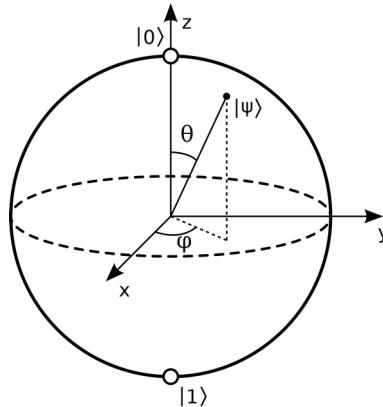


Polynomial approximation of non-Gaussian unitaries by counting one photon at a time

Francesco Arzani, Nicolas Treps, Giulia Ferrini

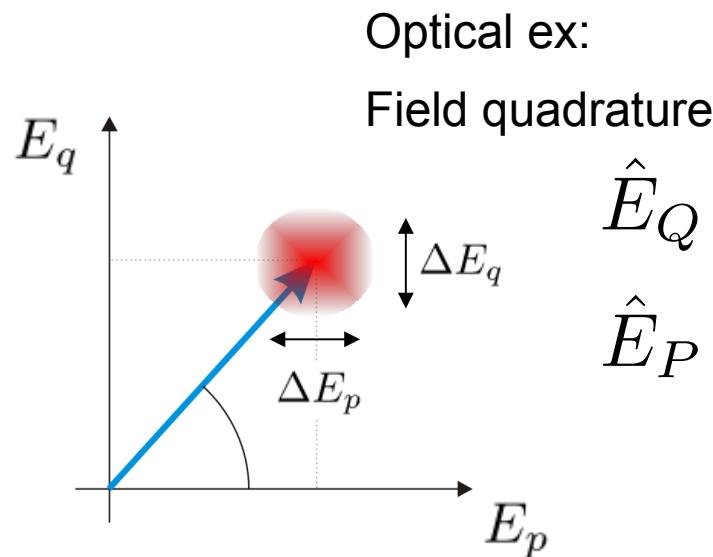


DV : information encoded in qu-bits



Optical ex:
Polarization of
single photon

CV : information encoded in observables with continuous spectrum, e.g. : \hat{q} , \hat{p}



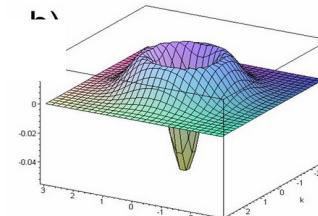
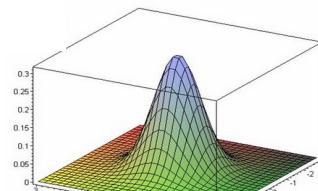
$$\hat{E}_Q \propto \hat{a} + \hat{a}^\dagger$$

$$\hat{E}_P \propto i(\hat{a}^\dagger - \hat{a})$$

Wigner function

$$W(q, p) = \frac{1}{2\pi} \int dx e^{ipx} \left\langle q - \frac{x}{2} \right|_q \hat{\rho} \left| q + \frac{x}{2} \right\rangle_q$$

$$\langle \hat{O} \rangle = \text{Tr} [\hat{\rho} \hat{O}] = \int dq dp W_\rho(q, p) W_O(q, p)$$



For **pure states**:

Gaussian
iff
Non-negative

R. Hudson
Rep. Math. Phys. 6, 249 (1974)

c)

Discrete encoding $\vec{x} = 1001111010101001\dots$

bits $\vec{x} \mapsto b(\vec{x})$ For any boolean function

qubits $|\vec{x}\rangle|\phi\rangle \mapsto U_b|\vec{x}\rangle|\phi\rangle = |\vec{x}\rangle|b(\vec{x})\rangle$

Continuous Variables

qumodes $|\psi\rangle \in (L^2(\mathbb{R}, \mathbb{C}))^{\otimes n} \mapsto e^{-itH(\hat{a}, \hat{a}^\dagger)}|\psi\rangle$

General polynomial

Universal set:

$$\left\{ e^{i\hat{q}s}, e^{i\hat{q}^2 s}, e^{i\frac{\pi}{4}(\hat{q}^2 + \hat{p}^2)}, e^{i\hat{q}_1 \otimes \hat{q}_2}, e^{i\hat{q}^3 s} \right\}$$

Single-mode, Gaussian

Non-Gaussian

Two-modes C_Z

Computation with arbitrary encoding

No quantum advantage without non-Gaussianity !

*A. Mari and J. Eisert
PRL 109, 230503 (2012)*

Actually, it's negativity of the WF...

Continuous Variables

qumodes $|\psi\rangle \in (L^2(\mathbb{R}, \mathbb{C}))^{\otimes n} \mapsto e^{-itH(\hat{a}, \hat{a}^\dagger)} |\psi\rangle$

General polynomial

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Single-mode, Gaussian

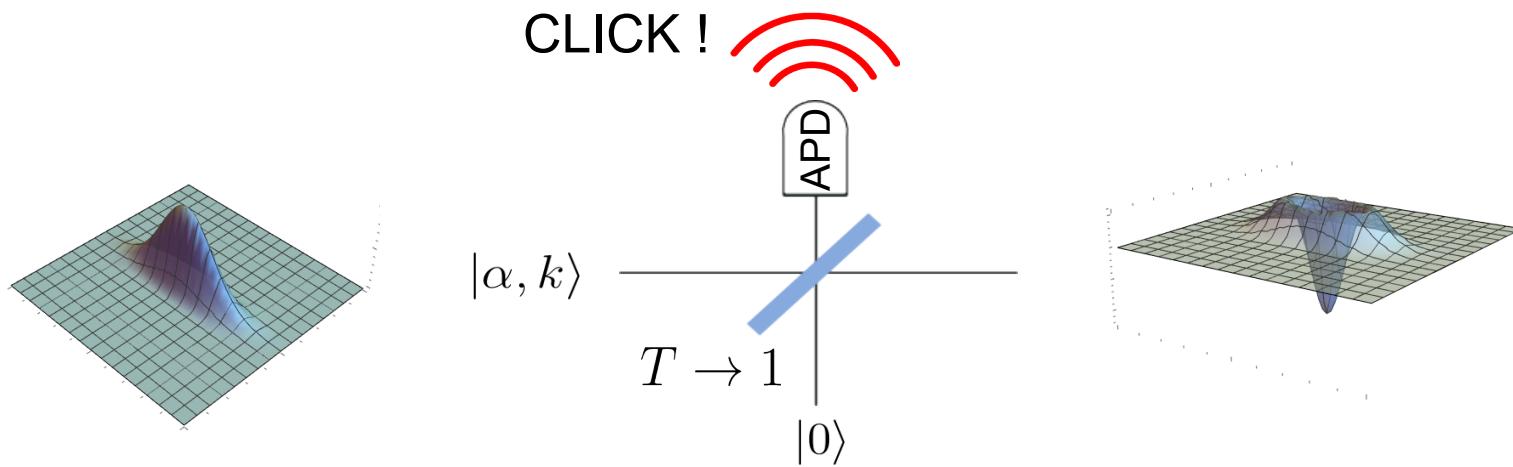
*S. Lloyd and S. L. Braunstein
PRL 82, 1784 (1999)*

Non-Gaussian

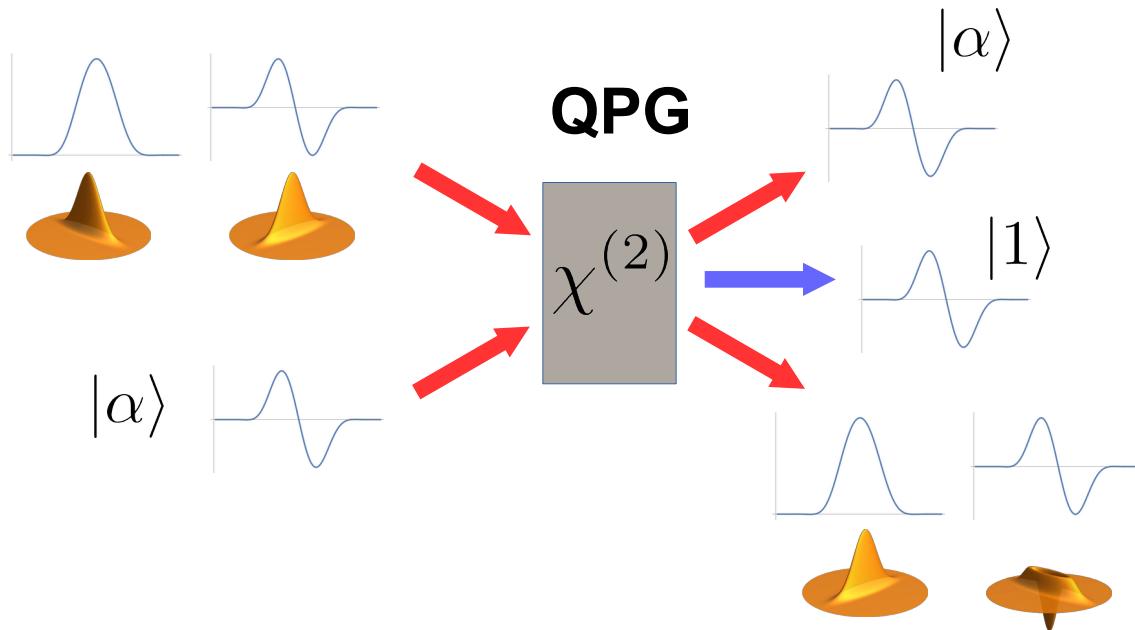
Two-modes C_Z



Computation with arbitrary encoding



Motivation: experiments in LKB!

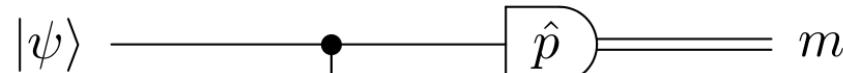
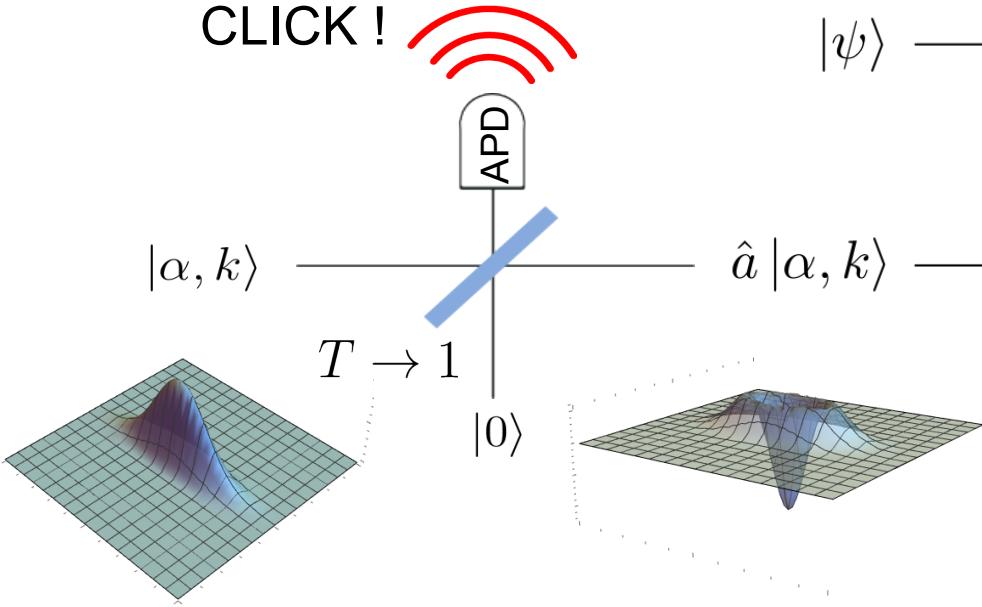


A. Eckstein et al,
Opt. Expr. 19;15 (2011)

J. S. Ra et al,
PRX 7, 031012 (2017)

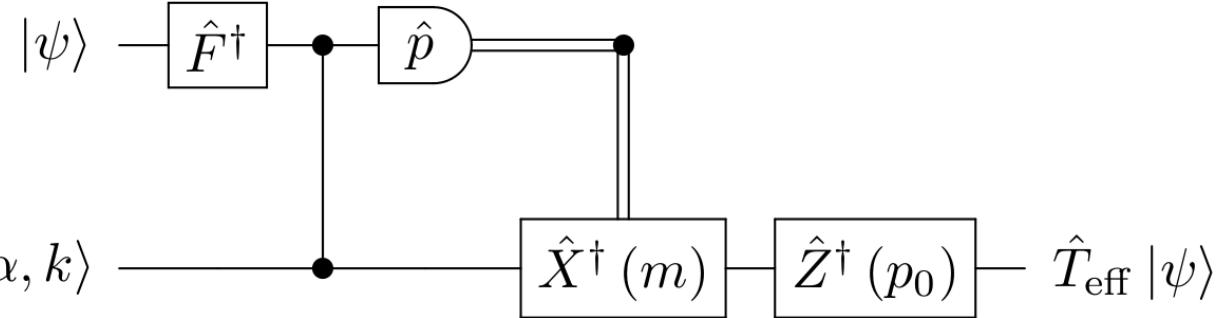
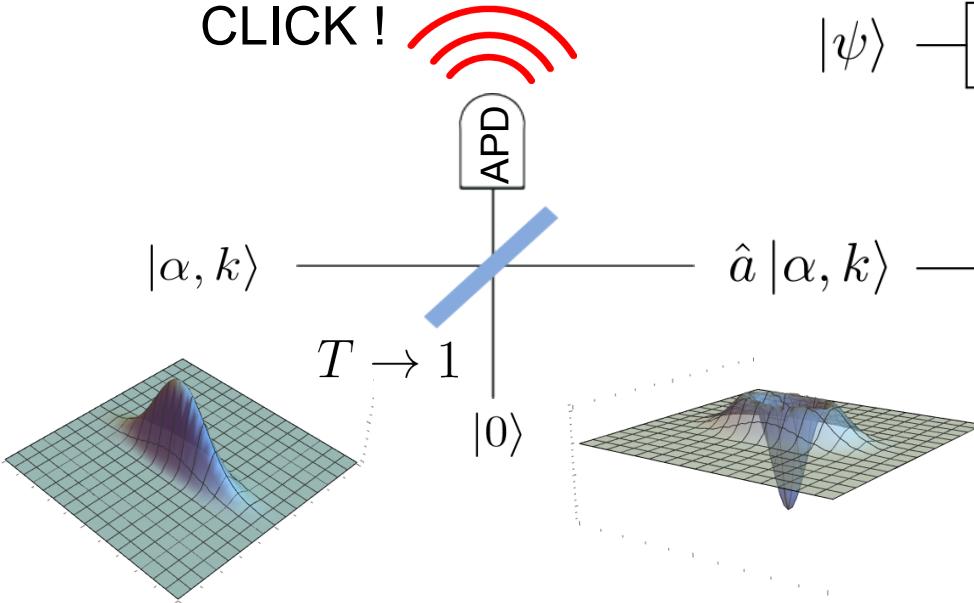
Measurement-based approach

CLICK !



*F. Arzani et al,
PRA 95 (5), 052352*

CLICK !



*F. Arzani et al,
PRA 95 (5), 052352*

$$\hat{T}_{\text{eff}} = \mathcal{N} \exp \left\{ -\frac{(\hat{q} - q_0 + m)^2}{k^2} \right\} (\hat{q} - \lambda(\alpha, k, m))$$

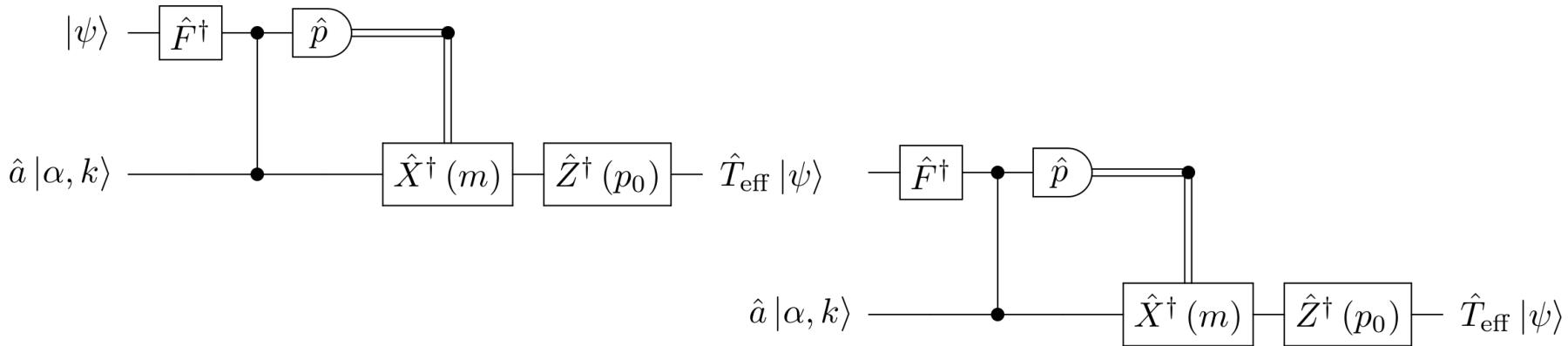
Normalization envelope Monomial in q

$$\lambda(\alpha, k, m) = - \left(\frac{2}{k^2 - 2} \right) q_0 - i \left(\frac{k^2}{k^2 - 2} \right) p_0 - m$$

Depends on

- Experimental parameters
- Measurement (!!?)

Repeated application: **polynomial** in the quadratures



Three times: get something like

$$e^{i\nu\hat{q}^3} \approx \mathbb{I} + i\nu\hat{q}^3 = (\hat{q} - \lambda_1)(\hat{q} - \lambda_2)(\hat{q} - \lambda_3)$$

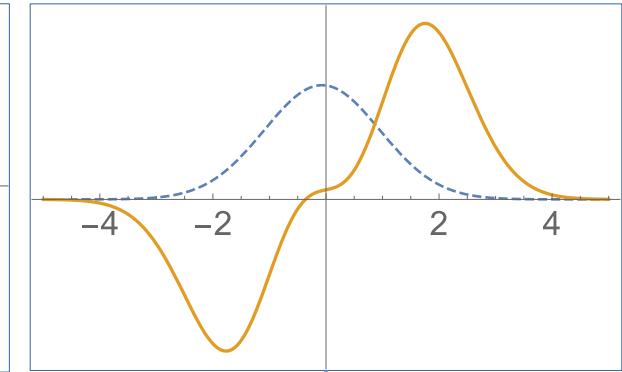
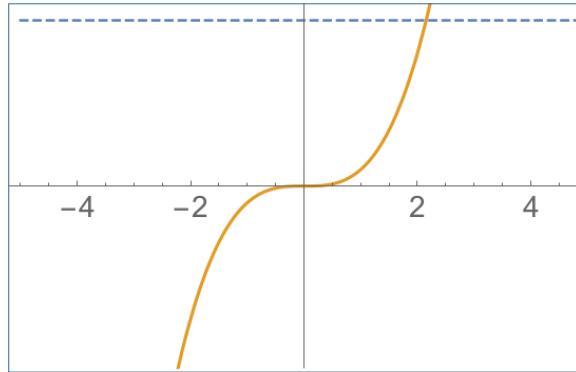
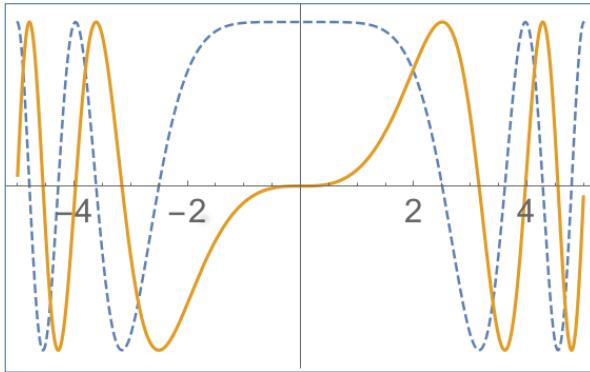
P. Marek et al
PRA 84(5), 053802 (2011)

K. Marshall et al
PRA 91, 032321 (2015)

Approximating Unitary operators

$$e^{0.1ix^3}$$

$$1 + 0.1ix^3$$



$$\mathcal{T}_{\text{eff}}(x, \vec{m}) \propto \prod_{i=1}^3 [\mathcal{G}(x, \alpha_i, k, m_i) (x - \lambda(\alpha_i, k, m_i))]$$

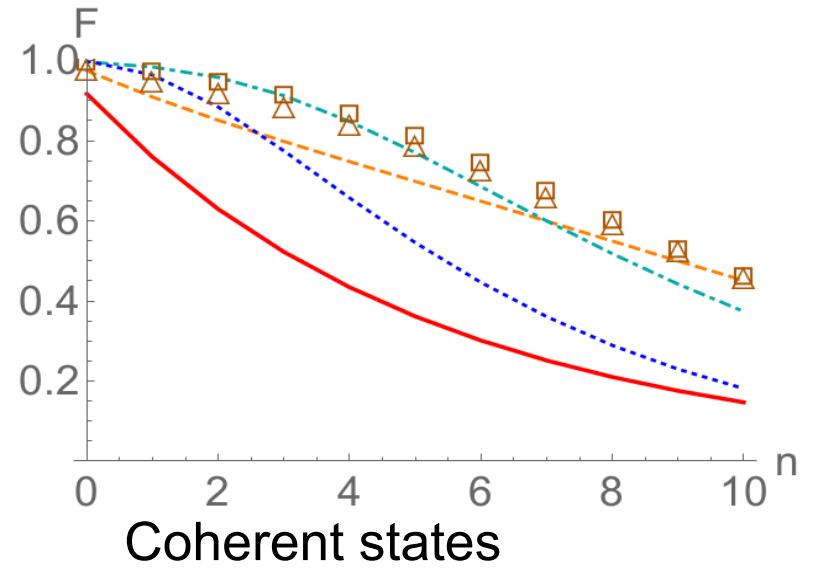
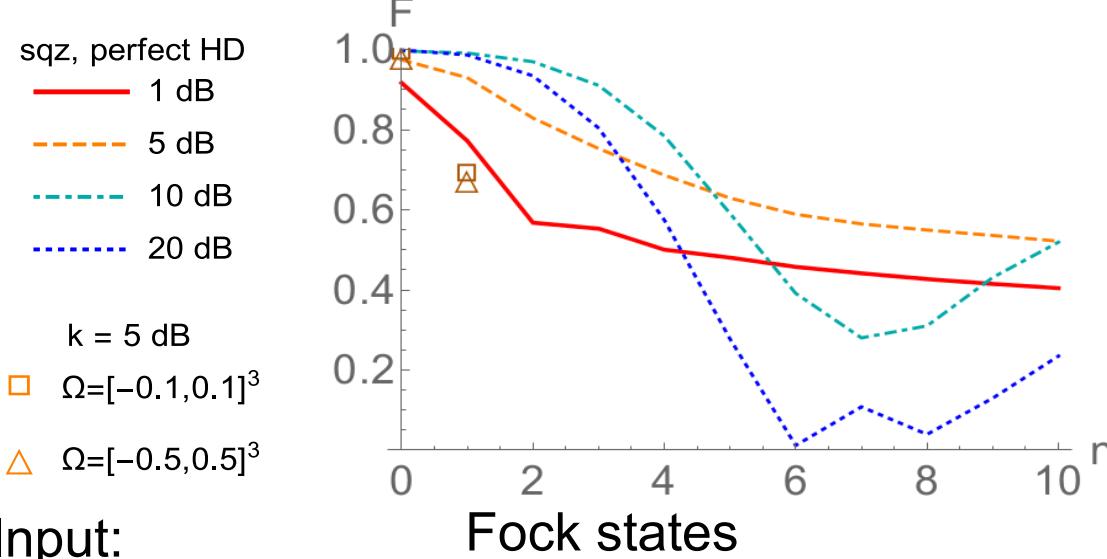
Finite success probability: acceptance region Ω
 $\approx 10^{-9} - 10^{-12}$

Average output state

$$\rho_\Omega = \int_\Omega d^n m \frac{p(\vec{m})}{p_\Omega} \hat{\mathcal{T}}_{\text{eff}}(\vec{m}) |\psi\rangle \langle \psi| \hat{\mathcal{T}}_{\text{eff}}^\dagger(\vec{m})$$

$$\rho_{\Omega} = \int_{\Omega} d^n m \frac{p(\vec{m})}{p_{\Omega}} \hat{\mathcal{T}}_{\text{eff}}(\vec{m}) |\psi\rangle \langle \psi| \hat{\mathcal{T}}_{\text{eff}}^{\dagger}(\vec{m})$$

$$\mathcal{F} = \sqrt{\langle \psi | \hat{U}^{\dagger} \rho \hat{U} | \psi \rangle} \quad \hat{U} = e^{i\nu \hat{q}^3}$$



Know the input: increase success probability

$$\left(\hat{q} - \lambda(\alpha, k, m) \right) \rightsquigarrow \lambda(\alpha, k, m) = - \left(\frac{2}{k^2 - 2} \right) q_0 - i \left(\frac{k^2}{k^2 - 2} \right) p_0 - m$$

Ex: Cubic phase state

$$|\gamma(\nu)\rangle = \hat{\gamma}(\nu)|0\rangle_p \rightarrow \hat{\gamma}_{\text{appr}}(\nu)|k_{in}\rangle_p$$

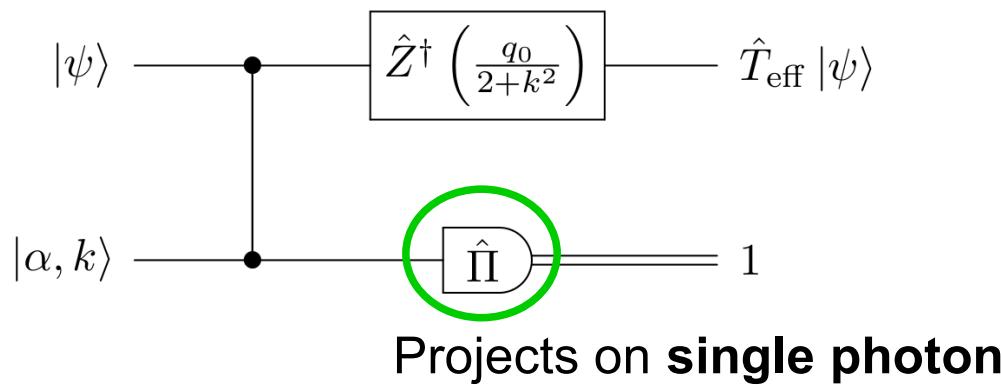
Fidelity = 0.9 with

Success probability $\sim 10^{-4}$

$$\hat{\gamma}(\nu)|k_{in}\rangle_p = e^{i\nu\hat{q}^3}|k_{in}\rangle_p$$

for $k = 5 \text{ dB}$, $k_{in} = 5 \text{ dB}$, $\nu = 0.1$

Alternative Scheme: Single-Photon Counter



Still post-selection but no binning of HD outcomes

~ Conjugate process of

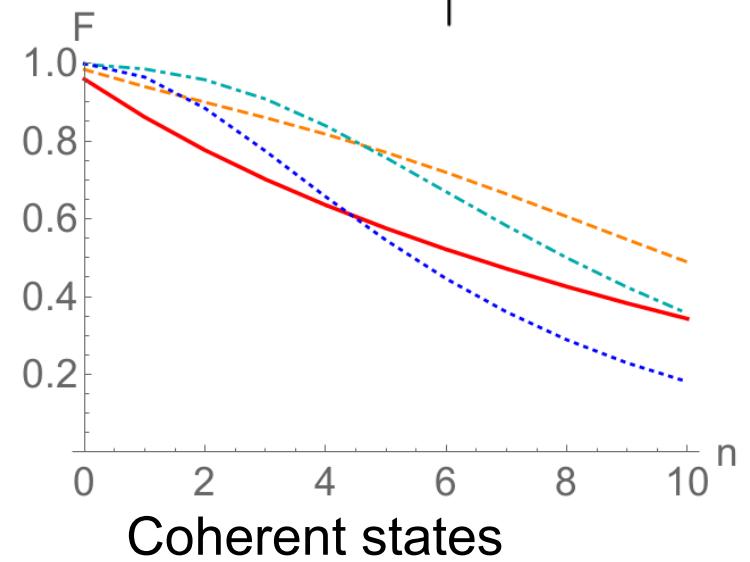
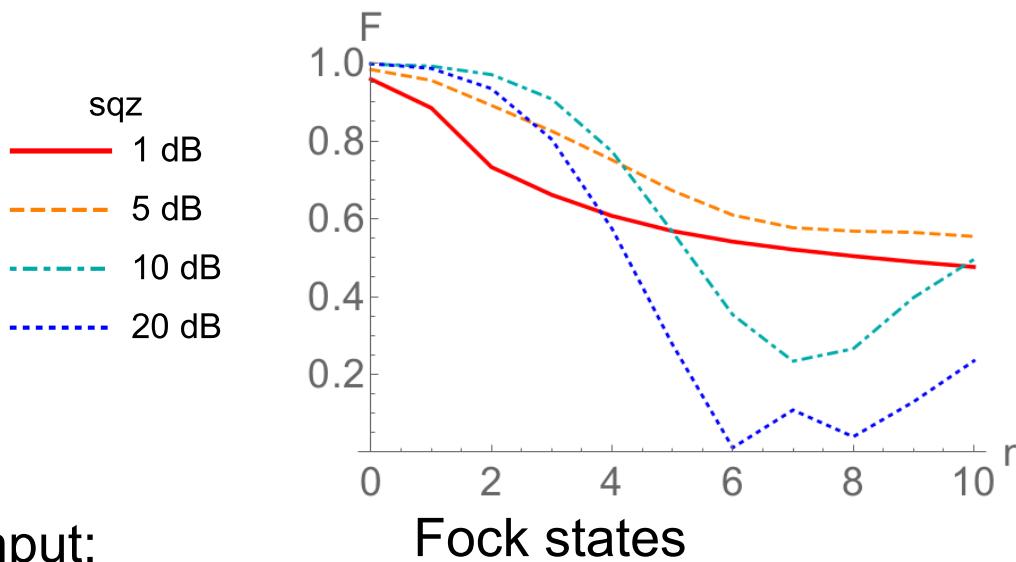
K. Park et al
PRA 90, 013804 (2014)

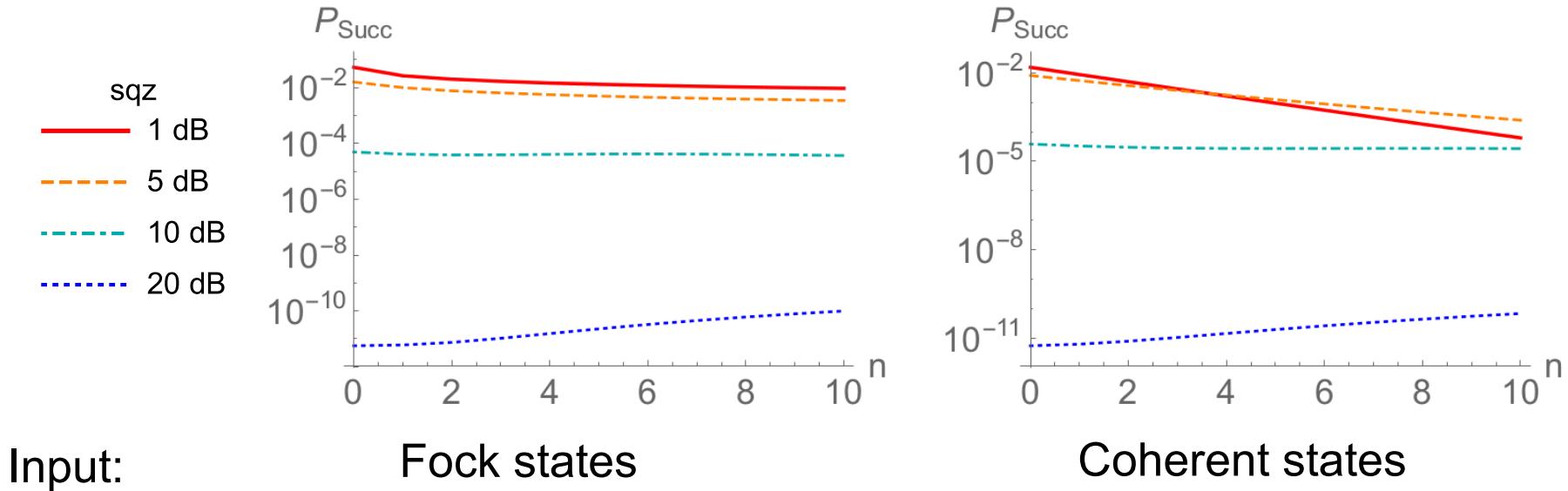
$$\hat{T}_{\text{eff}} = \tilde{\mathcal{N}} \exp \left\{ - \left(\frac{k^2}{4 + 2k^2} \right) (\hat{q} + p_0)^2 \right\} \left(\hat{q} - \lambda(\alpha, k) \right)$$

$$\lambda(\alpha, k) = \frac{2i}{k^2} q_0 - p_0$$

$$\hat{U} = e^{i\nu \hat{q}^3}$$

$$\mathcal{F} = \left| \langle \psi | \hat{U}^\dagger \hat{\mathcal{T}} | \psi \rangle \right|$$





Input:

Fock states

Coherent states

- Probability of detecting exactly one photon in each step
- Lower for high squeezing: many photons in the ancilla

- CV quantum computing
- Non-gaussianity through photon-subtraction
- Polynomial approximation of unitaries
- Alternative scheme : single photon counter

Thank you !

