$$\frac{\Gamma \vdash i : \operatorname{Int}}{\Gamma \vdash i : \operatorname{Int}} \quad \frac{\Gamma \vdash e_1 : \operatorname{Int}}{\Gamma \vdash e_1 \; \Box \; e_2 : \operatorname{Int}} \quad \frac{\Gamma \vdash e_1 : \operatorname{Int}}{\Gamma \vdash e_1 \; \Box \; e_2 : \operatorname{Int}} \quad T\text{-Arith}$$

$$\frac{\Gamma \vdash i : \operatorname{Int}}{\Gamma \vdash \operatorname{true} : \operatorname{Bool}} \quad T\text{-True} \quad \frac{\Gamma \vdash e_1 : \operatorname{Bool}}{\Gamma \vdash \operatorname{false} : \operatorname{Bool}} \quad T\text{-False}$$

$$\frac{\Gamma \vdash e_1 : \operatorname{Bool}}{\Gamma \vdash e_1 : \operatorname{Int}} \quad \frac{\Gamma \vdash e_2 : \operatorname{Int}}{\Gamma \vdash e_1 : \operatorname{then} \; e_2 \operatorname{else} \; e_3 : \tau_1} \quad T\text{-If}$$

$$\frac{\Gamma \vdash e_1 : \operatorname{Int}}{\Gamma \vdash e_1 : \operatorname{Int}} \quad \frac{\Gamma \vdash e_2 : \operatorname{Int}}{\Gamma \vdash e_1 : e_2 : \operatorname{Bool}} \quad T\text{-ReL}$$

$$\frac{\Gamma(x) = \mathsf{T}}{\Gamma \vdash e_1 : \mathsf{T}} \quad T\text{-Var} \quad \frac{\Gamma \vdash e_1 : \mathsf{T}_1 \quad x : \mathsf{T}_1, \Gamma \vdash e_2 : \mathsf{T}_2}{\Gamma \vdash \operatorname{let} \; x = e_1 \operatorname{in} \; e_2 : \mathsf{T}_2} \quad T\text{-Let}$$

$$\frac{x : \mathsf{T}_1, \Gamma \vdash e : \mathsf{T}_2}{\Gamma \vdash (\operatorname{lambda} \; x : \mathsf{T}_1 : e) : \mathsf{T}_1 \to \mathsf{T}_2} \quad T\text{-Lambda}$$

$$\frac{\Gamma \vdash e_1 : \mathsf{T}_1 \to \mathsf{T}_2 \quad \Gamma \vdash e_2 : \mathsf{T}_3 \quad (\mathsf{T}_1 = \mathsf{T}_3)}{\Gamma \vdash (e_1 \; e_2) : \mathsf{T}_2} \quad T\text{-App}$$

$$\frac{f : \mathsf{T}, \Gamma \vdash e : \mathsf{T}}{\Gamma \vdash (\operatorname{fix} \; f : \mathsf{T} \; \operatorname{is} \; e) : \mathsf{T}} \quad T\text{-Fix}$$

$$\begin{array}{c} \Gamma \vdash (\mathsf{fix}\, f : \mathsf{T} \; \mathsf{is} \; e) : \mathsf{T} \\ \\ \frac{}{\Gamma \vdash \mathsf{Nil}[\mathsf{T}] : \mathsf{List}[\mathsf{T}]} \; \mathsf{T-NIL} & \frac{\Gamma \vdash e_1 : \mathsf{T} \quad \Gamma \vdash e_2 : \mathsf{List}[\mathsf{T}]}{\Gamma \vdash e_1 : : e_2 : \mathsf{List}[\mathsf{T}]} \; \mathsf{T-Cons} \\ \\ \Gamma \vdash e_1 : \mathsf{List}[\mathsf{T}_1] & \Gamma \vdash e_2 : \mathsf{T}_2 \end{array}$$

$$\frac{x:\mathsf{T}_1,y:\mathsf{List}[\mathsf{T}_1],\Gamma\vdash e_3:\mathsf{T}_3}{\Gamma\vdash\mathsf{match}\ e_1\ \mathsf{with}\ \mathsf{Nil}\to e_2\mid x::y\to e_3\ \mathsf{end}:\mathsf{T}_2}\ \mathsf{T-MATCH}}$$

$$\frac{\Gamma\vdash e:\mathsf{T}_2}{\Gamma\vdash (e\@\mathsf{T}_1):\mathsf{T}_1}\ \mathsf{T-Annor}$$

Problem 5 (20 points)

Consider the following λ^+ expression:

```
(lambda n: Int.
  if n = 0 then (lambda k: Int. true) else (lambda m: Int. r (n-1) (m-1))) 7
```

Denote the above expression as e. We claim e is well-typed under the environment $1 = r : \text{Int} \to \text{Int} \to \text{Bool}$. Below is the skeleton of the derivation tree that shows $1 \vdash e : 1$ for some type 1. The derivation starts from the bottom of the page with the double bar, and proceeds upwards. Complete the derivation tree by filling in the table on the next page. A solid box should be filled with an λ^+ expression or type, a dashed box with the name of a typing rule, and a circled number n with a typing environment. Note that type equalities of the form 1 = 1 are omitted from the tree, but you should mentally check that they hold.