

15. The SPIN Model Checker



Computer-Aided Verification

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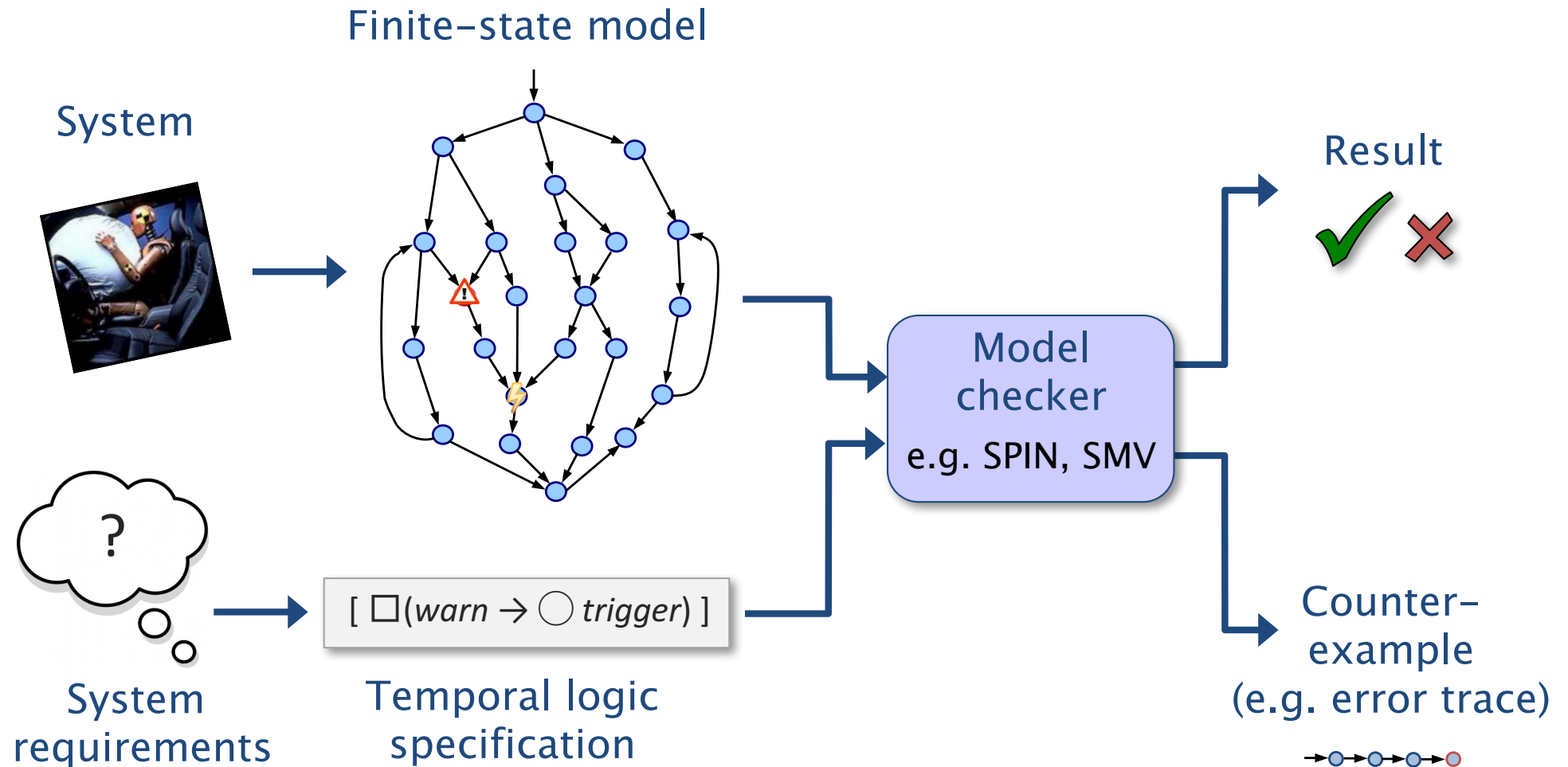
Module syllabus

- Modelling sequential and parallel systems
 - labelled transitions systems, parallel composition
- Temporal logic
 - LTL, CTL and CTL*, etc.
- Model checking
 - CTL model checking algorithms
 - automata-theoretic model checking (LTL)
- Verification tools: SPIN
- Advanced verification techniques
 - bounded model checking via propositional satisfiability
 - (symbolic execution), (symbolic model checking)
- Quantitative verification
 - (real-time systems), probabilistic systems

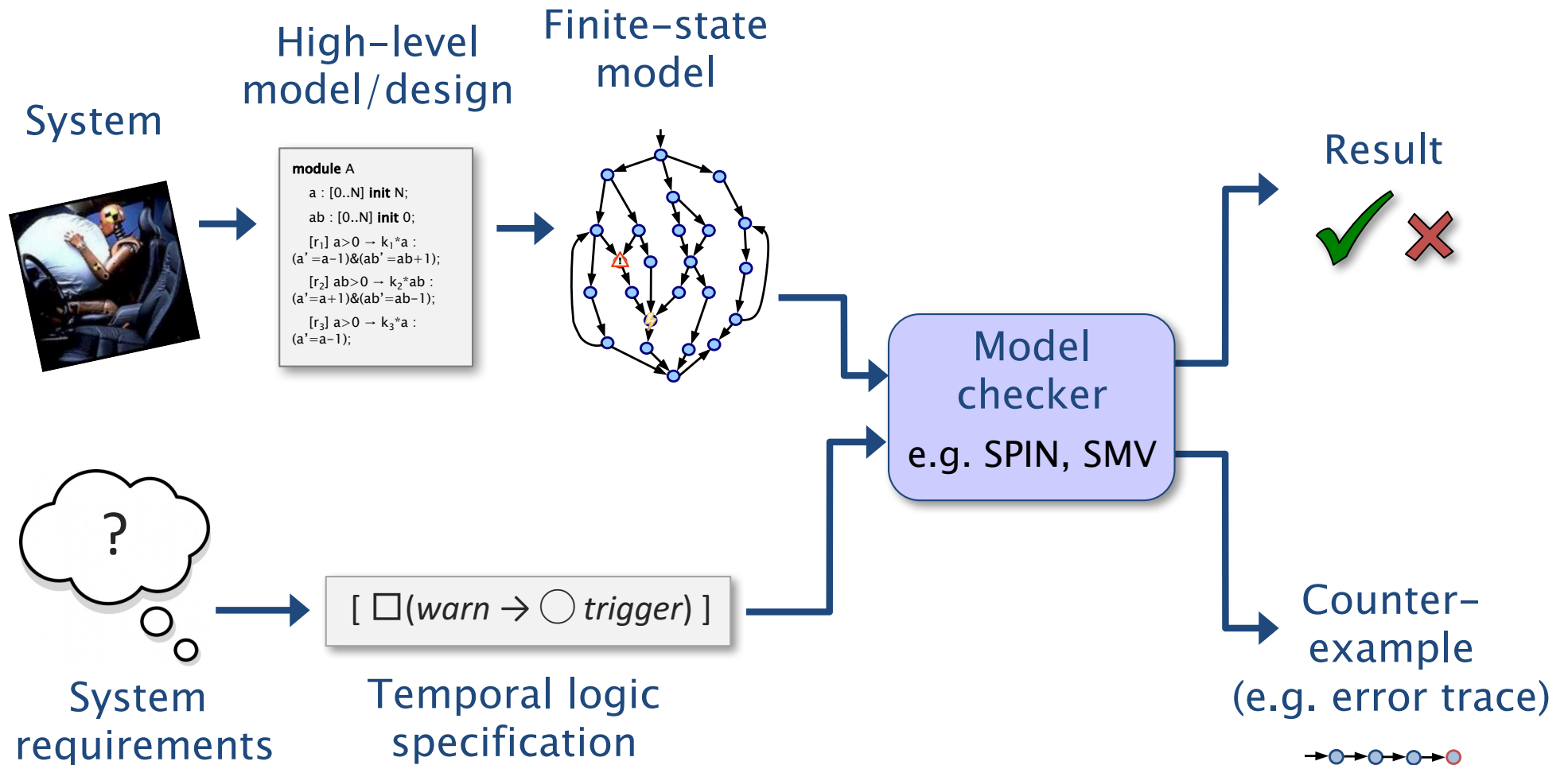
Overview

- The SPIN model checker
 - overview
 - ProMeLa
 - demo
 - examples
- Background reading & reference:
 - basic manual: <http://spinroot.com/spin/Man/Manual.html>
 - tool options: <http://spinroot.com/spin/Man/Spin.html>

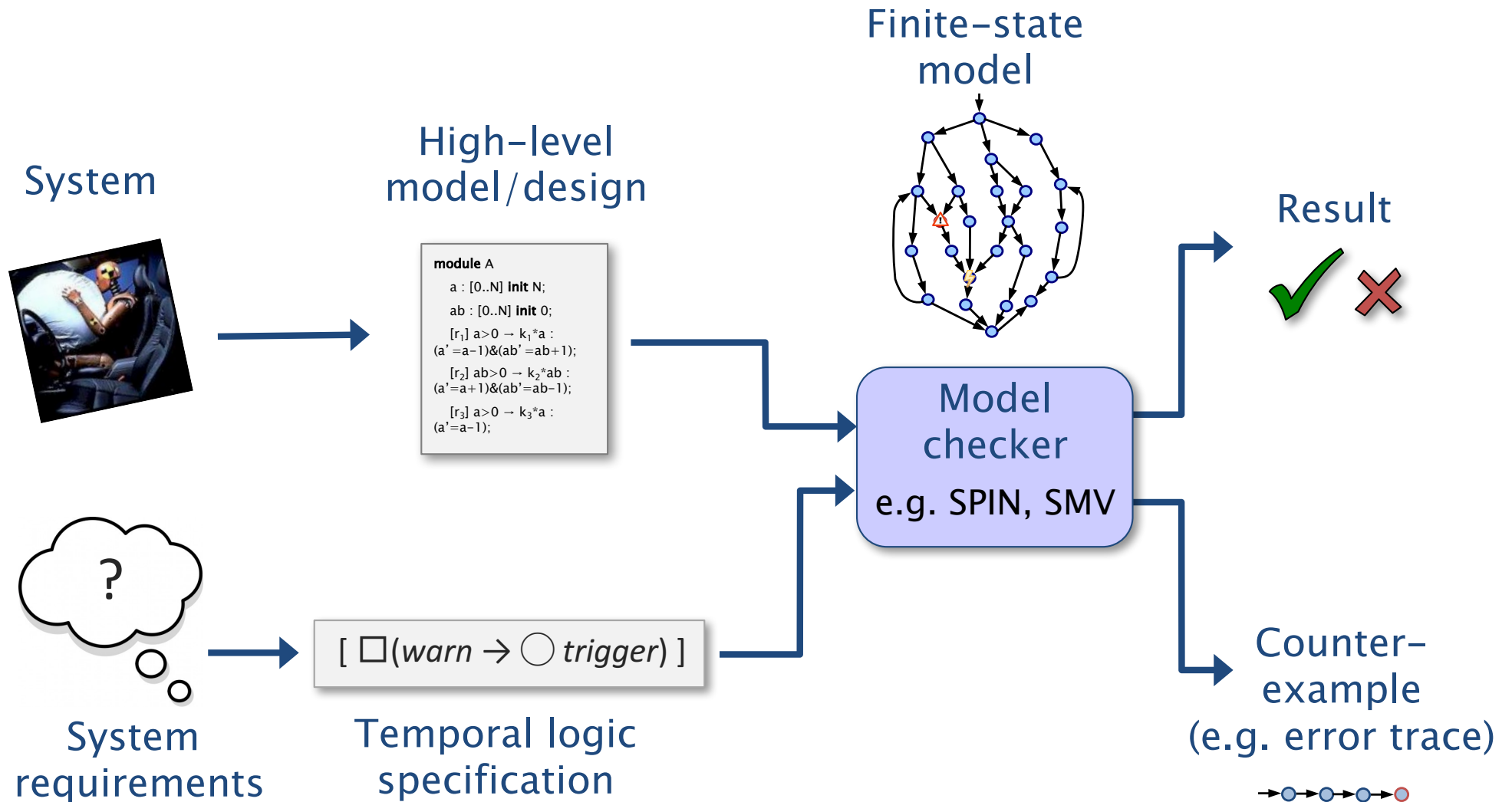
Verification via model checking



Verification via model checking



Verification via model checking



High-level model description

- Example concurrent program:

```
process Inc = while true do if x < 200 then x := x + 1 fi od
```

```
process Dec = while true do if x > 0 then x := x - 1 fi od
```

```
process Reset = while true do if x = 200 then x := 0 fi od
```

- ProMeLa/SPIN model (fragment):

```
int x=0;  
proctype Inc () {  
    do :: true -> if :: (x < 200) -> x = x + 1 fi od  
}  
proctype Dec() {  
    do :: true -> if :: (x > 0) -> x = x - 1 fi od  
}  
...
```

SPIN

- SPIN model checker
 - prominent, widely used verification tool
 - open source, freely available (<http://spinroot.com/>)
 - developed by Gerald Holzmann at Bell Labs
 - many success stories: NASA mission software (Deep Space 1, Mars Exploration Rovers), Toyota control software investigation
- Key features
 - custom modelling language: ProMeLa
 - on-the-fly model checking for safety, liveness, LTL
 - verification vs falsification (bug hunting)
 - state storage using hash table of lists of states
 - simulator (random, interactive, guided)
 - separate user interface (iSpin)

ProMeLa

- ProMeLa: Process Meta Language
 - modelling language for the SPIN tool
- Key ingredients
 - **processes** (one or more, in parallel)
 - typed **data** variables (can be shared, for communication)
 - **channels** (synchronous/asynchronous communication)
- Language notation
 - **guarded commands** (nondeterministic choice)
 - plus **imperative**-style control flow (and embedded C)
- Semantics
 - labelled transition systems (LTSs), via program graphs

ProMeLa – Guarded commands

- Guarded commands

```
:: guard -> statement
```

- execute statement if guard is true

- Conditionals

```
if :: guard -> statement1  
   :: !guard -> statement2  
fi
```

- if (guard) then statement₁ else statement₂

- Loops

```
do :: guard -> statement od
```

- while (guard) do statement

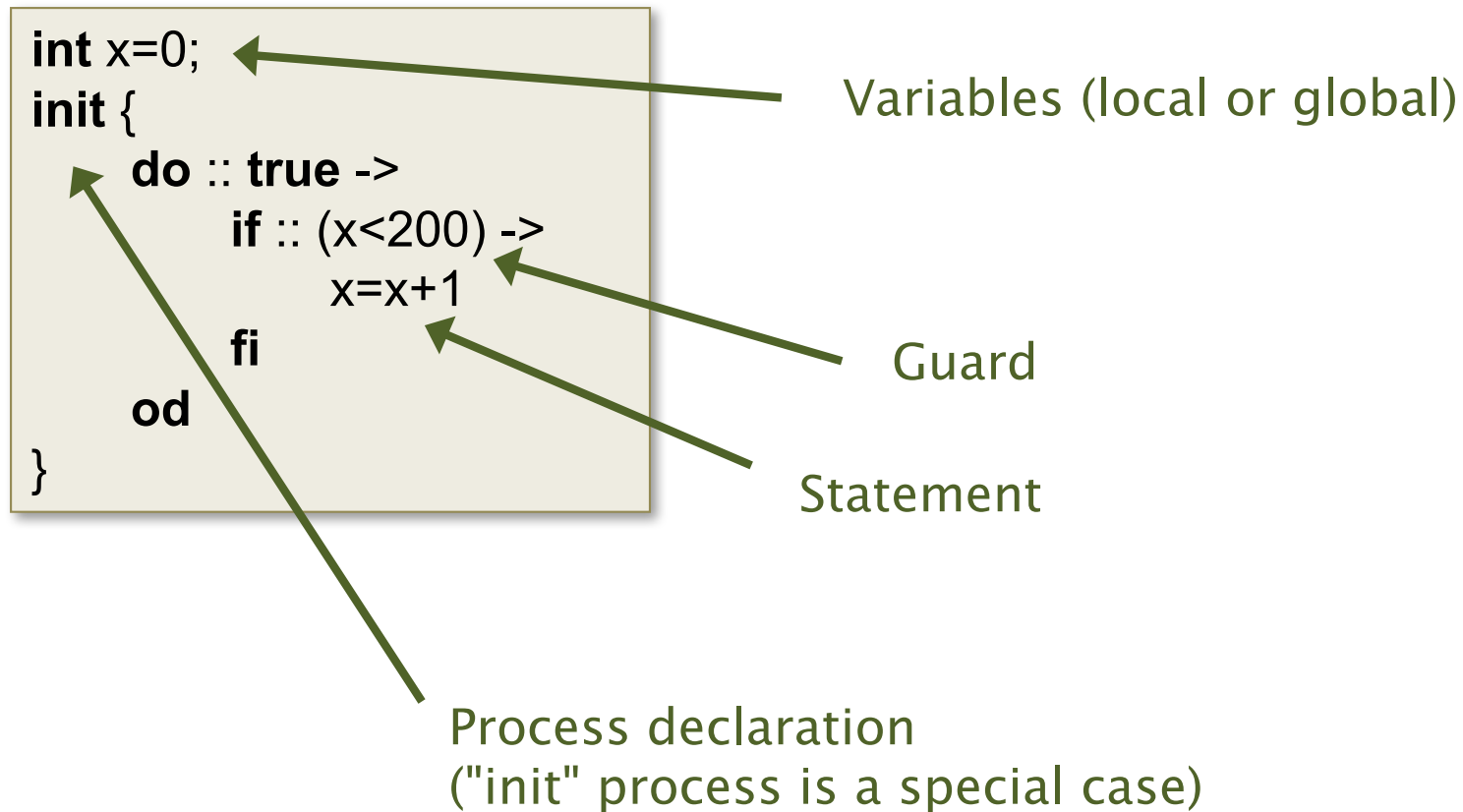
```
do :: guard1 -> statement1  
   :: guard2 -> statement2  
od
```

Nondeterminism



ProMeLa – Example

- One process, one variable (integer x)



Demo

ProMeLa – Concurrent program

- Earlier example:

```
int x = 0;

proctype Inc() {
    do :: true -> if :: (x < 200) -> x = x + 1 fi od
}
proctype Dec() {
    do :: true -> if :: (x > 0) -> x = x - 1 fi od
}
proctype Reset() {
    do :: true -> if :: (x==200) -> x=0 fi od
}

init {
    run Inc() ; run Dec() ; run Reset()
}
```

ProMeLa – Concurrent program

- Add a "monitor" process that checks whether $0 \leq x \leq 200$

```
int x = 0;
proctype Inc() {
    do :: true -> if :: (x < 200) -> x = x + 1 fi od
}
proctype Dec() {
    do :: true -> if :: (x > 0) -> x = x - 1 fi od
}
proctype Reset() {
    do :: true -> if :: (x==200) -> x=0 fi od
}
proctype Check () {
    assert (x >= 0 && x <= 200)
}
init {
    run Inc() ; run Dec() ; run Reset() ; run Check()
}
```