

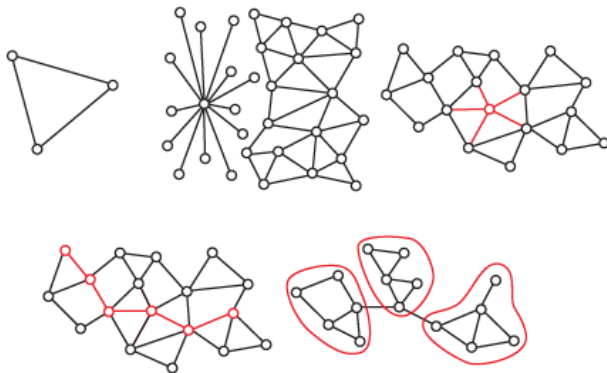
Introduction to data science & artificial intelligence (INF7100)

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#281 Networks

été 2020

Networks



Networks

Let $V = \{1, \dots, n_V\}$ denote either **nodes**, or **vertices** (n_V is the **order**). Let $E \in \{0, 1\}^{n_V \times n_V}$ represents the **relationships**, through an adjacency matrix \mathbf{A} , $A_{i,j} = 1$ indicates a **link** - or **edge** - between i and j , or a collection of links $\{e_1, \dots, e_{n_E}\}$. Let $n_E = |E|$ denote the number of edges, called **size**.

The **degree** $d(\cdot)$ of a vertex v is its number of incident edges.

A **network** is a pair (V, E)

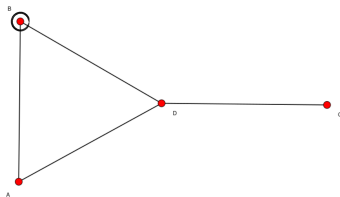
- ▶ $G = \text{Internet}$, $V = \text{computers}$, $E = \text{IP network adjacency}$
- ▶ $G = \text{World Wide Web}$, $V = \text{web pages}$, $E = \text{hyperlink}$
- ▶ $G = \text{Articles}$, $V = \text{authors}$, $E = \text{citations}$
- ▶ $G = \text{Friendship Network}$, $V = \text{persons}$, $E = \text{friendship}$
- ▶ $G = \text{Airport Network}$, $V = \text{airports}$, $E = \text{non-stop flight}$

Networks

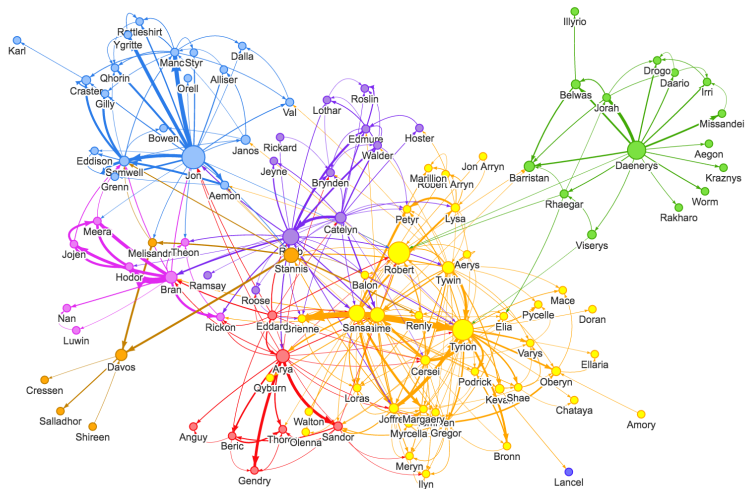
Adjacency matrix \mathbf{A} , $A_{i,j} = 1$ indicates a link between i and j
The **degree** of a vertex v is $d(v)$, the number of vertices in V incident to v (i.e. the number of neighbors of v)
Row i contains list of vertices connected to vertex i .

$$d(i) = \sum_{j=1}^{n_V} A_{i,j} = \mathbf{A}_{i,\cdot}^T \mathbf{1}.$$

$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ \boxed{1} & \boxed{0} & \boxed{0} & \boxed{1} \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

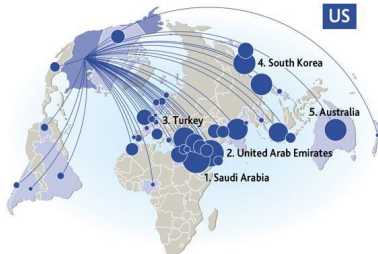


Networks



Networks

CÉCILE MARIN

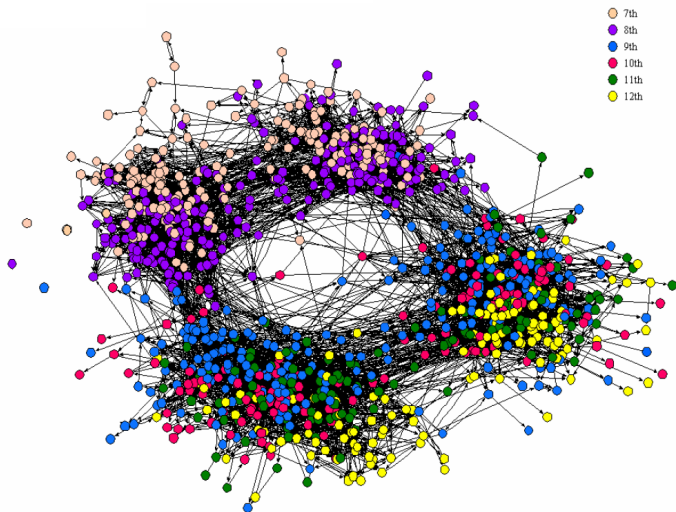


Total arms exports 2011-5 (in excess of \$100m)

Millions of dollars (constant 1990 PPP)

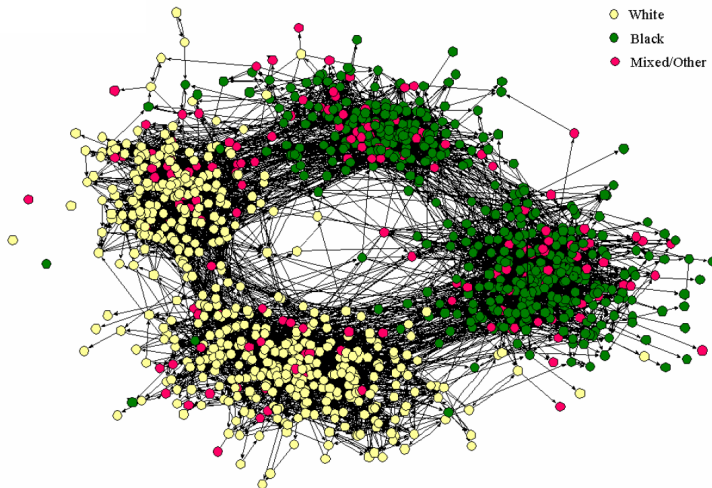
Source: Sipri, Arms Transfers Database, 2016

Networks



Race, School Integration and Friendship Segregation in America

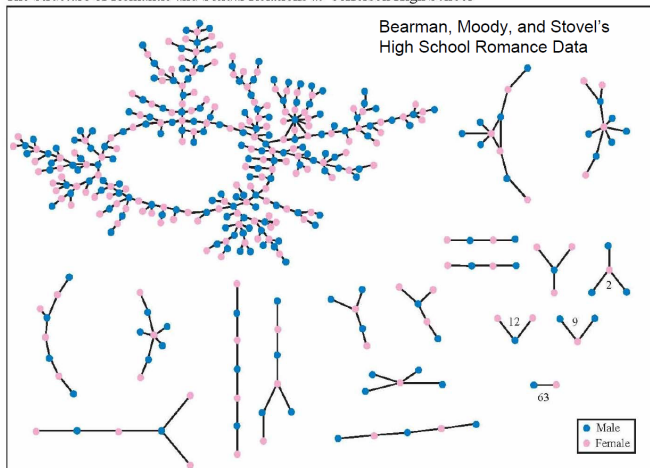
Networks



Race, School Integration and Friendship Segregation in America

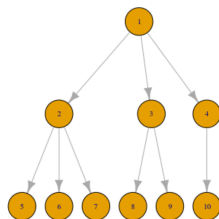
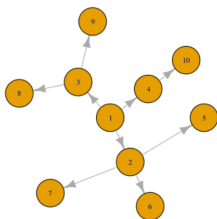
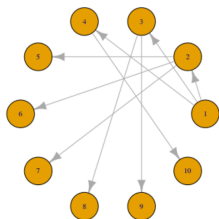
Networks

The Structure of Romantic and Sexual Relations at "Jefferson High School"



See [Chains of Affection: The Structure of Adolescent Romantic and Sexual Networks](#)

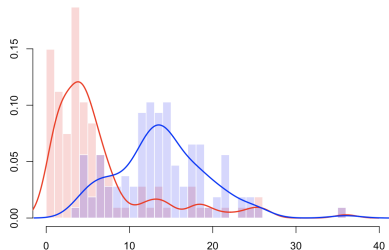
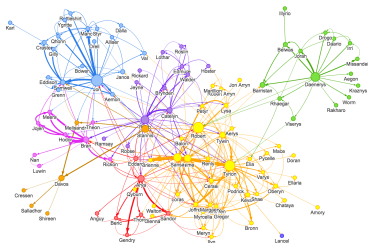
Networks



(same network, different representations)

Friendship Paradox

The average degree of a friend is strictly greater than the average degree of a random node. See [Feld \(1991\)](#) and [Zuckerman & Jost \(2001\)](#).



The average number of friends of a random person in the graph is

$$\mu = \frac{1}{n_V} \sum_{v \in V} d(v) = \frac{2n_E}{n_V} = \frac{1}{n_V} \|\mathbf{d}\|_1$$

where \mathbf{d} is the vector of $d(v)$'s (or $\mathbf{d} = \mathbf{A}\mathbf{1}$).

Friendship Paradox

The average number of friends that a typical friend has is

$$\frac{1}{n_V} \sum_{v \in V} \left(\frac{1}{d(v)} \sum_{v': (v, v') \in E} d(v') \right) = \frac{1}{\|\mathbf{d}\|_1} \mathbf{d}^\top \mathbf{d}$$

We can prove that

$$\frac{1}{\|\mathbf{d}\|_1} \mathbf{d}^\top \mathbf{d} \geq \frac{1}{n_V} \|\mathbf{d}\|_1$$

See also

$$\frac{\mathbb{E}[D^2]}{\mathbb{E}[D]} = \mathbb{E}[D] + \frac{\text{Var}[D]}{\mathbb{E}[D]} \geq \mathbb{E}[D]$$

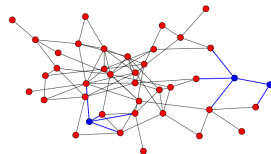
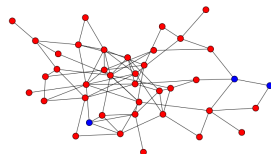
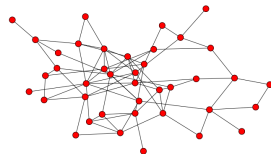
because $\text{Var}[D] \geq 0$.

Sampling & Networks

Surveying hard-to-reach groups through sampled respondents in a social network

Use **Snowball Sampling** (or **Star Sampling**)

Idea: start with an initial vertex sample V_0 and observe all incident edges, as well as all vertices sharing those edges. Iterate to have n edges.



Small Word Property

“I read somewhere that everybody on this planet is separated by only six other people. Six degrees of separation. Between us and everybody else on this planet. The president of the United States. A gondolier in Venice. Fill in the names. . . . How every person is a new door, opening up into other worlds. Six degrees of separation between me and everyone else on this planet. But to find the right six people...” from [Six Degrees of Separation](#)

The [small world](#) property is mathematically intuitive: if the number of vertices within a given distance of a specific node growth exponentially with the distance, then the average path length increases a $\log n_V$.

Formally, we talk here about the [average path length](#)

$$\bar{\ell} = \binom{n_V}{2}^{-1} \sum_{u \neq v} d(u, v) = O(\log n_V)$$

Small Word Property

Intuition: if $d_v = d$, if we reach everyone after k hops $n_V \sim d^k$ i.e. $k = O(\log n_V)$

“Each person in the world (at least among the 1.59 billion people active on Facebook) is connected to every other person by an average of three and a half other people. The average distance we observe is 4.57, corresponding to 3.57 intermediaries or ‘degrees of separation’”, from Three and a half degrees of separation

An alternative is based on the harmonic mean, related to graph efficiency

$$\ell^{-1} = \frac{2}{n_V(n_V - 1)} \sum_{i > j} d_{i,j}^{-1}$$