UC Berkeley – Computer Science
CS188: Introduction to Artificial Intelligence
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Midterm I, Fall 2016

This test has **8** questions worth a total of **100** points, to be completed in 110 minutes. The exam is closed book, except that you are allowed to use a single two-sided hand written cheat sheet. No calculators or other electronic devices are permitted. Give your answers and show your work in the space provided.

Write the statement out below in the blank provided and sign. You may do this before the exam begins. Any plagiarism, no matter how minor, will result in an F.

"I have neither given nor received any assistance in the taking of this exam."

_____

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Signature: _____

Name: _____ **Your EdX Login:** _____
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Exam Room: _____ **Name of person to right:** _____
Primary TA: _____

- ◯ indicates that only one circle should be filled in.
- ▫ indicates that more than one box may be filled in.
- Be sure to fill in the ◯ and ▫ boxes completely and erase fully if you change your answer.
- There may be partial credit for incomplete answers. Write as much of the solution as you can, but bear in mind that we may deduct points if your answers are much more complicated than necessary.
- There are a lot of problems on this exam. Work through the ones with which you are comfortable first. **Do not get overly captivated by interesting problems or complex corner cases you're not sure about.** Fun can come after the exam: There will be time after the exam is try them again.
- Not all information provided in a problem may be useful.
- Write the last four digits of your SID on each page in case pages get shuffled during scanning.

Problem	1	2	3	4	5	6	7	8
Points	11.5	9.5	14.5	12	8.5	16	19	9

Optional: Mark along the line to show your feelings on the spectrum between :(and ☺.

Before exam: [:(_____ ☺]

After exam: [:(_____ ☺]

1. Munching (11.5 pts)

i) (7.5 pts) Suppose Pac-Man is given a map of a maze, which shows the following: it is an ***M*-by-*N*** grid with a single entrance, and there are **W** walls in various locations in the maze that block movement between empty spaces. Pac-Man also knows that there are **K** pellets in the maze, but the pellets are not on the map, i.e. **Pac-Man doesn't know their locations in advance**. Pac-Man wants to model this problem as a search problem. Pac-Man's goal is to plan a path, before entering the maze, that lets him eat the hidden pellets as quickly as possible. **Pac-Man can only tell he has eaten a pellet when he moves on top of one.**

Which of the following features should be included in a minimal correct state space representation? For each feature, if it should be included in the search state, calculate the size of this feature (the number of values it can take on). If it should **not** be included, **very briefly** justify why not.

Feature	Include?	Size (if included) or Justification (if not included)
Current location	◉ Yes / ○ No	mxn
Dimensions of maze	○ Yes / ◉ No	already in the domain of the position
Locations of walls	○ Yes / ◉ No	we don't need for every state to store this information
Places explored	◉ Yes / ○ No	2^mxn
Number of pellets found	○ Yes / ◉ No	we don't have this information
Time elapsed	○ Yes / ◉ No	we will codify this in the algorithm

ii) (1 pt) Which approach is best suited for solving the K hidden pellets problem? Completely fill in the circle next to one of the following:

◉ **State Space Search** ○ **CSP** ○ **Minimax** ○ **Expectimax** ○ **MDP** ○ **RL**

iii) (2 pts) Now assume that there are K_R red pellets, K_G green pellets, and K_B blue pellets, that Pac-Man **knows the locations and colors of all the pellets**, and that Pac-Man wants **the shortest path that lets him eat <u>at</u> <u>least</u> one pellet of each color,** i.e. one red, one green, and one blue. What is the total size of the search state space, assuming a minimal correct state space representation?

mxn x 2^3

iv) (1 pt) Which approach is best suited for solving the one-pellet-of-each-color problem? Completely fill in the circle next to one of the following:

◉ **State Space Search** ○ **CSP** ○ **Minimax** ○ **Expectimax** ○ **MDP** ○ **RL**

2. Uninformed Search (9.5 pts)

Part A: Depth Limited Search (DLS) is a variant of depth first search where nodes at a given depth L are treated as if they have no successors. We call L the **depth limit.** Assume all edges have uniform weight.

i) (1 pt) Briefly describe a major advantage of depth limited search over a simple depth first tree search that explores the entire search tree.

ii) (1 pt) Briefly describe a major disadvantage of depth limited search compared to a simple depth first tree search that explores the entire search tree.

iii) (2 pts) Is DLS optimal? ○ Yes / ◉ No Is DLS complete? ○ Yes / ◉ No

Part B: IDDFS. The iterative deepening depth first search (IDDFS) algorithm repeatedly applies depth limited search with increasing depth limits L: **First it performs DLS with L=1, then L=2, and so forth, until the goal is found.** IDDFS is neither BFS nor DFS, but rather an entirely new search algorithm. Suppose we are running IDDFS on a tree with branching factor B, and **the tree has a goal node located at depth D**. Assume all edges have uniform weight.

i) (2 pts) Is IDDFS optimal? ◉ Yes / ○ No Is IDDFS complete? ◉ Yes / ○ No

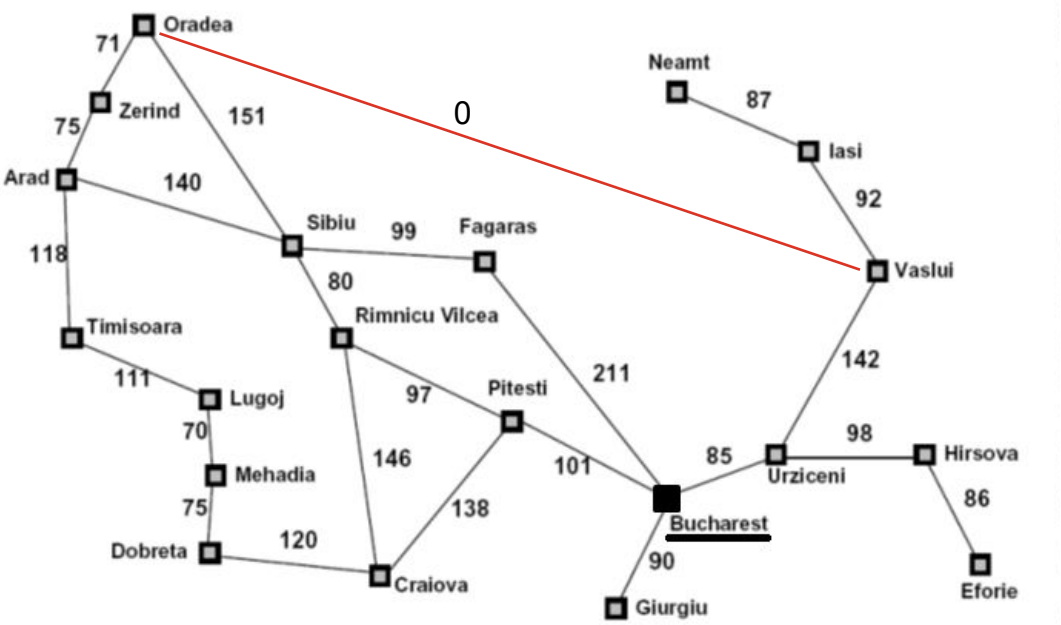
ii) (1 pt) In the worst case asymptotically, which algorithm will **use more memory** looking for the goal?

○ IDDFS / ◉ BFS

iii) (2.5 pts) In project 2, you used DLS to implement the minimax algorithm to play Pac-Man. Suppose we are instead writing a minimax AI to play in a speed chess tournament. Unlike the Pac-Man project, there are many pieces that you or your opponent might move, and they may move in more complicated ways than Pac-Man. Furthermore, you are given only 60 seconds to select a move, otherwise you are given a random move. **Explain why** we'd prefer to implement minimax using IDDFS instead of DLS for this problem, and **explain how** would we need to modify IDDFS (as described above) so that it works for this problem.

3. Aliens in Romania (14.5 pts)

Here's our old friend, a road map of Romania. The **edges** represent **travel time along a road**. Our goal in this problem is to find the shortest path from some arbitrary city to Bucharest (underlined in black). You may assume that the heuristic at Bucharest is always 0. In this question, **L(X, Y) represents the travel time if we went in a straight line from city X to city Y**, ignoring roads, teleporters, wind, birds, etc. You should make no other assumptions about the values of L(X, Y).



Part A: Warm up.

i) (6 pts) For each of the heuristics below, completely fill in the circle next to "Yes" or "No" to indicate whether the heuristic is admissible and/or consistent for **A* search**. You should fill in 6 answers (one has been provided for you). **Reminder, h(Bucharest) is always zero.**

Heuristic h(C) = **[except C = Bucharest]**	**Admissible**	**Consistent**
0	● Yes / ○ No	◉ Yes / ○ No
90	○ Yes / ◉ No	○ Yes / ◉ No
L(C, Bucharest)	◉ Yes / ○ No	◉ Yes / ○ No
min(L(C, Bucharest), 146)	◉ Yes / ○ No	

ii) (2.5 pts) For what values of x is h(C) = max(L(C, Bucharest), x) guaranteed to be admissible? Your answer should be in the form of an expression on x (e.g. "x = 0 OR x > 102").

x = 0

Part B: Teleporters

Aliens build a teleporter from Oradea to Vaslui. This means that we can now get from Oradea (top left) to Vaslui (near the top right) with zero cost.

i) (1 pt) Update the drawing on the **previous page** so that your graph now reflects the existence of the teleporter.

ii) (5 pts) Completely fill in the circle next to "Yes" or "No" to indicate which of the following are guaranteed admissible for A* search, given the existence of the alien teleporter. **Reminder, h(Bucharest) is always zero, and L(A, B) is the travel time between A and B assuming a straight line between A and B ("as a bird flies").**

Heuristic h(C) = [except C = Bucharest]	**Admissible**
77	◉ Yes / ○ No
L(C, Bucharest)	○ Yes / ◉ No
max(L(C, Bucharest), 80)	○ Yes / ◉ No
max(L(C, Bucharest), L(C, Oradea))	○ Yes / ◉ No
min(L(C, Bucharest), L(C, Oradea))	◉ Yes / ○ No

Survey Completion

Enter the secret code for survey completion: _____

If you took the survey but don't have the secret code, explain:

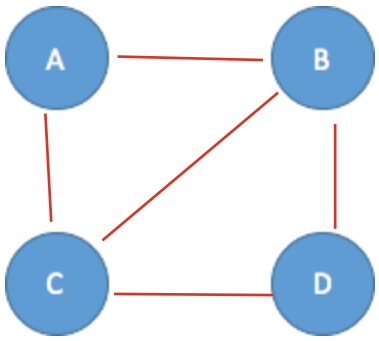
_____

4. CSPs (12 pts)

In a combined 3rd-5th grade class, students can be 8, 9, 10, or 11 years old. We are trying to solve for the ages of Ann, Bob, Claire, and Doug. Consider the following constraints:

- No student is older in years than **C**laire (but may be the same age).
- **B**ob is two years older than **A**nn.
- **B**ob is younger in years than **D**oug.

The figure below shows these four students schematically ("A" for **A**nn, "B" for **B**ob, etc).



i) (2 pts) In the figure, draw the arcs that represent the binary constraints described in the problem.

ii) (1 pt) Suppose we're using the AC-3 algorithm for arc consistency. How many total arcs will be enqueued when the algorithm begins execution? Completely fill in the circle next to one of the following:

○ **0** / ○ **5** / ○ **6** / ○ **8** / ◉ **10** / ○ **12**

iii) Assuming all ages {8, 9, 10, or 11} are possible for each student before running arc consistency, manually run arc consistency on **only** the arc *from* A *to* B.

A. (2 pts) What **values on A** remain viable after this operation? Fill in all that apply.
☑ 8 / ☑ 9 / ☐ 10 / ☐ 11

B. (2 pts) What **values on B** remain viable after this operation? Fill in all that apply.
☑ 8 / ☑ 9 / ☑ 10 / ☑ 11

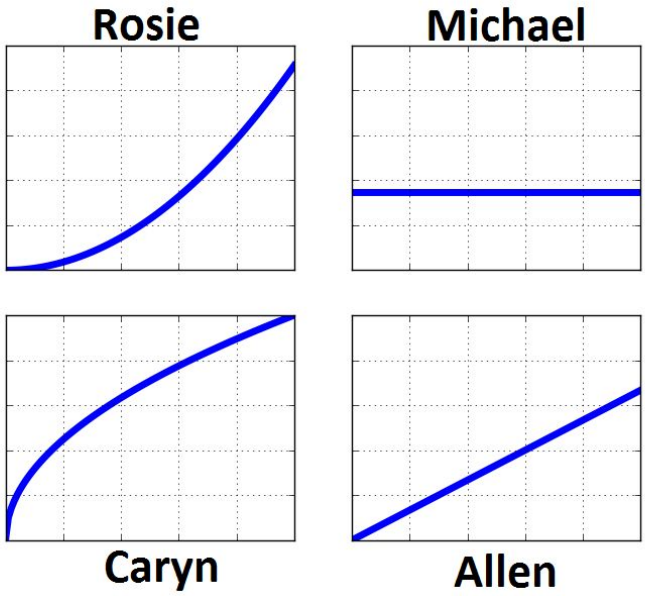
C. (1 pt) Assuming there were no arcs left in the list of arcs to be processed, which arc(s) would be added to the queue for processing after this operation?
Everyone but A->B

iv) (4 pts) Suppose we enforce arc consistency on all arcs. What ages remain in each person's domain?

Ann: 8 **B**ob: 10 **C**laire: 11 **D**oug: 11

5. Utilities (8.5 pts)

Part A: Interpreting Utility. (3.5 pts) Suppose Rosie, Michael, Caryn, and Allen are four people with the utility functions for marshmallows as described below (e.g. Michael doesn't care how many marshmallows he gets). The x-axis is the number of marshmallows, and the y-axis is the utility. Rank our four TAs from **least risk-seeking (at the bottom)** to most risk-seeking (at the top) using the four lines provided. **If two people have the same degree of risk-seeking, put them on the same line**. You may not need all four blanks (e.g. if two people are on the same line).



Rank in terms of risk-seeking behavior:

Most: _____

_____

_____

Least: _____

Part B: Of Two Minds. (5 pts) Some utility functions are sometimes risk-seeking, and sometimes risk-averse. Show that $U(x) = x^3$ is one such function, by providing two lotteries: one for which it is risk-seeking, and one for which it is risk-averse. Provide justification by filling out the table below. Averse is just the opposite of "seeking". Specifically, it means "having a strong dislike for something," in this case, risk.

	Risk-seeking:	Risk-averse:
Lottery		
Utility of lottery		
Utility of expected value of lottery		

6. Games / Mini-Max: The Potato Game (16 pts)

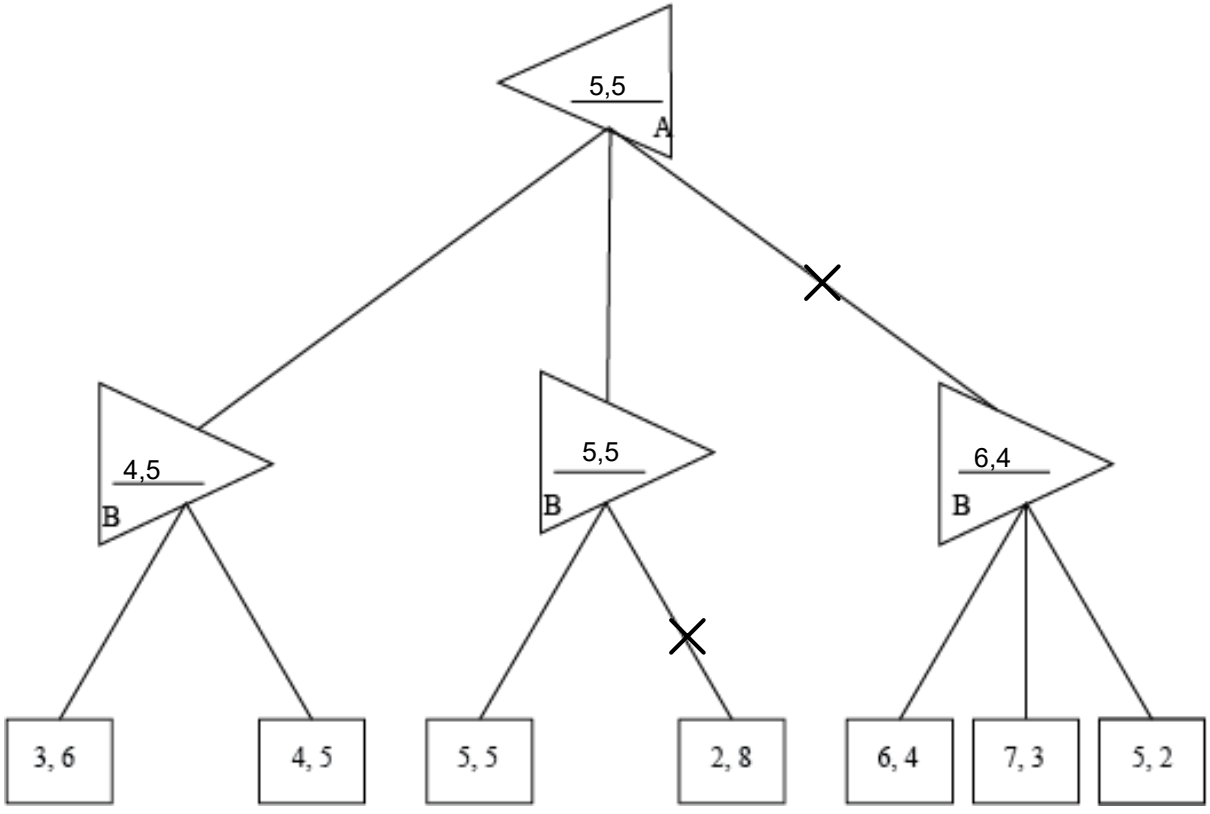
Aldo and Becca are playing the game below, where the left value of each node is the number of potatoes that Aldo gets, and the right value of each node is the number of potatoes that Becca gets (i.e. Aldo: left, Becca: right). Unlike prior scenarios where having more potatoes results in more utility, Becca and Aldo will have a more complex view of what makes a "good" distribution of potatoes.

Becca will use the following thought process to decide which move to take:

- Among all choices where she gets at least as many potatoes as Aldo, she'll pick the one that maximizes Aldo's number of potatoes (very nice!).
- If there are no choices where she gets at least as many potatoes as Aldo, she'll simply maximize her own potato count (ignoring Aldo's value).

Aldo will do the same thing, but substituting "he" for "she", "her" for "his", and "Becca" for "Aldo". The rules above effectively tell us how Becca (and Aldo) rank the utility of any set of choices.

i) (4 pts) Fill in the blanks below with the choice that each player would make at each stage. Assume that both Aldo and Becca are trying to maximize their own utility as described above.

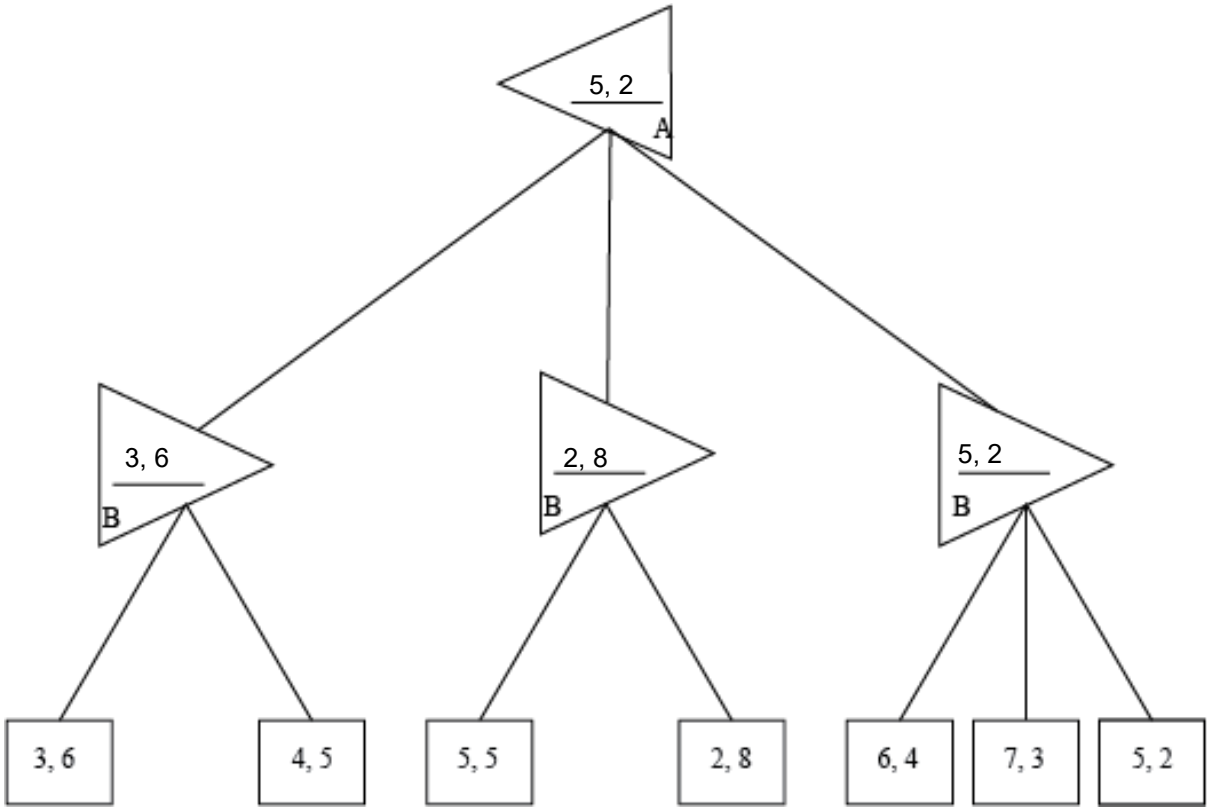


ii) (3 pts) Assume that Aldo and Becca know that **the sum at any leaf node is no more than 10**. Cross out the edges to any nodes that can be pruned (if none can be pruned then write "no pruning possible" below the tree).

iii) (3 pts) Describe a general rule for pruning edges, assuming that the sum at any leaf node is no more than 10.

Prune when 5,5 is found

iv) (4 pts) Suppose that **Becca attempts to minimize Aldo's utility instead of maximizing her own utility, and that Aldo knows this, and tries to maximize his own utility -- taking into account Becca's strategy**. Repeat part ii. There is no need to cross out edges that are pruned.



v) (2 pts) Suppose that **Becca chooses uniformly randomly and Aldo knows this and tries to maximize his own utility -- taking into account Becca's strategy**. Which move will Aldo make to maximize his expected utility, assuming he treats Becca as a chance node: left, middle, or right? If there is not enough information, pick "Not Enough Information".

○ Left ○ Middle ○ Right ◉ Not Enough Information

7. The Spice Must Flow (19 pts)

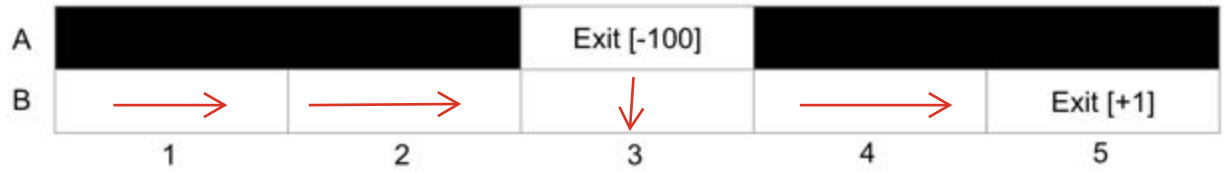
i) (3 pts) Suppose we have a robot that has 5 possible actions: {↑, ↓, ←, →, Exit}. If the robot picks a direction, it has an **80%** chance of going in that direction, a **10%** chance of going 90 degrees counterclockwise of the direction it picked (e.g. picked →, but actually goes ↑), and a **10%** chance of going 90 degrees clockwise of the direction it picks.

If the robot hits a wall, it does not move this timestep, but time still elapses. For example, if the robot picks ↓ in state B1 and is 'successful' (80% chance), it will hit the wall and not move this timestep. As another example, if it picks → in state B1, but gets the 10% chance of going down, it will hit the wall and not move this timestep

In states B1, B2, B3, and B4, the set of possible actions is {↑, ↓, ←, →}. In states A3 and B5, the only possible action is {Exit}. Exit is always successful. The blackened areas (A1, A2, A4, and A5) represent walls.

All transition rewards are 0, except R(A3, Exit) = -100, R(B5, Exit) = 1.

In the boxes below, draw an optimal policy for states B1, B2, B3, and B4, if there is no discounting, i.e. $\gamma = 1$. Use the symbols {↑, ↓, ←, →}.



ii) (2 pts) Give the expected utility for states B1, B2, B3, and B4 under the policy you gave above.

V*(B1): 1 V*(B2): 1 V*(B3): 1 V*(B4): 1

iii) (2.5 pts)What is V*(B4), the expected utility of state B4, if we have a discount $0 < \gamma < 1$ and use the optimal policy? Give your answer in terms of γ.

tau = 1/gamma

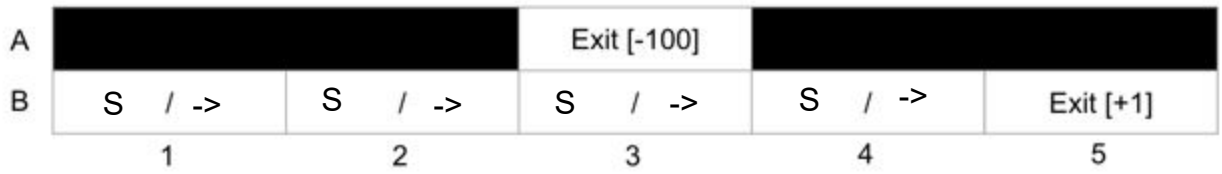
V*(B4): 0.8/(tau - 0.2)

Part B: Spice. Now suppose we add a sixth special action **SPICE**. This action has a special property that it always takes the robot **back to its current state**, with a reward of 0. This action might seem useless, except for one thing - the robot will always get its top preference on its next action (instead of being subject to noise).

For example, suppose that the robot is in state **s1** and picks action **a** which has a chance **p2** of moving to **s2**, and a chance **p3** of moving to **s3**. If the robot takes the **SPICE** action at timestep t, and **a** at timestep $t + 1$, the robot will not end up randomly in **s2** or **s3**, but will instead go to its **preference** of **s2** or **s3**. It will still collect R(**s1**, **a**, **s'**), which is unchanged. Preferring **s2** or **s3** does not take an additional time step, e.g. if the robot is in B1 at timestep 0, chooses SPICE, then chooses ↓ but prefers the outcome B2, the robot arrives in B2 at timestep 2.

The powers granted by **SPICE** last one timestep. **SPICE** may be used any number of times.

i) (2 pts) Give an optimal policy for B1, B2, B3, and B4, assuming no discounting, i.e. $\gamma = 1$. **On the left of each slash,** write the optimal action if the previous action was not **SPICE**, and on the **right side of each slash**, write the optimal action if the previous action was **SPICE, assuming the robot always "prefers" the outcome that points in the same direction as its action (e.g. if it picks →, it prefers going right)**. In total, you should write 8 actions. Use the symbols {↑, ↓, ←, →, S}, where "S" represents indicate the spice action.



ii) (2.5 pts) Assuming we have a discount $0 < \gamma < 1$, what is V*(B4) if the previous action was not **SPICE**? Give your answer in terms of γ. It is OK to leave your answer in terms of a max operation.

V*(B4): max(gamma^2, 0.8gamma/(1 -0.2gamma)

OFFICIAL RELAXATION SPACE. Draw or write anything here:

Part C: Bellman Equations. (You can do this problem even if you didn't do B, but it'll be harder)

In class, we derived the Bellman Equations for an MDP given below.

$$Q^*(s,a) = \sum_{s' \in S_{(s,a)}} T(s,a,s')[R(s,a,s') + \gamma V^*(s')]$$

$$V^*(s) = \max_{a \in A_s} Q^*(s,a)$$

You'll notice these look slightly different than the version on your cheat sheet, but these equations represent exactly the same idea. The only difference is that to improve the clarity of these equations, we've specified the sets over which we are summing and maximizing. Specifically, $S_{(s, a)}$ is the set of states that might result from the action (s, a), and A_s is the set of actions that one can take in state s. **For this problem, assume that A_s does not include the special SPICE action.**

i) (3.5 pts) Derive Q*(s, SPICE). Assume that the previous action was not **SPICE**. Your answer should have two max operators in it. Hint: Consider drawing the diagram we used to derive the Bellman Equation.

ii) (3.5 pts) Derive V*(s), the value of a state in this MDP. You should account for the the fact that the **SPICE** action is available to the robot. Assume that the previous action was not **SPICE**.

8. The Reinforcement of Brotherly Love (9 pts)

Consider two brother ghosts playing a game in an M-by-N grid world, Ghost Y is the Younger brother of Ghost O. Each ghost takes 100 steps, trying to pet cats which randomly appear in the maze. The reward for the game is equal to the number of cats petted. Each ghost has a policy for how to run through the maze. Ghost O being the older brother has a policy that has a **higher expected cumulative reward** (i.e. $V_Y(s) \leq V_O(s) \ \forall s \in S$). However, Ghost O is **not necessarily optimal** (i.e. $V_O(s) \leq V^*(s) \ \forall s \in S$). Ghost Y wants to learn from Ghost O instead of starting from scratch.

Part A: Q-Learning

We will now explore how Ghost Y can achieve this, while performing **Q-Learning**. Ghost Y starts out with a Q function **initialized randomly** and wants to catch up to Ghost O's performance. Denote the Q functions for each Ghost as Q_Y , Q_O . In order for Ghost Y to converge to the optimal policy, he must properly balance **exploring** and **exploiting**. Consider the following exploration strategy: At state s with probability 1- ε choose $max_a \, Q_Y(s,a)$ otherwise choose $max_a \, Q_O(s,a)$.

(2 pts) Is this guaranteed to converge to the optimal value function V*? Briefly justify.

◯ Yes / ◯ No _____

Part B: Model-Based RL

Ghost Y decides model free learning is scary, and decides to try to learn a model instead. Ghost Y will watch Ghost O try to catch Pac-Man and estimate the **transition** probabilities and **rewards** from Ghost O's demonstrations (i.e. learn a model for T(s,a,s') and R(s,a,s')).

i) (2 pts) For Ghost Y to learn the correct and exact values of R(s,a,s') for all state/action/next_state transitions, what must be true about the episodes he watches Ghost O perform?

ii) (2 pts) For Ghost Y to learn the correct and exact values of T(s,a,s') for all state/action/next_state transitions, what must be true about the episodes he watches Ghost O perform?

iii) (2 pts) Assuming T(s,a,s') and R(s,a,s') are learned **exactly**, what algorithm/s below could be helpful for Ghost Y to decide which actions to take from each state, based **only** on information in his learned model?

▢ **Value Iteration** ▢ **Policy Iteration** ▢ **Q-Learning** ▢ **TD-Learning**

iv) (1 pt) Let π be any of the policies learned in part iii. Does following π result in maximized expected utility for every state, i.e. is $V^{\pi}(s) = V^*(s)$ for all states? [Consider all policies]

◯ **Yes** ◯ **No** ◯ **Not enough information**

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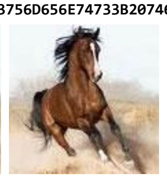
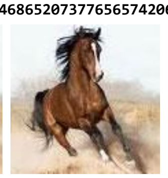
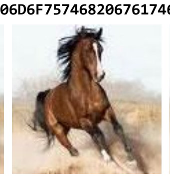
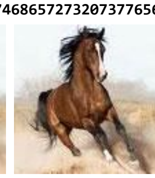
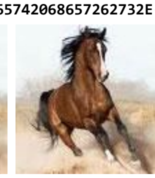

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