

Lesson 5

Higher-order functions

Higher order functions

- functions that return functions as result

obvious with curryfication

- functions that take functions as parameters

$\text{twice} :: (a \rightarrow a) \rightarrow a \rightarrow a$

$\text{twice } f = f . f$

$\text{twice } (*2) 3$

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```
twice reverse [1,2,3]  
[1,2,3]
```

```
map :: (a -> b) -> [a] -> [b]  
map f xs = [ fx | x <- xs]
```

```
map (+1) [1,2,3]  
[2,3,4]
```

```
map reverse ["abc", "def"]  
["cba", "fed"]
```

two maps to work on lists of lists :: `[[a]]`

```
map (map (+1)) [[1,2,3],[4,5]]
```

={applying outer map}

```
[map (+1) [1,2,3], map (+1) [4,5]]
```

={applying inner maps}

```
[[2,3,4],[5,6]]
```

```
filter :: (a -> Bool) -> [a] -> [a]
```

```
filter p xs = [x | x <- xs, p x]
```

```
filter even [1..10]
```

```
[2,4,6,8,10]
```

```
filter (\= ` `) "abc def ghi"  
"abcdefghi"
```

```
filter p [] = []  
filter p (x:xs) | px = x : filter p xs  
                | otherwise = filter p xs
```

From the article «Why functional programming matters» by John Hughes

The importance of modularity and «gluing» functions together


```
sum [] = 0  
sum (x:xs) = x + sum xs
```

sum can be modularized by gluing together a general recursive pattern and the boxed parts

the recursive pattern is called foldr

$\text{sum} = \text{foldr } (+) 0$ generalized to $\text{foldr } f v$

parameterizing sum:

$\text{sum } [] = 0$		$(\text{foldr } f v) [] = v$
$\text{sum } (x:xs) = x + \text{sum } xs$		$\text{foldr } f v (x:xs) = f x ((\text{foldr } f v) xs)$

$\text{foldr} :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b$

`product = foldr (*) 1`

`anyTrue = foldr (||) False`

`allTrue = foldr (&&) True`

it replaces all `:` in a list by `f` and `[]` with `0/1/False/True`

`foldr (+) 0 (1:2:3:[])`

`1+(2+(3+0))` e in generale:

`foldr (#) v [x0,x1,...xn] = x0 # (x1 # (...(xn # v)...))`

foldr (:) [] simply copies a list

append a b = foldr (:) b a

(:) replaced by (:) and [] by b

length = foldr count 0

count _ n = n + 1

surprising: reverse a list with foldr

$\text{snoc } x \text{ } xs = xs ++ [x]$

$\text{reverse} = \text{foldr snoc } []$

$\text{foldr } (\#) \text{ } v \text{ } [x_0, x_1, \dots, x_n] = x_0 \# (x_1 \# (\dots (x_n \# v) \dots))$

if we want to double all elements of a list:

```
doubleAll:: Num a => [a] -> [a]
```

```
doubleAll = foldr doubleAndCons []
```

where

```
doubleAndCons n xs = (n*2) : xs
```

doubleAndCons can be modularized further:

```
doubleAndCons = fAndCons double
```

```
where double n = 2 * n    and
```

```
fAndCons f x xs = (f x) : xs = (:) (fx) xs
```

e alla fine:

```
fAndCons f = (:) . f
```

```
doubleAll= foldr (:).double []
```

but we know that `doubleAll` can be defined also with `map`:

```
doubleAll = map double
```

this is not by chance

```
map f = foldr (:).f Nil
```

`map` always has `(:)`, `foldr` does not, is more general

the same idea applies to other data structures (not just to lists)

`data Tree a = Node a [Tree a] -- Node is a constructor, i.e. a function that builds the values of type Tree a:`

`Node :: a -> [Tree a] -> Tree a`

leafs have [], i.e. empty list of subtrees

trees are built with Node, (:) and [], 3 things, hence:

`foldtree f g a Node / [Trees] / []`

$\text{foldtree } f \ g \ a \ (\text{Node } x \ tx) = f \ x \ (\text{foldtree } f \ g \ a \ tx)$

$\text{foldtree } f \ g \ a \ (t:tx) = g \ (\text{foldtree } f \ g \ a \ t) \ (\text{foldtree } f \ g \ a \ tx)$

$\text{foldtree } f \ g \ a \ [] = a$

$\text{sumtree} = \text{foldtree } (+) \ (+) \ 0$

$\text{labels} = \text{foldtree } (:) \ \text{append} \ \text{Nil}$

$\text{maptree} = \text{foldr } (\text{Node} \ . \ f) \ (:) \ \text{Nil}$

$\text{foldl} :: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow [b] \rightarrow a$

fold from left to right

$\text{foldl } f \ v \ [] = v$

$\text{foldl } f \ v \ (x:xs) = \text{foldl } f \ (f \ v \ x) \ xs$

$\text{foldl } (\#) \ v \ [x_0, x_1, \dots, x_n] = (\dots((v \# x_0) \# x_1) \dots) \# x_n$

associates to the left

$\text{sum} = \text{foldl } (+) \ 0$

we can do reverse also with foldl:

$\text{reverse} = \text{foldl } (\backslash xs \rightarrow \backslash x \rightarrow x:xs) \ []$

$\text{map } f = \text{foldl } (\backslash v \rightarrow (\backslash x \rightarrow v ++ [f \ x])) \ []$

the function composition operator

$$(\cdot) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$$
$$f \cdot g = \lambda x \rightarrow f (g x)$$

identity for composition:

$$\text{id} : a \rightarrow a$$
$$\text{id} = \lambda x \rightarrow x$$
$$\text{compose} :: [a \rightarrow a] \rightarrow (a \rightarrow a)$$
$$\text{compose} = \text{foldr } (\cdot) \text{id}$$

Exercise 2 c

takeWhile :: (a -> Bool) -> [a] -> [a]

takeWhile p [] = []

```
takeWhile p (x:xs) = if p x then x : takeWhile p xs
                    else []
```

Exercise 2 d: dropWhile

Exercise 3 redefine filter p using foldr