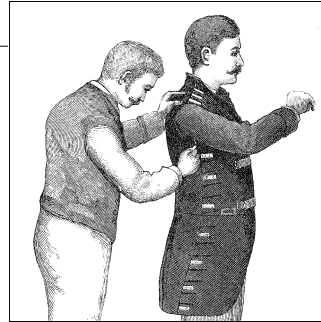


CHAPTER 14

The Linux Device Model



One of the stated goals for the 2.5 development cycle was the creation of a unified device model for the kernel. Previous kernels had no single data structure to which they could turn to obtain information about how the system is put together. Despite this lack of information, things worked well for some time. The demands of newer systems, with their more complicated topologies and need to support features such as power management, made it clear, however, that a general abstraction describing the structure of the system was needed.

The 2.6 device model provides that abstraction. It is now used within the kernel to support a wide variety of tasks, including:

Power management and system shutdown

These require an understanding of the system's structure. For example, a USB host adaptor cannot be shut down before dealing with all of the devices connected to that adaptor. The device model enables a traversal of the system's hardware in the right order.

Communications with user space

The implementation of the sysfs virtual filesystem is tightly tied into the device model and exposes the structure represented by it. The provision of information about the system to user space and knobs for changing operating parameters is increasingly done through sysfs and, therefore, through the device model.

Hotpluggable devices

Computer hardware is increasingly dynamic; peripherals can come and go at the whim of the user. The hotplug mechanism used within the kernel to handle and (especially) communicate with user space about the plugging and unplugging of devices is managed through the device model.

Device classes

Many parts of the system have little interest in how devices are connected, but they need to know what kinds of devices are available. The device model includes a mechanism for assigning devices to *classes*, which describe those

devices at a higher, functional level and allow them to be discovered from user space.

Object lifecycles

Many of the functions described above, including hotplug support and sysfs, complicate the creation and manipulation of objects created within the kernel. The implementation of the device model required the creation of a set of mechanisms for dealing with object lifecycles, their relationships to each other, and their representation in user space.

The Linux device model is a complex data structure. For example, consider Figure 14-1, which shows (in simplified form) a tiny piece of the device model structure associated with a USB mouse. Down the center of the diagram, we see the part of the core “devices” tree that shows how the mouse is connected to the system. The “bus” tree tracks what is connected to each bus, while the subtree under “classes” concerns itself with the functions provided by the devices, regardless of how they are connected. The device model tree on even a simple system contains hundreds of nodes like those shown in the diagram; it is a difficult data structure to visualize as a whole.

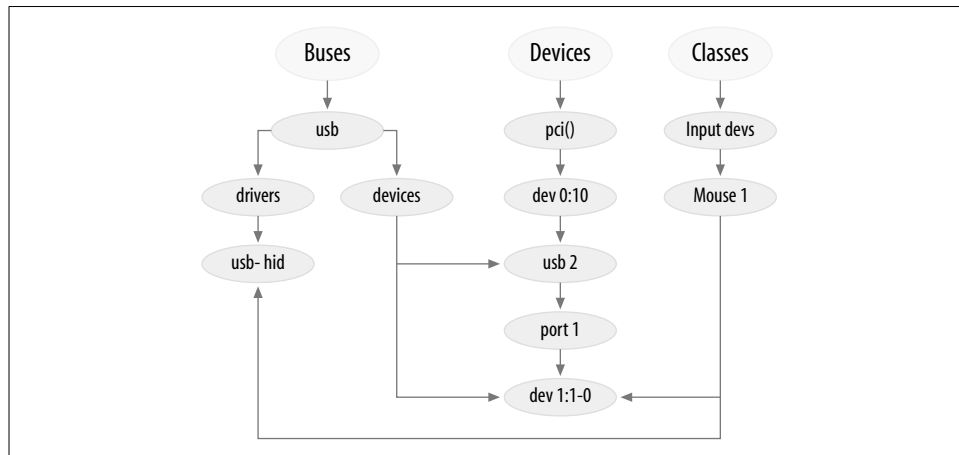


Figure 14-1. A small piece of the device model

For the most part, the Linux device model code takes care of all these considerations without imposing itself upon driver authors. It sits mostly in the background; direct interaction with the device model is generally handled by bus-level logic and various other kernel subsystems. As a result, many driver authors can ignore the device model entirely, and trust it to take care of itself.

There are times, however, when an understanding of the device model is a good thing to have. There are times when the device model “leaks out” from behind the other layers; for example, the generic DMA code (which we encounter in

Chapter 15) works with struct device. You may want to use some of the capabilities provided by the device model, such as the reference counting and related features provided by kobjects. Communication with user space via sysfs is also a device model function; this chapter explains how that communication works.

We start, however, with a bottom-up presentation of the device model. The complexity of the device model makes it hard to understand by starting with a high-level view. Our hope is that, by showing how the low-level device components work, we can prepare you for the challenge of grasping how those components are used to build the larger structure.

For many readers, this chapter can be treated as advanced material that need not be read the first time through. Those who are interested in how the Linux device model works are encouraged to press ahead, however, as we get into the low-level details.

Kobjects, Ksets, and Subsystems

The *kobject* is the fundamental structure that holds the device model together. It was initially conceived as a simple reference counter, but its responsibilities have grown over time, and so have its fields. The tasks handled by struct kobject and its supporting code now include:

Reference counting of objects

Often, when a kernel object is created, there is no way to know just how long it will exist. One way of tracking the lifecycle of such objects is through reference counting. When no code in the kernel holds a reference to a given object, that object has finished its useful life and can be deleted.

Sysfs representation

Every object that shows up in sysfs has, underneath it, a kobject that interacts with the kernel to create its visible representation.

Data structure glue

The device model is, in its entirety, a fiendishly complicated data structure made up of multiple hierarchies with numerous links between them. The kobject implements this structure and holds it together.

Hotplug event handling

The kobject subsystem handles the generation of events that notify user space about the comings and goings of hardware on the system.

One might conclude from the preceding list that the kobject is a complicated structure. One would be right. By looking at one piece at a time, however, it is possible to understand this structure and how it works.

Kobject Basics

A kobject has the type `struct kobject`; it is defined in `<linux/kobject.h>`. That file also includes declarations for a number of other structures related to kobjects and, of course, a long list of functions for manipulating them.

Embedding kobjects

Before we get into the details, it is worth taking a moment to understand how kobjects are used. If you look back at the list of functions handled by kobjects, you see that they are all services performed on behalf of other objects. A kobject, in other words, is of little interest on its own; it exists only to tie a higher-level object into the device model.

Thus, it is rare (even unknown) for kernel code to create a standalone kobject; instead, kobjects are used to control access to a larger, domain-specific object. To this end, kobjects are found embedded in other structures. If you are used to thinking of things in object-oriented terms, kobjects can be seen as a top-level, abstract class from which other classes are derived. A kobject implements a set of capabilities that are not particularly useful by themselves but that are nice to have in other objects. The C language does not allow for the direct expression of inheritance, so other techniques—such as embedding one structure in another—must be used.

As an example, let's look back at `struct cdev`, which we encountered in Chapter 3. That structure, as found in the 2.6.10 kernel, looks like this:

```
struct cdev {
    struct kobject kobj;
    struct module *owner;
    struct file_operations *ops;
    struct list_head list;
    dev_t dev;
    unsigned int count;
};
```

As we can see, the `cdev` structure has a kobject embedded within it. If you have one of these structures, finding its embedded kobject is just a matter of using the `kobj` field. Code that works with kobjects often has the opposite problem, however: given a `struct kobject` pointer, what is the pointer to the containing structure? You should avoid tricks (such as assuming that the kobject is at the beginning of the structure), and, instead, use the `container_of` macro (introduced in the section “The open Method” in Chapter 3). So the way to convert a pointer to a `struct kobject` called `kp` embedded within a `struct cdev` would be:

```
struct cdev *device = container_of(kp, struct cdev, kobj);
```

Programmers often define a simple macro for “back-casting” kobject pointers to the containing type.

Kobject initialization

This book has presented a number of types with simple mechanisms for initialization at compile or runtime. The initialization of a kobject is a bit more complicated, especially when all of its functions are used. Regardless of how a kobject is used, however, a few steps must be performed.

The first of those is to simply set the entire kobject to 0, usually with a call to `memset`. Often this initialization happens as part of the zeroing of the structure into which the kobject is embedded. Failure to zero out a kobject often leads to very strange crashes further down the line; it is not a step you want to skip.

The next step is to set up some of the internal fields with a call to `kobject_init()`:

```
void kobject_init(struct kobject *kobj);
```

Among other things, `kobject_init` sets the kobject's reference count to one. Calling `kobject_init` is not sufficient, however. Kobject users must, at a minimum, set the name of the kobject; this is the name that is used in sysfs entries. If you dig through the kernel source, you can find the code that copies a string directly into the kobject's name field, but that approach should be avoided. Instead, use:

```
int kobject_set_name(struct kobject *kobj, const char *format, ...);
```

This function takes a *printf*-style variable argument list. Believe it or not, it is actually possible for this operation to fail (it may try to allocate memory); conscientious code should check the return value and react accordingly.

The other kobject fields that should be set, directly or indirectly, by the creator are `ktype`, `kset`, and `parent`. We will get to these later in this chapter.

Reference count manipulation

One of the key functions of a kobject is to serve as a reference counter for the object in which it is embedded. As long as references to the object exist, the object (and the code that supports it) must continue to exist. The low-level functions for manipulating a kobject's reference counts are:

```
struct kobject *kobject_get(struct kobject *kobj);
void kobject_put(struct kobject *kobj);
```

A successful call to `kobject_get` increments the kobject's reference counter and returns a pointer to the kobject. If, however, the kobject is already in the process of being destroyed, the operation fails, and `kobject_get` returns NULL. This return value must always be tested, or no end of unpleasant race conditions could result.

When a reference is released, the call to `kobject_put` decrements the reference count and, possibly, frees the object. Remember that `kobject_init` sets the reference count to one; so when you create a kobject, you should make sure that the corresponding `kobject_put` call is made when that initial reference is no longer needed.

Note that, in many cases, the reference count in the kobject itself may not be sufficient to prevent race conditions. The existence of a kobject (and its containing structure) may well, for example, require the continued existence of the module that created that kobject. It would not do to unload that module while the kobject is still being passed around. That is why the `cdev` structure we saw above contains a `struct module` pointer. Reference counting for `struct cdev` is implemented as follows:

```
struct kobject *cdev_get(struct cdev *p)
{
    struct module *owner = p->owner;
    struct kobject *kobj;

    if (owner && !try_module_get(owner))
        return NULL;
    kobj = kobject_get(&p->kobj);
    if (!kobj)
        module_put(owner);
    return kobj;
}
```

Creating a reference to a `cdev` structure requires creating a reference also to the module that owns it. So `cdev_get` uses `try_module_get` to attempt to increment that module's usage count. If that operation succeeds, `kobject_get` is used to increment the kobject's reference count as well. That operation could fail, of course, so the code checks the return value from `kobject_get` and releases its reference to the module if things don't work out.

Release functions and kobject types

One important thing still missing from the discussion is what happens to a kobject when its reference count reaches 0. The code that created the kobject generally does not know when that will happen; if it did, there would be little point in using a reference count in the first place. Even predictable object life cycles become more complicated when sysfs is brought in; user-space programs can keep a reference to a kobject (by keeping one of its associated sysfs files open) for an arbitrary period of time.

The end result is that a structure protected by a kobject cannot be freed at any single, predictable point in the driver's lifecycle, but in code that must be prepared to run at whatever moment the kobject's reference count goes to 0. The reference count is not under the direct control of the code that created the kobject. So that code must be notified asynchronously whenever the last reference to one of its kobjects goes away.

This notification is done through a kobject's *release* method. Usually, this method has a form such as:

```
void my_object_release(struct kobject *kobj)
{
    struct my_object *mine = container_of(kobj, struct my_object, kobj);
```

```

        /* Perform any additional cleanup on this object, then... */
        kfree(mine);
    }

```

One important point cannot be overstated: every kobject must have a *release* method, and the kobject must persist (in a consistent state) until that method is called. If these constraints are not met, the code is flawed. It risks freeing the object when it is still in use, or it fails to release the object after the last reference is returned.

Interestingly, the *release* method is not stored in the kobject itself; instead, it is associated with the type of the structure that contains the kobject. This type is tracked with a structure of type `struct kobj_type`, often simply called a “ktype.” This structure looks like the following:

```

struct kobj_type {
    void (*release)(struct kobject *);
    struct sysfs_ops *sysfs_ops;
    struct attribute **default_attrs;
};

```

The `release` field in `struct kobj_type` is, of course, a pointer to the *release* method for this type of kobject. We will come back to the other two fields (`sysfs_ops` and `default_attrs`) later in this chapter.

Every kobject needs to have an associated `kobj_type` structure. Confusingly, the pointer to this structure can be found in two different places. The kobject structure itself contains a field (called *ktype*) that can contain this pointer. If, however, this kobject is a member of a `kset`, the `kobj_type` pointer is provided by that `kset` instead. (We will look at `ksets` in the next section.) Meanwhile, the macro:

```

struct kobj_type *get_ktype(struct kobject *kobj);

```

finds the `kobj_type` pointer for a given kobject.

Kobject Hierarchies, Ksets, and Subsystems

The kobject structure is often used to link together objects into a hierarchical structure that matches the structure of the subsystem being modeled. There are two separate mechanisms for this linking: the parent pointer and `ksets`.

The `parent` field in `struct kobject` is a pointer to another kobject—the one representing the next level up in the hierarchy. If, for example, a kobject represents a USB device, its `parent` pointer may indicate the object representing the hub into which the device is plugged.

The main use for the `parent` pointer is to position the object in the `sysfs` hierarchy. We’ll see how this works in the section “Low-Level `sysfs` Operations.”

Ksets

In many ways, a kset looks like an extension of the `kobj_type` structure; a kset is a collection of kobjects embedded within structures of the same type. However, while `struct kobj_type` concerns itself with the type of an object, `struct kset` is concerned with aggregation and collection. The two concepts have been separated so that objects of identical type can appear in distinct sets.

Therefore, the main function of a kset is containment; it can be thought of as the top-level container class for kobjects. In fact, each kset contains its own kobject internally, and it can, in many ways, be treated the same way as a kobject. It is worth noting that ksets are always represented in sysfs; once a kset has been set up and added to the system, there will be a sysfs directory for it. Kobjects do not necessarily show up in sysfs, but every kobject that is a member of a kset is represented there.

Adding a kobject to a kset is usually done when the object is created; it is a two-step process. The kobject's `kset` field must be pointed at the kset of interest; then the kobject should be passed to:

```
int kobject_add(struct kobject *kobj);
```

As always, programmers should be aware that this function can fail (in which case it returns a negative error code) and respond accordingly. There is a convenience function provided by the kernel:

```
extern int kobject_register(struct kobject *kobj);
```

This function is simply a combination of `kobject_init` and `kobject_add`.

When a kobject is passed to `kobject_add`, its reference count is incremented. Containment within the kset is, after all, a reference to the object. At some point, the kobject will probably have to be removed from the kset to clear that reference; that is done with:

```
void kobject_del(struct kobject *kobj);
```

There is also a `kobject_unregister` function, which is a combination of `kobject_del` and `kobject_put`.

A kset keeps its children in a standard kernel linked list. In almost all cases, the contained kobjects also have pointers to the kset (or, strictly, its embedded kobject) in their parent's fields. So, typically, a kset and its kobjects look something like what you see in Figure 14-2. Bear in mind that:

- All of the contained kobjects in the diagram are actually embedded within some other type, possibly even other ksets.
- It is not required that a kobject's parent be the containing kset (although any other organization would be strange and rare).

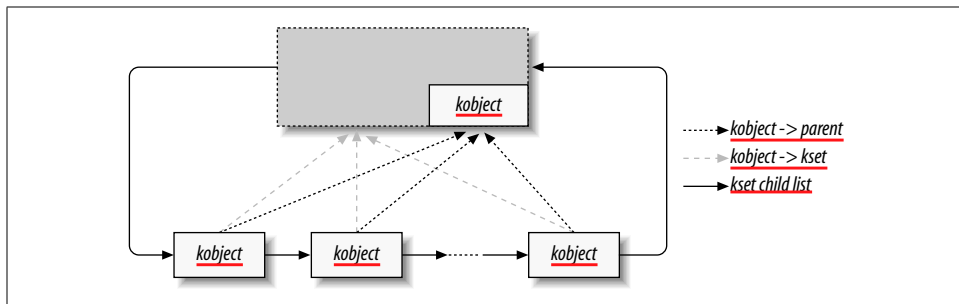


Figure 14-2. A simple kset hierarchy

Operations on ksets

For initialization and setup, ksets have an interface very similar to that of kobjects. The following functions exist:

```

void kset_init(struct kset *kset);
int kset_add(struct kset *kset);
int kset_register(struct kset *kset);
void kset_unregister(struct kset *kset);
    
```



For the most part, these functions just call the analogous *kobject* function on the kset's embedded kobject.

To manage the reference counts of ksets, the situation is about the same:

```

struct kset *kset_get(struct kset *kset);
void kset_put(struct kset *kset);
    
```

A kset also has a name, which is stored in the embedded kobject. So, if you have a kset called *my_set*, you would set its name with:

```

kobject_set_name(&my_set->kobj, "The name");
    
```

Ksets also have a pointer (in the *ktype* field) to the *kobj* type structure describing the kobjects it contains. This type is used in preference to the *ktype* field in a kobject itself. As a result, in typical usage, the *ktype* field in struct *kobject* is left NULL, because the same field within the kset is the one actually used.



Finally, a kset contains a subsystem pointer (called *subsys*). So it's time to talk about subsystems.

Subsystems

A subsystem is a representation for a high-level portion of the kernel as a whole. Subsystems usually (but not always) show up at the top of the sysfs hierarchy. Some example subsystems in the kernel include *block_subsys* (*/sys/block*, for block devices), *devices_subsys* (*/sys/devices*, the core device hierarchy), and a specific subsystem for every bus type known to the kernel. A driver author almost never needs to

create a new subsystem; if you feel tempted to do so, think again. What you probably want, in the end, is to add a new class, as discussed in the section “Classes.”

A subsystem is represented by a simple structure:

```
struct subsystem {
    struct kset kset;
    struct rw_semaphore rwsem;
};
```

A subsystem, thus, is really just a wrapper around a kset, with a semaphore thrown in.

Every kset must belong to a subsystem. The subsystem membership helps establish the kset's position in the hierarchy, but, more importantly, the subsystem's rwsem semaphore is used to serialize access to a kset's internal-linked list. This membership is represented by the subsys pointer in struct kset. Thus, one can find each kset's containing subsystem from the kset's structure, but one cannot find the multiple ksets contained in a subsystem directly from the subsystem structure.

Subsystems are often declared with a special macro:

```
decl_subsys(name, struct kobj_type *type,
            struct kset_hotplug_ops *hotplug_ops);
```

This macro creates a struct subsystem with a name formed by taking the name given to the macro and appending _subsys to it. The macro also initializes the internal kset with the given type and hotplug_ops. (We discuss hotplug operations later in this chapter.)

Subsystems have the usual list of setup and teardown functions:

```
void subsystem_init(struct subsystem *subsys);
int subsystem_register(struct subsystem *subsys);
void subsystem_unregister(struct subsystem *subsys);
struct subsystem *subsys_get(struct subsystem *subsys)
void subsys_put(struct subsystem *subsys);
```

Most of these operations just act upon the subsystem's kset.

Low-Level Sysfs Operations

Kobjects are the mechanism behind the sysfs virtual filesystem. For every directory found in sysfs, there is a kobject lurking somewhere within the kernel. Every kobject of interest also exports one or more *attributes*, which appear in that kobject's sysfs directory as files containing kernel-generated information. This section examines how kobjects and sysfs interact at a low level.

Code that works with sysfs should include `<linux/sysfs.h>`.

Getting a kobject to show up in sysfs is simply a matter of calling `kobject_add`. We have already seen that function as the way to add a kobject to a kset; creating entries

in sysfs is also part of its job. There are a couple of things worth knowing about how the sysfs entry is created:

- Sysfs entries for kobjects are always directories, so a call to *kobject_add* results in the creation of a directory in sysfs. Usually that directory contains one or more attributes; we see how attributes are specified shortly.
- The name assigned to the kobject (with *kobject_set_name*) is the name used for the sysfs directory. Thus, kobjects that appear in the same part of the sysfs hierarchy must have unique names. Names assigned to kobjects should also be reasonable file names: they cannot contain the slash character, and the use of white space is strongly discouraged.
- The sysfs entry is located in the directory corresponding to the kobject's parent pointer. If parent is NULL when *kobject_add* is called, it is set to the kobject embedded in the new kobject's kset; thus, the sysfs hierarchy usually matches the internal hierarchy created with ksets. If both parent and kset are NULL, the sysfs directory is created at the top level, which is almost certainly not what you want.

Using the mechanisms we have described so far, we can use a kobject to create an empty directory in sysfs. Usually, you want to do something a little more interesting than that, so it is time to look at the implementation of attributes.

Default Attributes



When created, every kobject is given a set of default attributes. These attributes are specified by way of the *kobj_type* structure. That structure, remember, looks like this:

```
struct kobj_type {
    void (*release)(struct kobject *);
    struct sysfs_ops *sysfs_ops;
    struct attribute **default_attrs;
};
```

The default attrs field lists the attributes to be created for every kobject of this type, and sysfs_ops provides the methods to implement those attributes. We start with *default_attrs*, which points to an array of pointers to attribute structures:

```
struct attribute {
    char *name;
    struct module *owner;
    mode_t mode;
};
```

In this structure, *name* is the name of the attribute (as it appears within the kobject's sysfs directory), *owner* is a pointer to the module (if any) that is responsible for the implementation of this attribute, and *mode* is the protection bits that are to be applied to this attribute. The mode is usually S_IRUGO for read-only attributes; if the attribute

is writable, you can toss in `S_IWUSR` to give write access to root only (the macros for modes are defined in `<linux/stat.h>`). The last entry in the `default_attrs` list must be zero-filled.

The `default_attrs` array says what the attributes are but does not tell sysfs how to actually implement those attributes. That task falls to the `kobj_type->sysfs_ops` field, which points to a structure defined as:

```
struct sysfs_ops {
    ssize_t (*show)(struct kobject *kobj, struct attribute *attr,
                    char *buffer);
    ssize_t (*store)(struct kobject *kobj, struct attribute *attr,
                    const char *buffer, size_t size);
};
```

Whenever an attribute is read from user space, the *show* method is called with a pointer to the kobject and the appropriate attribute structure. That method should encode the value of the given attribute into buffer, being sure not to overrun it (it is `PAGE_SIZE` bytes), and return the actual length of the returned data. The conventions for sysfs state that each attribute should contain a single, human-readable value; if you have a lot of information to return, you may want to consider splitting it into multiple attributes.

The same *show* method is used for all attributes associated with a given kobject. The *attr* pointer passed into the function can be used to determine which attribute is being requested. Some *show* methods include a series of tests on the attribute name. Other implementations embed the attribute structure within another structure that contains the information needed to return the attribute's value; in this case, *container_of* may be used within the *show* method to obtain a pointer to the embedding structure.

The *store* method is similar; it should decode the data stored in buffer (size contains the length of that data, which does not exceed `PAGE_SIZE`), store and respond to the new value in whatever way makes sense, and return the number of bytes actually decoded. The *store* method can be called only if the attribute's permissions allow writes. When writing a *store* method, never forget that you are receiving arbitrary information from user space; you should validate it very carefully before taking any action in response. If the incoming data does not match expectations, return a negative error value rather than possibly doing something unwanted and unrecoverable. If your device exports a *self_destruct* attribute, you should require that a specific string be written there to invoke that functionality; an accidental, random write should yield only an error.

Nondefault Attributes

In many cases, the kobject type's `default_attrs` field describes all the attributes that kobject will ever have. But that's not a restriction in the design; attributes can be

added and removed to kobjects at will. If you wish to add a new attribute to a kobject's sysfs directory, simply fill in an attribute structure and pass it to:

```
int sysfs_create_file(struct kobject *kobj, struct attribute *attr);
```

If all goes well, the file is created with the name given in the attribute structure, and the return value is 0; otherwise, the usual negative error code is returned.

Note that the same *show()* and *store()* functions are called to implement operations on the new attribute. Before you add a new, nondefault attribute to a kobject, you should take whatever steps are necessary to ensure that those functions know how to implement that attribute.

To remove an attribute, call:

```
int sysfs_remove_file(struct kobject *kobj, struct attribute *attr);
```

After the call, the attribute no longer appears in the kobject's sysfs entry. Do be aware, however, that a user-space process could have an open file descriptor for that attribute and that *show* and *store* calls are still possible after the attribute has been removed.

Binary Attributes

The sysfs conventions call for all attributes to contain a single value in a human-readable text format. That said, there is an occasional, rare need for the creation of attributes that can handle larger chunks of binary data. That need really only comes about when data must be passed, untouched, between user space and the device. For example, uploading firmware to devices requires this feature. When such a device is encountered in the system, a user-space program can be started (via the hotplug mechanism); that program then passes the firmware code to the kernel via a binary sysfs attribute, as is shown in the section “The Kernel Firmware Interface.”

Binary attributes are described with a `bin_attribute` structure:

```
struct bin_attribute {
    struct attribute attr;
    size_t size;
    ssize_t (*read)(struct kobject *kobj, char *buffer,
                    loff_t pos, size_t size);
    ssize_t (*write)(struct kobject *kobj, char *buffer,
                    loff_t pos, size_t size);
};
```

Here, `attr` is an attribute structure giving the name, owner, and permissions for the binary attribute, and `size` is the maximum size of the binary attribute (or 0 if there is no maximum). The *read* and *write* methods work similarly to the normal char driver equivalents; they can be called multiple times for a single load with a maximum of one page worth of data in each call. There is no way for sysfs to signal the last of a set

of write operations, so code implementing a binary attribute must be able to determine the end of the data some other way.

Binary attributes must be created explicitly; they cannot be set up as default attributes. To create a binary attribute, call:

```
int sysfs_create_bin_file(struct kobject *kobj,
                        struct bin_attribute *attr);
```

Binary attributes can be removed with:

```
int sysfs_remove_bin_file(struct kobject *kobj,
                        struct bin_attribute *attr);
```

Symbolic Links

The sysfs filesystem has the usual tree structure, reflecting the hierarchical organization of the kobjects it represents. The relationships between objects in the kernel are often more complicated than that, however. For example, one sysfs subtree (/sys/devices) represents all of the devices known to the system, while other subtrees (under /sys/bus) represent the device drivers. These trees do not, however, represent the relationships between the drivers and the devices they manage. Showing these additional relationships requires extra pointers which, in sysfs, are implemented through symbolic links.

Creating a symbolic link within sysfs is easy:

```
int sysfs_create_link(struct kobject *kobj, struct kobject *target,
                    char *name);
```

This function creates a link (called name) pointing to target's sysfs entry as an attribute of kobj. It is a relative link, so it works regardless of where sysfs is mounted on any particular system.

The link persists even if target is removed from the system. If you are creating symbolic links to other kobjects, you should probably have a way of knowing about changes to those kobjects, or some sort of assurance that the target kobjects will not disappear. The consequences (dead symbolic links within sysfs) are not particularly grave, but they are not representative of the best programming style and can cause confusion in user space.

Symbolic links can be removed with:

```
void sysfs_remove_link(struct kobject *kobj, char *name);
```

Hotplug Event Generation

A hotplug event is a notification to user space from the kernel that something has changed in the system's configuration. They are generated whenever a kobject is created or destroyed. Such events are generated, for example, when a digital camera is

plugged in with a USB cable, when a user switches console modes, or when a disk is repartitioned. Hotplug events turn into an invocation of `/sbin/hotplug`, which can respond to each event by loading drivers, creating device nodes, mounting partitions, or taking any other action that is appropriate.


The last major kobject function we look at is the generation of these events. The actual event generation takes place when a kobject is passed to `kobject_add` or `kobject_del`. Before the event is handed to user space, code associated with the kobject (or, more specifically, the kset to which it belongs) has the opportunity to add information for user space or to disable event generation entirely.

Hotplug Operations

Actual control of hotplug events is exercised by way of a set of methods stored in the `kset_hotplug_ops` structure:

```
struct kset_hotplug_ops {
    int (*filter)(struct kset *kset, struct kobject *kobj);
    char *(*name)(struct kset *kset, struct kobject *kobj);
    int (*hotplug)(struct kset *kset, struct kobject *kobj,
                  char **envp, int num_envp, char *buffer,
                  int buffer_size);
};
```

A pointer to this structure is found in the `hotplug_ops` field of the `kset` structure. If a given kobject is not contained within a kset, the kernel searches up through the hierarchy (via the parent pointer) until it finds a kobject that *does* have a kset; that kset's hotplug operations are then used.

The `filter` hotplug operation is called whenever the kernel is considering generating an event for a given kobject. If `filter` returns 0, the event is not created. This method, therefore, gives the kset code an opportunity to decide which events should be passed on to user space and which should not. 

As an example of how this method might be used, consider the block subsystem. There are at least three types of kobjects used there, representing disks, partitions, and request queues. User space may want to react to the addition of a disk or a partition, but it does not normally care about request queues. So the `filter` method allows event generation only for kobjects representing disks and partitions. It looks like this:

```
static int block_hotplug_filter(struct kset *kset, struct kobject *kobj)
{
    struct kobj_type *ktype = get_ktype(kobj);

    return ((ktype == &ktype_block) || (ktype == &ktype_part));
}
```

Here, a quick test on the type of kobject is sufficient to decide whether the event should be generated or not.

When the user-space hotplug program is invoked, it is passed to the name of the relevant subsystem as its one and only parameter. The *name* hotplug method is charged with providing that name. It should return a simple string suitable for passing to user space.

Everything else that the hotplug script might want to know is passed in the environment. The final hotplug method (*hotplug*) gives an opportunity to add useful environment variables prior to the invocation of that script. Again, this method's prototype is:

```
int (*hotplug)(struct kset *kset, struct kobject *kobj,
               char **envp, int num_envp, char *buffer,
               int buffer_size);
```

As usual, *kset* and *kobject* describe the object for which the event is being generated. The *envp* array is a place to store additional environment variable definitions (in the usual *NAME=value* format); it has *num_envp* entries available. The variables themselves should be encoded into *buffer*, which is *buffer_size* bytes long. If you add any variables to *envp*, be sure to add a *NULL* entry after your last addition so that the kernel knows where the end is. The return value should normally be 0; any nonzero return aborts the generation of the hotplug event.

The generation of hotplug events (like much of the work in the device model) is usually handled by logic at the bus driver level.

Buses, Devices, and Drivers

So far, we have seen a great deal of low-level infrastructures and a relative shortage of examples. We try to make up for that in the rest of this chapter as we get into the higher levels of the Linux device model. To that end, we introduce a new virtual bus, which we call *lddbus*,* and modify the *scullp* driver to “connect” to that bus.

Once again, much of the material covered here will never be needed by many driver authors. Details at this level are generally handled at the bus level, and few authors need to add a new bus type. This information is useful, however, for anybody wondering what is happening inside the PCI, USB, etc. layers or who needs to make changes at that level.

Buses

A bus is a channel between the processor and one or more devices. For the purposes of the device model, all devices are connected via a bus, even if it is an internal, virtual, “platform” bus. Buses can plug into each other—a USB controller is usually a

* The logical name for this bus, of course, would have been “sbus,” but that name was already taken by a real, physical bus.

PCI device, for example. The device model represents the actual connections between buses and the devices they control.

In the Linux device model, a bus is represented by the `bus_type` structure, defined in `<linux/device.h>`. This structure looks like:

```
struct bus_type {
    char *name;
    struct subsystem subsys;
    struct kset drivers;
    struct kset devices;
    int (*match)(struct device *dev, struct device_driver *drv);
    struct device *(*add)(struct device *parent, char *bus_id);
    int (*hotplug)(struct device *dev, char **envp,
                   int num_envp, char *buffer, int buffer_size);
    /* Some fields omitted */
};
```



The `name` field is the name of the bus, something such as `pci`. You can see from the structure that each bus is its own subsystem; these subsystems do not live at the top level in `sysfs`, however. Instead, they are found underneath the bus subsystem. A bus contains two ksets, representing the known drivers for that bus and all devices plugged into the bus. Then, there is a set of methods that we will get to shortly.

Bus registration

As we mentioned, the example source includes a virtual bus implementation called `lddbus`. This bus sets up its `bus_type` structure as follows:

```
struct bus_type ldd_bus_type = {
    .name = "ldd",
    .match = ldd_match,
    .hotplug = ldd_hotplug,
};
```

Note that very few of the `bus_type` fields require initialization; most of that is handled by the device model core. We do have to specify the name of the bus, however, and any methods that go along with it.

Inevitably, a new bus must be registered with the system via a call to `bus_register`. The `lddbus` code does so in this way:

```
ret = bus_register(&ldd_bus_type);
if (ret)
    return ret;
```

This call can fail, of course, so the return value must always be checked. If it succeeds, the new bus subsystem has been added to the system; it is visible in `sysfs` under `/sys/bus`, and it is possible to start adding devices.

Should it be necessary to remove a bus from the system (when the associated module is removed, for example), `bus_unregister` should be called:

```
void bus_unregister(struct bus_type *bus);
```

Bus methods

There are several methods defined for the `bus_type` structure; they allow the bus code to serve as an intermediary between the device core and individual drivers. The methods defined in the 2.6.10 kernel are:

```
int (*match)(struct device *device, struct device_driver *driver);
```

This method is called, perhaps multiple times, whenever a new device or driver is added for this bus. It should return a nonzero value if the given device can be handled by the given driver. (We get to the details of the device and device_driver structures shortly). This function must be handled at the bus level, because that is where the proper logic exists; the core kernel cannot know how to match devices and drivers for every possible bus type.

```
int (*hotplug) (struct device *device, char **envp, int num_envp, char
                *buffer, int buffer_size);
```

This method allows the bus to add variables to the environment prior to the generation of a hotplug event in user space. The parameters are the same as for the `kset hotplug` method (described in the earlier section “Hotplug Event Generation”).

The `lddbus` driver has a very simple `match` function, which simply compares the driver and device names:

```
static int ldd_match(struct device *dev, struct device_driver *driver)
{
    return !strcmp(dev->bus_id, driver->name, strlen(driver->name));
}
```

When real hardware is involved, the `match` function usually makes some sort of comparison between the hardware ID provided by the device itself and the IDs supported by the driver.

The `lddbus hotplug` method looks like this:

```
static int ldd_hotplug(struct device *dev, char **envp, int num_envp,
                      char *buffer, int buffer_size)
{
    envp[0] = buffer;
    if (snprintf(buffer, buffer_size, "LDBBUS_VERSION=%s",
                Version) >= buffer_size)
        return -ENOMEM;
    envp[1] = NULL;
    return 0;
}
```

Here, we add in the current revision number of the `lddbus` source, just in case anybody is curious.

Iterating over devices and drivers

If you are writing bus-level code, you may find yourself having to perform some operation on all devices or drivers that have been registered with your bus. It may be

tempting to dig directly into the structures in the `bus_type` structure, but it is better to use the helper functions that have been provided.

To operate on every device known to the bus, use:

```
int bus_for_each_dev(struct bus_type *bus, struct device *start,
                   void *data, int (*fn)(struct device *, void *));
```

This function iterates over every device on bus, passing the associated device structure to `fn`, along with the value passed in as `data`. If `start` is `NULL`, the iteration begins with the first device on the bus; otherwise iteration starts with the first device after `start`. If `fn` returns a nonzero value, iteration stops and that value is returned from `bus_for_each_dev`.

There is a similar function for iterating over drivers:

```
int bus_for_each_drv(struct bus_type *bus, struct device_driver *start,
                   void *data, int (*fn)(struct device_driver *, void *));
```

This function works just like `bus_for_each_dev`, except, of course, that it works with drivers instead.

It should be noted that both of these functions hold the bus subsystem's reader/writer semaphore for the duration of the work. So an attempt to use the two of them together will deadlock—each will be trying to obtain the same semaphore. Operations that modify the bus (such as unregistering devices) will also lock up. So, use the `bus_for_each` functions with some care.

Bus attributes

Almost every layer in the Linux device model provides an interface for the addition of attributes, and the bus layer is no exception. The `bus_attribute` type is defined in `<linux/device.h>` as follows:

```
struct bus_attribute {
    struct attribute attr;
    ssize_t (*show)(struct bus_type *bus, char *buf);
    ssize_t (*store)(struct bus_type *bus, const char *buf,
                    size_t count);
};
```

We have already seen `struct attribute` in the section “Default Attributes.” The `bus_attribute` type also includes two methods for displaying and setting the value of the attribute. Most device model layers above the `kobject` level work this way.

A convenience macro has been provided for the compile-time creation and initialization of `bus_attribute` structures:

```
BUS_ATTR(name, mode, show, store);
```

This macro declares a structure, generating its name by prepending the string `bus_attr` to the given name.

Any attributes belonging to a bus should be created explicitly with *bus_create_file*:

```
int bus_create_file(struct bus_type *bus, struct bus_attribute *attr);
```

Attributes can also be removed with:

```
void bus_remove_file(struct bus_type *bus, struct bus_attribute *attr);
```

The *lddbus* driver creates a simple attribute file containing, once again, the source version number. The *show* method and *bus_attribute* structure are set up as follows:

```
static ssize_t show_bus_version(struct bus_type *bus, char *buf)
{
    return snprintf(buf, PAGE_SIZE, "%s\n", Version);
}
```

```
static BUS_ATTR(version, S_IRUGO, show_bus_version, NULL);
```

Creating the attribute file is done at module load time:

```
if (bus_create_file(&ldd_bus_type, &bus_attr_version))
    printk(KERN_NOTICE "Unable to create version attribute\n");
```

This call creates an attribute file (*/sys/bus/ldd/version*) containing the revision number for the *lddbus* code.

Devices

At the lowest level, every device in a Linux system is represented by an instance of *struct device*:

```
struct device {
    struct device *parent;
    struct kobject kobj;
    char bus_id[BUS_ID_SIZE];
    struct bus_type *bus;
    struct device_driver *driver;
    void *driver_data;
    void (*release)(struct device *dev);
    /* Several fields omitted */
};
```

There are many other *struct device* fields that are of interest only to the device core code. These fields, however, are worth knowing about:

*struct device *parent*

The device's "parent" device—the device to which it is attached. In most cases, a parent device is some sort of bus or host controller. If *parent* is *NULL*, the device is a top-level device, which is not usually what you want.

struct kobject kobj;

The *kobject* that represents this device and links it into the hierarchy. Note that, as a general rule, *device->kobj->parent* is equal to *&device->parent->kobj*.

```
char bus_id[BUS_ID_SIZE];
```

A string that uniquely identifies this device on the bus. PCI devices, for example, use the standard PCI ID format containing the domain, bus, device, and function numbers.

```
struct bus_type *bus;
```

Identifies which kind of bus the device sits on.

```
struct device_driver *driver;
```

The driver that manages this device; we examine struct device_driver in the next section.

```
void *driver_data;
```

A private data field that may be used by the device driver.

```
void (*release)(struct device *dev);
```

The method is called when the last reference to the device is removed; it is called from the embedded kobject's *release* method. All device structures registered with the core must have a *release* method, or the kernel prints out scary complaints.

At a minimum, the parent, bus_id, bus, and release fields must be set before the device structure can be registered.



Device registration

The usual set of registration and unregistration functions exists:

```
int device_register(struct device *dev);
void device_unregister(struct device *dev);
```

We have seen how the *lddbus* code registers its bus type. However, an actual bus is a device and must be registered separately. For simplicity, the *lddbus* module supports only a single virtual bus, so the driver sets up its device at compile time:

```
static void ldd_bus_release(struct device *dev)
{
    printk(KERN_DEBUG "lddbus release\n");
}

struct device ldd_bus = {
    .bus_id = "ldd0",
    .release = ldd_bus_release
};
```

This is a top-level bus, so the parent and bus fields are left NULL. We have a simple, no-op *release* method, and, as the first (and only) bus, its name is ldd0. This bus device is registered with:

```
ret = device_register(&ldd_bus);
if (ret)
    printk(KERN_NOTICE "Unable to register ldd0\n");
```

Once that call is complete, the new bus can be seen under */sys/devices* in sysfs. Any devices added to this bus then shows up under */sys/devices/ldd0/*.

Device attributes

Device entries in sysfs can have attributes. The relevant structure is:

```
struct device_attribute {
    struct attribute attr;
    ssize_t (*show)(struct device *dev, char *buf);
    ssize_t (*store)(struct device *dev, const char *buf,
                    size_t count);
};
```

These attribute structures can be set up at compile time with this macro:

```
DEVICE_ATTR(name, mode, show, store);
```

The resulting structure is named by prepending `dev_attr_` to the given name. The actual management of attribute files is handled with the usual pair of functions:

```
int device_create_file(struct device *device,
                     struct device_attribute *entry);
void device_remove_file(struct device *dev,
                      struct device_attribute *attr);
```

The `dev_attrs` field of `struct bus_type` points to a list of default attributes created for every device added to that bus.

Device structure embedding

The device structure contains the information that the device model core needs to model the system. Most subsystems, however, track additional information about the devices they host. As a result, it is rare for devices to be represented by bare device structures; instead, that structure, like `kobject` structures, is usually embedded within a higher-level representation of the device. If you look at the definitions of `struct pci_dev` or `struct usb_device`, you will find a `struct device` buried inside. Usually, low-level drivers are not even aware of that `struct device`, but there can be exceptions.

The `lddbus` driver creates its own device type (`struct ldd_device`) and expects individual device drivers to register their devices using that type. It is a simple structure:

```
struct ldd_device {
    char *name;
    struct ldd_driver *driver;
    struct device dev;
};

#define to_ldd_device(dev) container_of(dev, struct ldd_device, dev);
```

This structure allows the driver to provide an actual name for the device (which can be distinct from its bus ID, stored in the device structure) and a pointer to driver information. Structures for real devices usually also contain information about the vendor, device model, device configuration, resources used, and so on. Good examples can be

found in struct `pci_dev` (`<linux/pci.h>`) or struct `usb_device` (`<linux/usb.h>`). A convenience macro (`to_ldd_device`) is also defined for struct `ldd_device` to make it easy to turn pointers to the embedded device structure into `ldd_device` pointers.

The registration interface exported by `lddbus` looks like this:

```
int register_ldd_device(struct ldd_device *ldddev)
{
    ldddev->dev.bus = &ldd_bus_type;
    ldddev->dev.parent = &ldd_bus;
    ldddev->dev.release = ldd_dev_release;
    strncpy(ldddev->dev.bus_id, ldddev->name, BUS_ID_SIZE);
    return device_register(&ldddev->dev);
}
EXPORT_SYMBOL(register_ldd_device);
```

Here, we simply fill in some of the embedded device structure fields (which individual drivers should not need to know about), and register the device with the driver core. If we wanted to add bus-specific attributes to the device, we could do so here.

To show how this interface is used, let us introduce another sample driver, which we have called `sculld`. It is yet another variant on the `sculdp` driver first introduced in Chapter 8. It implements the usual memory area device, but `sculld` also works with the Linux device model by way of the `lddbus` interface.

The `sculld` driver adds an attribute of its own to its device entry; this attribute, called `dev`, simply contains the associated device number. This attribute could be used by a module loading the script or the hotplug subsystem to automatically create device nodes when the device is added to the system. The setup for this attribute follows the usual patterns:

```
static ssize_t sculld_show_dev(struct device *ddev, char *buf)
{
    struct sculld_dev *dev = ddev->driver_data;

    return print_dev_t(buf, dev->cdev.dev);
}

static DEVICE_ATTR(dev, S_IRUGO, sculld_show_dev, NULL);
```

Then, at initialization time, the device is registered, and the `dev` attribute is created through the following function:

```
static void sculld_register_dev(struct sculld_dev *dev, int index)
{
    sprintf(dev->devname, "sculld%d", index);
    dev->ldev.name = dev->devname;
    dev->ldev.driver = &sculld_driver;
    dev->ldev.dev.driver_data = dev;
    register_ldd_device(&dev->ldev);
    device_create_file(&dev->ldev.dev, &dev_attr_dev);
}
```

Note that we make use of the `driver_data` field to store the pointer to our own, internal device structure.

Device Drivers

The device model tracks all of the drivers known to the system. The main reason for this tracking is to enable the driver core to match up drivers with new devices. Once drivers are known objects within the system, however, a number of other things become possible. Device drivers can export information and configuration variables that are independent of any specific device, for example.

Drivers are defined by the following structure:

```
struct device_driver {
    char *name;
    struct bus_type *bus;
    struct kobject kobj;
    struct list_head devices;
    int (*probe)(struct device *dev);
    int (*remove)(struct device *dev);
    void (*shutdown) (struct device *dev);
};
```

Once again, several of the structure's fields have been omitted (see `<linux/device.h>` for the full story). Here, `name` is the name of the driver (it shows up in `sysfs`), `bus` is the type of bus this driver works with, `kobj` is the inevitable `kobject`, `devices` is a list of all devices currently bound to this driver, `probe` is a function called to query the existence of a specific device (and whether this driver can work with it), `remove` is called when the device is removed from the system, and `shutdown` is called at shutdown time to quiesce the device.

The form of the functions for working with `device_driver` structures should be looking familiar by now (so we cover them very quickly). The registration functions are:

```
int driver_register(struct device_driver *drv);
void driver_unregister(struct device_driver *drv);
```

The usual attribute structure exists:

```
struct driver_attribute {
    struct attribute attr;
    ssize_t (*show)(struct device_driver *drv, char *buf);
    ssize_t (*store)(struct device_driver *drv, const char *buf,
                    size_t count);
};
DRIVER_ATTR(name, mode, show, store);
```

And attribute files are created in the usual way:

```
int driver_create_file(struct device_driver *drv,
                     struct driver_attribute *attr);
void driver_remove_file(struct device_driver *drv,
                      struct driver_attribute *attr);
```


The bus type structure contains a field (`drv attrs`) that points to a set of default attributes, which are created for all drivers associated with that bus.

Driver structure embedding

As is the case with most driver core structures, the `device_driver` structure is usually found embedded within a higher-level, bus-specific structure. The *lddbus* subsystem would never go against such a trend, so it has defined its own `ldd_driver` structure:

```
struct ldd_driver {
    char *version;
    struct module *module;
    struct device_driver driver;
    struct driver_attribute version_attr;
};

#define to_ldd_driver(drv) container_of(drv, struct ldd_driver, driver);
```

Here, we require each driver to provide its current software version, and *lddbus* exports that version string for every driver it knows about. The bus-specific driver registration function is:

```
int register_ldd_driver(struct ldd_driver *driver)
{
    int ret;

    driver->driver.bus = &ldd_bus_type;
    ret = driver_register(&driver->driver);
    if (ret)
        return ret;
    driver->version_attr.attr.name = "version";
    driver->version_attr.attr.owner = driver->module;
    driver->version_attr.attr.mode = S_IRUGO;
    driver->version_attr.show = show_version;
    driver->version_attr.store = NULL;
    return driver_create_file(&driver->driver, &driver->version_attr);
}
```

The first half of the function simply registers the low-level `device_driver` structure with the core; the rest sets up the version attribute. Since this attribute is created at runtime, we can't use the `DRIVER_ATTR` macro; instead, the `driver_attribute` structure must be filled in by hand. Note that we set the owner of the attribute to the driver module, rather than the *lddbus* module; the reason for this can be seen in the implementation of the *show* function for this attribute:

```
static ssize_t show_version(struct device_driver *driver, char *buf)
{
    struct ldd_driver *ldriver = to_ldd_driver(driver);

    sprintf(buf, "%s\n", ldriver->version);
    return strlen(buf);
}
```

One might think that the attribute owner should be the *lddbus* module, since the function that implements the attribute is defined there. This function, however, is working with the *ldd_driver* structure created (and owned) by the driver itself. If that structure were to go away while a user-space process tried to read the version number, things could get messy. Designating the driver module as the owner of the attribute prevents the module from being unloaded, while user-space holds the attribute file open. Since each driver module creates a reference to the *lddbus* module, we can be sure that *lddbus* will not be unloaded at an inopportune time.

For completeness, *sculld* creates its *ldd_driver* structure as follows:

```
static struct ldd_driver sculld_driver = {
    .version = "$Revision: 1.1 $",
    .module = THIS_MODULE,
    .driver = {
        .name = "sculld",
    },
};
```

A simple call to *register_ldd_driver* adds it to the system. Once initialization is complete, the driver information can be seen in *sysfs*:

```
$ tree /sys/bus/ldd/drivers
/sys/bus/ldd/drivers
|-- sculld
|   |-- sculld0 -> ../../../../devices/ldd0/sculld0
|   |-- sculld1 -> ../../../../devices/ldd0/sculld1
|   |-- sculld2 -> ../../../../devices/ldd0/sculld2
|   |-- sculld3 -> ../../../../devices/ldd0/sculld3
|   |-- version
|   `--
```

Classes

The final device model concept we examine in this chapter is the *class*. A class is a higher-level view of a device that abstracts out low-level implementation details. Drivers may see a SCSI disk or an ATA disk, but, at the class level, they are all simply disks. Classes allow user space to work with devices based on what they do, rather than how they are connected or how they work.

Almost all classes show up in *sysfs* under */sys/class*. Thus, for example, all network interfaces can be found under */sys/class/net*, regardless of the type of interface. Input devices can be found in */sys/class/input*, and serial devices are in */sys/class/tty*. The one exception is block devices, which can be found under */sys/block* for historical reasons.

Class membership is usually handled by high-level code without the need for explicit support from drivers. When the *sbull* driver (see Chapter 16) creates a virtual disk device, it automatically appears in */sys/block*. The *snulld* network driver (see Chapter 17) does not have to do anything special for its interfaces to be represented

in `/sys/class/net`. There will be times, however, when drivers end up dealing with classes directly.

In many cases, the class subsystem is the best way of exporting information to user space. When a subsystem creates a class, it owns the class entirely, so there is no need to worry about which module owns the attributes found there. It also takes very little time wandering around in the more hardware-oriented parts of sysfs to realize that it can be an unfriendly place for direct browsing. Users more happily find information in `/sys/class/some-widget` than under, say, `/sys/devices/pci0000:00/0000:00:10.0/usb2/2-0:1.0`.

The driver core exports two distinct interfaces for managing classes. **The `class_simple` routines** are designed to make it as easy as possible to add new classes to the system; their main purpose, usually, is to expose attributes containing device numbers to enable the automatic creation of device nodes. **The regular class interface** is more complex but offers more features as well. We start with the simple version.

The `class_simple` Interface

The `class_simple` interface was intended to be so easy to use that nobody would have any excuse for not exporting, at a minimum, an attribute containing a device's assigned number. Using this interface is simply a matter of a couple of function calls, with little of the usual boilerplate associated with the Linux device model.

The first step is to create the class itself. That is accomplished with a call to `class_simple_create`:

```
struct class_simple *class_simple_create(struct module *owner, char *name);
```

This function creates a class with the given name. The operation can fail, of course, so the return value should always be checked (using `IS_ERR`, described in the section “Pointers and Error Values” in Chapter 1) before continuing.

A simple class can be destroyed with:

```
void class_simple_destroy(struct class_simple *cs);
```

The real purpose of creating a simple class is to add devices to it; that task is achieved with:

```
struct class_device *class_simple_device_add(struct class_simple *cs,
                                             dev_t devnum,
                                             struct device *device,
                                             const char *fmt, ...);
```

Here, `cs` is the previously created simple class, `devnum` is the assigned device number, `device` is the `struct device` representing this device, and the remaining parameters are a *printf*-style format string and arguments to create the device name. This call adds an entry to the class containing one attribute, `dev`, which holds the device number. If the device parameter is not `NULL`, a symbolic link (called `device`) points to the device's entry under `/sys/devices`.

It is possible to add other attributes to a device entry. It is just a matter of using *class_device_create_file*, which we discuss in the next section with the rest of the full class subsystem.

Classes generate hotplug events when devices come and go. If your driver needs to add variables to the environment for the user-space event handler, it can set up a hot-plug callback with:

```
int class_simple_set_hotplug(struct class_simple *cs,
                           int (*hotplug)(struct class_device *dev,
                                           char **envp, int num_envp,
                                           char *buffer, int buffer_size));
```

When your device goes away, the class entry should be removed with:

```
void class_simple_device_remove(dev_t dev);
```

Note that the *class_device* structure returned by *class_simple_device_add* is not needed here; the device number (which should certainly be unique) is sufficient.

The Full Class Interface

The *class_simple* interface suffices for many needs, but sometimes more flexibility is required. The following discussion describes how to use the full class mechanism, upon which *class_simple* is based. It is brief: the class functions and structures follow the same patterns as the rest of the device model, so there is little that is truly new here.

Managing classes

A class is defined by an instance of struct *class*:

```
struct class {
    char *name;
    struct class_attribute *class_attrs;
    struct class_device_attribute *class_dev_attrs;
    int (*hotplug)(struct class_device *dev, char **envp,
                  int num_envp, char *buffer, int buffer_size);
    void (*release)(struct class_device *dev);
    void (*class_release)(struct class *class);
    /* Some fields omitted */
};
```

Each class needs a unique name, which is how this class appears under /sys/class. When the class is registered, all of the attributes listed in the (NULL-terminated) array pointed to by class_attrs is created. There is also a set of default attributes for every device added to the class; *class_dev_attrs* points to those. There is the usual *hot-plug* function for adding variables to the environment when events are generated. There are also two release methods: release is called whenever a device is removed from the class, while class_release is called when the class itself is released.

The registration functions are:

```
int class_register(struct class *cls);
void class_unregister(struct class *cls);
```

The interface for working with attributes should not surprise anybody at this point:

```
struct class_attribute {
    struct attribute attr;
    ssize_t (*show)(struct class *cls, char *buf);
    ssize_t (*store)(struct class *cls, const char *buf, size_t count);
};

CLASS_ATTR(name, mode, show, store);

int class_create_file(struct class *cls,
                     const struct class_attribute *attr);
void class_remove_file(struct class *cls,
                      const struct class_attribute *attr);
```

Class devices

The real purpose of a class is to serve as a container for the devices that are members of that class. A member is represented by struct `class_device`:

```
struct class_device {
    struct kobject kobj;
    struct class *class;
    struct device *dev;
    void *class_data;
    char class_id[BUS_ID_SIZE];
};
```

The class `id` field holds the name of this device as it appears in `sysfs`. The class pointer should point to the class holding this device, and `dev` should point to the associated device structure. Setting `dev` is optional; if it is non-NULL, it is used to create a symbolic link from the class entry to the corresponding entry under `/sys/devices`, making it easy to find the device entry in user space. The class can use `class_data` to hold a private pointer.

The usual registration functions have been provided:

```
int class_device_register(struct class_device *cd);
void class_device_unregister(struct class_device *cd);
```

The class device interface also allows the renaming of an already registered entry:

```
int class_device_rename(struct class_device *cd, char *new_name);
```

Class device entries have attributes:

```
struct class_device_attribute {
    struct attribute attr;
    ssize_t (*show)(struct class_device *cls, char *buf);
    ssize_t (*store)(struct class_device *cls, const char *buf,
                    size_t count);
};
```

```

CLASS_DEVICE_ATTR(name, mode, show, store);

int class_device_create_file(struct class_device *cls,
                           const struct class_device_attribute *attr);
void class_device_remove_file(struct class_device *cls,
                             const struct class_device_attribute *attr);

```

A default set of attributes, in the class's `class_dev_attrs` field, is created when the class device is registered; `class_device_create_file` may be used to create additional attributes. Attributes may also be added to class devices created with the `class_simple` interface.

Class interfaces

The class subsystem has an additional concept not found in other parts of the Linux device model. This mechanism is called an *interface*, but it is, perhaps, best thought of as a sort of trigger mechanism that can be used to get notification when devices enter or leave the class.

An interface is represented by:

```

struct class_interface {
    struct class *class;
    int (*add) (struct class_device *cd);
    void (*remove) (struct class_device *cd);
};

```

Interfaces can be registered and unregistered with:

```

int class_interface_register(struct class_interface *intf);
void class_interface_unregister(struct class_interface *intf);

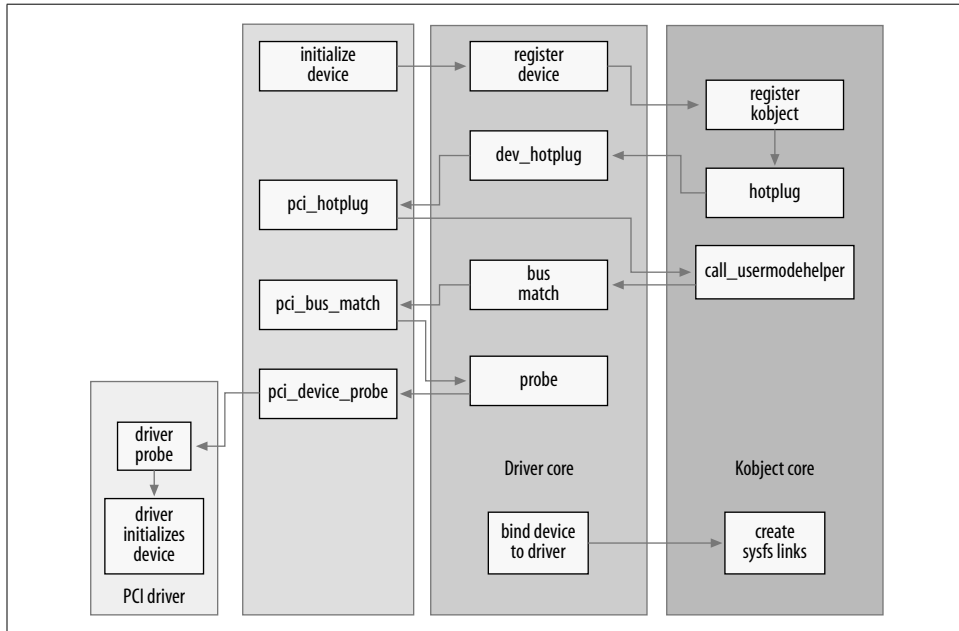
```

The functioning of an interface is straightforward. Whenever a class device is added to the class specified in the `class_interface` structure, the interface's `add` function is called. That function can perform any additional setup required for that device; this setup often takes the form of adding more attributes, but other applications are possible. When the device is removed from the class, the `remove` method is called to perform any required cleanup.

Multiple interfaces can be registered for a class.

Putting It All Together

To better understand what the driver model does, let us walk through the steps of a device's lifecycle within the kernel. We describe how the PCI subsystem interacts with the driver model, the basic concepts of how a driver is added and removed, and how a device is added and removed from the system. These details, while describing the PCI kernel code specifically, apply to all other subsystems that use the driver core to manage their drivers and devices.



and how to write a PCI driver, see Chapter 12). That structure contains a struct `device_driver` that is then initialized by the PCI core when the PCI driver is registered:

```
/* initialize common driver fields */
drv->driver.name = drv->name;
drv->driver.bus = &pci_bus_type;
drv->driver.probe = pci_device_probe;
drv->driver.remove = pci_device_remove;
drv->driver.kobj.ktype = &pci_driver_kobj_type;
```

This code sets up the bus for the driver to point to the `pci_bus_type` and points the *probe* and *remove* functions to point to functions within the PCI core. The *ktype* for the driver's kobject is set to the variable `pci_driver_kobj_type`, in order for the PCI driver's attribute files to work properly. Then the PCI core registers the PCI driver with the driver core:

```
/* register with core */
error = driver_register(&drv->driver);
```

The driver is now ready to be bound to any PCI devices it supports.

The PCI core, with help from the architecture-specific code that actually talks to the PCI bus, starts probing the PCI address space, looking for all PCI devices. When a PCI device is found, the PCI core creates a new variable in memory of type struct `pci_dev`. A portion of the struct `pci_dev` structure looks like the following:

```
struct pci_dev {
    /* ... */
    unsigned int    devfn;
    unsigned short vendor;
    unsigned short device;
    unsigned short subsystem_vendor;
    unsigned short subsystem_device;
    unsigned int    class;
    /* ... */
    struct pci_driver *driver;
    /* ... */
    struct device dev;
    /* ... */
};
```

The bus-specific fields of this PCI device are initialized by the PCI core (the `devfn`, `vendor`, `device`, and other fields), and the struct device variable's parent variable is set to the PCI bus device that this PCI device lives on. The bus variable is set to point at the `pci_bus_type` structure. Then the `name` and `bus_id` variables are set, depending on the name and ID that is read from the PCI device.

After the PCI device structure is initialized, the device is registered with the driver core with a call to:

```
device_register(&dev->dev);
```


Within the *device_register* function, the driver core initializes a number of the device's fields, registers the device's kobject with the kobject core (which causes a hotplug event to be generated, but we discuss that later in this chapter), and then adds the device to the list of devices that are held by the device's parent. This is done so that all devices can be walked in the proper order, always knowing where in the hierarchy of devices each one lives.

The device is then added to the bus-specific list of all devices, in this example, the *pci_bus_type* list. Then the list of all drivers that are registered with the bus is walked, and the *match* function of the bus is called for every driver, specifying this device. For the *pci_bus_type* bus, the *match* function was set to point to the *pci_bus_match* function by the PCI core before the device was submitted to the driver core.

The *pci_bus_match* function casts the struct device that was passed to it by the driver core, back into a struct *pci_dev*. It also casts the struct *device_driver* back into a struct *pci_driver* and then looks at the PCI device-specific information of the device and driver to see if the driver states that it can support this kind of device. If the match is not successful, the function returns 0 back to the driver core, and the driver core moves on to the next driver in its list.

If the match is successful, the function returns 1 back to the driver core. This causes the driver core to set the driver pointer in the struct device to point to this driver, and then it calls the *probe* function that is specified in the struct device driver.

Earlier, before the PCI driver was registered with the driver core, the probe variable was set to point at the *pci_device_probe* function. This function casts (yet again) the struct device back into a struct *pci_dev* and the struct driver that is set in the device back into a struct *pci_driver*. It again verifies that this driver states that it can support this device (which seems to be a redundant extra check for some unknown reason), increments the reference count of the device, and then calls the PCI driver's *probe* function with a pointer to the struct *pci_dev* structure it should bind to.

If the PCI driver's *probe* function determines that it can not handle this device for some reason, it returns a negative error value, which is propagated back to the driver core and causes it to continue looking through the list of drivers to match one up with this device. If the *probe* function can claim the device, it does all the initialization that it needs to do to handle the device properly, and then it returns 0 back up to the driver core. This causes the driver core to add the device to the list of all devices currently bound by this specific driver and creates a symlink within the driver's directory in *sysfs* to the device that it is now controlling. This symlink allows users to see exactly which devices are bound to which devices. This can be seen as:

```
$ tree /sys/bus/pci
/sys/bus/pci/
|-- devices
|   |-- 0000:00:00.0 -> ../../../../devices/pci0000:00/0000:00:00.0
|   |-- 0000:00:00.1 -> ../../../../devices/pci0000:00/0000:00:00.1
|   |-- 0000:00:00.2 -> ../../../../devices/pci0000:00/0000:00:00.2
```

```

| -- 0000:00:02.0 -> ../../../../devices/pci0000:00/0000:00:02.0
| -- 0000:00:04.0 -> ../../../../devices/pci0000:00/0000:00:04.0
| -- 0000:00:06.0 -> ../../../../devices/pci0000:00/0000:00:06.0
| -- 0000:00:07.0 -> ../../../../devices/pci0000:00/0000:00:07.0
| -- 0000:00:09.0 -> ../../../../devices/pci0000:00/0000:00:09.0
| -- 0000:00:09.1 -> ../../../../devices/pci0000:00/0000:00:09.1
| -- 0000:00:09.2 -> ../../../../devices/pci0000:00/0000:00:09.2
| -- 0000:00:0c.0 -> ../../../../devices/pci0000:00/0000:00:0c.0
| -- 0000:00:0f.0 -> ../../../../devices/pci0000:00/0000:00:0f.0
| -- 0000:00:10.0 -> ../../../../devices/pci0000:00/0000:00:10.0
| -- 0000:00:12.0 -> ../../../../devices/pci0000:00/0000:00:12.0
| -- 0000:00:13.0 -> ../../../../devices/pci0000:00/0000:00:13.0
| -- 0000:00:14.0 -> ../../../../devices/pci0000:00/0000:00:14.0
|-- drivers
| -- ALI15x3_IDE
|   |-- 0000:00:0f.0 -> ../../../../devices/pci0000:00/0000:00:0f.0
| -- ehci_hcd
|   |-- 0000:00:09.2 -> ../../../../devices/pci0000:00/0000:00:09.2
| -- ohci_hcd
|   |-- 0000:00:02.0 -> ../../../../devices/pci0000:00/0000:00:02.0
|   |-- 0000:00:09.0 -> ../../../../devices/pci0000:00/0000:00:09.0
|   |-- 0000:00:09.1 -> ../../../../devices/pci0000:00/0000:00:09.1
| -- orinoco_pci
|   |-- 0000:00:12.0 -> ../../../../devices/pci0000:00/0000:00:12.0
| -- radeonfb
|   |-- 0000:00:14.0 -> ../../../../devices/pci0000:00/0000:00:14.0
| -- serial
| -- trident
|   |-- 0000:00:04.0 -> ../../../../devices/pci0000:00/0000:00:04.0

```

Remove a Device

A PCI device can be removed from a system in a number of different ways. All Card-Bus devices are really PCI devices in a different physical form factor, and the kernel PCI core does not differentiate between them. Systems that allow the removal or addition of PCI devices while the machine is still running are becoming more popular, and Linux supports them. There is also a fake PCI Hotplug driver that allows developers to test to see if their PCI driver properly handles the removal of a device while the system is running. This module is called `fakephp` and causes the kernel to think the PCI device is gone, but it does not allow users to physically remove a PCI device from a system that does not have the proper hardware to do so. See the documentation with this driver for more information on how to use it to test your PCI drivers.

The PCI core exerts a lot less effort to remove a device than it does to add it. When a PCI device is to be removed, the `pci_remove_bus_device` function is called. This function does some PCI-specific cleanups and housekeeping, and then calls the `device_unregister` function with a pointer to the struct `pci_dev`'s struct `device` member.

In the *device_unregister* function, the driver core merely unlinks the sysfs files from the driver bound to the device (if there was one), removes the device from its internal list of devices, and calls *kobject_del* with a pointer to the struct *kobject* that is contained in the struct *device* structure. That function makes a hotplug call to user space stating that the kobject is now removed from the system, and then it deletes all sysfs files associated with the *kobject* and the sysfs directory itself that the *kobject* had originally created.

The *kobject_del* function also removes the *kobject* reference of the device itself. If that reference was the last one (meaning no user-space files were open for the sysfs entry of the device), then the *release* function for the PCI device itself, *pci_release_dev*, is called. That function merely frees up the memory that the struct *pci_dev* took up.

After this, all sysfs entries associated with the device are removed, and the memory associated with the device is released. The PCI device is now totally removed from the system.

Add a Driver

A PCI driver is added to the PCI core when it calls the *pci_register_driver* function. This function merely initializes the struct *device_driver* structure that is contained within the struct *pci_driver* structure, as previously mentioned in the section about adding a device. Then the PCI core calls the *driver_register* function in the driver core with a pointer to the struct *device_driver* structure contained in the struct *pci_driver* structure.

The *driver_register* function initializes a few locks in the struct *device_driver* structure, and then calls the *bus_add_driver* function. This function does the following steps:

- Looks up the bus that the driver is to be associated with. If this bus is not found, the function instantly returns.
- The driver's sysfs directory is created based on the name of the driver and the bus that it is associated with.
- The bus's internal lock is grabbed, and then all devices that have been registered with the bus are walked, and the match function is called for them, just like when a new device is added. If that match function succeeds, then the rest of the binding process occurs, as described in the previous section.

Remove a Driver

Removing a driver is a very simple action. For a PCI driver, the driver calls the *pci_unregister_driver* function. This function merely calls the driver core function *driver_unregister*, with a pointer to the struct *device_driver* portion of the struct *pci_driver* structure passed to it.

The *driver_unregister* function handles some basic housekeeping by cleaning up some sysfs attributes that were attached to the driver's entry in the sysfs tree. It then iterates over all devices that were attached to this driver and calls the *release* function for it. This happens exactly like the previously mentioned *release* function for when a device is removed from the system.

After all devices are unbound from the driver, the driver code does this unique bit of logic:

```
down(&drv->unload_sem);  
up(&drv->unload_sem);
```

This is done right before returning to the caller of the function. This lock is grabbed because the code needs to wait for all reference counts on this driver to be dropped to 0 before it is safe to return. This is needed because the *driver_unregister* function is most commonly called as the exit path of a module that is being unloaded. The module needs to remain in memory for as long as the driver is being referenced by devices and by waiting for this lock to be freed, this allows the kernel to know when it is safe to remove the driver from memory.

Hotplug

There are two different ways to view hotplugging. The kernel views hotplugging as an interaction between the hardware, the kernel, and the kernel driver. Users view hotplugging as the interaction between the kernel and user space through the program called */sbin/hotplug*. This program is called by the kernel when it wants to notify user space that some type of hotplug event has just happened within the kernel.

Dynamic Devices

The most commonly used meaning of the term “hotplug” happens when discussing the fact that most all computer systems can now handle devices appearing or disappearing while the system is powered on. This is very different from the computer systems of only a few years ago, where the programmers knew that they needed to scan for all devices only at boot time, and they never had to worry about their devices disappearing until the power was turned off to the whole machine. Now, with the advent of USB, CardBus, PCMCIA, IEEE1394, and PCI Hotplug controllers, the Linux kernel needs to be able to reliably run no matter what hardware is added or removed from the system. This places an added burden on the device driver author, as they must now always handle a device being suddenly ripped out from underneath them without any notice.

Each different bus type handles the loss of a device in a different way. For example, when a PCI, CardBus, or PCMCIA device is removed from the system, it is usually a while before the driver is notified of this action through its *remove* function. Before that happens, all reads from the PCI bus return all bits set. This means that drivers

need to always check the value of the data they read from the PCI bus and properly be able to handle a 0xff value.

An example of this can be seen in the *drivers/usb/host/ehci-hcd.c* driver, which is a PCI driver for a USB 2.0 (high-speed) controller card. It has the following code in its main handshake loop to detect if the controller card has been removed from the system:

```
result = readl(ptr);
if (result == ~(u32)0)    /* card removed */
    return -ENODEV;
```

For USB drivers, when the device that a USB driver is bound to is removed from the system, any pending urbs that were submitted to the device start failing with the error `-ENODEV`. The driver needs to recognize this error and properly clean up any pending I/O if it occurs.

Hotpluggable devices are not limited only to traditional devices such as mice, keyboards, and network cards. There are numerous systems that now support removal and addition of entire CPUs and memory sticks. Fortunately the Linux kernel properly handles the addition and removal of such core “system” devices so that individual device drivers do not need to pay attention to these things.

The /sbin/hotplug Utility

As alluded to earlier in this chapter, whenever a device is added or removed from the system, a “hotplug event” is generated. This means that the kernel calls the user-space program */sbin/hotplug*. This program is typically a very small bash script that merely passes execution on to a list of other programs that are placed in the */etc/hotplug.d/* directory tree. For most Linux distributions, this script looks like the following:

```
DIR="/etc/hotplug.d"
for I in "${DIR}/${1}/${*}.hotplug "${DIR}/default/*.hotplug ; do
    if [ -f $I ]; then
        test -x $I && $I $1 ;
    fi
done
exit 1
```

In other words, the script searches for all programs bearing a *.hotplug* suffix that might be interested in this event and invokes them, passing to them a number of different environment variables that have been set by the kernel. More details about how the */sbin/hotplug* script works can be found in the comments in the program and in the *hotplug(8)* manpage.

As mentioned previously, */sbin/hotplug* is called whenever a kobject is created or destroyed. The hotplug program is called with a single command-line argument providing a name for the event. The core kernel and specific subsystem involved also set a series of environment variables (described below) with information on what has just occurred. These variables are used by the hotplug programs to determine what has just happened in the kernel, and if there is any specific action that should take place.

The command-line argument passed to */sbin/hotplug* is the name associated with this hotplug event, as determined by the kset assigned to the kobject. This name can be set by a call to the *name* function that is part of the kset's *hotplug_ops* structure described earlier in this chapter; if that function is not present or never called, the name is that of the kset itself.

The default environment variables that are always set for the */sbin/hotplug* program are:

ACTION

The string *add* or *remove*, depending on whether the object in question was just created or destroyed.

DEVPATH

A directory path, within the *sysfs* filesystem, that points to the kobject that is being either created or destroyed. Note that the mount point of the *sysfs* filesystem is not added to this path, so it is up to the user-space program to determine that.

SEQNUM

The sequence number for this hotplug event. The sequence number is a 64-bit number that is incremented for every hotplug event that is generated. This allows user space to sort the hotplug events in the order in which the kernel generates them, as it is possible for a user-space program to be run out of order.

SUBSYSTEM

The same string passed as the command-line argument as described above.

A number of the different bus subsystems all add their own environment variables to the */sbin/hotplug* call, when devices associated with the bus are added or removed from the system. They do this in their *hotplug* callback that is specified in the struct *kset_hotplug_ops* assigned to their bus (as described in the section “Hotplug Operations”). This allows user space to be able to automatically load any necessary module that might be needed to control the device that has been found by the bus. Here is a list of the different bus types and what environment variables they add to the */sbin/hotplug* call.

IEEE1394 (FireWire)

Any devices on the IEEE1394 bus, also known as Firewire, have the */sbin/hotplug* parameter name and the `SUBSYSTEM` environment variable set to the value `ieee1394`. The `ieee1394` subsystem also always adds the following four environment variables:

`VENDOR_ID`

The 24-bit vendor ID for the IEEE1394 device

`MODEL_ID`

The 24-bit model ID for the IEEE1394 device

`GUID`

The 64-bit GUID for the device

`SPECIFIER_ID`

The 24-bit value specifying the owner of the protocol spec for this device

`VERSION`

The value that specifies the version of the protocol spec for this device

Networking

All network devices create a hotplug event when the device is registered or unregistered in the kernel. The */sbin/hotplug* call has the parameter name and the `SUBSYSTEM` environment variable set to the value `net`, and just adds the following environment variable:

`INTERFACE`

The name of the interface that has been registered or unregistered from the kernel. Examples of this are `lo` and `eth0`.

PCI

Any devices on the PCI bus have the parameter name and the `SUBSYSTEM` environment variable set to the value `pci`. The PCI subsystem also always adds the following four environment variables:

`PCI_CLASS`

The PCI class number for the device, in hex.

`PCI_ID`

The PCI vendor and device IDs for the device, in hex, combined in the format `vendor:device`.

`PCI_SUBSYS_ID`

The PCI subsystem vendor and subsystem device IDs, combined in the format `subsys_vendor:subsys_device`.

`PCI_SLOT_NAME`

The PCI slot “name” that is given to the device by the kernel. It is in the format `domain:bus:slot:function`. An example might be `0000:00:0d.0`.

Input

For all input devices (mice, keyboards, joysticks, etc.), a hotplug event is generated when the device is added and removed from the kernel. The `/sbin/hotplug` parameter and the `SUBSYSTEM` environment variable are set to the value `input`. The input subsystem also always adds the following environment variable:

PRODUCT

A multivalue string listing values in hex with no leading zeros. It is in the format `bustype:vendor:product:version`.

The following environment variables may be present, if the device supports it:

NAME

The name of the input device as given by the device.

PHYS

The device's physical address that the input subsystem gave to this device. It is supposed to be stable, depending on the bus position into which the device was plugged.

EV

KEY

REL

ABS

MSC

LED

SND

FF

These all come from the input device descriptor and are set to the appropriate values if the specific input device supports it.

USB

Any devices on the USB bus have the parameter name and the `SUBSYSTEM` environment variable set to the value `usb`. The USB subsystem also always adds the following environment variables:

PRODUCT

A string in the format `idVendor/idProduct/bcdDevice` that specifies those USB device-specific fields

TYPE

A string in the format `bDeviceClass/bDeviceSubClass/bDeviceProtocol` that specifies those USB device-specific fields

If the `bDeviceClass` field is set to 0, the following environment variable is also set:

INTERFACE

A string in the format `bInterfaceClass/bInterfaceSubClass/bInterfaceProtocol` that specifies those USB device-specific fields.

If the kernel build option, `CONFIG_USB_DEVICEFS`, which selects the `usbfs` filesystem to be built in the kernel, is selected, the following environment variable is also set:

DEVICE

A string that shows where in the `usbfs` filesystem the device is located. This string is in the format `/proc/bus/usb/USB_BUS_NUMBER/USB_DEVICE_NUMBER`, in which `USB_BUS_NUMBER` is the three-digit number of the USB bus that the device is on, and `USB_DEVICE_NUMBER` is the three-digit number that has been assigned by the kernel to that USB device.

SCSI

All SCSI devices create a hotplug event when the SCSI device is created or removed from the kernel. The `/sbin/hotplug` call has the parameter name and the `SUBSYSTEM` environment variable set to the value `scsi` for every SCSI device that is added or removed from the system. There are no additional environment variables added by the SCSI system, but it is mentioned here because there is a SCSI-specific user-space script that can determine what SCSI drivers (disk, tape, generic, etc.) should be loaded for the specified SCSI device.

Laptop docking stations

If a Plug-and-Play-supported laptop docking station is added or removed from the running Linux system (by inserting the laptop into the station, or removing it), a hotplug event is created. The `/sbin/hotplug` call has the parameter name and the `SUBSYSTEM` environment variable set to the value `dock`. No other environment variables are set.

S/390 and zSeries

On the S/390 architecture, the channel bus architecture supports a wide range of hardware, all of which generate `/sbin/hotplug` events when they are added or removed from the Linux virtual system. These devices all have the `/sbin/hotplug` parameter name and the `SUBSYSTEM` environment variable set to the value `dasd`. No other environment variables are set.

Using `/sbin/hotplug`

Now that the Linux kernel is calling `/sbin/hotplug` for every device added and removed from the kernel, a number of very useful tools have been created in user space that take advantage of this. Two of the most popular tools are the Linux Hotplug scripts and `udev`.

Linux hotplug scripts

The Linux hotplug scripts started out as the very first user of the `/sbin/hotplug` call. These scripts look at the different environment variables that the kernel sets to describe the device that was just discovered and then tries to find a kernel module that matches up with that device.

As has been described before, when a driver uses the `MODULE_DEVICE_TABLE` macro, the program, `depmod`, takes that information and creates the files located in `/lib/module/KERNEL_VERSION/modules.*map`. The `*` is different, depending on the bus type that the driver supports. Currently, the module map files are generated for drivers that work for devices that support the PCI, USB, IEEE1394, INPUT, ISAPNP, and CCW subsystems.

The hotplug scripts use these module map text files to determine what module to try to load to support the device that was recently discovered by the kernel. They load all modules and do not stop at the first match, in order to let the kernel work out what module works best. These scripts do not unload any modules when devices are removed. If they were to try to do that, they could accidentally shut down devices that were also controlled by the same driver of the device that was removed.

Note, now that the `modprobe` program can read the `MODULE_DEVICE_TABLE` information directly from the modules without the need of the module map files, the hotplug scripts might be reduced to a small wrapper around the `modprobe` program.

udev

One of the main reasons for creating the unified driver model in the kernel was to allow user space to manage the `/dev` tree in a dynamic fashion. This had previously been done in user space with the implementation of `devfs`, but that code base has slowly rotted away, due to a lack of an active maintainer and some unfixable core bugs. A number of kernel developers realized that if all device information was exported to user space, it could perform all the necessary management of the `/dev` tree.

`devfs` has some very fundamental flaws in its design. It requires every device driver to be modified to support it, and it requires that device driver to specify the name and location within the `/dev` tree where it is placed. It also does not properly handle dynamic major and minor numbers, and it does not allow user space to override the naming of a device in a simple manner, forcing the device naming policy to reside within the kernel and not in user space. Linux kernel developers really hate having policy within the kernel, and since the `devfs` naming policy does not follow the Linux Standard Base specification, it really bothers them.

As the Linux kernel started to be installed on huge servers, a lot of users ran into the problem of how to manage very large numbers of devices. Disk drive arrays of over 10,000 unique devices presented the very difficult task of ensuring that a specific disk

was always named with the same exact name, no matter where it was placed in the disk array or when it was discovered by the kernel. This same problem also plagued desktop users who tried to plug two USB printers into their system and then realized that they had no way of ensuring that the printer known as */dev/lpt0* would not change and be assigned to the other printer if the system was rebooted.

So, *udev* was created. It relies on all device information being exported to user space through *sysfs* and on being notified by */sbin/hotplug* that a device was added or removed. Policy decisions, such as what name to give a device, can be specified in user space, outside of the kernel. This ensures that the naming policy is removed from the kernel and allows a large amount of flexibility about the name of each device.

For more information on how to use *udev* and how to configure it, please see the documentation that comes included with the *udev* package in your distribution.

All that a device driver needs to do, for *udev* to work properly with it, is ensure that any major and minor numbers assigned to a device controlled by the driver are exported to user space through *sysfs*. For any driver that uses a subsystem to assign it a major and minor number, this is already done by the subsystem, and the driver doesn't have to do any work. Examples of subsystems that do this are the *tty*, *misc*, *usb*, *input*, *scsi*, *block*, *i2c*, *network*, and *frame buffer* subsystems. If your driver handles getting a major and minor number on its own, through a call to the *cdev_init* function or the older *register_chrdev* function, the driver needs to be modified in order for *udev* to work properly with it.

udev looks for a file called *dev* in the */class/* tree of *sysfs*, in order to determine what major and minor number is assigned to a specific device when it is called by the kernel through the */sbin/hotplug* interface. A device driver merely needs to create that file for every device it controls. The *class_simple* interface is usually the easiest way to do this.

As mentioned in the section “The *class_simple* Interface,” the first step in using the *class_simple* interface is to create a struct *class_simple* with a call to the *class_simple_create* function:

```
static struct class_simple *foo_class;
...
foo_class = class_simple_create(THIS_MODULE, "foo");
if (IS_ERR(foo_class)) {
    printk(KERN_ERR "Error creating foo class.\n");
    goto error;
}
```

This code creates a directory in *sysfs* in */sys/class/foo*.

Whenever a new device is found by your driver, and you assign it a minor number as described in Chapter 3, the driver should call the *class_simple_device_add* function:

```
class_simple_device_add(foo_class, MKDEV(FOO_MAJOR, minor), NULL, "foo%d", minor);
```

This code causes a subdirectory under `/sys/class/foo` to be created called `fooN`, where `N` is the minor number for this device. There is one file created in this directory, `dev`, which is exactly what `udev` needs in order to create a device node for your device.

When your driver is unbound from a device, and you give up the minor number that it was attached to, a call to `class_simple_device_remove` is needed to remove the sysfs entries for this device:

```
class_simple_device_remove(MKDEV(FOO_MAJOR, minor));
```

Later, when your entire driver is being shut down, a call to `class_simple_destroy` is needed to remove the class that you created originally with the call to `class_simple_create`:

```
class_simple_destroy(foo_class);
```

The `dev` file that is created by the call to `class_simple_device_add` consists of the major and minor numbers, separated by a `:` character. If your driver does not want to use the `class_simple` interface because you want to provide other files within the class directory for the subsystem, use the `print_dev_t` function to properly format the major and minor number for the specific device.

Dealing with Firmware

As a driver author, you may find yourself confronted with a device that must have firmware downloaded into it before it functions properly. The competition in many parts of the hardware market is so intense that even the cost of a bit of EEPROM for the device's controlling firmware is more than the manufacturer is willing to spend. So the firmware is distributed on a CD with the hardware, and the operating system is charged with conveying the firmware to the device itself.

You may be tempted to solve the firmware problem with a declaration like this:

```
static char my_firmware[] = { 0x34, 0x78, 0xa4, ... };
```

That approach is almost certainly a mistake, however. Coding firmware into a driver bloats the driver code, makes upgrading the firmware hard, and is very likely to run into licensing problems. It is highly unlikely that the vendor has released the firmware image under the GPL, so mixing it with GPL-licensed code is usually a mistake. For this reason, drivers containing wired-in firmware are unlikely to be accepted into the mainline kernel or included by Linux distributors.

The Kernel Firmware Interface

The proper solution is to obtain the firmware from user space when you need it. Please resist the temptation to try to open a file containing firmware directly from kernel space, however; that is an error-prone operation, and it puts policy (in the

form of a file name) into the kernel. Instead, the correct approach is to use the firmware interface, which was created just for this purpose:

```
#include <linux/firmware.h>
int request_firmware(const struct firmware **fw, char *name,
                    struct device *device);
```

A call to *request_firmware* requests that user space locate and provide a firmware image to the kernel; we look at the details of how it works in a moment. The name should identify the firmware that is desired; the normal usage is the name of the firmware file as provided by the vendor. Something like *my_firmware.bin* is typical. If the firmware is successfully loaded, the return value is 0 (otherwise the usual error code is returned), and the *fw* argument is pointed to one of these structures:

```
struct firmware {
    size_t size;
    u8 *data;
};
```

That structure contains the actual firmware, which can now be downloaded to the device. Be aware that this firmware is unchecked data from user space; you should apply any and all tests you can think of to convince yourself that it is a proper firmware image before sending it to the hardware. Device firmware usually contains identification strings, checksums, and so on; check them all before trusting the data.

After you have sent the firmware to the device, you should release the in-kernel structure with:

```
void release_firmware(struct firmware *fw);
```

Since *request_firmware* asks user space to help, it is guaranteed to sleep before returning. If your driver is not in a position to sleep when it must ask for firmware, the asynchronous alternative may be used:

```
int request_firmware_nowait(struct module *module,
                           char *name, struct device *device, void *context,
                           void (*cont)(const struct firmware *fw, void *context));
```

The additional arguments here are *module* (which will almost always be *THIS_MODULE*), *context* (a private data pointer that is not used by the firmware subsystem), and *cont*. If all goes well, *request_firmware_nowait* begins the firmware load process and returns 0. At some future time, *cont* will be called with the result of the load. If the firmware load fails for some reason, *fw* is NULL.

How It Works

The firmware subsystem works with sysfs and the hotplug mechanism. When a call is made to *request_firmware*, a new directory is created under */sys/class/firmware* using your device's name. That directory contains three attributes:

loading

This attribute should be set to one by the user-space process that is loading the firmware. When the load process is complete, it should be set to 0. Writing a value of -1 to loading aborts the firmware loading process.

data

data is a binary attribute that receives the firmware data itself. After setting loading, the user-space process should write the firmware to this attribute.

device

This attribute is a symbolic link to the associated entry under */sys/devices*.

Once the sysfs entries have been created, the kernel generates a hotplug event for your device. The environment passed to the hotplug handler includes a variable *FIRMWARE*, which is set to the name provided to *request_firmware*. The handler should locate the firmware file, and copy it into the kernel using the attributes provided. If the file cannot be found, the handler should set the *loading* attribute to -1.

If a firmware request is not serviced within 10 seconds, the kernel gives up and returns a failure status to the driver. That time-out period can be changed via the sysfs attribute */sys/class/firmware/timeout*.

Using the *request_firmware* interface allows you to distribute the device firmware with your driver. When properly integrated into the hotplug mechanism, the firmware loading subsystem allows devices to simply work “out of the box.” It is clearly the best way of handling the problem.

Please indulge us as we pass on one more warning, however: device firmware should not be distributed without the permission of the manufacturer. Many manufacturers will agree to license their firmware under reasonable terms when asked politely; some others can be less cooperative. Either way, copying and distributing their firmware without permission is a violation of copyright law and an invitation for trouble.

Quick Reference

Many functions have been introduced in this chapter; here is a quick summary of them all.

Kobjects

`#include <linux/kobject.h>`

The include file containing definitions for kobjects, related structures, and functions.

`void kobject_init(struct kobject *kobj);`

`int kobject_set_name(struct kobject *kobj, const char *format, ...);`

Functions for kobject initialization.

`struct kobject *kobject_get(struct kobject *kobj);`

`void kobject_put(struct kobject *kobj);`

Functions that manage reference counts for kobjects.

`struct kobj_type;`

`struct kobj_type *get_ktype(struct kobject *kobj);`

Represents the type of structure within which a kobject is embedded. Use *get_ktype* to get the *kobj_type* associated with a given kobject.

`int kobject_add(struct kobject *kobj);`

`extern int kobject_register(struct kobject *kobj);`

`void kobject_del(struct kobject *kobj);`

`void kobject_unregister(struct kobject *kobj);`

kobject_add adds a kobject to the system, handling kset membership, sysfs representation, and hotplug event generation. *kobject_register* is a convenience function that combines *kobject_init* and *kobject_add*. Use *kobject_del* to remove a kobject or *kobject_unregister*, which combines *kobject_del* and *kobject_put*.

`void kset_init(struct kset *kset);`

`int kset_add(struct kset *kset);`

`int kset_register(struct kset *kset);`

`void kset_unregister(struct kset *kset);`

Initialization and registration functions for ksets.

`decl_subsys(name, type, hotplug_ops);`

A macro that makes it easier to declare subsystems.

`void subsystem_init(struct subsystem *subsys);`

`int subsystem_register(struct subsystem *subsys);`

`void subsystem_unregister(struct subsystem *subsys);`

`struct subsystem *subsys_get(struct subsystem *subsys)`

`void subsys_put(struct subsystem *subsys);`

Operations on subsystems.

Sysfs Operations

`#include <linux/sysfs.h>`

The include file containing declarations for sysfs.

```
int sysfs_create_file(struct kobject *kobj, struct attribute *attr);
int sysfs_remove_file(struct kobject *kobj, struct attribute *attr);
int sysfs_create_bin_file(struct kobject *kobj, struct bin_attribute *attr);
int sysfs_remove_bin_file(struct kobject *kobj, struct bin_attribute *attr);
int sysfs_create_link(struct kobject *kobj, struct kobject *target, char
                    *name);
void sysfs_remove_link(struct kobject *kobj, char *name);
```

Functions for creating and removing attribute files associated with a kobject.

Buses, Devices, and Drivers

```
int bus_register(struct bus_type *bus);
void bus_unregister(struct bus_type *bus);
```

Functions that perform registration and unregistration of buses in the device model.

```
int bus_for_each_dev(struct bus_type *bus, struct device *start, void *data,
                    int (*fn)(struct device *, void *));
int bus_for_each_drv(struct bus_type *bus, struct device_driver *start, void
                    *data, int (*fn)(struct device_driver *, void *));
```

Functions that iterate over each of the devices and drivers, respectively, that are attached to the given bus.

```
BUS_ATTR(name, mode, show, store);
int bus_create_file(struct bus_type *bus, struct bus_attribute *attr);
void bus_remove_file(struct bus_type *bus, struct bus_attribute *attr);
```

The *BUS_ATTR* macro may be used to declare a *bus_attribute* structure, which may then be added and removed with the above two functions.

```
int device_register(struct device *dev);
void device_unregister(struct device *dev);
```

Functions that handle device registration.

```
DEVICE_ATTR(name, mode, show, store);
int device_create_file(struct device *device, struct device_attribute *entry);
void device_remove_file(struct device *dev, struct device_attribute *attr);
```

Macros and functions that deal with device attributes.


```
int driver_register(struct device_driver *drv);
void driver_unregister(struct device_driver *drv);
```

Functions that register and unregister a device driver.

```
DRIVER_ATTR(name, mode, show, store);
int driver_create_file(struct device_driver *drv, struct driver_attribute
                      *attr);
void driver_remove_file(struct device_driver *drv, struct driver_attribute
                       *attr);
```

Macros and functions that manage driver attributes.

Classes

```
struct class_simple *class_simple_create(struct module *owner, char *name);
void class_simple_destroy(struct class_simple *cs);
struct class_device *class_simple_device_add(struct class_simple *cs, dev_t
      devnum, struct device *device, const char *fmt, ...);
void class_simple_device_remove(dev_t dev);
int class_simple_set_hotplug(struct class_simple *cs, int (*hotplug)(struct
      class_device *dev, char **envp, int num_envp, char *buffer, int
      buffer_size));
```

Functions that implement the class_simple interface; they manage simple class entries containing a dev attribute and little else.

```
int class_register(struct class *cls);
void class_unregister(struct class *cls);
```

Registration and unregistration of classes.

```
CLASS_ATTR(name, mode, show, store);
int class_create_file(struct class *cls, const struct class_attribute *attr);
void class_remove_file(struct class *cls, const struct class_attribute *attr);
```

The usual macros and functions for dealing with class attributes.

```
int class_device_register(struct class_device *cd);
void class_device_unregister(struct class_device *cd);
int class_device_rename(struct class_device *cd, char *new_name);
CLASS_DEVICE_ATTR(name, mode, show, store);
int class_device_create_file(struct class_device *cls, const struct
      class_device_attribute *attr);
void class_device_remove_file(struct class_device *cls, const struct
      class_device_attribute *attr);
```

Functions and macros that implement the class device interface.

```
int class_interface_register(struct class_interface *intf);
void class_interface_unregister(struct class_interface *intf);
```

Functions that add an interface to a class (or remove it).

Firmware

```
#include <linux/firmware.h>
int request_firmware(const struct firmware **fw, char *name, struct device
                    *device);
int request_firmware_nowait(struct module *module, char *name, struct device
    *device, void *context, void (*cont)(const struct firmware *fw, void
    *context));
void release_firmware(struct firmware *fw);
```

Functions that implement the kernel firmware-loading interface.