

Q 6 Little E's Monster Challenge

Solution ① key Insight

find the longest increasing subsequence (LIS)

② Approach

(1) find the monster cannot be defeated

(2) both h and a values
strictly increasing

(3) Dynamic Programming

1. DP array $dp[i]$
represents the length of the
longest valid increasing subsequence
end at monster i

2. Initialize $dp[i] = 1$ for all i

because monster 1 can end at itself

3. for each pair (i, j) with $i < j$
update $dp[j] = \max(dp[j], dp[i] + 1)$ if $a[i] < a[j]$

4. answer the max val in dp

Example 1

$n=3$ $H=4$ $A=5$ $h=[1, 2, 3]$ $a=[3, 2, 1]$

Solution ① monster M_1 1 3 (defeatable?)
 M_2 2 2 ✓
 M_3 3 1 ✓

② potential subsequence

M_1 (1, 3) $dp[0] = 1$

M_2 (2, 2) $dp[1] = 1$

$$M_3(3, 1) \quad dp[2] = 1$$

③ try to chain them

fail loop monster
loop after monster
 $dp[after] = \max(after, before + 1)$

$$M_1(1, 3) \quad (2, 2) \quad (3, 1)$$

$$M_2(2, 2) \quad (1, 3) \quad (3, 1)$$

$$M_3(3, 1) \quad (1, 3) \quad (2, 2)$$

④ output : 1

Example 2

$$M_1(6, 8)$$

$$M_2(9, 9)$$

$$M_4(4, 2)$$

$$M_5(7, 5)$$

$$DP = [1 \quad 1 \quad 1 \quad 1]$$

$$\text{loop monster } [M_1 \quad M_2 \quad M_4 \quad M_5]$$

$$\text{loop after monster } [M_2 \quad M_4 \quad M_5]$$

$$i = 0 \quad j = 1$$

$$(6, 8) \rightarrow (9, 9)$$

$$dp[j] = \max(dp[j], dp[i] + 1)$$

$$DP = [1 \quad 2 \quad 1 \quad 1]$$

$$i = 3 \quad j = 4$$

$$(4, 2) \rightarrow (7, 5)$$

$$dp[4] = \max(1, 1+1) = 2$$

$$dp = [1 \quad 2 \quad 1 \quad 2]$$

Answer 1