University of Vienna, Sommersemester 2015

INTEROPERABILITY - Information Integration-

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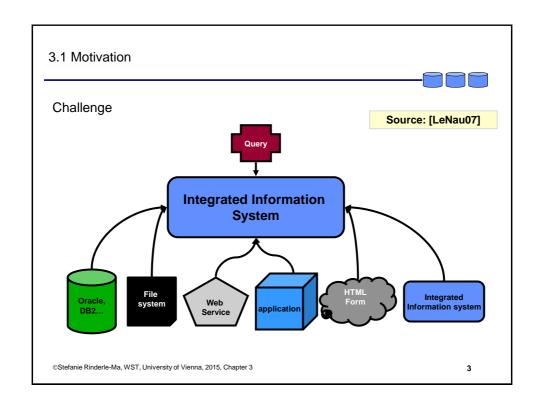
Outline

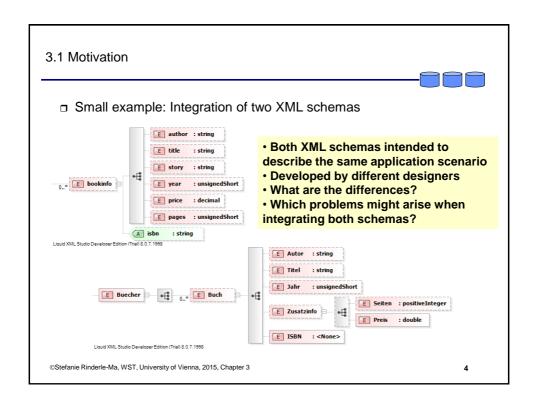


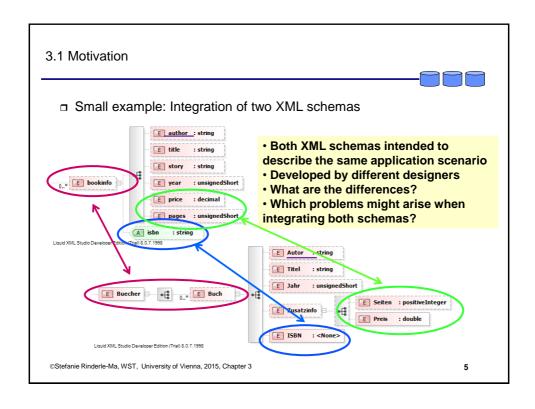
- 3.1 Motivation
- 3.2 Data Extraction
- 3.3 Schema and data integration
- 3.4 Summary and outlook

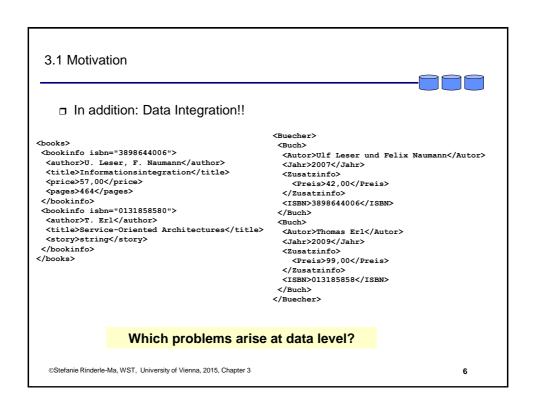
References

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3.1 Motivation □ In addition: Data Integration!! Different data format h> <books> Aut rollf Leser und Felix Naumann /Autor> <Jahr>2007</Jahr> <Zusatzinfo> Orais 42,00 Preis> (/Zusatzinto/ (/SBN>3898644006<//SBN> </bookinfo> <bookinfo isbn="0131858580"> </Buch> <author>T. Erl</author> <Buch> <title>Service-Oriented Architectures</title> <Autor>Thomas Erl</Autor> <story>string</story> <Jahr>2009</Jahr> </bookinfo> <Zusatzinfo> <Preis>99,00</Preis> </Zusatzinfo> <ISBN>013185858</ISBN> </Ruch> </Buecher> ☐ These structural problems can be SOMEHOW solved. □ Even harder: semantical problems → Example? ©Stefanie Rinderle-Ma, WST, University of Vienna, 2015, Chapter 3 7

3.1 Motivation Before schema and data integration Where is the data extracted? How can the data be extracted? Typically: heterogeneous data sources Examples: Databases File systems XML documents Legacy applications ETL process (Extraction → Transformation → Loading) Typically applied in Data Warehouse Systems (DWH) Goal of DWH: Integration and Analysis

3.1 Motivation



Teaching objectives:

- Understanding the challenges of information integration
- □ Learning different techniques for schema and data integration
- □ Particularly being able to distinguish different terminology, for example
 - o Data cleasing
 - o Data scrubbing
 - o Data fusion
- □ Schema mapping
- Schema matching
- □ Understanding "Pre-integration" process (→ ETL)
- Getting insights into specific IS application by means of Data Warehouses

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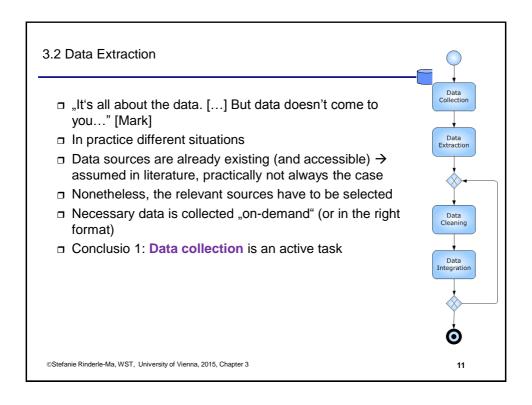
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Outline



- 3.1 Motivation
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Conclusio 1: Data collection is an active task

- □ Identification of relevant data sources
- □ Clarification of issues such as data access (particularly, if external data sources are to be accessed)

Use Case 1: Patient treatment processes

- ☐ EBMC2 project: http://ebmc2.univie.ac.at/
 - o Co-funded by University of Vienna and Medical University of Vienna
 - o Formalizing medical guidelines for skin cancer treatment
 - o Mining and analysis of real-world treatment processes
 - o In particular regarding their compliance with the guidelines

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Balance between:

- □ What data sources do we need (to fulfill a certain analysis goal) and
- □ Which data sources are actually available and accessible (privacy, data ownership, data access costs, etc.)

Available data sources [Dunkl12]:

- □ detailed data collection of clinical Cutaneous Melanoma (CM) stage IV protocols (Stage IV Melanoma Database, **S4MDB**, for short)
- administrative data of the Main Association of Austrian Social Security Institutions comprising a billing-oriented view of medical patient treatments (GAP-DRG)

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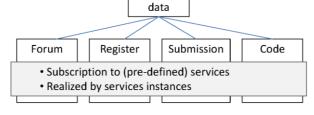
3.2 Data Extraction



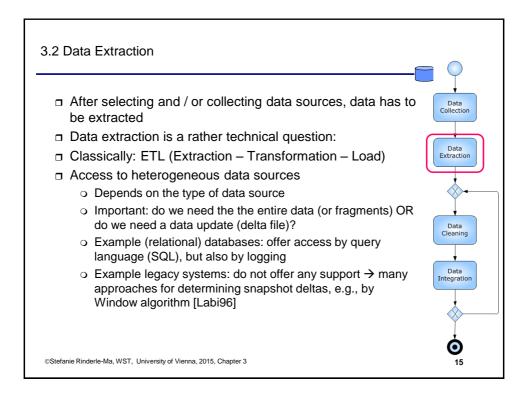
Use Case 2: Higher-Education Data (HEP)

- www.wst.univie.ac.at/communities/hep/
- □ Collected from service-oriented learning platform CEWebs [Dern04]

Figure taken from [Ly12]



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Differential snapshot → basic problem [LaGa96]

- □ Data extraction from legacy systems → snapshots, e.g., item list
- ☐ However: Possibly expensive (why?)
- Challenge:
 - o How can we detect the actual changes?
 - → Calculating delta between old and new snapshot

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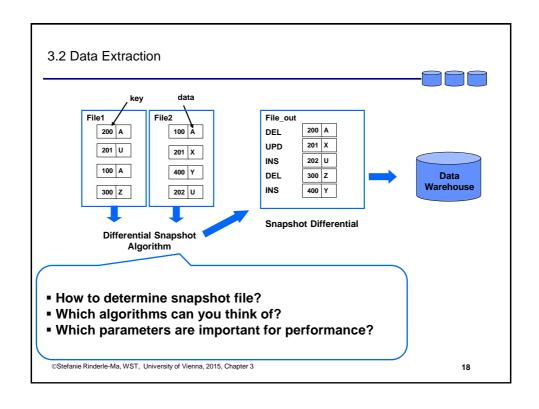


Differential snapshot problem - basic scenario

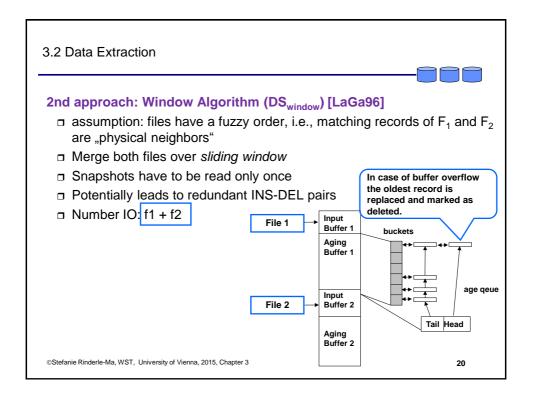
- □ Snapshot F extracted from source:
 - \circ File F: set of records (K, A₁, ..., A_n) where K denotes the key and A₁, ..., A_n denote the data fields of the record
- Let F₁ and F₂ be two snapshot files where F₁ has been generated before F₂
- \square Goal: determine the smallest set of UPDATE /INSERT / DELETE operation (within a so called **snapshot differential file**) that transforms F_1 into F_2 , formally:
 - o For two files F_1 and F_2 calculate smallest set $O = \{INS, DEL, UPD\}^*$ with $O(F_1) = F_2$
- $\quad \ \ \, \square \ \, O \text{ is not unique, e.g., } O_1 = \{INS(X), \varnothing, DEL(X)\} \equiv O_2 = \{\varnothing, \varnothing, \varnothing\}$
- \Box \rightarrow differential snapshot problem

[LaGa96]

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WINDOW ALGORITHM [LaGa96]

INPUT: F1, F2, n

OUTPUT: F_{out} /* the snapshot differential */

- (1) Input Buffer₁ \leftarrow Read n blocks from F₁
- (2) Input Buffer₂ \leftarrow Read n blocks form F_2
- (3) while ((Input Buffer $_1 \neq \emptyset$) and (Input Buffer $_2 \neq \emptyset$))
- (4) Match Input Buffer₁ against Input Buffer₂
- (5) Match Input $Buffer_1$ against $Aging Buffer_2$
- (6) Match Input Buffer₂ against Aging Buffer₁
- (7) Put contents of Input Buffer, to Aging Buffer,
- (8) Put contents of Input Buffer₂ to Aging Buffer₂
- (9) Input Buffer₁ \leftarrow Read n blocks from F_1
- (10) Input Buffer $_2 \leftarrow \text{Read n blocks from } \text{F}_2$
- (11) Report records in Aging Buffer_1 as deletes
- (12) Report records in Aging Buffer₂ as inserts

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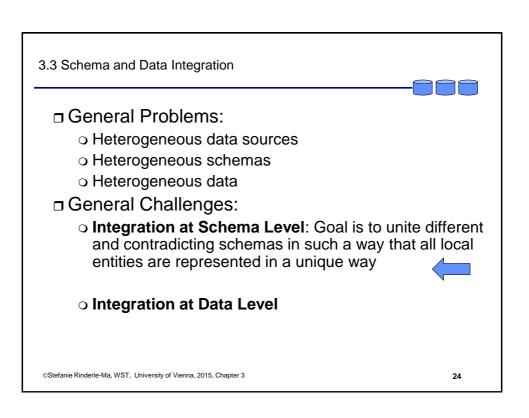
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3.2 Data Extraction

Operation pair	Effect	Example
INSERT(R)/ DELETE(R)	Neutral	INSERT(2, "HI"), DELETE(2, NULL)
UPDATE(R)/ DELETE(R)	delete	UPDATE(2, "HI"), DELETE(2, NULL) → DELETE(2, NULL)
DELETE(R)/ INSERT(R)	a) Same attribute v → neutral b) Different attribut value → UPDATE	INSERT(2,"HI") DELETE(2, NULL),INSERT(2, "BYE")
UPDATE(R)/ UPDATE(R)		
INSERT(R)/ UPDATE(R)		Redundant operations in DS ([RJR07])

Remark: the first entry in each tuple corresponds to the primary key of the underlying relation.

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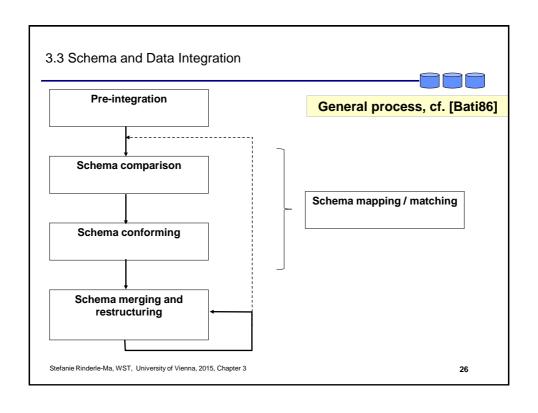




Characteristics of integrated schema – Quality Criteria for measuring the success of the integration process, cf. [Bati86]

- Completeness: all objects within the participating local schemas have to be present in the global schema as well
- □ Correctness:
 - o Semantic correctness (content)
 - o No contradictions
- □ Minimality: no redundancies, every object just once
- Understandability: Documentation of processes, transformation steps

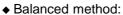
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Integration Strategies → [Bati86]

- ☐ Analysis and documentation of source schemas with respect to structure
- Determining integration strategy
 - o Binary approaches: in each step two schemas are merged
 - ♦ left deep tree:
 - o Local schemas are successively merged into developed schema
 - o Global schema is completed or extended with new entities



- o Level-wise iterative merging of two schemas
- o Build starting point for next integration cycle





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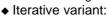
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3.3 Schema and Data Integration



Integration Strategies ctd. → [Bati86]

- o n-ary approaches
 - ♦ One-Shot strategy:
 - o All existing schemas are integrated in one step



- Arbitrary combination of schemas
- o Choice of integration strategy corresponds to an implicit preference (schema that are integrated in the beginning have higher influence than those that are brought in later)



☐ Find out about commonalities and differences between the affected schemas



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☐ Frequently used taxonomy for schema integration conflicts (4 conflict classes) [SPD92]:

Semantic conflicts

"The two designers do not perceive exactly the same set of real world objects, but instead they visualize overlapping sets (included or intersecting sets). [...] "Student" object class in one schema, while a more restrictive "CS-Student" object class in another schema."

Descriptive conflicts

"When describing related sets of real-world objects, two designers do not perceive exactly the same set of properties (e.g., because of their different interest in the many available pieces of information on car models in the real World). "

- ◆ Schema S1: Expensive_car (modelname, manufacturer, maximumspeed, price)
- ◆ Schema S2: Car_model (name, horsepower, fuelconsumption, price)

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3.3 Schema and Data Integration



- ☐ Frequently used taxonomy for schema integration conflicts (4 conflict classes) [SPD92] continued:
 - o Heterogeneity conflicts
 - "The designers use different data models, for example a relational one and an object-oriented one."
 - Structural conflicts
 - "Lastly, even if they use the same data model, they can choose different constructs to represent common real-world objects."
- Often: combination of different conflicts
- □ When looking at the different kinds of conflicts, obviously transformation of semantics and content is necessary!
 - o Semantics → identify and apply precise semantics
 - Structure → Schema integration
 - Data → Data integration and Data cleansing

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A few more words on descriptive conflicts:

- Same real-world objects within different systems are described in a different way
- □ Detecting and eliminating such semantic conflicts has often be done by the (human) experts (maybe partly supported by ontologies)
- Important and typical semantic conflicts:
 - o Synonyms
 - ◆ Example: vehicle ← → car
 - o Homonyms
 - Example: bank means either an institution for receiving, lending, exchanging, and safeguarding money or a long pile or heap; mass: a bank of earth; a bank of clouds (source: dictionary.com)
- Resolving synonyms and homonyms requires context knowledge (e.g., using "bank" in the context of a loan application scenario indicates the actual meaning)
- □ Furthermore: the real semantic can be described within by using meta data

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3.3 Schema and Data Integration

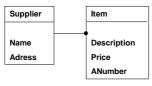


A few more words on Heterogeneity conflicts

- □ For a discussion of different heterogeneity conflicts and their reasons see [Lese2007]
- Different data models of the schemas to integrate
 - Examples: Network data model, hierarchical data model, OO data model, relational data model
 - o Using different modeling constructs
 - o Implying structural conflicts
- □ Transformation in a common global data model

Relational:

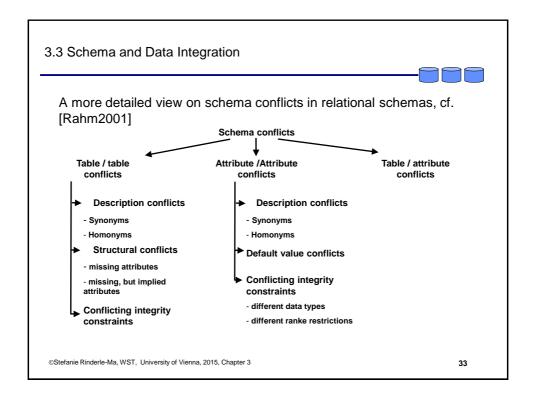
Object-oriented:



SID	Name	Adress

SID Description Price ANumber

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- ☐ Schema Integration Approaches cf. [Lese2007]
 - o Correspondence Assertions
 - o Generic Integration Model (GIM)
- □ Both approaches assume that all schemas are objectoriented
- □ [Lese2007], p. 122:
 - o Schema integration is an art
 - o Highly complex
 - Mostly manual

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Schema mapping, cf. [Rahm 2011]

- □ Part of schema comparison
- Goal: based on two schemas as input a mapping between elements of these schemas that are semantically corresponding should be found
- □ Formally [Bella2011, Chapter 4, pp. 82 ff.]:

Let S^* and T^* be two relational schemas. Then a mapping between S^* and T^* is defined as (S, T, m) where S is a relation in S^* and T is a relation in T^* and m is a set of attribute correspondences between S and T.

At instance level: Let D_S and D_T be instances of S and T. Then D_S and D_T **satisfy** mapping m if for \forall t_s in $D_S \exists$ t_t in D_T such that \forall attribute correspondences $(s,t) \in m$, the value of attribute s in t_s is the same value of attribute t in t_t .

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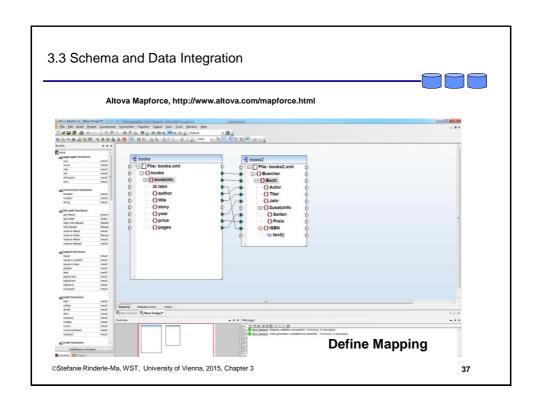
3.3 Schema and Data Integration

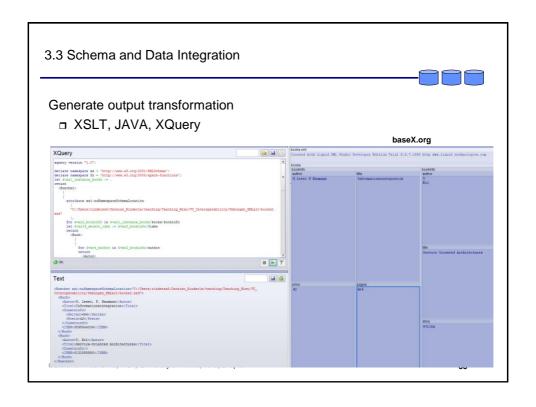


Schema mapping, cf. [Rahm 2011]

- □ Manual task → errorprone and tedious
- General algorithm:
 - o Given two schemas S and T (two relations) with attribute sets A and B
 - o Core idea:
 - ◆ Build cross product A × B between all attributes from A and B
 - For each pair calculate similarity
 - o E.g., regarding attribute name
 - o E.g., regarding stored data
 - ◆ Choose a mapping
 - ◆ Most similar pairs until threshold
 - ◆ In addition: consider constraints

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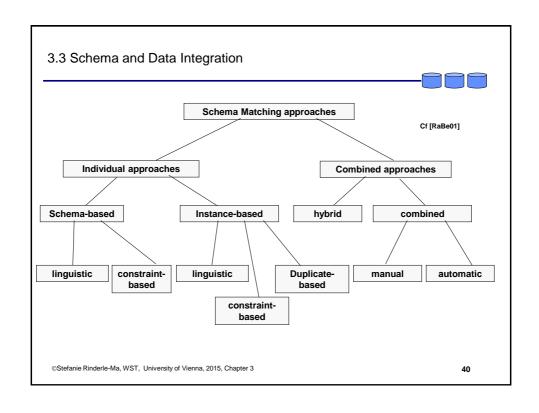




Schema matching, cf. [Lese2007]

- ☐ Trying to find mapping, i.e., set of correspondences between to schemas, automatically
- □ Automatic finding (supported by tools) is mostly confined to find similar labels
 - o String edit distance
 - o Homonyms / synonyms (→ontologies)
- □ Tools (commercial), cf. [Bella2011, p. 20]
 - o Microsoft Biztalk
 - o SAP Netweaver Process Integration
 - o Altova MapForce
- □ Tools (open source, research)
 - o COMA 3.0 (http://dbs.uni-leipzig.de/de/Research/coma.html)
 - Protègè (http://protege.stanford.edu/) and Prompt (http://protege.stanford.edu/plugins/prompt/prompt.html)
- □ Protègè, for example, provides support for ontology matching

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Schema Matching approaches ctd. (cf. [Lese2007]

- □ Label-based:
 - o Considers distances between the attribute labels
 - O Not always very helpful, e.g., d(year, Jahr) = 0
- □ Instance-based:
 - o Is based on (manually defined) properties on the database instances
 - o Example: length of the attribute name + attribute type
 - o Limited success
- ☐ Structure-based, take into consideration:
 - o Belonging of attributes to relations
 - Neighborhood of attributes
 - Newer approaches: model managament → represent relations as graphs and transfer similarity calculations to graph

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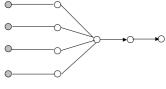
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3.3 Schema and Data Integration



Schema merging and restructuring

- □ The adapted schemas are overlapped and merged such that a global schema is created
- Main steps:
 - o Overlapping of adapted schemas
 - Quality tests along quality criteria (completeness, correctness, minimality, understandability)
 - o Further transformations of resulting schemas



Integration Transformation(s)

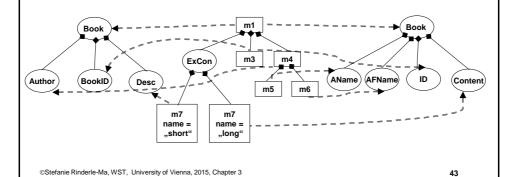
Cf. [BaLe86]

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- ☐ Fundamentals on schema merging have been proposed in [Bune1992]
- □ As for schema matching, newer approaches use model management, i.e., graph representations, as basis for the merge
- □ [Pott2003]: example adapted from [Bella2011, p. 240]



3.3 Schema and Data Integration



- □ Approaches based on 3-Ways Merge
- ☐ Typically applied in version control systems in software development, e.g., [Mens2002], e.g., SmartGit
- ☐ Prerequisite: merging two schemas where common ancestor of both schemas is known

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Remember:

- □ General Problems:
 - o Heterogeneous data sources
 - o Heterogeneous schemas
 - Heterogeneous data
- □ General Challenges:
 - Integration at Schema Level: Goal is to unite different and contradicting schemas in such a way that all local entities are represented in a unique way
 - o Integration at Data Level



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3.3 Schema and Data Integration

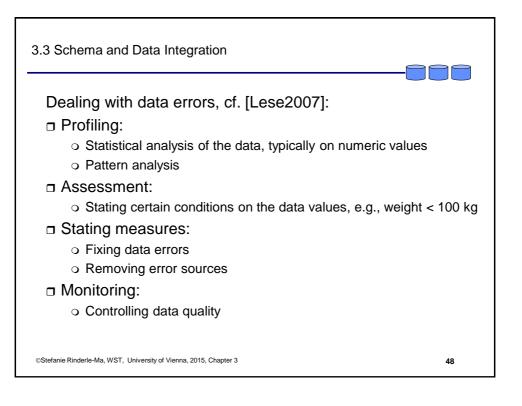


Problems with data (cf. [Lese2007])

- Data errors
 - o Different formats
 - o Errors (e.g., typos)
 - o Inconsistencies (e.g., zip code does not match city)
- Duplicates
- Data quality
 - o Credibility
 - o Relevance
- □ Completeness
 - o Are all real world objects considered?
 - o Do all attributes have values?

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3.3 Schema and Data Integration Data cleaning (cf. [Lese2007]) □ Simple data errors: can be detected by analyzing one tuple □ Difficult data errors: can only be detected by analyzing several tuples **Data Quality Problems** Multi Source Problems Single Source Problems Schema Level Instance Level Schema Level Instance Level (Heterogeneous (overlapping, (Missing integrity (Data entry errors) Data models contradicting, Constraints, poor → Misspellings and schema conflicting data) schema design) → redundancies → Missing integrity designs) → Referential \rightarrow contradictory > naming conflicts constraints integrity values → structural → Poor schema → uniqueness conflicts design





Data normalization (cf. [Lese2007])

- □ (De-)capitalization
- □ Abbreviations / spelling: Str., street, Straße ←→ strasse, ...
- □ Stemming
- Names
- □ Formats:
 - o Date: 18 February 2005, 18.02.2005, 2/18/05
 - o Coding: 1: female, 2: male
 - o Preciseness, field length, digits
 - o Scales: grades, temperature, currency, etc.
- ☐ Good support by commercial systems (SQL)
- □ In addition: outlier detection, detection of duplicates

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Outline



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3.4 Summary & Outlook



- Dealing with data is THE prerequisite for many applications
- Often a complex and expensive task
- Calculate enough time and manpower
- □ Include the domain experts
- □ Document every step of the integration process (→ tool support)
- ☐ Always keep an eye on your analysis goals
- □ Also keep in mind maintenance issues (your data sources might be changing!)

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