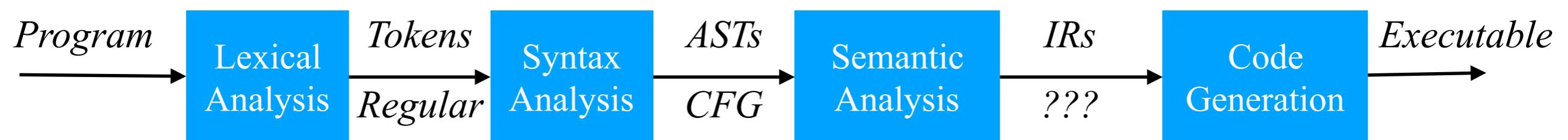


CS 162 Programming languages

Lecture 15: Course Review & Software Robustness

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Winter 2021

A typical flow of a compiler



Want to learn compiler? Check out my CS160 in Fall

What we have learned

- Core PL
 - Lambda calculus (alpha-renaming, beta-reduction)
 - Operational semantics
 - Type checking
 - Type inference
- Functional programming
 - Higher-order functions, recursions
 - Data types

WHO CARES?

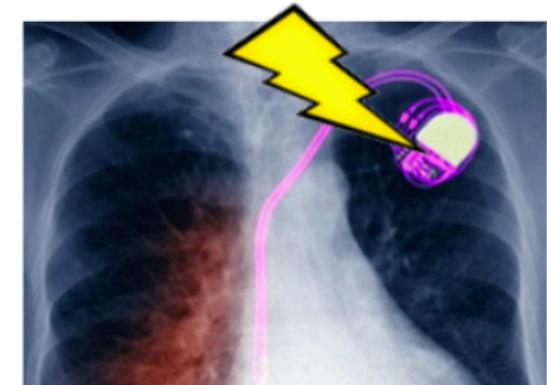
Software is great



Software is NOT so great



```
0101010101101010110101101  
0110101 NAME ADRES  
01101001010010101101001001  
0110101 LOGIN PASSWORD 1  
01101001010010101101001001  
01101010 NAME ADRES  
01101001010010101101001001  
01101010101101010110101010  
011010010100101011010010011010  
0110101010110101011010011010
```



How to make a robust bridge



Make it thicker!

How to make robust software

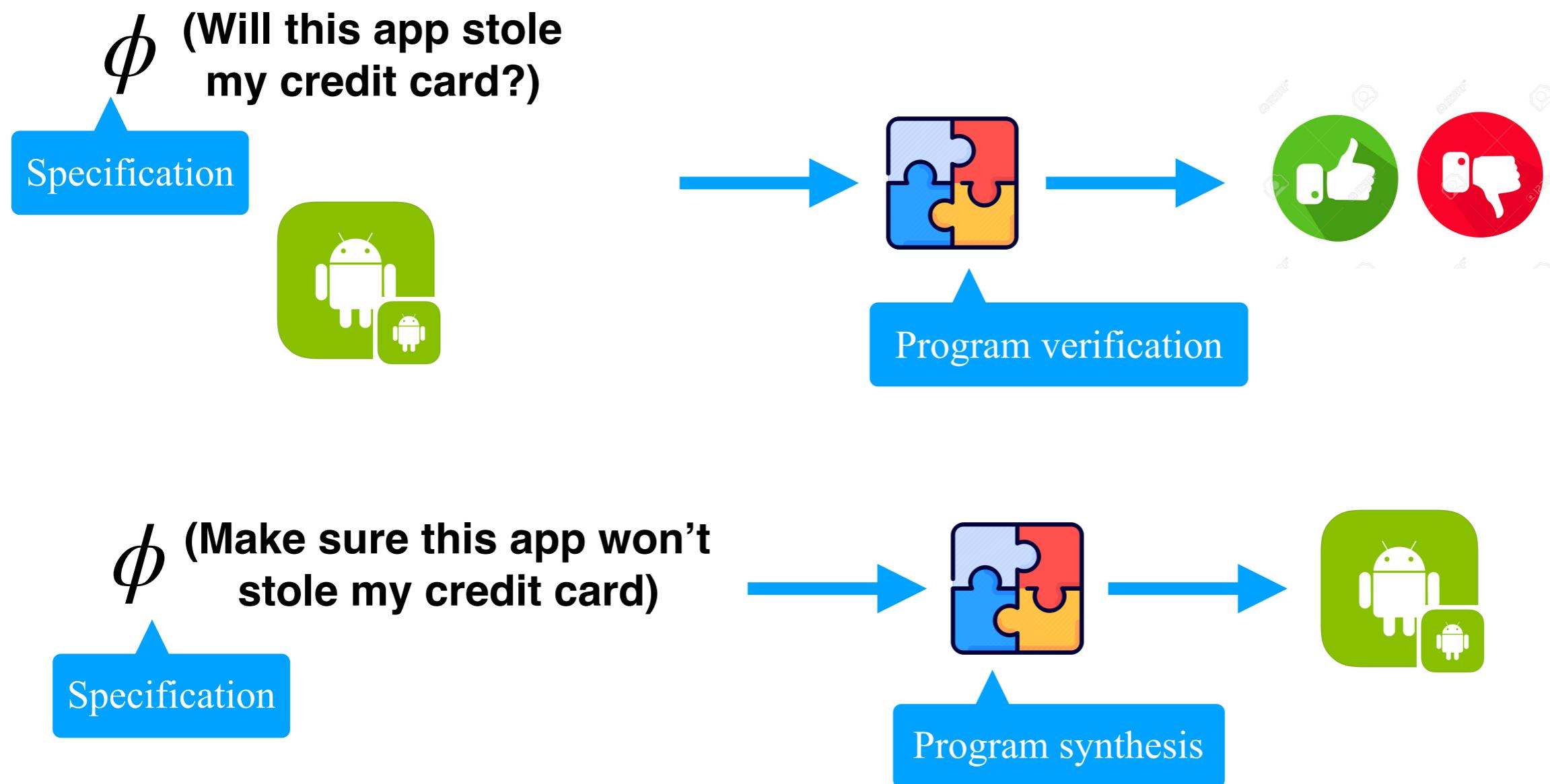


It is undecidable!

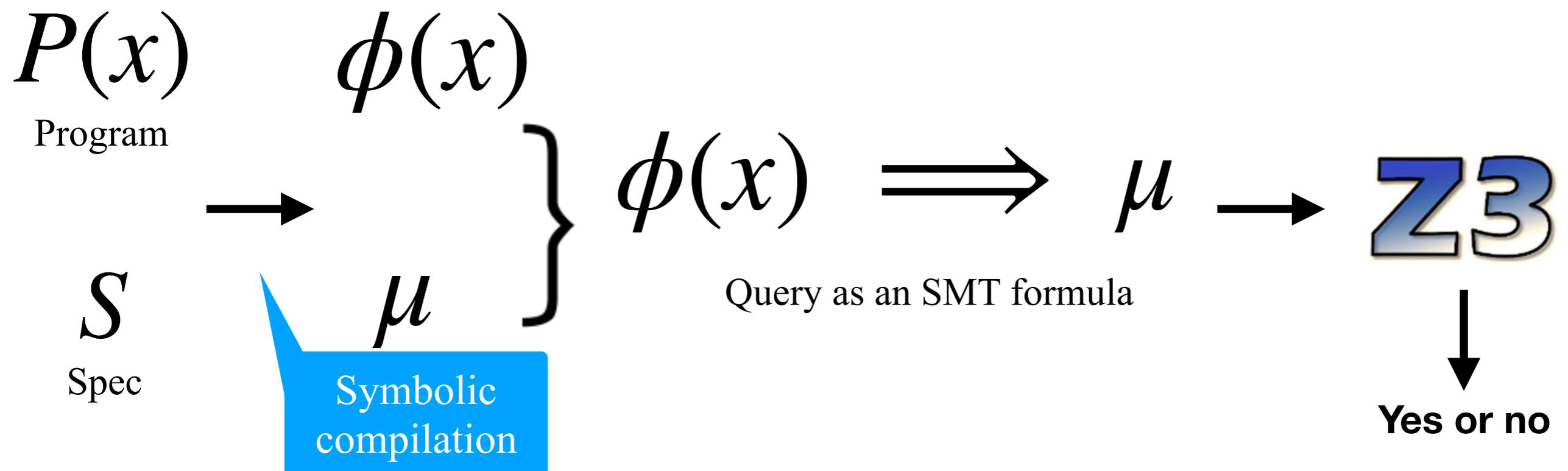


Let me show you a demo with a concrete example

Ensure robustness with PL



Checking correctness



Symbolic compilation can take years of effort!

<https://github.com/Z3Prover/z3>

Checking correctness

```
foo () {  
    x = 10;  
    y = 5;  
}
```

$$x = 10 \wedge y = 5$$

```
foo (int a) {  
    if (a > 0)  
        x = 10;  
    else  
        y = 5;  
}
```

$$a > 0 \implies x = 10 \wedge a \leq 0 \implies y = 5$$

```
foo (int a) {  
    if (a > 0)  
        x = 10;  
    else  
        y = 5;  
    assert y > 4  
}
```

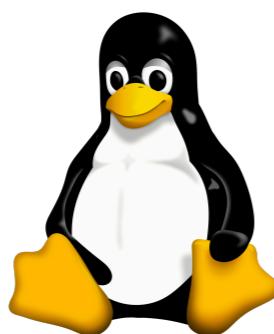
$$(a > 0 \implies x = 10 \wedge a \leq 0 \implies y = 5)$$

$$\implies y > 4$$

Checking correctness



??



??

How to deal with complex systems?

It is undecidable!

A programming model that integrates solvers into the language, providing constructs for program verification, synthesis, and debugging.

Solver-aided programming



```
p(x) {  
    v = 12  
  
    p(x) {  
        v = ??  
        ...  
    }  
    assert safe(x, p(x))
```

- Find an input on which the program fails.
- Localize bad parts of the program.
- Find values that repair the failing run.
- Find code that repairs the program.

Racket

- Like ML, functional focus with imperative features:
Anonymous functions, closures, pattern-matching, etc
- No static type system: accepts more programs, but most errors do not occur until run-time
- Advanced features like macros, modules, quoting/eval, continuations, contracts, ... (Will do only macros)

File structure

- Start every file with a line containing only
`#lang racket`
- A comment starts with semicolon “;”
- A file is a module containing a *collection of definitions* (bindings)

Example

```
#lang racket

(define x 3)
(define y (+ x 2))

(define cube ; function
  (lambda (x)
    (* x (* x x)))))

(define pow ; recursive function
  (lambda (x y)
    (if (= y 0)
        1
        (* x (pow x (- y 1)))))))
```

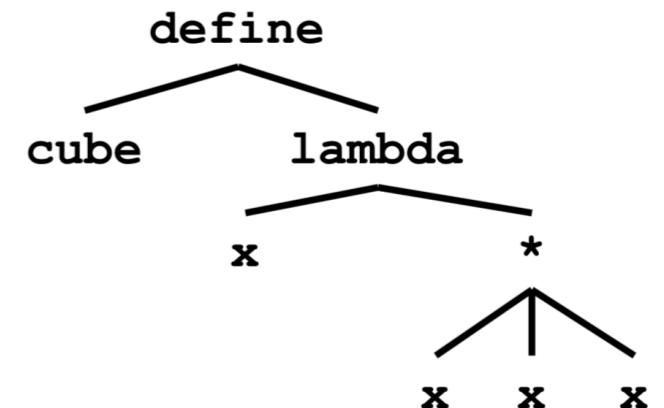
Why is this good?

- By parenthesizing everything, converting the program text into a tree representing the program (*parsing*) is trivial and unambiguous

- Atoms are leaves
- Sequences are nodes with elements as children

- (No other rules)

```
(define cube  
  (lambda (x)  
    (* x x x)))
```



- Also makes indentation easy
- No need to discuss “operator precedence” (e.g., $x + y * z$)

Our old friend: currying/ β -reduction

```
#lang racket

(define pow
  (lambda (x)
    (lambda (y)
      (if (= y 0)
          1
          (* x ((pow x) (- y 1)))))))

(define three-to-the (pow 3))
(define eightyone (three-to-the 4))
(define sixteen ((pow 2) 4))
```

Another old friend: list

- Empty list: **null**
- Cons constructor: **cons**
- Access head of list: **car**
- Access tail of list: **cdr**
- Append two lists: **append**
- Check for empty: **null?**
- Notes
 - Empty list is represented as '**()**' or **(list)** or **null**
 - **(list e₁ ... e_n)** for building lists

Another old friend: list

```
#lang racket

(define (sum xs)
  (if (null? xs)
      0
      (+ (car xs) (sum (cdr xs)))))

(define (my-append xs ys)
  (if (null? xs)
      ys
      (cons (car xs) (my-append (cdr xs) ys))))

(define (my-map f xs)
  (if (null? xs)
      null
      (cons (f (car xs)) (my-map f (cdr xs)))))
```