Abstract

Categories and Subject Descriptors CR-number [*subcategory*]: third-level

Keywords gradual types, Clojure

Tobin-Hochstadt and Felleisen [1, 2]

References

[1] S. Tobin-Hochstadt and M. Felleisen. The design and implementation of typed scheme. *SIGPLAN Not.*, 43 (1):395–406, Jan. 2008. ISSN 0362-1340. URL `http://doi.acm.org/10.1145/1328897.1328486`.

[2] S. Tobin-Hochstadt and M. Felleisen. Logical types for untyped languages. *SIGPLAN Not.*, 45(9):117–128, Sept. 2010. ISSN 0362-1340. URL `http://doi.acm.org/10.1145/1932681.1863561`.

$$
\begin{array}{lll}
v & ::= \#\mathsf{f} \mid \#\mathsf{t} & \text{Values} \\
e & ::= x \mid v \mid (e\ e) \mid \lambda x^{\tau}.e \mid (\mathbf{if}\ e\ e\ e) \mid (\mathbf{let}\ (x\ e)\ e) \mid (\mathbf{begin}\ e\ e) \mid (\mathit{cons}\ e\ e) & \text{Expressions} \\
 & & \\
\mathcal{R} & ::= \{\omega \mid \psi\} & \text{Refinement Types} \\
 & & \\
\sigma, \tau & ::= \mathbf{\#t} \mid \mathbf{\#f} \mid \mathbf{N} \mid x{:}\mathcal{R} \to \mathcal{R} \mid \langle \mathcal{R}, \mathcal{R}' \rangle & \text{Types} \\
 & & \\
\kappa & ::= \omega \mid x & \text{Object Variables} \\
 & & \\
\psi & ::= \tau_{\pi(\kappa)} \mid \overline{\tau}_{\pi(\kappa)} \mid \psi \supset \psi \mid \psi \wedge \psi \mid \psi \vee \psi \mid \mathbb{tt} \mid \mathbb{ff} \mid \pi(\kappa) \equiv \pi(\kappa) & \text{Propositions} \\
o & ::= \pi(\kappa) \mid \emptyset & \text{Objects} \\
\pi & ::= \overrightarrow{pe} & \text{Paths} \\
pe & ::= \mathbf{car} \mid \mathbf{cdr} & \text{Path Elements} \\
 & & \\
\Gamma & ::= \overrightarrow{\psi} & \text{Proposition Environment}
\end{array}
$$

Figure 1. Syntax of Terms, Types, Propositions, and Objects

$$
\begin{array}{lcl}
\top & = & \{\omega \mid \mathbb{tt}\} \\
\bot & = & \{\omega \mid \mathbb{ff}\} \\
\mathbf{\#t} & = & \{\omega \mid \mathbf{\#t}_{\omega}\} \\
\mathbf{\#f} & = & \{\omega \mid \mathbf{\#f}_{\omega}\} \\
(\bigcup\ \overrightarrow{\tau}) & = & \{\omega \mid \bigvee \overrightarrow{\tau_{\omega}}\} \\
\mathbf{N} & = & \{\omega \mid \mathbf{N}_{\omega}\} \\
\mathbf{B} & = & (\bigcup\ \mathbf{\#t}\ \mathbf{\#f}) \\
\langle \mathcal{R}, \mathcal{R}' \rangle & = & \{\omega \mid \langle \mathcal{R}, \mathcal{R}' \rangle_{\omega}\} \\
x{:}\tau \to \tau & = & \{\omega \mid (x{:}\mathcal{R} \to \mathcal{R})_{\omega}\}
\end{array}
$$

Figure 2. Refinement Type Abbreviations

$$
\textsc{T-Var}\quad \frac{\omega \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash \tau_x}{\Gamma \vdash x : \{\omega \mid \tau_{\omega} \wedge \omega \equiv x\}}
\qquad
\textsc{T-True}\quad \Gamma \vdash \#\mathsf{t} : \mathbf{\#t}
\qquad
\textsc{T-False}\quad \Gamma \vdash \#\mathsf{f} : \mathbf{\#f}
$$

$$
\textsc{T-If}\quad \frac{\omega_1 \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash e_1 : \{\omega_1 \mid \psi_1\} \qquad \Gamma, \overline{\mathbf{\#f}}_{\omega_1}, \psi_1 \vdash e_2 : \mathcal{R} \qquad \Gamma, \mathbf{\#f}_{\omega_1}, \psi_1 \vdash e_3 : \mathcal{R}}{\Gamma \vdash (\mathbf{if}\ e_1\ e_2\ e_3) : \mathcal{R}}
$$

$$
\textsc{T-App}\quad \frac{\omega \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash e : x{:}\{\omega \mid \psi\} \to \mathcal{R} \qquad \Gamma \vdash e' : \{\omega \mid \psi\}}{\Gamma \vdash (e\ e') : \mathcal{R}[\omega/x]}
$$

$$
\textsc{T-Subsume}\quad \frac{\Gamma \vdash e : \mathcal{R} \qquad \Gamma \vdash \mathcal{R} <: \mathcal{R}'}{\Gamma \vdash e : \mathcal{R}'}
$$

$$
\textsc{T-Let}\quad \frac{x, \omega_1 \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash e_1 : \{\omega_1 \mid \psi_1\} \qquad \Gamma, \psi_1[\omega_1/x] \vdash e : \mathcal{R}}{\Gamma \vdash (\mathbf{let}\ (x\ e_1)\ e) : \mathcal{R}}
$$

$$
\textsc{T-Cons}\quad \frac{\Gamma \vdash e_1 : \mathcal{R}_1 \qquad \Gamma \vdash e_2 : \mathcal{R}_2}{\Gamma \vdash (\mathit{cons}\ e_1\ e_2) : \langle \mathcal{R}_1, \mathcal{R}_2 \rangle}
$$

$$
\textsc{T-Car}\quad \frac{\omega, \omega_1 \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash e : \{\omega_1 \mid \psi_1\} \qquad \psi_1 \vdash \langle \top, \top \rangle_{\omega_1}}{\Gamma \vdash (\mathit{car}\ e) : \{\omega \mid \psi_1 \wedge \mathbf{car}(\omega_1) \equiv \omega\}}
$$

$$
\textsc{T-Cdr}\quad \frac{\omega, \omega_1 \notin \mathrm{fv}(\Gamma) \qquad \Gamma \vdash e : \{\omega_1 \mid \psi_1\} \qquad \psi_1 \vdash \langle \top, \top \rangle_{\omega_1}}{\Gamma \vdash (\mathit{cdr}\ e) : \{\omega \mid \psi_1 \wedge \mathbf{cdr}(\omega_1) \equiv \omega\}}
$$

$$
\textsc{T-Abs}\quad \frac{x \notin \mathrm{fv}(\Gamma) \qquad \mathcal{R} = \{\omega \mid \psi\} \qquad \Gamma, \psi[x/\omega] \vdash e : \mathcal{R}'}{\Gamma \vdash \lambda x^{\mathcal{R}}.e : x{:}\mathcal{R} \to \mathcal{R}'}
$$

Figure 3. Typing Rules

$$
\textsc{S-Refl}\quad \Gamma \vdash \tau <: \tau
\qquad
\textsc{S-Pair}\quad \frac{\Gamma \vdash \mathcal{R}_1 \sqsubseteq \mathcal{R}_1' \qquad \Gamma \vdash \mathcal{R}_2 \sqsubseteq \mathcal{R}_2'}{\Gamma \vdash \langle \mathcal{R}_1, \mathcal{R}_2 \rangle <: \langle \mathcal{R}_1', \mathcal{R}_2' \rangle}
\qquad
\textsc{S-Refine}\quad \frac{\Gamma \vdash \psi_1 \supset \psi_2}{\Gamma \vdash \{\omega \mid \psi_1\} \sqsubseteq \{\omega \mid \psi_2\}}
$$

$$
\textsc{S-Fun}\quad \frac{\Gamma \vdash \mathcal{R}_1' \sqsubseteq \mathcal{R}_1 \qquad \Gamma \vdash \mathcal{R}_2 \sqsubseteq \mathcal{R}_2'}{\Gamma \vdash x{:}\mathcal{R}_1 \to \mathcal{R}_2 <: x{:}\mathcal{R}_1' \to \mathcal{R}_2'}
$$

Figure 4. Subtyping rules

$$\frac{\psi \in \Gamma}{\Gamma \vdash \psi}\ \text{L-Atom} \qquad \Gamma \vdash \mathtt{tt}\ \text{L-True} \qquad \frac{\Gamma \vdash \mathtt{ff}}{\Gamma \vdash \psi}\ \text{L-False} \qquad \frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \wedge \psi_2}\ \text{L-AndI}$$

$$\frac{\Gamma, \psi_1 \vdash \psi \text{ or } \Gamma, \psi_2 \vdash \psi}{\Gamma, \psi_1 \wedge \psi_2 \vdash \psi}\ \text{L-AndE} \qquad \frac{\Gamma, \psi_1 \vdash \psi_2}{\Gamma \vdash \psi_1 \supset \psi_2}\ \text{L-ImplI} \qquad \frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_1 \supset \psi_2}{\Gamma \vdash \psi_2}\ \text{L-ImplE}$$

$$\frac{\Gamma \vdash \psi_1 \text{ or } \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \vee \psi_2}\ \text{L-OrI} \qquad \frac{\Gamma, \psi_1 \vdash \psi \quad \Gamma, \psi_2 \vdash \psi}{\Gamma, \psi_1 \wedge \psi_2 \vdash \psi}\ \text{L-OrE} \qquad \frac{\Gamma \vdash \tau_{\pi(\kappa)} \quad \vdash \tau <: \sigma}{\Gamma \vdash \sigma_{\pi(\kappa)}}\ \text{L-Sub} \qquad \frac{\Gamma \vdash \overline{\sigma}_{\pi(\kappa)} \quad \vdash \tau <: \sigma}{\Gamma \vdash \overline{\tau}_{\pi(\kappa)}}\ \text{L-SubNot}$$

$$\frac{\Gamma \vdash \psi \quad \Gamma \vdash \pi(\kappa_1) \equiv \kappa_2 \text{ or } \Gamma \vdash \kappa_2 \equiv \pi(\kappa_1)}{\Gamma \vdash \psi[\pi(\kappa_1)/\kappa_2]}\ \text{L-Alias} \qquad \frac{\Gamma \vdash \bot_{\pi(\kappa)}}{\Gamma \vdash \psi}\ \text{L-Bot} \qquad \frac{\Gamma \vdash \tau_{\pi'(\kappa)} \quad \Gamma \vdash \nu_{\pi(\pi'(\kappa))}}{\Gamma \vdash \text{update}(\tau, \nu, \pi)_{\pi'(\kappa)}}\ \text{L-Update}$$

Figure 5. Proof System

$$\begin{array}{lcll}
\Gamma[o/\kappa] & = & \overrightarrow{\psi[o/\kappa]} & \Gamma = \overrightarrow{\psi} \\
\psi_+|\psi_-[o/\kappa] & = & \psi_+[o/\kappa]|\psi_-[o/\kappa] & \\
\\
\nu_{\pi(\kappa)}[\pi'(\kappa')/\kappa] & = & (\nu[\pi(\kappa')/\kappa])_{\pi(\pi'(\kappa'))} & \\
\nu_{\pi(\kappa)}[\emptyset/\kappa]_+ & = & \mathtt{tt} & \\
\nu_{\pi(\kappa)}[\emptyset/\kappa]_- & = & \mathtt{ff} & \\
\nu_{\pi(\kappa)}[o/\kappa'''] & = & \nu_{\pi(\kappa)} & \kappa \neq \kappa''' \text{ and } \kappa''' \notin \text{fv}(\nu) \\
\nu_{\pi(\kappa)}[o/\kappa''']_+ & = & \mathtt{tt} & \kappa \neq \kappa''' \text{ and } \kappa''' \in \text{fv}(\nu) \\
\nu_{\pi(\kappa)}[o/\kappa''']_- & = & \mathtt{ff} & \kappa \neq \kappa''' \text{ and } \kappa''' \in \text{fv}(\nu) \\
\mathtt{tt}[o/\kappa] & = & \mathtt{tt} & \\
\mathtt{ff}[o/\kappa] & = & \mathtt{ff} & \\
(\psi_1 \supset \psi_2)[o/\kappa]_+ & = & \psi_1[o/\kappa]_- \supset \psi_2[o/\kappa]_+ & \\
(\psi_1 \supset \psi_2)[o/\kappa]_- & = & \psi_1[o/\kappa]_+ \supset \psi_2[o/\kappa]_- & \\
(\psi_1 \vee \psi_2)[o/\kappa] & = & \psi_1[o/\kappa] \vee \psi_2[o/\kappa] & \\
(\psi_1 \wedge \psi_2)[o/\kappa] & = & \psi_1[o/\kappa] \wedge \psi_2[o/\kappa] & \\
(\pi'(\kappa') \equiv \pi''(y'))[\pi(\kappa''')/\kappa] & = & \pi'(\kappa')[\pi(\kappa''')/\kappa] \equiv \pi''(y')[\pi(\kappa''')/\kappa] & \\
(\pi(\kappa'') \equiv \pi'(\kappa'''))[\emptyset/\kappa] & = & \pi(\kappa'') \equiv \pi'(\kappa''') & \kappa \neq \kappa'' \text{ and } \kappa \neq \kappa''' \\
(\pi(\kappa) \equiv \pi'(\kappa''))[\emptyset/\kappa'''] & = & \mathtt{tt} & \\
\\
\pi(\kappa)[\pi'(\kappa'')/\kappa] & = & \pi(\pi'(\kappa'')) & \\
\pi(\kappa)[\emptyset/\kappa] & = & \emptyset & \\
\pi(\kappa)[o/\kappa'''] & = & \pi(\kappa) & \kappa \neq \kappa''' \\
\emptyset[o/\kappa] & = & \emptyset & \\
\\
\{\omega \mid \psi\}[o/\kappa] & = & \{\omega \mid \psi[o/\kappa]\} & \\
(\kappa{:}\sigma \to \tau)[o/\kappa'''] & = & \kappa{:}\sigma[o/\kappa'''] \to \tau[o/\kappa'''] & \kappa \neq \kappa'''
\end{array}$$

Figure 6. Substitution

$$\begin{array}{lcl}
\delta_\tau(number?) & = & x{:}\top \to \{\omega \mid \mathbf{B}_\omega \wedge \overline{\mathbf{\#f}}_\omega \supset \mathbf{N}_x \wedge \mathbf{\#f}_\omega \supset \overline{\mathbf{N}}_x\} \\
\delta_\tau(boolean?) & = & x{:}\top \to \{\omega \mid \mathbf{B}_\omega \wedge \overline{\mathbf{\#f}}_\omega \supset \mathbf{B}_x \wedge \mathbf{\#f}_\omega \supset \overline{\mathbf{B}}_x\} \\
\delta_\tau(cons?) & = & x{:}\top \to \{\omega \mid \mathbf{B}_\omega \wedge \overline{\mathbf{\#f}}_\omega \supset \langle\top,\top\rangle_x \wedge \mathbf{\#f}_\omega \supset \overline{\langle\top,\top\rangle}_x\} \\
\delta_\tau(add1) & = & x{:}\mathbf{N} \to \mathbf{N} \\
\delta_\tau(zero?) & = & x{:}\mathbf{N} \to \mathbf{B}
\end{array}$$

Figure 7. Constant Typing