```
NÚMEROS GRANDES
#include <cmath> //log10,pow,floor
#include <cstring> //memset
#include <string> //string class
#include <sstream> //stringstream class
#include <algorithm> //max and min
using namespace std;
const int NUM=25; //number of elementes in array
typedef long long int base_t;
const base_t BASE=pow(10,floor(log10(2)*4*sizeof(base_t)));
const int zeros=log10(BASE);
struct longint{
    longint(long long int i=0) :used(0),sign((i>=0)?1:-1){
        memset(inner, 0, NUM*sizeof(base t));
        for(;i>=BASE;i=i/BASE)
            inner[used++]=i%BASE;
        inner[used++]=i;
    }
    longint operator+(longint b) const{
        if(sign!=b.sign) return b.sign*=-1,((sign<0)?*this-b: b-*this);</pre>
        b.used=min(max(b.used,used)+1 , NUM); base_t carry=0;
        for(int i=0;i<b.used;i++){</pre>
            carry=(b.inner[i]+=inner[i]+carry)/BASE;
            b.inner[i]%=BASE;
        while(b.used && !b.inner[b.used-1]) --b.used;
        return b;
    }
    longint operator-(longint b) const{
        longint const *m=this,*M=&b;
        if(used>b.used || used==b.used && inner[used-1]>b.inner[used-1])
            M=this, m=&b;
        if(b.sign!=sign) return b.sign*=-1,((*M)+(*m));
        longint res(*M); base_t carry=0;
        for(int i=0;i<=min(m->used,NUM-1);i++)
            carry=(res.inner[i]-=(m->inner[i]+carry))<0?res.inner[i]+=BASE,1:0;</pre>
        while(res.used && !res.inner[res.used-1]) --res.used;
        return res;
    }
    longint operator*(const longint& b) const{
        longint res;
        longint const *m, *M= (used>b.used)? (m=&b,this) : (m=this,&b);
        for(int i=0;i<m->used; i++)
            for(base t j=0, carry=0; j <= (min(M->used, NUM-i-1)); j++){
                carry=(res.inner[j+i]=(M->inner[j]*(m->inner[i]))+carry)/BASE;
                res.inner[j+i]%=BASE;
        res.used=min(M->used+m->used,NUM);
        while(res.used && !res.inner[res.used-1]) --res.used;
        res.sign=sign*b.sign;
        return res;
    }
    string str() const{
        stringstream ss;
        ss << sign*inner[used-1];</pre>
        for(int i=used-2;i>=0;i--){
            int cfr= inner[i] ? floor(log10(inner[i]))+1 : 1;
            for(int j=cfr;j<zeros;j++) ss << 0;</pre>
            ss << inner[i];</pre>
        return ss.str();
    }
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char sign; //sign bit
    int used; //number of base t in used
    base t inner[NUM]; //base array
};
longint operator+(long long int i, const longint& l){ return l+i; }
longint operator*(long long int i,const longint& l){ return l*i; }
longint operator-(long long int i,const longint& l){ return longint(i)-l;}
                            DPLETARIO
#include <iostream>
#include <vector>
#include <stack>
#include <utility>
using namespace std;
typedef pair<int,int> ii;
typedef vector<int> vi;
//calcula la primera LIS que vale. Tuneese como corresponda.
void dplis(){
        int n=8;
        int sec[]={-7,10,9,2,3,8,8,1};
        int tam[8]; //tamaño de la LIS que termina en i
        int prev[8]; //contiene el predecesor de i en la LIS que termina en i
        vector<int> lis;
        stack<int> st;
        int bestEnd=0;
        tam[0]=1;
        prev[0]=-1; //casos base;
        for(int i=1; i<n; i++){</pre>
                tam[i]=1;
                prev[i]=-1;
                for(int j=0; j<i; j++) if(sec[j]<sec[i] && tam[j]+1>tam[i]){
                        tam[i]=tam[j]+1;
                        prev[i]=j;
                }
        for(int k=0; k<n; k++) if(tam[k]>tam[bestEnd]) bestEnd=k;
        //deshacer el camino de la LIS empezando en bestEnd.
        while(bestEnd!=-1){
                st.push(bestEnd);
                bestEnd=prev[bestEnd];
        while(!st.empty()){
                cout<<st.top()<<" ";
                lis.push back(st.top());
                st.pop();
        }//enhorabuena, en lis tienes tu LIS
//calcula la primera LCS que vale. (como subsecuencia)
void lcd() {
        const int n=9, m=10;
        int a[]={1,4,5,3,7,9,4,7,0};
        int b[]={7,5,3,7,9,1,4,2,0,5};
        int tabla[100][100]; //tamaño de la maxima LCS entre las subcadenas que empiezan en i y j
        ii next[100][100]; //principio de la LCS en ambas subcadenas
        int i,j;
        memset(tabla,0, sizeof tabla);
        for(i=0; i< n; i++) tabla[i][m-1]=(a[i]==b[m-1])?1:0;
        for(j=0; j<m; j++) tabla[n-1][j]=(a[n-1]==b[j])?1:0;</pre>
        for(i=n-1; i>=0; i--) for(j=m-1; j>=0; j--) if(a[i]==b[j]){
                tabla[i][j]=tabla[i+1][j+1]+1;
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next[i][j]=make_pair(i+1,j+1);
        }else if(tabla[i][j+1]>tabla[i+1][j]){
                tabla[i][j]=tabla[i][j+1];
                next[i][j]=next[i][j+1];
        }else{
                tabla[i][j]=tabla[i+1][j];
                next[i][j]=next[i+1][j];
        cout<<tabla[0][0]<<endl;</pre>
        for(i=0; i<n; i++) {for(j=0; j<m; j++)</pre>
                cout<<'('<<next[i][j].first<<','<<next[i][j].second<<") ";
        }cout<<endl;}</pre>
        ii indice=(a[0]==b[0])?make_pair(0,0):next[0][0];
        while(indice.first || indice.second)
                cout<<indice.first<<" "<<indice.second<<endl;</pre>
                indice=next[indice.first][indice.second];
        }
}
//calcula la primera LCSS que vale. (subcadena, no subsecuencia)
//(en realidad, se resuelve en O((n+m)log(n+m)) usando un suffix array)
void dplcss()
{
        const int n=9, m=10;
        int a[]={1,4,5,3,7,9,4,7,0};
        int b[]={7,5,3,7,9,1,4,2,0,5};
        int tabla[100][100]; //tamaño de la coincidencia entre a[i] y b[j]
        int i,j;
        int bestA=0, bestB=0;
        for(i=0; i<n; i++) tabla[i][m-1]=(a[i]==b[m-1])?1:0;</pre>
        for(j=0; j<m; j++) tabla[n-1][j]=(a[n-1]==b[j])?1:0;</pre>
        for(i=n-1; i>=0; i--) for (j=m-1; j>=0; j--)
                tabla[i][j]=(a[i]==b[j])?tabla[i+1][j+1]+1:0;
        for(i=0; i< n; i++) for(j=0; j< m; j++) if(tabla[bestA][bestB]< tabla[i][j]){}
                bestA=i; bestB=j;
        cout<<bestA<<" "<<bestB<<" "<<tabla[bestA][bestB];</pre>
        //tu LCS empieza en bestA,bestB y tiene tamaño tabla[bestA][bestB];
}
                                  MATESLETARIO
#include <bitset>
                     // compact STL for Sieve, more efficient than vector<bool>!
#include <cmath>
#include <cstdio>
#include <map>
#include <vector>
using namespace std;
typedef long long ll;
typedef vector<int> vi;
typedef pair<int,int> ii;
typedef map<int, int> mii;
ll _sieve_size;
                      // 10^7 should be enough for most cases
bitset<10000010> bs;
           // compact list of primes in form of vector<int>
vi primes;
// llamar a este metodo nada mas empezar.
void sieve(ll upperbound) {
                                      // create list of primes in [0..upperbound]
  _sieve_size = upperbound + 1;
                                                    // add 1 to include upperbound
                                                               // set all bits to 1
  bs.set();
 bs[0] = bs[1] = 0;
                                                            // except index 0 and 1
  for (ll i = 2; i <= _sieve_size; i++) if (bs[i]) {</pre>
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// cross out multiples of i starting from i * i!
    for (ll j = i * i; j <= _sieve_size; j += i) bs[j] = 0;
primes.push_back((int)i); // also add this vector containing list of primes</pre>
                                                 // call this method in main method
bool isPrime(ll N) {
                                       // a good enough deterministic prime tester
  if (N <= _sieve_size) return bs[N];</pre>
                                                            // 0(1) for small primes
  for (int i = 0; i < (int)primes.size() && i*i <= N; i++)
    if (N % primes[i] == 0) return false;
  return true;
                                    // it takes longer time if N is a large prime!
                        // note: only work for N <= (last prime in vi "primes")^2</pre>
vi primeFactors(ll N) { // remember: vi is vector of integers, ll is long long
                                   // vi `primes' (generated by sieve) is optional
  vi factors;
 ll PF_idx = 0, PF = primes[PF_idx]; // using PF = 2, 3, 4, ..., is also ok while (N != 1 && (PF * PF <= N)) { // stop at sqrt(N), but N can get smaller
    while (N % PF == 0) { N /= PF; factors.push_back(PF); } // remove this PF
    PF = primes[++PF idx];
                                                            // only consider primes!
  if (N != 1) factors.push_back(N);
                                         // special case if N is actually a prime
                          // if pf exceeds 32-bit integer, you have to change vi
  return factors;
}
ll EulerPhi(ll N) {
  ll PF_idx = 0, PF = primes[PF_idx], ans = N;
                                                               // start from ans = N
 while (N != 1 \&\& (PF * PF <= \overline{N}))  {
    if (N \% PF == 0) ans -= ans / PF;
                                                        // only count unique factor
    while (N % PF == 0) N /= PF;
    PF = primes[++PF_idx];
 if (N != 1) ans -= ans / N;
                                                                       // last factor
 return ans;
//Algoritmo de Floyd de busqueda de ciclos.
//pongase la funcion f de la que calcular ciclos aqui
int f(int x) { return 2; }
ii floydCycleFinding(int x0) { // function "int f(int x)" must be defined earlier}
  // 1st part: finding v, hare's speed is 2x tortoise's
  int tortoise = f(x0), hare = f(f(x0)); // f(x0) is the element/node next to x0
 while (tortoise != hare) { tortoise = f(tortoise); hare = f(f(hare)); }
  // 2nd part: finding mu, hare and tortoise move at the same speed
  int mu = 0; hare = x0;
 while (tortoise != hare) { tortoise = f(tortoise); hare = f(hare); mu++; }
  // 3rd part: finding lambda, hare moves, tortoise stays
  int lambda = 1; hare = f(tortoise);
 while (tortoise != hare) { hare = f(hare); lambda++; }
  return ii(mu, lambda);
}
//Algoritmo de Euclides extendidio
//ax+by=g=gcd(a,b)
void eea (int a, int b,
        int& gcd, int& x, int& y) {
        x=0, y=1;
        int u=1, v=0, m, n, q, r;
        gcd = b;
        while (a!=0) {
                 q=gcd/a; r=gcd%a;
                m=x-u*q; n=y-v*q;
                 gcd=a; a=r; x=u; y=v; u=m; v=n;
        }
//ecuacion diofantica
//con d=mcd(A,B)
x=x0 + L*(B/d)
y=y0 - L*(A/d)
                     ____ MISCLETARIO
```

```
#include <vector>
using namespace std;
typedef pair<int, int> ii;
typedef vector<ii> vii;
typedef vector<int> vi;
//mascaras de bits
#define is0n(S, j) (S & (1 << j))
#define setBit(S, j) (S \mid= (1 << j))
#define clearBit(S, j) (S &= \sim(1 << j))
#define toggleBit(S, j) (S ^= (1 << ^-j))
#define lowBit(S) (S & (-S))
#define setAll(S, n) (S = (1 << n) - 1)
//Arbol de segmentos. Permite buscar el minimo elemento de un array LEERLO
// Segment Tree Library: The segment tree is stored like a heap array
void st_build(vi &st, const vi &A, int vertex, int L, int R) {
  if (L == R)
                                             // as L == R, either one is fine
                                                             // store the index
    st[vertex] = L;
    se { // recursively compute the values in the left and right subtrees
int nL = 2 * vertex, nR = 2 * vertex + 1;
  else {
                                         (L + R) / 2);
    st_build(st, A, nL, L
    st_build(st, A, nR, (L + R) / 2 + 1, R
    int lContent = st[nL] , rContent = st[nR];
int lValue = A[lContent], rValue = A[rContent];
    st[vertex] = (lValue <= rValue) ? lContent : rContent;</pre>
} }
void st_create(vi &st, const vi &A) {
                                              // if original array size is N,
     // the required segment tree array length is 2*2^{(floor(log2(N)) + 1)};
  int len = (int)(2*pow(2.0, floor((log((double)A.size())/log(2.0)) + 1)));
  st.assign(len, 0); // create vector of size `len' and fill it with zeroes
  st_build(st, A, 1, 0, (int)A.size() - 1);
                                                            // recursive build
int st_rmq(vi &st, const vi &A, int vertex, int L, int R, int i, int j) {
  if (i > R \mid | j < L) return -1; // current segment outside query range
  if (L >= i && R <= j) return st[vertex];</pre>
                                                        // inside query range
     // compute the min position in the left and right part of the interval
  int p1 = st_rmq(st, A, 2 * vertex , L
                                                           (L+R) / 2, i, j);
  int p2 = st_rmq(st, A, 2 * vertex + 1, (L+R) / 2 + 1,
                                                                       , i, j);
                         // return the position where the overall minimum is
  if (p1 == -1) return p2;
                                // if we try to access segment outside query
  if (p2 == -1) return p1;
                                                              // same as above
  return (A[p1] <= A[p2]) ? p1 : p2; }
int st_rmq(vi &st, const vi& A, int i, int j) {      // function overloading
  return st_rmq(st, A, 1, 0, (int)A.size() - 1, i, j); }
int st_update_point(vi &st, vi &A, int node, int b, int e, int idx, int new_value) {
  // this update code is still preliminary, i == j
  // must be able to update range in the future!
  int i = idx, j = idx;
  // if the current interval does not intersect
  // the update interval, return this st node value!
  if (i > e || j < b)
    return st[node];
  // if the current interval is included in the update range,
  // update that st[node]
  if (b == i && e == j) {
    A[i] = new_value; // update the underlying array
    return st[node] = b; // this index
  // compute the minimum position in the
  // left and right part of the interval
  int p1, p2;
```

```
, b
                                                            , (b + e) / 2, idx, new_value);
 p1 = st_update_point(st, A, 2 * node
 p2 = st update point(st, A, 2 * node + 1, (b + e) / 2 + 1, e
                                                                       , idx, new value);
  // return the position where the overall minimum is
  return st[node] = (A[p1] <= A[p2]) ? p1 : p2;</pre>
int st_update_point(vi &st, vi &A, int idx, int new_value) {
  return st_update_point(st, A, 1, 0, (int)A.size() - 1, idx, new_value); }
//Arbol de Fenwick. Permite resolver el problema de la suma en un rango
//con actualizaciones. (RSQ(a,b)=suma de los a[i], i entre a y b)
// initialization: n + 1 zeroes, ignoring index 0, just using index [1..n]
void ft_create(vi &ft, int n) { ft.assign(n + 1, 0); }
int ft_rsq(const vi &ft, int b) {
                                                        // returns RSQ(1, b)
  int sum = \theta; for (; b; b -= LSOne(b)) sum += ft[b];
  return sum; }
int ft_rsq(const vi &ft, int a, int b) {
                                                        // returns RSQ(a, b)
  return ft_rsq(ft, b) - (a == 1 ? 0 : ft_rsq(ft, a - 1)); }
// adjusts value of the k-th element by v (v can be +ve/inc or -ve/dec)
void ft_adjust(vi &ft, int k, int v) {
                                                 // note: n = ft.size() - 1
  for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v; }</pre>
//Backtracking
Funcion Backtracking (Etapai) devuelve: boolean
Inicio
        Exito = falso;
        IniciarOpciones(i, GrupoOpciones o);
        Repetir
                SeleccionarnuevaOpcion(o, Opcion n);
                Si (Aceptable(n)) entonces
                        AnotarOpcion(i, n);
                        SiSolucionCompleta(i) entonces
                                 Exito = verdadero;
                        Sino
                                 Éxito = Backtracking(i+1);
                                 Si Éxito = false entonces
                                         cancelamosAnotacion(i, n);
                                 finsi;
                        Finsi;
                Finsi;
        Hasta (éxito = verdadero) o (NoQuedanOpciones(o));
        Retorna Éxito;
Fin;
                        STRINGLETARIO
#include <cstdio>
#include <cstring>
#include<iostream>
#include<string>
#include<algorithm>
using namespace std;
#define MAXN 100010
//Knuth-Morris-Pratt: sirve para buscar la cadena P
//en la cadena T. Ojo, hay que preprocesar P usando
//kmpPreprocess(). MAXN<=100010</pre>
char T[MAXN], P[MAXN]; // T = text, P = pattern
int b[MAXN], n, m; // b = back table, n = length of T, m = length of P
void kmpPreprocess() { // call this before calling kmpSearch()
  int i = 0, j = -1; b[0] = -1; // starting values
 while (i < m) { // pre-process the pattern string P
   while (j \ge 0 \&\& P[i] != P[j]) j = b[j]; // if different, reset j using b
   i++; j++; // if same, advance both pointers b[i]=j; // observe i=8,\ 9,\ 10,\ 11,\ 12 with j=0,\ 1,\ 2,\ 3,\ 4
              // in the example of P = "SEVENTY SEVEN" above
} }
```

```
void kmpSearch() { // this is similar as kmpPreprocess(), but on string T
  int i = 0, j = 0; // starting values
  while (i < n) { // search through string T</pre>
    while (j \ge 0 \& T[i] != P[j]) j = b[j]; // if different, reset j using b
    i++; j++; // if same, advance both pointers
    if (j == m) \{ // a match found when <math>j == m
      printf("P is found at index %d in T\n", i - j);
      j = b[j]; // prepare j for the next possible match
} } }
/*ejemplo de uso
int main() {
  n = (int)strlen(gets(T));
  m = (int)strlen(gets(P));
  kmpPreprocess();
  kmpSearch();
  return 0;
}*/
//Array de sufijos. r contiene la posicion del iesimo sufijo ordenadamente de s
//pongo un ejemplo de uso mas abajo
//0jo: MAXN<1000005
//Importante: añadir un '$' al final de las cadenas para romper desempates en el
//count sort.
int n,t; //n es el tamaño de la cadena
int p[MAXN],r[MAXN],h[MAXN];
//p es el inverso del suffix array, no usa indices del suffix array ordenado
//h el el tamaño del lcp entre el i-esimo y el i+1-esimo elemento de suffix array ordenado
string s;
void fix_index(int *b, int *e) {
   int pkm1, pk, np, i, d, m;
   pkm1 = p[*b + t];
   m = e - b; d = 0;
   np = b - r;
   for(i = 0; i < m; i++) {
      if (((pk = p[*b+t]) != pkm1) && !(np <= pkm1 && pk < np+m)) {
         pkm1 = pk;
         d = i;
      p[*(b++)] = np + d;
   }
}
bool comp(int i, int j) {
   return p[i + t] < p[j + t];
void suff_arr() {
   int i, j, bc[256];
   t = 1;
   for(i = 0; i < 256; i++) bc[i] = 0; //alfabeto
   for(i = 0; i < n; i++) ++bc[int(s[i])]; //counting sort inicial del alfabeto
   for(i = 1; i < 256; i++) bc[i] += bc[i - 1];
   for(i = 0; i < n; i++) r[--bc[int(s[i])]] = i;
   for(i = n - 1; i >= 0; i--) p[i] = bc[int(s[i])];
   for(t = 1; t < n; t *= 2) {
      for(i = 0, j = 1; i < n; i = j++) {
         while(j < n && p[r[j]] == p[r[i]]) ++j;
         if (j - i > 1) {
            sort(r + i, r + j, comp);
            fix_index(r + i, r + j);
         }
      }
   }
}
void lcp() {
   int tam = 0, i, j;
for(i = 0; i < n; i++)if (p[i] > 0) {
      j = r[p[i] - 1];
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while(s[i + tam] == s[j + tam]) ++tam;
      h[p[i] - 1] = tam;
      if (tam > 0) --tam;
   h[n - 1] = 0;
/*ejemplo de uso
int main(){
   s="margarita$";
   n=s.size();
   suff_arr();
   lcp();
   for(int i=0;i<n;i++)cout<<r[i]<<" ";cout<<endl;</pre>
   for(int i=0;i<n;i++)cout<<h[i]<<" ";cout<<endl;</pre>
   return 0;
}*/_
                    __ IOMANIP
// FALGS DE FOMATO

    boolalpha/noboolalpha: muesta su valor verdadero/falso

         cout << boolalpha << b; :== true
- showbase/noshowbase: muesta el prefijo de la base (0x _hex_, 0 _oct_, nada _dec_).
        cout << hex << showbase << 14 << endl; :== 0x14
- showpoint/noshowponint: muestra el punto de _double/float_. Se puede combinar con setprecicion. cout.precision (5); cout << showpoint << 30.0 << 10000.0 << 3.1416; :== 30.000 10000. 3.1416
- showpos/noshowpos: muesta el signo + para numeros positivos
        cout << showpos << 1 << 0 << -1; :== +1 +0 -1
- skipws: elimina todos los espacios en blanco de la entrada
          istrinstream iss (" 123");
          iss >> skipws >> a >> b >> c;
          cout << a << b << c; :== 123
- uppercase/nouppercase: convierte el texto a mayusculas.
          cout << showbase << hex;</pre>
          cout << uppercase << 77; :== 0X4D
          cout << nouppercase << 77; :== 0x4d
- fixed: escribe _numeros flotantes_ en notacion de punto fijo
 scientific: escribe numeros flotantes en notación científica
          cout.setprecission(5);
                           << 3.1415926532 << 2006.0 << 1.0e-10;
          cout <<
                              2006
                3.1416
                                             1e-010
                              << 3.1415926532 << 2006.0 << 1.0e-10;
          cout <<
                      fixed
               3.14159
                              2006.00000 0.00000
          cout << scientific << 3.1415926532 << 2006.0 << 1.0e-10;</pre>
                3.14159e+000 2.00600e+003 1.00000e-10
- internal: espacia el texto separando el signo del valor
- left: muestra el texto alineado a la izquierda con el ancho indicdo
- rigth: muestra el texto a la derecha con el ancho indicado
         cout.width(6); cout << internal << -77 << endl; :== -</pre>
         cout.width(6); cout << left << -77 << endl; :== -77
         cout.width(6); cout << right << -77 << endl;</pre>
                                                            :==
// MANIPULADORES PARAMETRIZADOS

    setfill (char c): rellena el ancho indicado con el caracter pasado por parametro

          cout << setfill ('x') << setw (10;</pre>
          cout << 77 << endl; :== xxxxxxxx77
- setprecision: imprime el número según la precisión indicada. Si es fixed rellena con \theta
          cout << setprecision (5) << 3.14159 << endl; :== 3.1416
          cout << setprecision (9) << 3.14159 << endl; :== 3.14159
          cout << fixed;</pre>
          cout << setprecision (5) << 3.14159 << endl; :== 3.14159
          cout << setprecision (9) << 3.14159 << endl; :== 3.141590000

    setw(int n): setea el ancho de linea indicado

           cout \ll setw(10);
           cout << 77 << endl;
```