## **Abstract**

Categories and Subject Descriptors CR-number [subcategory]: third-level

*Keywords* gradual types, Clojure Tobin-Hochstadt and Felleisen [1, 2]

## References

- [1] S. Tobin-Hochstadt and M. Felleisen. The design and implementation of typed scheme. SIGPLAN Not., 43 (1):395–406, Jan. 2008. ISSN 0362-1340. URL http://doi.acm.org/10.1145/1328897.1328486.
- [2] S. Tobin-Hochstadt and M. Felleisen. Logical types for untyped languages. SIGPLAN Not., 45(9):117–128, Sept. 2010. ISSN 0362-1340. URL http://doi.acm.org/10.1145/1932681.1863561.

Figure 1. Syntax of Terms, Types, Propositions, and Objects

$$\begin{array}{rcl} \top &=& \{\omega \mid \mathtt{tt}\} \\ \bot &=& \{\omega \mid \mathtt{ff}\} \\ \mathbf{B} &=& (\bigcup \{\omega \mid \mathtt{ft}_\omega\} \; \{\omega \mid \mathtt{ff}_\omega\}) \end{array}$$

Figure 2. Refinement Type Abbreviations

Figure 3. Typing Rules

$$\begin{array}{lll} & \text{S-Pair} & \text{S-Fun} \\ & \Gamma \vdash \mathcal{R}_1 \sqsubseteq \mathcal{R}_1' & \Gamma \vdash \mathcal{R}_2 \sqsubseteq \mathcal{R}_2' \\ & \Gamma \vdash \langle \mathcal{R}_1, \mathcal{R}_2 \rangle <: \langle \mathcal{R}_1', \mathcal{R}_2' \rangle & \frac{\Gamma \vdash \mathcal{R}_1' \sqsubseteq \mathcal{R}_1 & \Gamma \vdash \mathcal{R}_2 \sqsubseteq \mathcal{R}_2'}{\Gamma \vdash x : \mathcal{R}_1 & \Gamma \vdash \mathcal{R}_2 \sqsubseteq \mathcal{R}_2'} \\ & \frac{\text{S-UnionSuper}}{\Gamma \vdash \tau <: \sigma_i} & \frac{\text{S-UnionSub}}{\Gamma \vdash \tau_i <: \sigma_i} \\ & \frac{\exists i. \ \Gamma \vdash \tau <: \sigma_i}{\Gamma \vdash \psi_1 \supset \psi_2} & \frac{\Gamma \vdash \psi_1 \supset \psi_2}{\Gamma \vdash \{\omega \mid \psi_1\} \sqsubseteq \{\omega \mid \psi_2\}} \end{array}$$

Figure 4. Subtyping rules

Figure 5. Proof System

$$\begin{array}{lll} \Gamma[o/\kappa] & = & \overline{\psi[o/\kappa]} \\ \psi_+|\psi_-|o/\kappa| & = & \psi_+|o/\kappa||\psi_-|o/\kappa| \\ \end{array} & = & \psi_+|o/\kappa||\psi_-|o/\kappa| \\ \psi_+|\psi_-|o/\kappa| & = & \psi_+|o/\kappa||\psi_-|o/\kappa| \\ \psi_-|\psi_-|o/\kappa| & = & \psi_+|o/\kappa||\psi_-|o/\kappa| \\ \psi_-|\psi_-|o/\kappa| & = & \psi_+|o/\kappa||\psi_-|o/\kappa|| \\ \psi_-|\psi_-|o/\kappa| & = & \pi(\kappa')|\pi(\kappa'')/\kappa|| \\ \pi(\kappa)|\pi'(\kappa'')/\kappa| & = & \pi(\kappa')|\pi(\kappa'')/\kappa| \\ \pi(\kappa)|\pi'(\kappa'')/\kappa| & = & \pi(\kappa)|\pi(\kappa'')/\kappa| \\ \pi(\kappa)|\pi'(\kappa'')/\kappa| & = & \pi(\kappa)|\pi(\kappa'')/\kappa| \\ \pi(\kappa)|\pi'(\kappa'')/\kappa| & = & \pi(\kappa)|\pi(\kappa'')/\kappa| \\ \pi(\kappa)|\pi'(\kappa'')/\kappa| & = & \pi(\kappa)|\pi(\kappa)/\kappa| \\ \pi(\kappa)|\pi'(\kappa)/\kappa|\pi| & = & \pi(\kappa)|\pi(\kappa)/\kappa|\pi| \\ \pi(\kappa)|\pi'(\kappa)/\kappa|\pi| & = & \pi(\kappa)|$$

Figure 6. Substitution

$$\begin{array}{lll} \delta_{\tau}(number?) & = & x: \top \to \{\omega \mid \mathbf{B}_{\omega} \wedge \overline{\#}\mathbf{f}_{\omega} \supset \mathbf{N}_{x} \wedge \#\mathbf{f}_{\omega} \supset \overline{\mathbf{N}}_{x} \} \\ \delta_{\tau}(boolean?) & = & x: \top \to \{\omega \mid \mathbf{B}_{\omega} \wedge \overline{\#}\mathbf{f}_{\omega} \supset \mathbf{B}_{x} \wedge \#\mathbf{f}_{\omega} \supset \overline{\mathbf{B}}_{x} \} \\ \delta_{\tau}(cons?) & = & x: \top \to \{\omega \mid \mathbf{B}_{\omega} \wedge \overline{\#}\mathbf{f}_{\omega} \supset \langle \top, \top \rangle_{x} \wedge \#\mathbf{f}_{\omega} \supset \overline{\langle \top, \top \rangle_{x}} \} \\ \delta_{\tau}(add1) & = & x: \mathbf{N} \to \mathbf{N} \\ \delta_{\tau}(zero?) & = & x: \mathbf{N} \to \mathbf{B} \end{array}$$

Figure 7. Constant Typing

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 \langle \{\omega \mid \psi \wedge \nu_{\pi(\omega)}\}, \mathcal{R} \rangle \\ \langle \mathcal{R}, \{\omega \mid \psi \wedge \nu_{\pi(\omega)}\} \rangle  restrict(\Gamma, \tau, \sigma)
\begin{array}{l} \mathsf{update}(\Gamma, \langle \{\omega \mid \psi\}, \mathcal{R} \rangle, \nu, \pi :: \mathbf{car}) \\ \mathsf{update}(\Gamma, \langle \mathcal{R}, \{\omega \mid \psi\} \rangle, \nu, \pi :: \mathbf{cdr}) \end{array}
\mathsf{update}(\Gamma,\tau,\sigma,\epsilon)
\mathsf{update}(\Gamma, \tau, \overline{\sigma}, \epsilon)
                                                                                                                                                                remove(\Gamma, \tau, \overline{\sigma})
\mathsf{restrict}(\Gamma, \tau, \sigma)
                                                                                                                                                                                                                                                       if \not\exists v.\ \Gamma \vdash v: \{\omega \mid \sigma_\omega\} and \Gamma \vdash v: \{\omega \mid \tau_\omega\}
                                                                                                                                                                \perp
 \begin{array}{l} \operatorname{restrict}(\Gamma, (\bigcup \overrightarrow{\tau}), \sigma) \\ \operatorname{restrict}(\Gamma, \tau, \sigma) \end{array} 
                                                                                                                                                                (\bigcup \ \overrightarrow{\mathsf{restrict}}(\Gamma, \tau, \sigma))
                                                                                                                                               =
                                                                                                                                                                                                                                                                                                                                                                           if \Gamma \vdash \tau <: \sigma
                                                                                                                                                             \tau
\mathsf{restrict}(\Gamma, \tau, \sigma)
                                                                                                                                                                                                                                                                                                                                                                                        otherwise
                                                                                                                                               =
                                                                                                                                                               \sigma
\mathsf{remove}(\Gamma,\tau,\sigma)
                                                                                                                                                                                                                                                                                                                                                                           \text{if }\Gamma \vdash \tau \mathop{<:} \sigma
      \mathsf{remove}(\Gamma, (\bigcup \overrightarrow{\tau}), \sigma)               \mathsf{remove}(\Gamma, \tau, \sigma)        
                                                                                                                                                               (\bigcup \operatorname{remove}(\Gamma, \tau, \sigma))
                                                                                                                                                =
                                                                                                                                                                                                                                                                                                                                                                                        otherwise
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Figure 8. Type Update