The Minimum Dilation Triangulation Problem

Introduction and Theoretical Aspects

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October 29, 2019







Triangulations

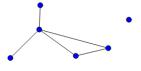
Let $S \subset \mathbb{R}^2$ finite.

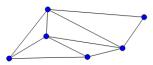
Definition

A planar network on S is a set of line segments with endpoints in S, where no two segments intersect nontrivially (they may only share an endpoint).

Definition

A triangulation of S is a planar network which is maximal for inclusion.





How good is a triangulation?

Let T be a triangulation of S. For $p, q \in S$, write $d_T(p, q)$ for the minimal possible sum of euclidean lengths of edges forming a path from p to q (i.e. the distance in the weighted graph (S, T)).

Diameter:

$$\operatorname{diam}(T) := \max_{p,q \in S} d_T(p,q) \quad \left(\text{or } \frac{\max_{p,q \in S} d_T(p,q)}{\max_{p,q \in S} \|p-q\|}, \cdots \right)$$

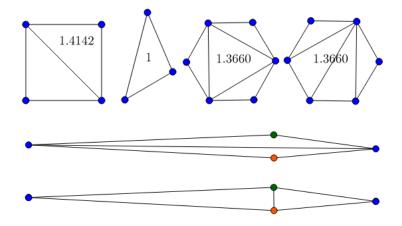
Dilation:

$$\operatorname{dil}(\mathcal{T}) := \max_{p
eq q \in \mathcal{S}} \left. rac{d_{\mathcal{T}}(p,q)}{\|p-q\|}
ight| \in [1,+\infty)$$

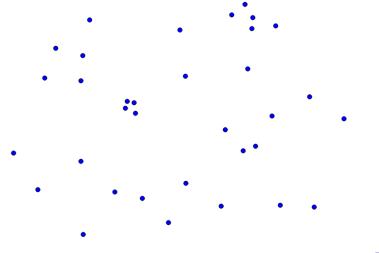
Goal

Find T such that dil(T) is minimal.

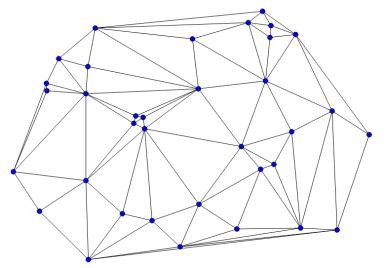
Examples



Can you always find a triangulation with bounded dilation?



Can you always find a triangulation with bounded dilation?



General upper bounds

Definition

The dilation of a finite set $S \subset \mathbb{R}^2$ is defined as

$$\operatorname{dil}(S) := \min_{T \in \mathcal{T}} \operatorname{dil}(T)$$

where \mathcal{T} is the set of all triangulations of S.

Theorem (1989; Chew)

For every finite $S \subset \mathbb{R}^2$,

$$\operatorname{dil}(S) \leq 2.$$

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For every finite $S \subset \mathbb{R}^2$,

$$\operatorname{dil}(S) \leq 2.$$

Theorem (2011; Xia)

For every finite $S \subset \mathbb{R}^2$,

$$dil(S) \leq 1.998$$
.

What is the true bound?

So, we know that every S has a triangulation with dilation ≤ 1.998 .

Proposition (Trivial)

If S is the set of vertices of a square, then every triangulation of S has dilation at least $\sqrt{2} \approx 1.4142$.

Proposition (2004; Multzer)

If S is the set of vertices of a regular 21-gon, then every triangulation of S has dilation at least 1.4161.

Proposition (2016; Dumitrescu)

If S is the set of vertices of a regular 23-gon, then every triangulation of S has dilation at least 1.4308.

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Proposition (2019; Galant, P.)

If S is the set of vertices of a regular 53-gon, then every triangulation of S has dilation at least 1.4336.

Convex case

Convex polygons seem to be special for the study of dilation:

- The known examples S with largest dilation are convex (even regular) polygons
- General algorithms often have much lower performance with convex polygons than with other (e.g. random) point sets

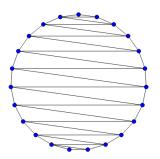


Figure: Output of the approximation algorithm based on Delaunay triangulations. This algorithm is proven to always output a triangulation with dilation less than 1.998. But in this case, the output is actually the worst triangulation (dilation $\approx \pi/2$, the actual answer is ≈ 1.43)

Upper Bounds in the convex case

Theorem (2011; Xia)

For every finite $S \subset \mathbb{R}^2$,

$$\operatorname{dil}(S) \leq 1.998.$$

Theorem (2016; Bose et al.)

For every finite $S \subset \mathbb{R}^2$ in convex position,

$$\operatorname{dil}(S) \leq 1.88.$$

Conversely, there is a set S in convex position with $dil(S) \ge 1.4336$.

Dilation of regular polygons

Let c_n be the dilation of the regular n-gon.

n	Cn	n	C _n	n	Cn	n	Cn
4	1.4142	12	1.3836	20	1.4142	28	1.4147
5	1.2360	13	1.3912	21	1.4161	29	1.4199
6	1.3660	14	1.4053	<u>22</u>	<u>1.4047</u>	30	1.4236
7	1.3351	15	1.4089	<u>23</u>	<u>1.4308</u>	31	1.4120
8	1.4142	16	1.4092	<u>24</u>	<u>1.4013</u>	32	1.4161
9	1.3472	17	1.4084	25	1.4050	33	1.4185
10	1.3968	18	1.3816	26	1.4169	34	1.4167
11	1.3770	19	1.4098	27	1.4185	35	1.4213

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Theorem (2019; P.)

The sequence c_n converges.

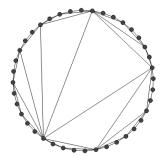
Proof sketch.

Let $\Delta := \liminf c_n$. Let $\varepsilon > 0$. It is sufficient to show that $c_n \le \Delta + \varepsilon$ for all n sufficiently large.

By definition of Δ , there exists a m for which $c_m \leq \Delta + \varepsilon/2$. Let T_m be an optimal triangulation of the regular m-gon (i.e. $\operatorname{dil}(T_m) = c_m$). Let n_0 be much larger than m. Let $n > n_0$.

We construct a triangulation T_n of the regular n-gon which is as similar to T_m as possible. It is possible to prove that $\operatorname{dil}(T_n) \leq \operatorname{dil}(T_m) + \varepsilon/2$, which concludes the proof.





Generalisation

Theorem (2019; P.)

Let C be a strictly convex differentiable closed curve. There is a real number Δ , with the following property.

Let $(S_n)_{n\in\mathbb{N}}$ be a sequence of finite subsets of \mathcal{C} such that $\forall \varepsilon > 0$, S_n is an ε -covering of \mathcal{C} for all but finitely many n.

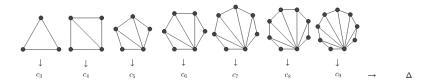
Then, $dil(S_n)$ converges to Δ .

Thus, the limit is independent of the sequence of polygons, and depends only on the curve $\mathcal{C}.$

Summary

To simplify, let S_n be the regular n-gon. For each n, there is at least one triangulation T_n of S_n with smallest possible dilation.

- We know that the sequence of dilations $(c_n) := (\operatorname{dil}(T_n))$ converges to some value Δ .
- ② What can be said about the triangulations T_n themselves?

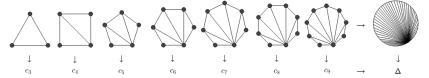


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Dream:



Infinite triangulations



Definition (Infinite triangulations)

- Let $\mathcal C$ be a strictly convex closed curve. A geometric network on $\mathcal C$ is a set of segments with both endpoints in $\mathcal C$, which are pairwise non-intersecting (except possibly at the endpoints).
- ullet A triangulation of ${\mathcal C}$ is a geometric network which is maximal for inclusion.

Definition (Distance in an infinite triangulation)

Let T be a triangulation of a strictly convex closed curve \mathcal{C} . The *distance* $d_T(p,q)$ between two points $p,q\in\mathcal{C}$ is the length of the shortest path^a from p to q which uses only edges of T and arcs of the curve \mathcal{C} .

^aInfimum over all finite paths.

Infinite triangulations

Definition (Dilation of an infinite triangulation)

Let T be a triangulation of a strictly convex closed curve C.

• The dilation of T isa

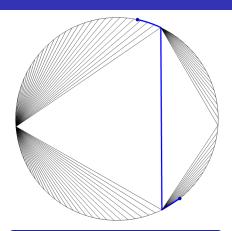
$$\operatorname{dil}(T) := \sup_{p,q \in \mathcal{C}} \frac{d_T(p,q)}{\|p-q\|}.$$

ullet The dilation of ${\mathcal C}$ is

$$\operatorname{dil}(\mathcal{C}) := \inf_{\mathcal{T} \in \mathcal{T}} \operatorname{dil}(\mathcal{T}),$$

where \mathcal{T} is the set of triangulations of \mathcal{C} .

^aAs C is compact, sup → max.

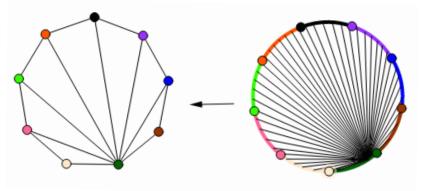


Question

Is there a triangulation T of C that achieves dil(T) = dil(C)?

The quotient operation

We need a way to transfer triangulations of "larger sets" (e.g. the circle or a regular polygon with many sides) to triangulations of "smaller sets" (e.g. a regular polygon with less sides).



The quotient operation

Lemma

Let T be a triangulation (of the circle or a regular N-gon with $N \ge n$). Let T' be the triangulation obtained after a quotient of T into n evenly spaced equivalence classes. Then

$$|\operatorname{dil}(T') - \operatorname{dil}(T)| = O\left(\frac{\log n}{n}\right).$$

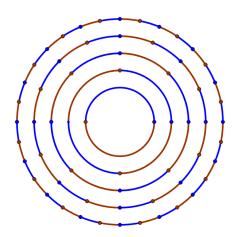
Corollary

Let T be a triangulation of the circle. Let $T_{\phi(n)}$ be a sequence of quotients of T into $\phi(n)$ evenly spaced equivalence classes (where $\phi: \mathbb{N} \to \mathbb{N}$ strictly increasing). Then $\lim \operatorname{dil}(T_{\phi(n)}) = \operatorname{dil}(T)$.

Profinite triangulations

Let T be a triangulation of the circle. We consider a special sequence of quotients of T. Let T_{2^n} be the quotient of T into 2^n classes as shown on the picture.

We have the following property: for all $n \ge k$, the quotient of T_{2^n} into the appropriate 2^k equivalence classes is exactly T_{2^k} .



Profinite triangulations

Definition

^a For $n \ge 1$, let \mathcal{T}_{2^n} be the set of triangulations of a regular 2^n -gon. We define

$$\varprojlim \mathcal{T}_{2^n}$$

to be the set of sequences of triangulations

$$(T_2, T_4, T_8, \ldots)$$

where

- $\bullet \ T_{2^n} \in \mathcal{T}_{2^n},$
- for all $n \ge k$, the quotient of T_{2^n} into the appropriate 2^k equivalence classes is exactly T_{2^k} .



^aThis definition can be generalised to any strictly convex closed curve.

Theorem (2019; P.)

Let $\mathcal C$ be a differentiable strictly convex closed curve and let $\mathcal T$ be the set of triangulations of $\mathcal C$.

• We have a natural bijection

$$\mathcal{T}\cong \varprojlim \mathcal{T}_{2^n}.$$

- ② The maps $\operatorname{dil}_1: \mathcal{T} \to \mathbb{R}, \ T \mapsto \operatorname{dil}(T) \ \text{and} \ \operatorname{dil}_2: \varprojlim \mathcal{T}_{2^n} \to \mathbb{R}, \ (\mathcal{T}_{2^n})_n \mapsto \lim_n \operatorname{dil}(\mathcal{T}_{2^n}) \ \text{agree under this bijection.}$
- **1** The space $\varprojlim \mathcal{T}_{2^n}$ is compact and dil_2 is continuous, so there is a triangulation $T \in \mathcal{T}$ such that $\operatorname{dil}(T) = \operatorname{dil}(\mathcal{C})$.

Corollary

Let c_n be the dilation of the regular n-gon, and $\Delta = \lim c_n$. Then $\Delta = \operatorname{dil}(\mathit{circle})$ and

$$|c_n - \Delta| = O\left(\frac{\log n}{n}\right).$$

Summary

For each n, let T_n be a triangulation of the regular n-gon with minimal dilation.



Dilations:

$$\lim_{n\to\infty}\operatorname{dil}(T_n)$$

exists.

• Triangulations: A subsequence $(T_{\phi(n)})$ of (T_n) converges to a triangulation T of the circle.¹

Furthermore:

$$\operatorname{dil}(T) = \operatorname{dil}(\operatorname{circle}) = \operatorname{dil}(T_n) + O\left(\frac{\log n}{n}\right).$$

¹The convergence is in the space \mathcal{T} (choose an arbitrary-lift for every $T_{\phi(n)}$).

Open problems

Question 1

How does the dilation of a curve change as the curve changes?

Question 2

Can we improve the bound 1.88 for the dilation of convex curves?

Question 3

For the regular polygons, is minimizing the dilation equivalent to minimizing the diameter of the graph? If not, how close are those two problems?

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Thank you for your attention!

