Efficient algorithms for factorization

Here are the implementation of several factorization algorithms, each of which individually can work as quickly or very slowly, but together they provide a very fast method.

Descriptions of these methods are, the more that they are well described on the Internet.

Method Pollard p-1

Probabilistic test gives fast response is not for all properties.

Returns either found divider, or 1 if the divisor was not found.

```
template <class T>
T pollard p 1 (T n)
       // Algorithm parameters significantly affect the performance and the quality of
search
       const T b = 13;
       const T q [] = {2, 3, 5, 7, 11, 13};
       // Several attempts algorithm
       T a = 5\% n;
       for (int j = 0; j < 10; j + +)
              // Look for is a, which is relatively prime to n
              while (\gcd(a, n)! = 1)
                      mulmod (a, a, n);
                      a + = 3;
                      a\% = n;
              }
              // Calculate a ^ M
              for (size t = 0; i < size of q / size of q [0]; <math>i + +)
                      T qq = q [i];
                      T = T floor (log ((double) b) / log ((double) qq));
```

Pollard's method of "Po"

}

Probabilistic test gives fast response is not for all properties.

Returns either found divider, or 1 if the divisor was not found.

```
template <class T>
T pollard_rho (T n, unsigned iterations_count = 100000)
      Τ
             b0 = rand ()\% n,
             b1 = b0,
             g;
       mulmod (b1, b1, n);
      if (+ + b1 == n)
              b1 = 0;
      g = gcd (abs (b1 - b0), n);
      for (unsigned count = 0; count <iterations count && (g == 1 | | g == n); count + +)
      {
             mulmod (b0, b0, n);
             if (+ + b0 == n)
                    b0 = 0;
             mulmod (b1, b1, n);
             + + B1;
             mulmod (b1, b1, n);
```

Bent method (modification of the method of Pollard "Po")

Probabilistic test gives fast response is not for all properties.

Returns either found divider, or 1 if the divisor was not found.

```
template <class T>
T pollard bent (T n, unsigned iterations count = 19)
       Т
              b0 = rand ()\% n,
              b1 = (b0 * b0 + 2)\% n,
              a = b1;
       for (unsigned iteration = 0, series len = 1; iteration <iterations count; iteration +
+, series_len * = 2)
       {
              T q = qcd (b1-b0, n);
              for (unsigned len = 0; len <series len && (g == 1 && g == n); len + +)
                     b1 = (b1 * b1 + 2)\% n;
                     g = gcd (abs (b1-b0), n);
              b0 = a;
              a = b1;
              if (g! = 1 \&\& g! = n)
                     return g;
       return 1;
}
```

Pollard's method of Monte Carlo

Probabilistic test gives fast response is not for all properties.

Returns either found divider, or 1 if the divisor was not found.

```
template <class T>
T pollard_monte_carlo (T n, unsigned m = 100)
      T b = rand()\% (m-2) 2 +;
       static std :: vector <T> primes;
       static T m max;
       if (primes.empty ())
              primes.push_back (3);
       if (m_max <m)
              m_max = m;
              for (T prime = 5; prime \leq m; + + + + prime)
                     bool is prime = true;
                     for (std :: vector <T> :: const_iterator iter = primes.begin (), end =
primes.end ();
                            iter! = end; + + iter)
                     {
                            T div = * iter;
                            if (div * div> prime)
                                   break;
                            if (prime% div == 0)
                                   is prime = false;
                                   break;
                            }
                     if (is prime)
                            primes.push back (prime);
              }
      }
       Tg = 1;
       for (size_t i = 0; i <pri>primes.size () && g == 1; i + +)
              T cur = primes [i];
              while (cur <= n)
                     cur * = primes [i];
```

Method Farm

This wide method, but it can be very slow if you have small number of divisors.

So run it costs only after all other methods.

```
template <class T, class T2>
T ferma (const T & n, T2 unused)
       T2
              x = sq root(n),
              y = 0,
              r = x * x - y * y - n;
       for (; ;)
              if (r == 0)
                      return x! = y? xy: x + y;
               else
                      if (r> 0)
                             r - = y + y + 1;
                             + + Y;
                      }
                      else
                             r + = x + x + 1;
                             + + X;
                      }
}
```

Trivial division

This basic method is useful to immediately process numbers with very small divisors.

```
template <class T, class T2>
T2 prime div trivial (const T & n, T2 m)
       // First check for trivial cases
       if (n == 2 | | n == 3)
               return 1;
       if (n <2)
               return 0;
       if (even (n))
              return 2;
       // Generate a simple 3 to m
       T2 pi;
       const vector <T2> & primes = get primes (m, pi);
       // Divide by all simple
       for (std :: vector <T2> :: const_iterator iter = primes.begin (), end = primes.end ();
              iter! = end; + + iter)
       {
              const T2 & div = * iter;
              if (div * div> n)
                      break;
              else
                      if (n\% \text{ div} == 0)
                             return div;
       }
       if (n <m * m)
              return 1;
       return 0;
```

Putting it all together

Combine all methods in a single function.

Also, the function uses the simplicity of the test, otherwise the factorization algorithms can work for very long. For example, you can select a test BPSW (read article BPSW).

```
template <class T, class T2>
void factorize (const T & n, std :: map <T,unsigned> & result, T2 unused)
       if (n == 1)
       else
              // Check if not a prime number
              if (isprime (n))
                     + + Result [n];
              else
                     // If the number is small enough that it expand the simple search
                     if (n <1000 * 1000)
                            T div = prime div trivial (n, 1000);
                            + + Result [div];
                            factorize (n / div, result, unused);
                     }
                     else
                     {
                            // Number of large, run it factorization algorithms
                            T div;
                            // First go fast algorithms Pollard
                            div = pollard monte carlo (n);
                            if (div == 1)
                                   div = pollard_rho (n);
                            if (div == 1)
                                   div = pollard p 1 (n);
                            if (div == 1)
                                   div = pollard bent (n);
                            // Need to run 100% algorithm Farm
                            if (div == 1)
                                   div = ferma (n, unused);
                            // Recursively Point Multipliers
                            factorize (div, result, unused);
                            factorize (n / div, result, unused);
                     }
```

Application

Download [5 KB] source program that uses all of these methods and test factorization BPSW on simplicity.