

# Advanced Euclidean algorithm

While the "normal" Euclidean algorithm simply finds the greatest common divisor of two numbers  $a$  and  $b$ , extended Euclidean algorithm finds the GCD also factors in addition to  $x$ , and  $y$  such that:

$$a \cdot x + b \cdot y = \gcd(a, b).$$

It finds the coefficients by which the GCD of two numbers expressed in terms of the numbers themselves.

## Algorithm

Make the calculation of these coefficients in the Euclidean algorithm is simple enough to derive formulas by which they change during the transition from couple  $(a, b)$  to couple  $(b \% a, a)$  (percent sign we denote modulo arithmetic).

Thus, suppose we have found a solution  $(x_1, y_1)$  of the problem for the new couple  $(b \% a, a)$ :

$$(b \% a) \cdot x_1 + a \cdot y_1 = g,$$

and want to get a solution  $(x, y)$  for our couples  $(a, b)$ :

$$a \cdot x + b \cdot y = g.$$

To do this, transform the quantity  $b \% a$ :

$$b \% a = b - \left\lfloor \frac{b}{a} \right\rfloor \cdot a.$$

Substituting this into the above expression  $x_1$  and  $y_1$  obtain:

$$g = (b \% a) \cdot x_1 + a \cdot y_1 = \left( b - \left\lfloor \frac{b}{a} \right\rfloor \cdot a \right) \cdot x_1 + a \cdot y_1,$$

and performing regrouping terms, we obtain:

$$g = b \cdot x_1 + a \cdot \left( y_1 - \left\lfloor \frac{b}{a} \right\rfloor \cdot x_1 \right).$$

Comparing this with the original expression of the unknown  $x$ , and  $y$  we obtain the desired expression:

$$\begin{cases} x = y_1 - \left\lfloor \frac{b}{a} \right\rfloor \cdot x_1, \\ y = x_1. \end{cases}$$

## Implementation

```
int gcd (int a, int b, int & x, int & y) {  
    if (a == 0) {  
        x = 0; y = 1;  
        return b;  
    }  
    int x1, y1;  
    int d = gcd (b%a, a, x1, y1);  
    x = y1 - (b / a) * x1;  
    y = x1;  
    return d;  
}
```

This is a recursive function, which still returns the GCD of the numbers  $a$  and  $b$ , but beyond that - as desired coefficients  $x$  and  $y$  parameters in the form functions are passed by reference.

Recursion base - case  $a = 0$ . Then the GCD is  $b$ , and obviously required coefficients  $x$  and  $y$  are 0 and 1, respectively. In other cases, the usual solution is working, and the coefficients are converted by the above formulas.

Advanced Euclidean algorithm in this implementation works correctly even for negative numbers.

## Literature

- Thomas Cormen, Charles Leiserson, Ronald Rivest, Clifford Stein. **Algorithms: Design and analysis of** [2005]