## Digital Signal Processing Ch8&9 DFS DFT FFT

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modulo N
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- 。若 $\widetilde{\chi}[n]$ 為週期信號,則滿足 $\widetilde{\chi}[n] = \widetilde{\chi}[n+N]$
- 。((n))n:將起出範圍的n折返到週期內
- · EX: 作意 蒙 究[n] = (1.2.3.4), N=4

$$X[((0))_{4}] = \hat{X}[0] = 1$$
  
 $X[((3)_{4}] = \hat{X}[3] = 4$   
 $X[((5))_{4}] = \hat{X}[1] = 2$   
 $X[((5))_{4}] = \hat{X}[1] = 2$   
5起出範圍

⇒ 
$$\chi(k) = \chi(k-N)$$
 皆為週期信號. 故範圍為  $(-\infty,\infty)$ 

$$\circ e^{j\left(\frac{2\pi}{N}\right)(k+\ell N)n} = e^{j\left(\frac{2\pi}{N}k\right) + 2\pi k \mathcal{I}}$$

$$\int_{N} \int_{n=0}^{N} e^{\frac{1}{2} \left(\frac{2\pi}{N}\right) \ln n} = \frac{1}{N} - \frac{1 - e^{\frac{1}{2} \left(\frac{2\pi}{N}\right) \ln n}}{1 - e^{\frac{1}{2} \frac{\pi}{N} \ln n}} = \begin{cases} 1 & \text{if } l = mN \\ 0 & \text{o.w.} \end{cases}$$

$$\begin{cases} X(e^{j\omega}) = \sum_{N=0}^{N-1} \chi(n) e^{-j\omega n} = \sum_{N=0}^{N-1} \chi(n) e^{-j\omega n} \\ X(k) = \sum_{N=0}^{N-1} \chi(n) W_N^{kn} = \sum_{N=0}^{N-1} \chi(n) e^{-j(2\pi) k} = X(e^{j\omega}) \\ \omega = \frac{2\pi}{N} k \end{cases}$$

由於 DFS 的  $\tilde{\chi}[n]$  秋  $\tilde{\chi}[k]$  皆為週期信號  $\rightarrow$  range為  $(-\infty.\infty)$ 而 DFT 則是選取 DFS 的其中一段, 使其為有限區間

## DFT

$$\lambda(n) = \begin{cases} \chi(n) : 0 \le n \le N-1 \\ 0 : otherwise \end{cases}$$

$$X(k) = \begin{cases} \widetilde{X(k)} ; & 0 \le N \le N-1 \\ 0 ; & \text{otherwise} \end{cases}$$

$$\chi(n) = \chi((n))_{N}$$

$$X[k] = X[((k))^{N}]$$

$$X[u] = X[((u))^{N}]$$

$$\chi(n) = \chi((n))_N$$
  $\chi(n), \chi(k)$  過期信號. 
$$\chi(k) = \chi((k))_N$$
 取長度 N 做為循環.

$$\Rightarrow \begin{cases} \chi[k] = \sum_{n=0}^{N-1} \chi[n] W_N^{kn} & ; 0 \le n \le N-1 \\ \chi[n] = \sum_{n=0}^{N-1} \chi[k] W_N^{-kn} & ; 0 \le k \le N-1 \end{cases}$$

 $x_1[n] = \begin{cases} \widetilde{x}_1[n], & 0 \le n \le N-1 \\ 0, & \text{otherwise} \end{cases}$ 

「平移後的信號」對應到「原本信號」 ↓

 $\chi_{\lfloor 0 \rfloor} = \chi \lceil ((0+2))_b \rceil = \chi[2]$ 

 $\chi_{1}[1] = \chi[((1+2))_{6}] = \chi[3]$ 

 $\chi_{\iota}[4] = \chi \left[ ((4+2))_{\iota} \right] = \chi[0]$ 

 $\chi_{i}[5] = \chi[((5+2))_{i}] = \chi[i]$ 

Sampling DFT.  $\chi(n) = \sum_{r=0}^{\infty} \chi(n-rN) : DFS : N > length of <math>\chi(n)$  Affin we can recover  $\chi(n)$  from  $\chi(n)$  $\circ \chi \widetilde{(k)} = \chi(e^{\widetilde{J}\omega})|_{\omega = \frac{2\pi}{N}k}$ = X(Z) | Z= ej 2/k  $\rightarrow \chi(n) = \frac{1}{N} \sum_{k=0}^{N-1} \chi(k) W_N^{-kn} = \frac{1}{N} \sum_{k=0}^{N-1} \sum_{m=-\infty}^{\infty} \chi(m) W_N^{-kn} \cdot W_N^{-kn}$  $\sum_{m=-\infty}^{\infty} \chi[m] \frac{1}{N} \sum_{k=0}^{N-1} W_{N}^{-k(n-m)} = p[n-m]$  $\frac{1}{2} p[n] = \sum_{k=0}^{\infty} S[n-rN] = \chi[n] * p[n] \frac{1}{N} \sum_{k=0}^{N-1} W_{N}^{+n} = p[n] = \sum_{k=0}^{\infty} S[n-rN]$   $\Rightarrow p(e^{j\omega}) = \frac{2\pi}{N} \sum_{k=0}^{\infty} S(\omega - \frac{2\pi}{N}k) = \sum_{k=0}^{\infty} \chi[n-rN]$