



LWPMFS: Light-Weight Persistent Memory Filesystem

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FreeBSD Developer Summit, 2018-10-18

Who am I? Why am I here?

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- **Panasas – High-Performance Storage Appliances**

- Founded in 1999; grew out of research from Carnegie Mellon University's Parallel Data Lab
- Clustered storage – data spread across lots individual servers
 - Custom hardware platform: blades, power-supplies, integrated network switches
- Native DirectFlow Protocol (kernel modules for Linux and macOS)
 - Direct parallel access to storage nodes
- Legacy Protocols (NFSv3, CIFSv2)
 - Access goes through protocol gateway, which does parallel access behind the scenes

Who am I? Why am I here?



- Ravi Pokala – rpokala@freebsd.org
 - Joined Panasas in late 2002; various hardware-facing hacking ever since
 - FreeBSD src-committer since 2015-11
- Caveat: I am not the original author of what is being presented
 - I inherited the bottom levels of the stack
 - Colleague now maintains the higher levels
 - Scheduling conflict – could not attend to present this himself
 - Take open questions to freebsd-hackers@freebsd.org mailing list

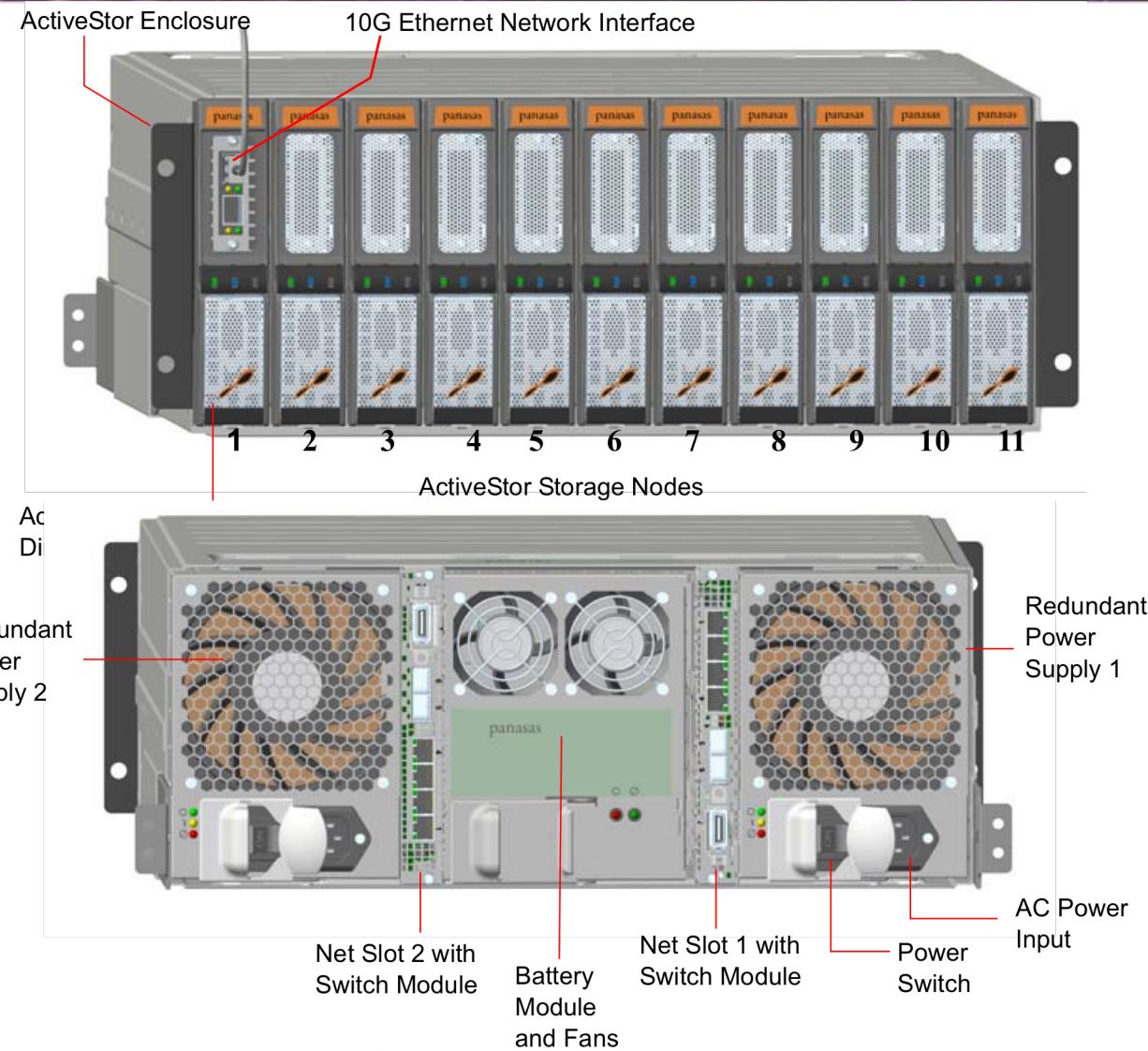
Power-failure and Panasas

Power-failure and Panasas: Problem



- Problem: High-Performance Storage requires caching
 - HDD caches and motherboard DRAM are both volatile
 - Loss of power => loss of data
 - First-generation Panasas hardware was designed about 17 years ago
 - Few good non-volatile memory options
 - PCI card w/ battery-backed DRAM
 - ???
 - PCI is slow compared to main system DRAM
 - PCI cards were too big for our custom blade form-factor
 - When the battery is drained, data is lost
 - Battery can't be hot-swapped
 - Does not help with drive cache contents

Power-failure and Panasas: Custom Hardware



Power-failure and Panasas: Solution



- **Solution: Shared enclosure-wide battery module**
 - When both power-supplies lose AC input, switch to DC input from the battery
 - De-assert AC_GOOD signal
 - Blades see un-interrupted DC input
 - Entire blade – CPU, system DRAM, HDDs, drive caches – stays online
 - Software detects that both AC_GOODs are de-asserted => running from battery
 - Expedited power-fail shutdown path
 - Stream out of DRAM to reserved space; defer proper placement until reboot after AC input is restored
 - One battery module per enclosure
 - Hot-swappable
 - Relatively cheap sealed lead-acid cells (like a motorcycle battery)

- Not without problems
 - Custom battery module requires custom monitoring and management
 - System has to perform safe, automated testing
 - Lead-acid cells are big and heavy
 - Expensive to ship
 - Hard to replace at top-of-rack
 - More recently: shipping restrictions
 - Lead is toxic
 - Batteries can be dangerous => more padding, more paperwork

Power-failure and Panasas: 2012 – COTS?



- Fast-forward: 2012
- Move protocol gateways and metadata management to off-the-shelf hardware
 - Faster, cheaper adoption of new CPUs and networking
- Still need a power-fail solution
 - Memory vendor pitched new idea called "NV-DIMM"
 - Connected to normal DDR3 memory bus
 - Looks like DRAM, because it's made from DRAM ...
 - Normal DDR3 speeds, normal DRAM latencies
 - Byte addressable
 - ... until power fails ...

Power-failure and Panasas: 2012 – NVDIMM?



- Theory of operation
 - Power-supply raises signal when AC is lost
 - NV-DIMM detaches from DDR3 bus, self-refreshes using supercap
 - Automatically copies DRAM contents into NAND
- Pros
 - Everything we like about main memory
- Cons
 - Only supported by specific motherboards and power-supplies
 - Required non-standard motherboard firmware
 - Supercaps are large, and their lifespan is unproven

Power-failure and Panasas: 2012 – No NVDIMM



- Too much science-experiment, not enough product
 - Plenty of other work to do while the technology solidified
 - COTS experiment was very educational, but never shipped

Power-failure and Panasas: 2015 – COTS Product



- 2015: Building on previous experiments, create a shippable COTS product
- DDR4 DRAM+NAND NVDIMMs are real products
 - Supercap packs are smaller
 - Supercap lifetimes better characterized
 - Standardized by JEDEC and ACPI
 - ... in mid-2015
 - Widely supported by motherboards, power-supplies, and firmware
 - ... by 2017

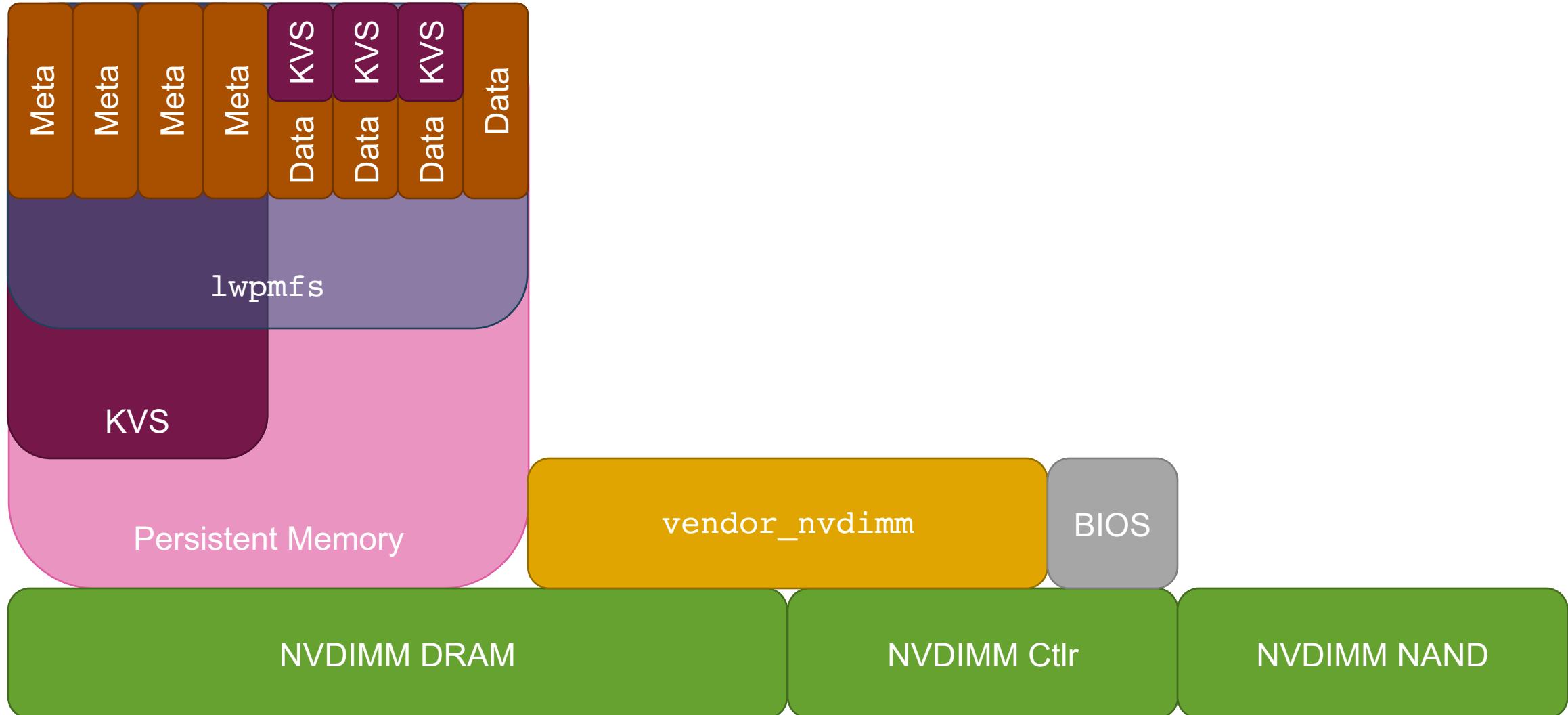
Power-failure and Panasas: 2015 – COTS Product



- In early 2015, NVDIMM products were real, but not-quite-standard
 - Board firmware supported proprietary APIs from two NVDIMM vendors
 - Several board and NVDIMM firmware upgrades required for stable functionality
- Long lead times for logistics, development, testing, etc
 - Couldn't just wait for JEDEC / ACPI to be adopted

NVDIMM Support

NVDIMM Support – Full Stack



NVDIMM Support – Low-level (Proprietary)



- Motherboard communicates with NVDIMM over SMBus
 - Just like SPD EEPROM or JEDEC temperature sensor
- Driver for SMBus controller (part of memory controller (part of CPU))
 - Eventually committed as `imcsmb(4)` ({Sandy,Ivy}bridge-Xeon and {Broad,Has}well-Xeon)

- Driver for NVDIMM vendor proprietary interface
 - Ported from Linux driver
 - Sources had a compatible license, but were only accessible under NDA
 - Only useful for this particular pre-standard NVDIMM, so no point in upstreaming anyway
 - Parts were extracted and committed as `jedec_dimm(4)`
- NVDIMM DRAM exposed as byte-addressable `/dev/nvdimm` device node
 - `read(2)` and `write(2)` do `uiomove(9)`
 - `write(2)` includes instructions to flush the appropriate cachelines
 - `CLFLUSHOPT` isn't available on {Broad,Has}well, so simulated with `MOVQ` / `MOVNTI` / `SFENCE`
- `ioctl(4)`s to arm and disarm auto-save, check supercap health, etc.

- Abstract interface
 - Byte-addressable
 - Memory-mapped
 - Contiguous virtual address space
 - Explicit write-back caching semantics
 - Explicit fencing around transactions
 - Persistence only guaranteed after synchronous commit
 - Writes between commits may be re-ordered

NVDIMM Support -- Persistent Memory

- **Three implementations**

- Direct – PM backed by /dev/nvdimm device – base class
 - Commit / fence: `s_fence` (if needed)
 - Writeback: Same `MOVQ` / `MOVNTI` / `SFENCE` tricks as NVDIMM driver
- `mmap` – PM semantics on mmap-ed user-space file
 - PM backed by file in `1wpmfs`
 - Commit: `fsync(2)`
 - Writeback: `msync(2)`
- `kmmmap` – PM semantics on vnode object
 - PM backed by file in `1wpmfs`, from inside kernel
 - Commit: `VOP_FSYNC(9)`
 - Writeback: `vm_map_sync(9)`



Buffer-cache bypass mode:
Explicit writeback not required

NVDIMM Support -- Persistent Memory



- Buffer-cache bypass
 - VOP_GETPAGES() adjusts mappings, does not do memcpy()
 - VOP_PUTPAGES() not called

- Key-Value Store
 - 64-bit keys, Arbitrary-sized values
 - Multiple K/V pairs can be updated atomically (transactional)
- Values are copy-on-write
 - Refcounted sparse tree of block pointers and checksums
 - Never modifies a referenced block in-place
- Create-time parameters
 - Maximum number of K/V pairs
 - Maximum number of modifications in a single transaction
 - Number and size (at least 16 bytes, power-of-two) of blocks for value data

- K/V pair metadata (record)
 - Key, value
 - Transaction ID
 - Transaction size
 - Flags
 - Checksum
- KVS reserves spaces for N (max pairs) + M (max in-flight txns) records

- Update single record
 - Find free record
 - Increment transaction ID
 - Update record data
 - Update record checksum
 - Power-fail during modification?
 - Record being modified has invalid checksum
 - Previous version of record is still present; use it

- **Update multiple records in a transaction**
 - Just like updating a single record, but all records being updated will share a transaction ID
 - Power-fail during modification?
 - Transaction has fewer valid records than specified
 - Throw out all records modified by the transaction
- **Reads & concurrent updates**
 - Increment refcount on value's radix tree
 - If value is modified w/ non-zero refcount, leaf blocks (and indirect blocks) are copied
 - Copies are modified and marked dirty
 - When modification is committed, only dirty blocks need to be committed

- LWPMFS is built on top of both KVS and PM
 - `newfs_lwpmfs` splits underlying device into a KVS for metadata, and blocks for data
 - Elaborate code to estimate proper KVS sizing – some assumptions based on our use-case
 - Remainder of device used for LWPMFS data blocks
- LWPMFS metadata stored in KVS
 - Key=inode
 - Value
 - Generic FS metadata (size, times, parent, permissions, etc)
 - List of FS data blocks
- LWPMFS data stored outside of KVS

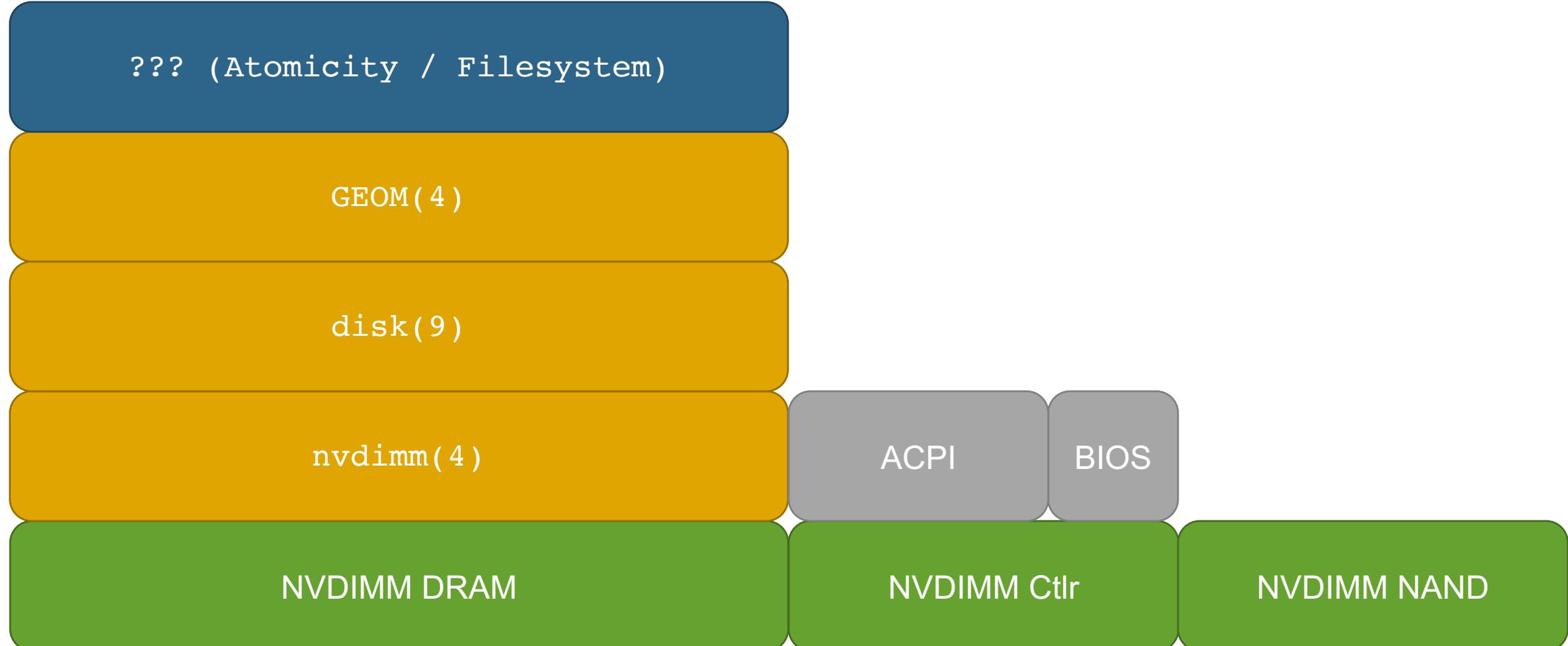
Standard NVDIMM Support

Standard NVDIMM Support – In-tree



- Konstantin Belousov (Kostik, aka kib@freebsd.org) has been working on this
- Two related commits earlier this week
 - r339386 – pmap_large_map() KPI
 - map very large contiguous physical regions
 - With the size of NVDIMM that we used, and our use-case, this was not an issue that we needed to address
 - Optimized cache flushing
 - CLFLUSHOPT when possible, falling back to CLFLUSH
 - Our MOVQ / MOVNTI / SFENCE trick might be better than CLFLUSH? Need to discuss
 - r339391 – JEDEC / ACPI NVDIMM device driver
 - Uses ACPI “NFIT” table to discover and enumerate NVDIMMs
 - Hooks into disk(9), geom(4)

Standard NVDIMM Support – Full Stack



Standard NVDIMM Support – Still needed



- Something to enforce atomicity and prevent partial updates
- DAX (Direct Access) filesystem (buffer-cache bypass)
- Some of these things are available through PMDK
 - Intel Persistent Memory Development Kit
 - Why didn't Panasas use that in the first place?
 - At the time, it didn't cover all our needs, and might not have supported FreeBSD at all?
 - Now PMDK does support FreeBSD (but not in ports yet...)
- Panasas PM / KVS / LWPMFS code is probably not directly useful upstream
 - But we're happy to share it, if it will help fill the gaps!

Q & A

Links



- <https://people.freebsd.org/~rpokala/2018-10-DeveloperSummitLWPMFS.pdf>