

# Cyber Planner Code Report v0.1.0 (C++ Version)

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Abstract—This report presents a new arm planning algorithm for team 8214 in the 2025 FRC season. The algorithm is based on a finite convex polytope set representation of objects, where path planning is built upon a Lipschitz continuous artificial potential field. In trajectory planning, constraints on current and voltage are incorporated, and time-optimal trajectories are derived using the Augmented Lagrangian method and the L-BFGS optimization algorithm. While traditional FRC teams commonly use A\* and forwardbackward traversal methods for trajectory planning, these approaches show limitations when applied to the configuration space of cascading robotic arms, where the L2 norm is not meaningful. The proposed method overcomes these challenges by ensuring efficient trajectory optimization for complex robotic systems.

## I. Introduction

IN FRC (FIRST Robotics Competition), effective trajectory planning is crucial for maximizing performance. Traditional AD\* and forward-backward traversal have been widely adopted for path planning and trajectory optimization in drivetrain. However, these methods face significant challenges when applied to more complex robotic systems, such as multi-joint robotic arms, where the conventional L2 norm does not apply effectively in configuration space.

We handle these limitations in our new arm planning algorithm for the 2025 FRC season. By leveraging finite convex polytope sets for object representation, the algorithm introduces a path planning strategy based on Lipschitz continuous artificial potential fields. Furthermore, the trajectory planning incorporates practical constraints related to current and voltage, ensuring the system's efficiency and safety. To optimize the trajectory and find time-optimal solutions, we use an advanced optimization approach combining the Augmented Lagrangian method and the L-BFGS algorithm. This method enhance the precision and performance of the robotic arm, overcoming

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the shortcomings of previous trajectory planning techniques, and offering significant advantages in dynamic and complex environments.

#### II. PATH PLANNING

## A. Improved A \* Algorithm

As Image 1a shown, the original A\* algorithm may always find a path that very close to the obstacles. This is often led by the heuristic function. However, if any errors occured in estimate or control, the robot may collide with the obstacles. To avoid this, we introduce artificial potential field to the cost function.

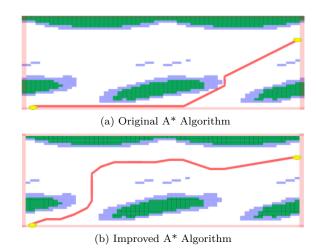


Fig. 1: Difference between original and improved  $\mathbf{A}^*$  algorithm

Every grid in the map has a potential value. The potential value is calculated by the distance to the obstacles in a dynamic programming way. It can find us a heuristic path, but it is still discrete. To make the path continuous, we need to interpolate the path.

### B. Polynomial Curve

The key idea of a polynomial curve is to make the discrete path continuous, not only in position but also in velocity and acceleration.

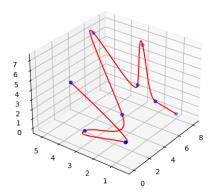


Fig. 2: 3D Cubic Spline

$$p_{i} \in \mathcal{P}$$

$$p_{i} = [x_{i}, y_{i}], \ \forall i \in [0, n]$$

$$q_{i,x}(s) = a_{i,x}s^{3} + b_{i,x}s^{2} + c_{i,x}s^{1} + d_{i,x}$$

$$\forall i \in [0, n - 1], s \in [0, 1)$$
(1)

s is the normalized representation of the arc length. The minimum order of polynomial to make the path continuous in position, velocity, and acceleration is 3. And to ensure the continuous, we need

$$\begin{cases} q_{i-1,x}(1) = q_{i,x}(0) = x_i \\ q'_{i-1,x}(1) = q'_{i,x}(0) \\ q''_{i-1,x}(1) = q''_{i,x}(0) \end{cases}, \forall i \in [1, n]$$
(2)

If the object is stationary in the beginning and at the end, the velocity should be  $q_0'(0) = q_n'(1) = 0$ . By solving these equations we can uniquely determine a polynomial to fit those points.

$$\begin{cases}
 a_i = x_{i+2} - 2x_{i+1} + x_i \\
 b_i = -x_{i+2} + 2x_{i+1,k} - x_i \\
 c_i = x_{i+1} - x_i \\
 d_i = x_i
\end{cases}, \forall i \in [1, n-1]$$
(3)

In natural form [12].

$$\begin{bmatrix} 2 & 1 & & & & \\ 1 & 4 & 1 & & & \\ & \ddots & \ddots & \ddots & \\ & & 1 & 4 & 1 \\ & & & & 1 & 2 \end{bmatrix} \begin{bmatrix} D_0 \\ D_1 \\ \vdots \\ D_{m-1} \\ D_m \end{bmatrix} = \begin{bmatrix} 3(x_1 - x_0) \\ 3(x_2 - x_0) \\ \vdots \\ 3(x_m - x_{m-2}) \\ 3(x_m - x_{m-1}) \end{bmatrix}$$

In case,

$$\begin{cases}
 a_i = 2(x_i - x_{i+1}) + D_i + D_{i+1} \\
 b_i = 3(x_{i+1} - x_i) - 2D_i - D_{i+1} \\
 c_i = D_i \\
 d_i = x_i
\end{cases}$$
(5)

#### C. Continuous Potential Field

In section II.A, we proposed a method to calculate the approximate and discrete potential field. However, the potential field often cause the path a severe shaking. Although you can set the potential field below a certain threshold to zero, this still does not make the system satisfactory.

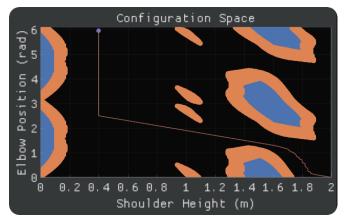


Fig. 3: The path is not smooth and has sharp jitters.

Any object can be represented by a set of convex polygons. The potential energy of a certain point is the minimum distance from the point to the edge of the polygon. In robot arm case, the certain point represent a kind of configuration of the robot arm. The potential energy is the minimum distance between the set of polygons that make up the current configuration of the robot and the set of polygons that make up the environmental obstacles.

If the point is inside the polygon, the distance is the minimum distance to the edge of the polygon. If the point is outside the polygon, the distance can be obtained by solving the Low Dimensional Quadratic Programming (LDQP) problem.

$$\min_{x} ||x - p||^2$$
s.t. $Ax \le b$  (6)

p is the given point, and x is the point inside the polygon.

If the polygon is intersect or overlap with the other polygon, the distance is the distance is defined as the minimum distance from a vertex of a polygon contained by another polygon to an edge of that polygon. If the two polygons are not intersect, the distance can alse be obtained by solving the LDQP problem.

$$\min_{x,y} ||x - y||^{2}$$
s.t. $Ax \le b$ 

$$Cy < d$$
(7)

By doing so, we can get a lipshitz continuous potential field. We can optimize the path by the gradient descent method. A probable object function is

Fig. 4: The difference between the discrete and continuous potential field.

$$\sum_{i=0}^{n} || \exp(-d_{x_{i}}) ||^{2}$$

$$+ \sum_{i=0}^{n} || x_{i} - \bar{x}_{i} ||_{\Sigma}^{2}$$

$$+ \sum_{i=0}^{n-1} || x_{i} - x_{i+1} ||_{\Sigma}^{2}$$

$$+ \sum_{i=1}^{n-1} || x_{i-1} - 2x_{i} + x_{i+1} ||_{\Sigma}^{2}$$
(8)

 $||x_i - \bar{x}_i||_1^2$  is the prior regularization term obtained by A\* search. The Mahalanobis distance is used because the units of each dimension are different, making the L2 norm meaningless. Squaring ensures differential flatness.

## III. TRAJECTORY PLANNING

Classical method to deal with the time parameterization problem in FRC is called forward and backward pass. Its essence is based on the trapezoidal curve or S-curve model, adjusting the time scale so that the motion planning meets the given constraints. However, it is only able to deal with 1 DOF planning problem. It has difficulty in multi-DOF trajectory planning, especially when it is not in Cartesian space [7], [5].

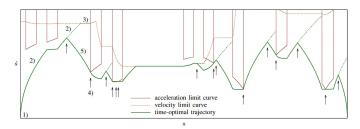


Fig. 5: Phase-plane trajectory[15]

Team 6328 proposed a Kairos solver in 2023. In their 23 season machine and code, they used an intake active avoidance strategy to avoid anti-collision. This allowed

their solver to complete planning in a simpler way. They use linear interpolation as the initial position of waypoints, and then adjusts the position of waypoints to make the speed, acceleration, and torque meet the constraints based on the assumption that the time interval between each waypoint remains unchanged[10].

3

Our approach is very different from that of Team 6328. Instead of globally adjusting the time scale by assuming that the time intervals between points are consistent, we express the function of time with respect to arc length in a discretized form and parameter.

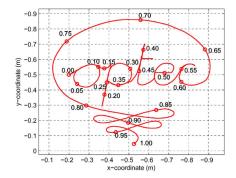


Fig. 6: The curve is uniquely determined by the arc length[2].

### A. Time Optimal Path Parameterization

The total time of the trajectory is given by the integration of the whole path.

$$T = \int_0^T 1 dt = \int_0^L \frac{1}{\frac{ds}{dt}} ds \tag{9}$$

Velocity and acceleration constraints at elevator position x and arm radian  $\theta$  are

$$-v_{x,\max} \le v_x \le v_{x,\max}$$

$$-a_{x,\max} \le a_x \le a_{x,\max}$$

$$-v_{\theta,\max} \le v_{\theta} \le v_{\theta,\max}$$

$$-a_{\theta,\max} \le a_{\theta} \le a_{\theta,\max}$$
(10)

Given by chain rules, we have

$$\frac{\mathrm{d}x}{\mathrm{d}t} = \frac{\mathrm{d}x}{\mathrm{d}s} \cdot \frac{\mathrm{d}s}{\mathrm{d}t} 
\frac{\mathrm{d}^2x}{\mathrm{d}t^2} = \frac{\mathrm{d}^2x}{\mathrm{d}s^2} \cdot \frac{\mathrm{d}s}{\mathrm{d}t} + \frac{\mathrm{d}x}{\mathrm{d}s} \cdot \frac{\mathrm{d}^2s}{\mathrm{d}t^2} 
\frac{\mathrm{d}\theta}{\mathrm{d}t} = \frac{\mathrm{d}\theta}{\mathrm{d}s} \cdot \frac{\mathrm{d}s}{\mathrm{d}t} 
\frac{\mathrm{d}^2\theta}{\mathrm{d}t^2} = \frac{\mathrm{d}^2\theta}{\mathrm{d}s^2} \cdot \frac{\mathrm{d}s}{\mathrm{d}t} + \frac{\mathrm{d}\theta}{\mathrm{d}s} \cdot \frac{\mathrm{d}^2s}{\mathrm{d}t^2}$$
(11)

Denote

$$a(s) = \frac{\mathrm{d}^2 s}{\mathrm{d}t^2}, b(s) = \left(\frac{\mathrm{d}s}{\mathrm{d}t}\right)^2 \tag{12}$$

The problem is described by

$$\min_{a(s),b(s)} \int_{0}^{L} \frac{1}{\sqrt{b(s)}} ds$$

$$b'(s) = 2a(s)$$

$$x'(0)\sqrt{b(0)} = 0$$

$$x'(L)\sqrt{b(L)} = 0$$

$$\theta'(0)\sqrt{b(0)} = 0$$

$$\theta'(L)\sqrt{b(L)} = 0$$
s.t.
$$-b(s) \le 0$$

$$[x'(s)]^{2}b(s) \le v_{x,\max}^{2}$$

$$[\theta'(s)]^{2}b(s) \le v_{\theta,\max}^{2}$$

$$x''(s)b(s) + x'(s)a(s) \le a_{x,\max}$$

$$-x''(s)b(s) + x'(s)a(s) \le a_{x,\max}$$

$$-x''(s)b(s) + \theta'(s)a(s) \le a_{\theta,\max}$$

$$-\theta''(s)b(s) + \theta'(s)a(s) \le a_{\theta,\max}$$

$$-\theta''(s)b(s) + \theta'(s)a(s) \le a_{\theta,\max}$$

 $1)\ Discrete\ Case:$  Considering the problem in discrete case, we have

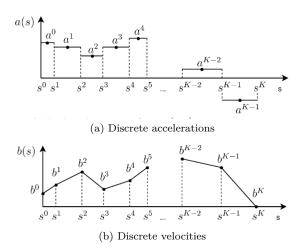


Fig. 7: Discrete function of time with respect to arc length[2].

2) Second-Order Cone: To rewrite the problem to a second-order conic programming form, we first try to bound nonlinear term  $\sqrt{b^k}$  with  $c^k$ , such that  $\sqrt{b^k} \ge c^k, k \in [0, K]$ . Going a step further, we introduce one more slack variable  $d^k$ , which satisfies  $\frac{1}{c^{k+1}+c^k} \le d^k, k \in [0, K-1]$ .

$$\left| \begin{vmatrix} b^{k} - 1 \\ 2c^{k} \end{vmatrix} \right|_{2} \le b^{k} + 1 \Leftrightarrow \begin{bmatrix} b^{k} + 1 \\ b^{k} - 1 \\ 2c^{k} \end{bmatrix} \in \mathcal{Q}^{3}$$

$$\left| \begin{vmatrix} c^{k} + c^{k+1} - d^{k} \\ 2 \end{vmatrix} \right|_{2} \le c^{k} + c^{k+1} + d^{k}$$

$$\Leftrightarrow \begin{bmatrix} c^{k} + c^{k+1} + d^{k} \\ c^{k} + c^{k+1} - d^{k} \\ 2 \end{vmatrix} \in \mathcal{Q}^{3}$$

$$(15)$$

$$\min_{a,b} \sum_{k=1}^{K} \frac{2(s^{k+1} - s^{k})}{\sqrt{b^{k+1}} + \sqrt{b^{k}}} \Leftrightarrow \\
\min_{a,b,c,d} \sum_{k=1}^{K} 2d^{k}(s^{k+1} - s^{k}) \\
\text{s.t.} \left\| \begin{vmatrix} c^{k} + c^{k+1} - d^{k} \\ 2 \end{vmatrix} \right\|_{2} \leq c^{k} + c^{k+1} + d^{k} \\
\left\| \begin{vmatrix} b^{k} - 1 \\ 2c^{k} \end{vmatrix} \right\|_{2} \leq b^{k} + 1$$
(16)

Rewriting the constraints of the optimization problem, especially the nonlinear terms  $\sqrt{b_k}$ , we get the final objective function [2].

$$\min_{a,b,c,d} \sum_{k=1}^{K} 2(s^{k+1} - s^k) d^k$$

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$$\begin{bmatrix} c^k + c^{k+1} + d^k \\ c^k + c^{k+1} - d^k \\ 2 \end{bmatrix} \in \mathcal{Q}^3 \qquad \forall k \in [0, K - 1]$$

$$\begin{bmatrix} b^k + 1 \\ b^k - 1 \\ 2c^k \end{bmatrix} \in \mathcal{Q}^3 \qquad \forall k \in [0, K]$$

$$2(s^{k+1} - s^k)a^k + b^k - b^{k+1} = 0 \qquad \forall k \in [0, K - 1]$$

$$[x'(s^0)]^2b^0 = 0$$

$$[x'(s^K)]^2b^K = 0$$

$$[\theta'(s^0)]^2b^0 = 0$$

$$[\theta'(s^K)]^2b^L = 0$$

$$-b^k \le 0, \qquad \forall k \in [0, K]$$

$$[x'(s^k)]^2b^k \le v_{x,\max}^2 \qquad \forall k \in [0, K]$$

$$[x'(s^k)]^2b^k \le v_{\theta,\max}^2 \qquad \forall k \in [0, K]$$

$$x''(s^k)b^k + x'(s^k)a^k \le a_{x,\max} \qquad \forall k \in [0, K - 1]$$

$$-x''(s^k)b^k + x'(s^k)a^k \le a_{\theta,\max} \qquad \forall k \in [0, K - 1]$$

$$\theta''(s^k)b^k + \theta'(s^k)a^k \le a_{\theta,\max} \qquad \forall k \in [0, K - 1]$$

$$\theta''(s^k)b^k + \theta'(s^k)a^k \le a_{\theta,\max} \qquad \forall k \in [0, K - 1]$$

$$-\theta''(s^k)b^k + \theta'(s^k)a^k \le a_{\theta,\max} \qquad \forall k \in [0, K - 1]$$

$$(17)$$

All constraints are divided into second-order cone constraints, equality constraints and inequality constraints [4].

$$\min_{x} c^{T} x$$
s.t.  $A_{i}x + b_{i} \in \mathcal{K}_{i}$ 

$$Gx = h$$

$$Px \leq q$$
(18)

3) Augmented Lagrangian Method: Therefore, we can reformulate the problem using the Augmented Lagrangian method for Symmetric Cones.

$$\mathcal{L}_{\rho}(x,\mu,\lambda,\eta) = c^{T}x + \frac{\rho}{2} \sum_{i=1}^{m} ||P_{\mathcal{K}_{i}}(\frac{\mu_{i}}{\rho} - A_{i}x - b_{i})||^{2} + \frac{\rho}{2} ||Gx - h + \frac{\lambda}{\rho}||^{2} + \frac{\rho}{2} ||\max[Px - q + \frac{\eta}{\rho}, 0]||^{2}$$
(19)

 $P_{\mathcal{K}}$  is the projection of a vector on a symmetric cone

$$P_{\mathcal{K}}(v) = \arg\min_{x \in \mathcal{K}} ||v - x||^2 \tag{20}$$

In second order cone case,

$$P_{\mathcal{K}=\mathcal{Q}^n}(v) = \begin{cases} 0, & v_0 \le -||v_1||_2\\ \frac{v_0 + ||v_1||_2}{2||v_1||_2} (||v_1||_2, v_1)^T, & |v_0| < ||v_1||_2\\ v, & v_0 \ge ||v_1||_2 \end{cases}$$

$$(21)$$

And the gradient of the augmented Lagrangian function is given by

$$\begin{bmatrix} c^{k} + c^{k+1} + d^{k} \\ c^{k} + c^{k+1} - d^{k} \\ 2 \end{bmatrix} \in \mathcal{Q}^{3} \qquad \forall k \in [0, K-1] \qquad \nabla \mathcal{L}_{\rho}(x, \mu, \lambda, \eta) = c - \rho \sum_{i=1}^{m} A_{i}^{T} P_{\mathcal{K}_{i}}(\frac{\mu_{i}}{\rho} - A_{i}x - b_{i}) + \frac{\rho G^{T}(Gx - h + \frac{\lambda}{\rho}) + \rho G^{T}(Gx - h + \frac{\lambda}{\rho})}{\rho P^{T}\{\max[Px - q + \frac{\eta}{\rho}, 0]\}}$$

$$(22)$$

So we can use a L-BFGS method to solve this convex and unconstrained function [3], [1], [16].

$$x \leftarrow \arg\min_{x} \mathcal{L}_{\rho}(x, \mu, \lambda, \eta)$$

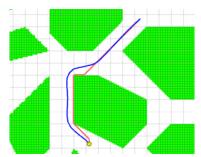
$$\mu_{i} \leftarrow P_{\mathcal{K}_{i}}(\mu_{i} - \rho(A_{i}x + b_{i}))$$

$$\lambda \leftarrow \lambda + \rho(Gx - h)$$

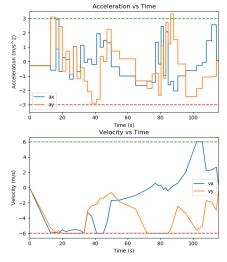
$$\eta \leftarrow \max[\eta + \rho(Px - q), 0]$$

$$\rho \leftarrow \min[(1 + \gamma)\rho, \beta]$$
(23)

 $\gamma$  is the growth rate of  $\rho$  and  $\beta$  is the upper bound of  $\rho$ , which is typically  $10^3$ .



(a) The corresponding path



(b) Velocity and acceleration under constraints

Fig. 8: Our previous work showing constraints on velocity and acceleration[19].

## B. Constraints on voltage and current

Can we go one step further? Although we can set constraints for speed and acceleration in segments instead of setting a global maximum value, it is difficult to exhaust all possibilities and enumerate all the required constraints. A conservative constraint will waste control resources.

On the contrary, under the guidance of the feedforward model, speed and acceleration together constitute a feasible control domain. Under the guidance of the theory of control allocation, we redesign the speed and acceleration constraints into voltage and current constraints. Doing so enables the model to dynamically adjust the energy distribution on speed and acceleration, maximize the use of control resources on the field.

Elevator feedforward model,  $i_e$  means the ratio from motor radian to elevator position meter:

$$V_E = K_g + K_S \cdot \operatorname{sign}(\dot{d}) + \frac{K_v}{i_e} \cdot \dot{d} + \frac{K_a}{i_e} \cdot \ddot{d}$$
 (24)

Arm feedforward model,  $i_a$  means the ratio from motor radian to arm radian:

$$V_A = K_g \cdot \cos(\theta) + K_S \cdot \operatorname{sign}(\dot{\theta}) + \frac{K_v}{i_a} \cdot \dot{\theta} + \frac{K_a}{i_a} \cdot \ddot{\theta} \quad (25)$$

The motor current at a given voltage and speed is:

$$I = -\frac{K_v}{R} \cdot \omega + \frac{1}{R} \cdot V \tag{26}$$

Notice that we used  $c^k$  to bound the lower side of  $\sqrt{b^k}$ . This is because

$$\frac{2(s^{k+1} - s^k)}{\sqrt{b^{k+1}} + \sqrt{b^k}} \le \frac{2(s^{k+1} - s^k)}{c^{k+1} + c^k} \le 2(s^{k+1} - s^k)d^k$$
 (27)

In case, we introduce a second order cone to handle the chain constraints. The equality stands only when  $\sqrt{b^k} = c^k$ . How about directly let  $\sqrt{b^k} = c^k$  instead of  $\sqrt{b^k} \ge c^k$ ? So that,

$$\frac{2(s^{k+1} - s^k)}{\sqrt{b^{k+1}} + \sqrt{b^k}} = \frac{2(s^{k+1} - s^k)}{c^{k+1} + c^k} \le 2(s^{k+1} - s^k)d^k \quad (28)$$

We introduce a dual variable  $\nu$ , to bound the quadratic term. The problem is modeled by the following.

$$\mathcal{L}_{\rho}(x,\mu,\nu,\lambda,\eta) = c^{T}x + \frac{\rho}{2} \sum_{i=1}^{m} ||P_{\mathcal{K}_{i}}(\frac{\mu_{i}}{\rho} - A_{i}x - b_{i})||^{2} + \frac{\rho}{2} \sum_{j=1}^{q} ||x^{T}J_{j}x - r_{j}^{T}x + \frac{\nu_{j}}{\rho}||^{2} + \frac{\rho}{2}||Gx - h + \frac{\lambda}{\rho}||^{2} + \frac{\rho}{2}||\max[Px - q + \frac{\eta}{\rho}, 0]||^{2}$$
(29)

$$\nabla \mathcal{L}_{\rho}(x,\mu,\nu,\lambda,\eta) = c - \rho \sum_{i=1}^{m} A_{i}^{T} P_{\mathcal{K}_{i}}(\frac{\mu_{i}}{\rho} - A_{i}x - b_{i}) +$$

$$\rho \sum_{j=1}^{q} (x^{T} J_{j} - r_{j}^{T} + \frac{\nu_{j}}{\rho}) [(J_{j}^{T} + J_{j})x - r_{j}] +$$

$$\rho G^{T}(Gx - h + \frac{\lambda}{\rho}) +$$

$$\rho P^{T}\{\max[Px - q + \frac{\eta}{\rho}, 0]\}$$
(30)

In case J is symmetric,  $J^T + J = 2J$ .

$$x \leftarrow \arg\min_{x} \mathcal{L}_{\rho}(x, \mu, \nu, \lambda, \eta)$$

$$\mu_{i} \leftarrow P_{\mathcal{K}_{i}}(\mu_{i} - \rho(A_{i}x + b_{i}))$$

$$\nu_{j} \leftarrow \nu_{j} + \rho(x^{T}J_{j}x - r_{j}^{T}x)$$

$$\lambda \leftarrow \lambda + \rho(Gx - h)$$

$$\eta \leftarrow \max[\eta + \rho(Px - q), 0]$$

$$\rho \leftarrow \min[(1 + \gamma)\rho, \beta]$$
(31)

Till now, we obtain all the math tools to solve our problem.

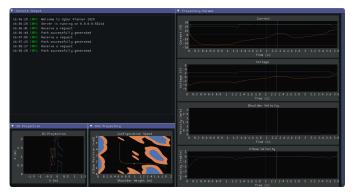


Fig. 9: Desktop APP Screenshot

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