

# Lesson 10

Chapter 12

Functors and Applicatives

we define some type classes that require the presence of  
general functions

these general functions abstract common programming patterns

we have already seen examples of such generalizations : foldr,  
foldl, map, maptree

we go further

Functor  $\rightarrow$  Applicative  $\rightarrow$  Monad

class Functors

$\text{inc} :: [\text{Int}] \rightarrow [\text{Int}]$

$\text{inc} [] = []$

$\text{inc} (n:ns) = n+1 : \text{inc} ns$

$\text{map} :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]$

$\text{map} f [] = []$

$\text{map} f (x:xs) = f x : \text{map} f xs$

$\text{inc} = \text{map} (+1)$

class Functor f where

fmap :: (a -> b) -> f a -> f b

examples of Functor's instances

instance Functor [] where

-- fmap :: (a -> b) -> [a] -> [b]

fmap = map

[a] is the same as [] a type list applied to the parameter type a

instances of Functor are always types T that have a type parameter

another example:

```
data Maybe a = Nothing | Just a
```

```
instance Functor Maybe where
```

```
--fmap :: (a -> b) -> Maybe a -> Maybe b
```

```
fmap _ Nothing = Nothing
```

```
fmap g (Just x) = Just (g x)
```

```
fmap (+1) Nothing = Nothing
```

```
fmap (+1) Just 3 = Just 4
```

```
fmap not Just True = Just False
```

```
data Tree a = Leaf a | Node (Tree a) (Tree a)
              deriving Show
```

```
instance Functor Tree where
```

```
--fmap :: (a -> b) -> Tree a -> Tree b
```

```
fmap g (Leaf x)      = Leaf (g x)
```

```
fmap g (Node l r)    = Node (fmap g l) (fmap g r)
```

```
fmap even Node (Leaf 1) (Leaf 2)
```

```
Node (Leaf false) (Leaf True)
```

in general a Functor  $T\ a$  is a container that contains value of type  $a$   
and `fmap g` applies `g` to each value in the container

instance Functor IO where

`--fmap :: (a -> b) -> IO a -> IO b`

`fmap g mx = do`

`x <- mx;`

`return (g x)`

`fmap g mx = do { x <- mx; return (g x) }`

example

the partially applied function type  $(a \rightarrow)$ , represented as  $((\rightarrow) a)$ ,  
is a Functor:

$$\text{fmap} :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$$
$$\text{fmap} = (.)$$



Functors => generalization

fmap can be used to process the values of any container (which is a Functor)

this generalization **propagates** to functions defined by mean of fmap

inc = fmap (+1)

inc (Just 1)   inc [1,2,3]   inc (Node (Leaf 1) (Leaf 2))

# Functor laws

- 1)  $\text{fmap id} = \text{id}$
- 2)  $\text{fmap } (g \cdot h) = \text{fmap } g \cdot \text{fmap } h$

(1) states that `fmap` preserves the identity

(2) states that `fmap` preserves function composition

*observe that  $\text{fmap id} :: f\ a \rightarrow f\ a$ , whereas  $\text{id} :: a \rightarrow a$*

together with the type of `fmap`, the laws guarantee that `fmap` is a mapping that does not reorder the values in the container that it is applied to

a def. of fmap that does not obey the laws  
instance Functor [] where

--fmap :: (a -> b) -> [a] -> [b ]

fmap g [] = []

fmap g (x: xs) = fmap g xs ++ [g x]

fmap id [1,2] = [2,1]

id [1,2] = [1,2]

In Chapter 16 we will see how one can show that a given Functor satisfies the Functor laws

**interesting fact:**

for any parameterized type in Haskell there is at most one function `fmap` that satisfies the laws

if we can make a parameterized type into a Functor (satisfying the laws), then `fmap` is unique

# Applicatives (are Functors)

Functors map functions with one argument over containers,  
Applicatives map functions with many parameters over  
containers

$\text{fmap0} :: a \rightarrow f\ a$

$\text{fmap1} :: (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

$\text{fmap2} :: (a \rightarrow b \rightarrow c) \rightarrow f\ a \rightarrow f\ b \rightarrow f\ c$

$\text{fmap3} :: (a \rightarrow b \rightarrow c \rightarrow d) \rightarrow f\ a \rightarrow f\ b \rightarrow f\ c \rightarrow f\ d$

.....

$\text{fmap2 } (+) (\text{Just } 1) (\text{Just } 2) = \text{Just } 3$

$\text{fmap} :: (a \rightarrow b) \rightarrow f a \rightarrow f b$

generalizes:

$\text{app} :: (a \rightarrow b) \rightarrow a \rightarrow b$

thanks to currying we don't need  $\text{app1}$ ,  $\text{app2}$ ,  $\text{app3}$ ,....

$\text{app1 } g x = g x$

$\text{app2 } g x y = \text{app1 } (\text{app1 } g x) y$

$\text{app3 } g x y z = \text{app1 } (\text{app1 } (\text{app1 } g x) y) z$

also here two basic functions are sufficient to define `fmap0`, `1`, `2`, `3`,....

`pure :: a -> f a`

`(<*>) :: f (a -> b) -> f a -> f b`

`pure` converts a value of type `a` into a structure `f a`

`<*>` is a generalized form of function application

`g <*> x <*> y <*> x` generalizes `g x y z`

typical use: `pure g <*> x <*> y <*> z` applicative style

$\text{fmap0} :: a \rightarrow f\ a$

$\text{fmap0} = \text{pure}$

$\text{fmap1} :: (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

$\text{fmap1}\ g\ x = \text{pure}\ g\ \langle * \rangle\ x$

$\text{fmap}\ 2 :: (a \rightarrow b \rightarrow c) \rightarrow f\ a \rightarrow f\ b \rightarrow f\ c$

$\text{fmap2}\ g\ x\ y = \text{pure}\ g\ \langle * \rangle\ x\ \langle * \rangle\ y$

$\text{fmap3} :: (a \rightarrow b \rightarrow c \rightarrow d) \rightarrow f\ a \rightarrow f\ b \rightarrow f\ c \rightarrow f\ d$

$\text{fmap3}\ g\ x\ y\ z = \text{pure}\ g\ \langle * \rangle\ x\ \langle * \rangle\ y\ \langle * \rangle\ z$



$g :: a \rightarrow b \rightarrow c \rightarrow d = (a \rightarrow (b \rightarrow (c \rightarrow d)))$       observe types

$\text{pure } g :: f (a \rightarrow (b \rightarrow (c \rightarrow d)))$

$(\langle * \rangle) :: f (a \rightarrow B) \rightarrow f a \rightarrow f B \quad \text{con } B = (b \rightarrow (c \rightarrow d))$

$\text{pure } g \langle * \rangle x :: f B = f (b \rightarrow (c \rightarrow d))$

$\text{pure } g \langle * \rangle x \langle * \rangle y :: f (c \rightarrow d)$

$\text{pure } g \langle * \rangle x \langle * \rangle y \langle * \rangle z :: f d$

definition

```
class Functor f => Applicative f where
```

```
pure :: a -> f a
```

```
(<*>) :: f (a -> b) -> f a -> f b
```

examples of Applicative:

```
instance Applicative Maybe where
```

```
--pure :: a -> Maybe a
```

```
pure = Just
```

```
--(<*>) :: Maybe (a -> b) -> Maybe a -> Maybe b
```

```
Nothing <*> _ = Nothing
```

```
(Just g) <*> mx = fmap g mx
```

pure (+1) <\*> Just 1

Just 2

pure (\*) <\*> Just 2 <\*> Just 3

Just 6

pure (\*) <\*> Just 2 <\*> Nothing

Nothing

lists are Applicatives

instance Applicative [] where

--pure :: a -> [a]

pure x = [x]

--(<\*>) :: [a -> b] -> [a] -> [b]

gs <\*> xs = [g x | g <- gs, x <- xs]

pure (+1) <\*> [1,2,3]

[2,3,4]

[(\*2),(+2)] <\*> [1,2,3]

[2,4,6,3,4,5]

`prods :: [Int] -> [Int] -> [Int]`

`prods [2,3] [3,4]`

`[6,8,9,12]`

`prods xs ys = pure (*) <*> xs <*> ys`

`xs` and `ys` may be results of non-deterministic functions

`failure = []`

`prods [2,3] [] = []`

# IO is Applicative

instance Applicative IO where

--pure :: a -> IO a

pure = return

--(<\*>) :: IO (a -> b) -> IO a -> IO b

mg <\*> mx = do { g <- mg; x <- mx; return (g x) }

getChars :: Int -> IO String

getChars 0 = []

getChars n = pure (:) <\*> getChar <\*> getChars (n-1)