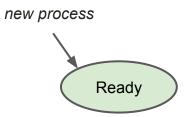
Process Lifecycle and Unix Process Creation



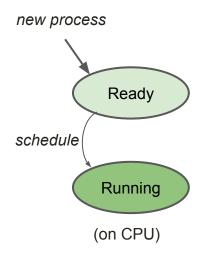
- To execute an instance of a program, new process must be created
 - Upon creation, process becomes active or *ready* for execution



- Multiple processes can be simultaneously ready
 - E.g., browser, music player, messaging app, etc. can be simultaneously open

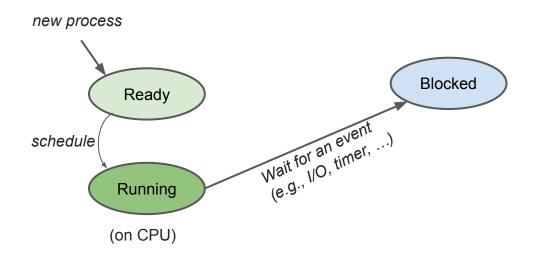


- When CPU is free, OS chooses / schedules a ready process to run on the CPU
 - The chosen / scheduled process can actively execute instructions on the CPU





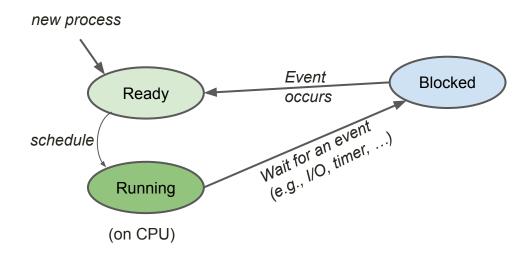
- Running process may sometimes need to pause & wait for some event
 - Process is put into blocked state so that it doesn't unnecessarily occupy the CPU



- In the meantime, OS schedules another *ready* process to *run*
 - I.e., OS switches between processes → concept of multiprogramming



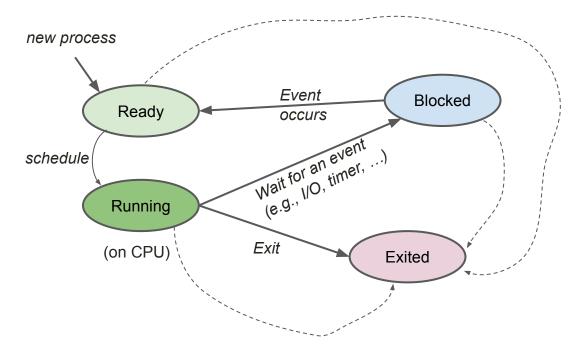
- Once event that blocked process is waiting on occurs...
 - ...process becomes *ready* again



• This *ready – running – blocked* cycle for a process can repeat

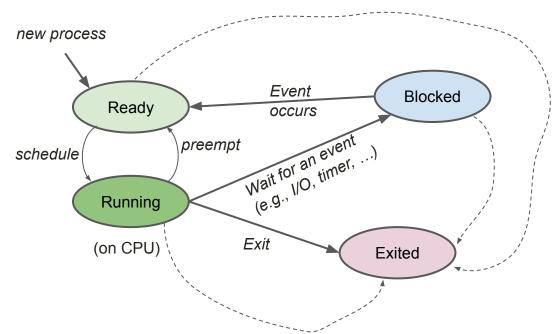


- Once process is done, it exits
- Process can actually exit from any state
 - E.g., user manually kills process, process gets killed due to error, etc.





- While process is running, OS may pause/preempt it and schedule another process (e.g., due to priority or simply for timesharing)
 - Preempted process goes back to ready state





Newly scheduled processes moves to running state

- Every time execution of running process is paused (due to blocking / preemption)
 - Its execution context needs to be saved so that it can resume later
- Every time a ready process is scheduled
 - Its execution context needs to be loaded / restored for it to run
- → Context Switching



To enable multiprogramming & context switching...

- ...OS must maintain information & execution context for each process
 - Process control struct / block used for this
 - Process ID
 - Process State (ready, running etc.)
 - Program Counter address of next instruction to be executed
 - Registers general purpose registers, stack pointer etc.
 - Scheduling information
 - Memory management information
 - Accounting information time limits, etc.
 - **■** ...



Now that we understand the process lifecycle in general, let's discuss process creation and termination in Unix-based systems



Unix process creation (fork)

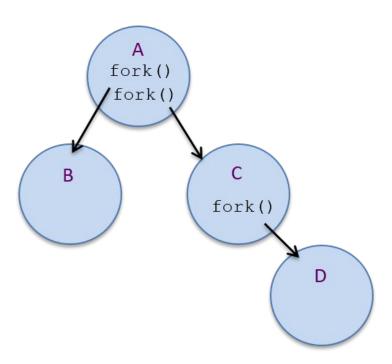
- Existing process makes fork system call to create new process
 - First process (*init*) gets created at system start

- Creating process is the parent & new process is the child
 - Data from parent process copied to memory of child process
 - Memory image, environment settings, I/O handles, etc.
 - Child gets its own ID, process control block, scheduling info, etc.
- Hierarchy of parent-child relationships exist between active processes



Unix process creation (fork)

- E.g.: If process A creates two new processes (B and C)...
 - ...and process C in turn creates another new process (D)





Differentiating between parent & child

- If process creation fails, fork returns -1 to parent & no child is created
- Successful fork call returns different values to parent and child
 - Fork returns child's process ID to parent process...
 - ...and returns 0 to child process
- Parent/child independently continue from execution point immediately after fork
 - Order of execution of the two processes may vary

```
pid_t id;
id = fork(); /* create new process */
if(id == -1) { /* error creating process */
    printf ("Error creating process\n");
} else if (id == 0) { /* only child process executes this */
    printf ("I'm a child process!\n");
} else { /* only parent process executes this */
    printf ("I just became a parent!\n");
}
```



Switching programs

- Child can *change* the program it executes
 - Invoke exec family of system calls to change program that it executes
 - Child stops executing code in parent program & starts executing new program

```
pid t id;
char *argv[2];
args[0] = "echo"; /* name of new program for child to execute */
args[1] = "hello"; /* parameter to pass to the above program */
args[2] = NULL; /* end of parameter list */
id = fork(); /* create new process */
if(id == -1) { /* error creating process */
    printf ("Error creating process\n");
} else if (id == 0) { /* only child process executes this */
    execvp(arqs[0], arqs); /* child process switches to program indicated by
                            args[0] with parameters in rest of args array */
} else { /* only parent process executes this */
   printf ("I just became a parent!\n");
```

Process termination

- As discussed earlier, process may terminate for many reasons
 - Voluntary exit upon task completion
 - Voluntary exit due to fatal error
 - Involuntary exit due to error/bug
 - Involuntary exit due to *kill* command by OS or other process who is authorized to do this



On Unix based systems...

- Parent must be allowed to read child's exit status or return value
 - Parent process can reap children by waiting for them to terminate
 - OS provides system call for this → wait
- If parent terminates before child...
 - Child is now an *orphan* process
 - Orphan processes are adopted by init process
- If a child process terminates before parent
 - System will still need to keep child's process control struct
 - Child process becomes a zombie process
 - "Dead", but not "reaped"





Waiting for child to terminate

```
pid t id, ret;
int status;
id = fork(); /* create new process */
if(id == -1) { /* error creating process */
    printf ("Error creating process\n");
} else if (id == 0) { /* only child process executes this */
    char * args[] = {"echo", "hello", NULL};
    execvp(args[0], args); /* child process switches to program indicated by
                            args[0] with parameters in rest of args array */
} else { /* only parent process executes this */
    printf ("I just became a parent!\n");
    ret = wait(&status); /* parent process waits for child to terminate */
```

