Explanation of the Frog-Star II Algorithm

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In the self-similar Poisson(1) frog model SFM(p, d), let V be the total number of visits to the root. Then V has the following RDE:

$$V \stackrel{\mathrm{D}}{=} \operatorname{Pois}(p_d^*) + \operatorname{Binom}(V^{(1)}, \hat{p}) + \sum_{i=2}^d A_i \operatorname{Binom}(V^{(i)}, \hat{p}), \tag{1}$$

where $A_i := \mathbf{1}_{\{\text{subtree at } v_i \text{ is activated}\}}, V^{(1)}, \dots, V^{(d)}$ are independent copies of V, $p_d^* = \frac{p(d-1)}{d-(d+1)p}$, and $\hat{p} = \frac{p}{1-p}$. Note that we will be abusing our notation here: the Poisson and Binomial's above refer to their independent random variables.

Since the root visits have $V_{\mathrm{SFM}(d,p)} \leq V_{\mathrm{nbFM}(d,p)} \leq V_{\mathrm{FM}(d,p)}$, proving $\mathrm{P}(V_{\mathrm{SFM}(d,p)} = \infty) = 1$ implies that $\mathrm{P}(V_{\mathrm{FM}(d,p)} = \infty) = 1$. Our idea is to show that for any $\lambda > 0$, given $V \succeq \mathrm{Pois}(\lambda)$, it follows that $V \succeq \mathrm{Pois}(\lambda + \epsilon)$ for a fixed $\epsilon > 0$. Since from (1) we have $V \succeq \mathrm{Pois}(p_d^*)$, by repeated application of what we mentioned we can show that $\mathrm{P}(V_{\mathrm{SFM}(d,p)} = \infty) = 1$.

Assume $V \succeq \operatorname{Pois}(\lambda)$, then we may replace independent $V^{(1)}, \ldots, V^{(d)}$ in (1) by independent $\operatorname{Pois}(\lambda)$ random variables and get

$$V \succeq \operatorname{Pois}(p_d^*) + \operatorname{Binom}(\operatorname{Pois}(\lambda), \hat{p}) + \sum_{i=2}^d A_i \operatorname{Binom}(\operatorname{Pois}(\lambda), \hat{p}),$$
$$= \operatorname{Pois}(p_d^*) + \operatorname{Pois}(\lambda \hat{p}) + \sum_{i=2}^d A_i \operatorname{Pois}(\lambda \hat{p}). \tag{2}$$

Now consider the auxilliary model $frog\text{-}star\ ii\ below^1$, which takes out the first three layers of the tree and replaces the distribution of each $v_j \to \varnothing'$ by $Pois(\lambda) \preceq V$. Now define W to be the number of activated vertices among v_2, \ldots, v_d in this frog-star ii. Note that the p.m.f. of W is now computable because we are in a finite model. By our setup we also have $W \preceq \sum_{i=2}^d A_i$. Hence from (2) onward we have

$$V \succeq \operatorname{Pois}(p_d^*) + \operatorname{Pois}(\lambda \hat{p}) + \sum_{i=2}^d A_i \operatorname{Pois}(\lambda \hat{p})$$
$$\succeq \operatorname{Pois}(p_d^*) + \operatorname{Pois}(\lambda \hat{p}) + W \cdot \operatorname{Pois}(\lambda \hat{p})$$
$$= \operatorname{Pois}(p_d^* + \lambda \hat{p}(1 + W)).$$

We now need to find under what λ is $\operatorname{Pois}(p_d^* + \lambda \hat{p}(1+W)) \succeq \operatorname{Pois}(\lambda + \epsilon)$ for some fixed $\epsilon > 0$.

To be specific, it has Poisson(1) frogs at \varnothing' and Poisson(λ) frogs at each of the leaves v_1, \ldots, v_d . \varnothing' and v_1 are activated at the beginning. The initial frogs at \varnothing' move to \varnothing with probability p_d^* , and to the leaves uniformly with total probability $1 - p_d^*$. Active frogs at leaf v_j moves to \varnothing' with probability 1, and then to \varnothing with probability \hat{p} , and to $v_1, \ldots, v_{j-1}, v_{j+1}, \ldots, v_d$ uniformly with total probability $1 - \hat{p}$.

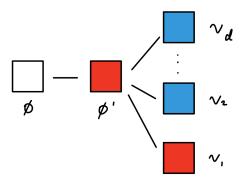


Figure 1: frog-star ii model

This is equivalent to proving that $\mathbf{E} \exp(p_d^* + \lambda \hat{p}(1+W)) \leq \exp(-\lambda - \epsilon)$. Simplification gives us

$$\mathbf{E} \exp(-p_d^* - \lambda \hat{p}(1+W)) = e^{-p_d^*} \sum_{w=0}^{d-1} \exp(-\lambda \hat{p}(1+w)) P(W=w)$$
$$= e^{-p_d^* - \lambda} \sum_{w=0}^{d-1} \exp(\lambda (1-\hat{p}(1+w))) P(W=w),$$

and thus

To remain consistent with the literature, let m = d and $\tilde{p} = p_d^* = \frac{p(m-1)}{m-(m+1)p}$. Under a fixed m, we want to numerically compute the minimal p such that

$$h(\lambda) = e^{-\tilde{p}} \sum_{w=0}^{m-1} \exp\left(\lambda(1 - \hat{p}(1+w))\right) P(W=w) \le 1 - \epsilon$$
(3)

holds for all $\lambda > 0$.

The p.m.f. of W can be computed via the following recursion algorithm². Define φ to be the probability that 0 of the initial frogs at \varnothing' moves to v_2 , and v to be the probability that 0 of the initial frogs at v_1 moves to v_2 . It is clear that $\varphi = \exp\left(-\frac{1-\tilde{p}}{m}\right)$ and $v = \exp\left(-\frac{\lambda(1-\hat{p})}{m-1}\right)$. Now let $p_{m,k} = P_m(W = k)$ be the probability of W = k when the frog-star ii has m leaves. With the initial

²c.f. Eric Yu

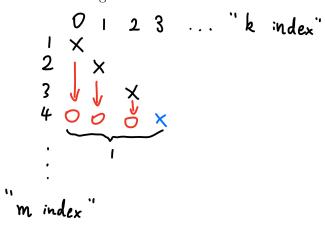
condition $p_{1,0} = 1$, the recursion is established by the following two equations:

$$p_{m,k} = {\binom{m-1}{k}} (\varphi v^{k+1})^{m-1-k} \cdot p_{k+1,k} \text{ when } 0 \le k \le m-2, \text{ and}$$
 (4)

$$p_{m,m-1} = 1 - \sum_{i=0}^{m-2} p_{m,i}.$$
 (5)

To visualize how to find the p.m.f. of W, consider the following table. Here the entry in the m-th row

Figure 2: recursion table



and k-th column corresponds to $p_{m,k}$. Take m=4 for example; $p_{4,0}, p_{4,1}$, and $p_{4,2}$ are all determined by the diagonal elements determined above [by equation (4)]. Since every row sums up to 1, the new diagonal element $p_{4,3}$ can be computed by $1-\sum_{i=0}^{2} p_{4,i}$ [equation (5)].

Hence it makes sense to implement the p.m.f. computation using a symbolic matrix. Be aware that Python-based SageMath is 0-indexed, and hence we have to shift the row index m here back by 1 (replaced by j in our code).

For numerical purposes we let p be an irreducible fraction $\frac{a}{b} < \frac{1}{2}$, and then

$$\hat{p} = \frac{a}{b-a}, 1 - \hat{p} = \frac{b-2a}{b-a}, \ \tilde{p} = \frac{am-a}{bm-am-a}, \ \text{and} \ 1 - \tilde{p} = \frac{bm-2am}{bm-am-a}.$$
 (6)

The way we want to approach equation (3) is that we want to convert $h(\lambda)$ into a polynomial g(y) by mapping $\lambda \mapsto -c \cdot \log(y)$ for some constant c. Through this change of variable it suffices to consider that

$$\sup_{\lambda \in (0,\infty)} h(\lambda) = \sup_{y \in (0,1)} g(y) = \max_{y \in [0,1]} g(y) \le 1 - \epsilon.$$

It turns out a working constant c can be chosen in general given m and $p = \frac{a}{b}$: let

$$c = (b-a)(m-1).$$

Now we verify our choice. In the p.m.f. of W, we have terms $\varphi = \exp(\frac{2a-b}{bm-am-a})$ and $v = y^{b-2a}$ by (6). Now we rewrite our objective function $h(\lambda)$ from (3) as follows:

$$g(y) = \exp\left(\frac{a - am}{bm - am - a}\right) \sum_{k=0}^{m-1} y^{[-(b-a)(m-1)] \cdot \left(\frac{b-2a}{b-a} - \frac{a}{b-a}k\right)} \cdot P(W = k)$$

$$= \exp\left(\frac{a - am}{bm - am - a}\right) \sum_{k=0}^{m-1} y^{(m-1)[(k+2)a-b]} \cdot P(W = k). \tag{7}$$

It remains to show that the exponents of y in (7) are indeed nonnegative integers. It is easy to show by induction that all $p_{m,k}$'s in the lower triangle of the recursion table only have nonnegative integer powers of y.

Then for $0 \le k \le m-2$, we know from (4) that the exponents of y in P(W=k) are all at least (b-2a)(k+1)(m-1-k). Thus in (7), the exponents in each summand from k=0 to m-2 are at least

$$(m-1)ak - (m-1)(b-2a) + (b-2a)(k+1)(m-1-k)$$
$$= (b-2a)(m-2-k)k + (m-1)ak \ge (m-1)ak \ge 0.$$

We now consider the summand when k = m - 1, i.e., $y^{(m-1)[(m+1)a-b]} \cdot P(W = m - 1)$. By (5) it suffices to show that $(m+1)a - b \ge 0$, i.e., $p = \frac{a}{b} \ge \frac{1}{m+1}$. This lower bound to the critical drift p_m was mentioned briefly in [BFJ+19]; basically we want to show that

assuming all frogs are initially awake in the original one-per-site model [meaning that this is a more recurrent model], if p < 1/(m+1), then there are only finite expected visits in total to the root.

Eric Yu wrote a proof of this in the general Discord channel early in June. The overall idea is that we wish to write the expected number of visits into a infinite sum of root visits E_n from vertices at each layer n to the root. It is quite clear that these random variables have their internal recursive relation, from which we can derive an explicit formula for these E_n 's. The summation will then become a geometric series. Since a more recurrent model is transient for p < 1/(m+1), $p_m \ge 1/(m+1)$ as desired.

It is really nice to prove that our substitution constant always gives a polynomial g(y), but ultimately we can use the computation algorithm below only up to a finite m. Showing that this substitution gives integer exponents and visually checking that g(y) is indeed a polynomial might be an alternative (albeit more tedious) approach.

Now we start explaining our algorithm. The only thing that needs to be adjusted is p = a/b based on our choice of m. The symbolic expression of the polynomial g(y) will then be computed accordingly. As preliminary exploration we shall visualize the graph of g(y) over the interval [0, 1].

It is almost always the case by visual inspection that g(y) attains its only peak on the interval (0,1) when y is somewhere close to 1, or equivalently, the original λ is not too far from 0.

The CountRoots method specifically from Mathematica is of great importance to us. In our context we use the Mathematica code CountRoots[dg[y],{y,0,1}] (which we imported into Sage-Math). This method uses Sturm's algorithm to count the number of roots (including multiplicities) on the designated closed interval [0,1]. Call it count. By our previous visual inspection we hope that this count is only 1 more than the multiplicity of the root 0. This multiplicity can be found easily by factoring g'(y) in SageMath, which gives

(some polynomial with nonzero constant) $\cdot y^t \cdot \text{(some constant)}$.

Hence this t is the multiplicity. If count = t + 1, then there is only one root of g'(y) in (0,1], which can be found numerically using $diff(g(y)).find_{root}(0.9,1)$. This finds a root of g'(y) on the closed interval $[0.9,1]^3$ using Brent's method, which in our case should be the only root y_0 in $[0.9,1] \subseteq (0,1]$.

Now there are two ways to check whether $y_0 = \arg\max_{y \in [0,1]} g(y)$ that I can think of. One approach is to use the second derivative test at y_0 : just verify $g''(y_0) < 0$ numerically. I am more in favor of the second approach, which is precise and avoids using g'' or g'. We know g(y) attains its maximum on the closed interval [0,1]. If we could check that $g(y_0) > g(0)$ and $g(y_0) > g(1)$, then the maximum of g(y) is attained in the open interval (0,1). By Fermat's theorem⁴ we know that this maximum must be attained at $y = y_0$, the only point in (0,1) such that g'(y) = 0.

We are only left to check if $g(y_0) = \max_{y \in (0,1)} g(y) \le 1 - \epsilon$. If this is correct, then our choice of p = a/b gives an upper bound to the critical drift $p_m = \inf\{p : \operatorname{FM}(m,p) \text{ is recurrent}\}$.

On the next page I will give a list of optimal upper bounds p that can be achieved numerically by the steps above. I only included m between 2 and 13. The Brent's method used in find_root starts to give error when $m \ge 14$ even for p with relatively small a and b. The length of the polynomial g(y) also becomes extremely long, and I doubt if we can go any further than m = 15.

³As I mentioned earlier, by visual inspection we know that $\max_{y \in [0,1]} g(y)$ should be achieved somewhere close to 1. Thus I choose 0.9 here.

⁴https://en.wikipedia.org/wiki/Fermat%27s theorem (stationary points)

m	upper bound $p =$
2	55/159 < 0.3460
3	31/107 < 0.2898
4	17/65 < 0.2616
5	23/94 < 0.2447
6	46/197 < 0.2336
7	23/102 < 0.2255
8	29/132 < 0.2197
9	20/93 < 0.2151
10	11/52 < 0.2116
11	5/24 < 0.2084
12	7/34 < 0.2059
13	11/54 < 0.2038

References

[BFJ+19] Erin Beckman, Natalie Frank, Yufeng Jiang, Matthew Junge, and Si Tang. "The frog model on trees with drift". In: *Electronic Communications in Probability* 24.26 (2019), pp. 1–10. DOI: 10.1214/19-ECP235.