

Randomized Algorithms: 2-dim Linear Programming

Deterministic algorithms: same behaviour on a fixed input

Randomized algorithms: random decisions

→ running time → random variable

→ output can be random

- With high probability, running time is small
- w.h.p., Output is correct.

Randomized Quicksort (randomly choose pivot)

Sampling, Avoid worst cases

- With some prob., hardware can fail.
- Algorithms : repeat multiple times
⇒ error probability very small

Many cases : randomized algorithms are the fastest known.

Also, simplest.

2-dim

Linear Programming

x, y

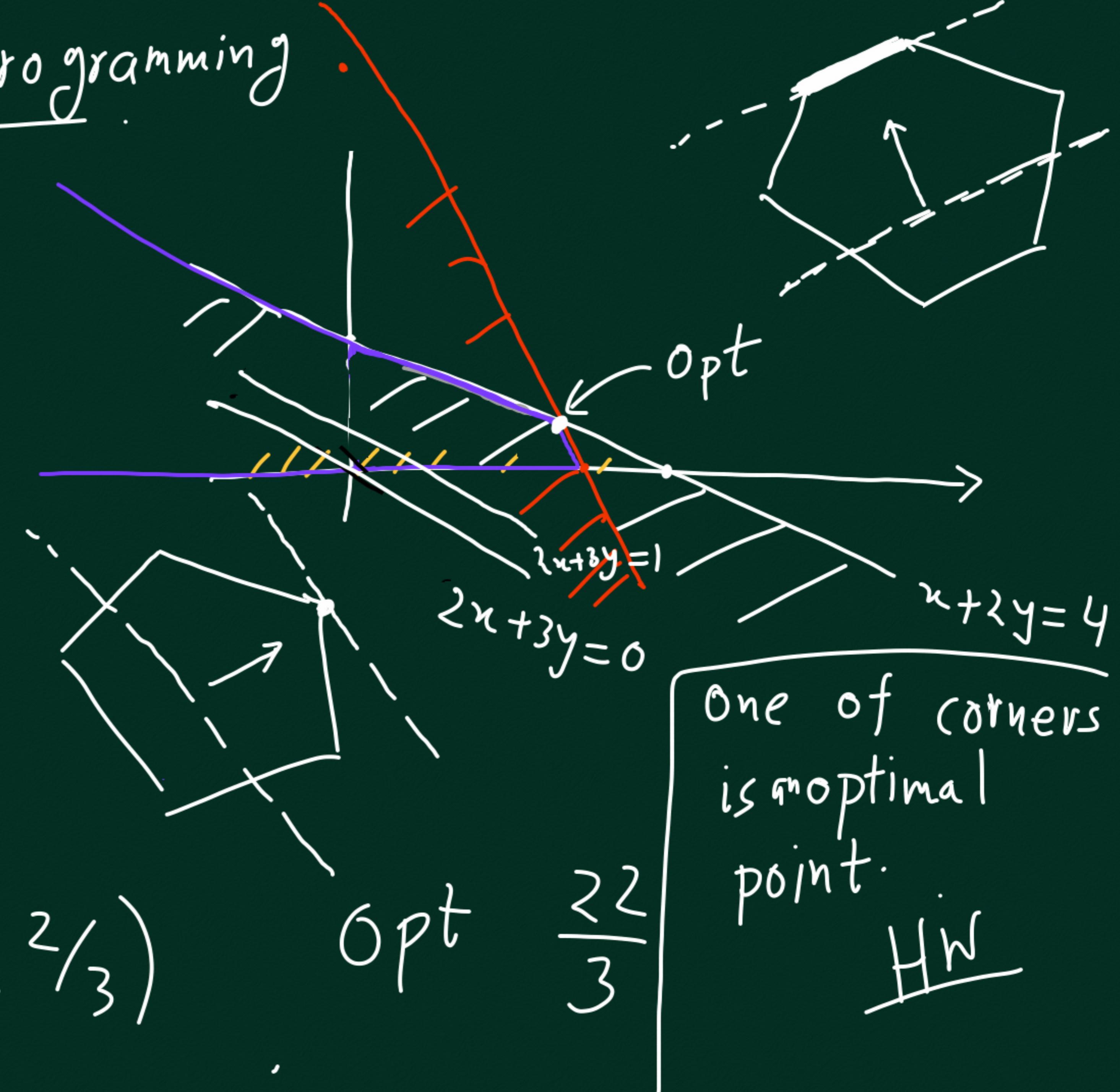
$$x + 2y \leq 4$$

$$2x + y \leq 6$$

$$y \geq 0$$

$$\max 2x + 3y$$

$$(8/3, 2/3)$$



Opt

$$\frac{22}{3}$$

One of corners
is optimal
point.

HW

n linear constraints.

check all the corners

$O(n^3)$

at most n



no. of intersection

points

$O(n^2)$



~~Hw~~ $\underline{\Theta(n \log n)}$ compute all corners.

Randomization

Find the optimal point

simple algorithm

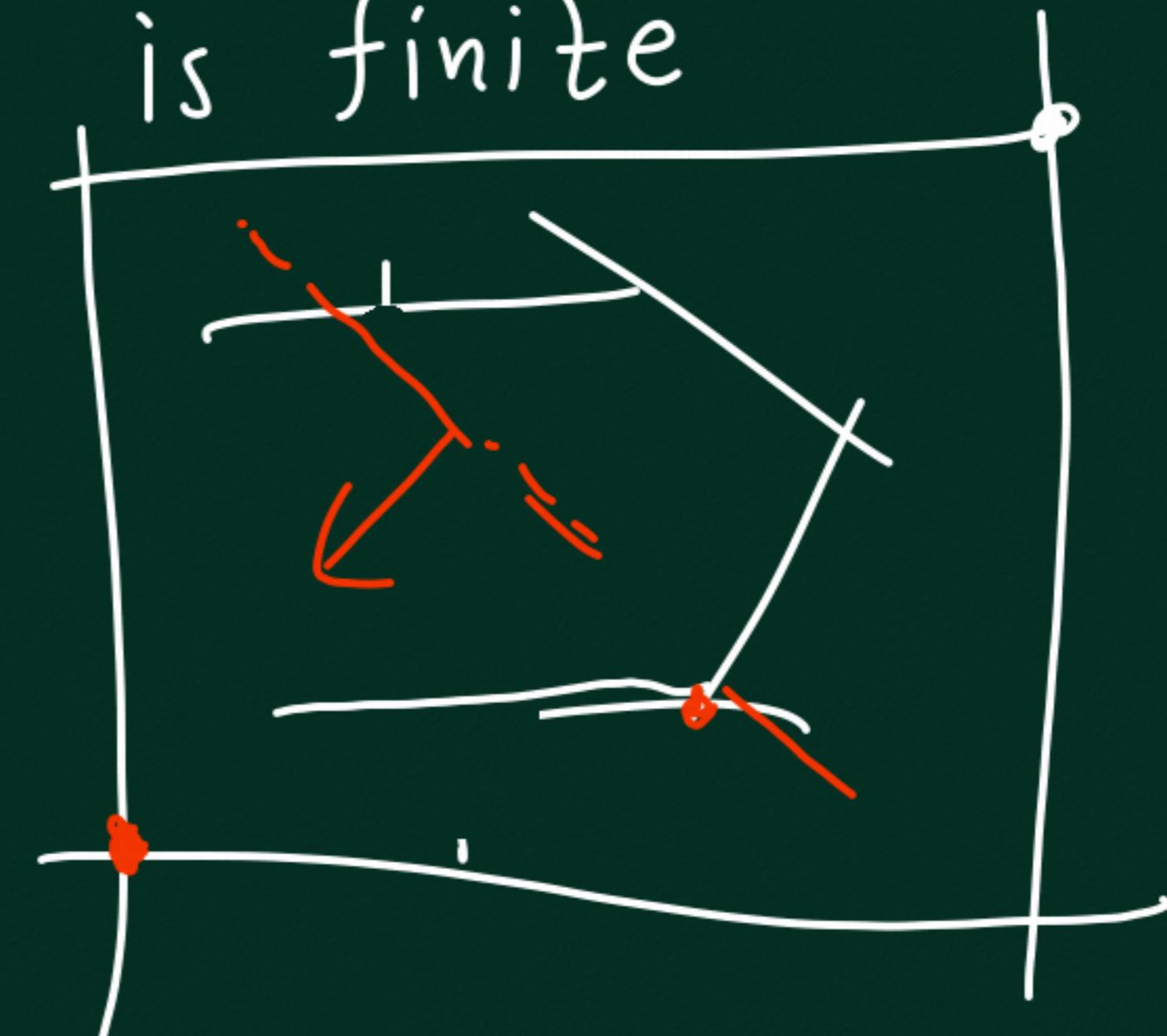
$\underline{\Theta(n)}$ time ($w \cdot h \cdot p$)

Assumptions: optimal value is finite

for every intersection point

$$-L < x\text{-coord} < L$$

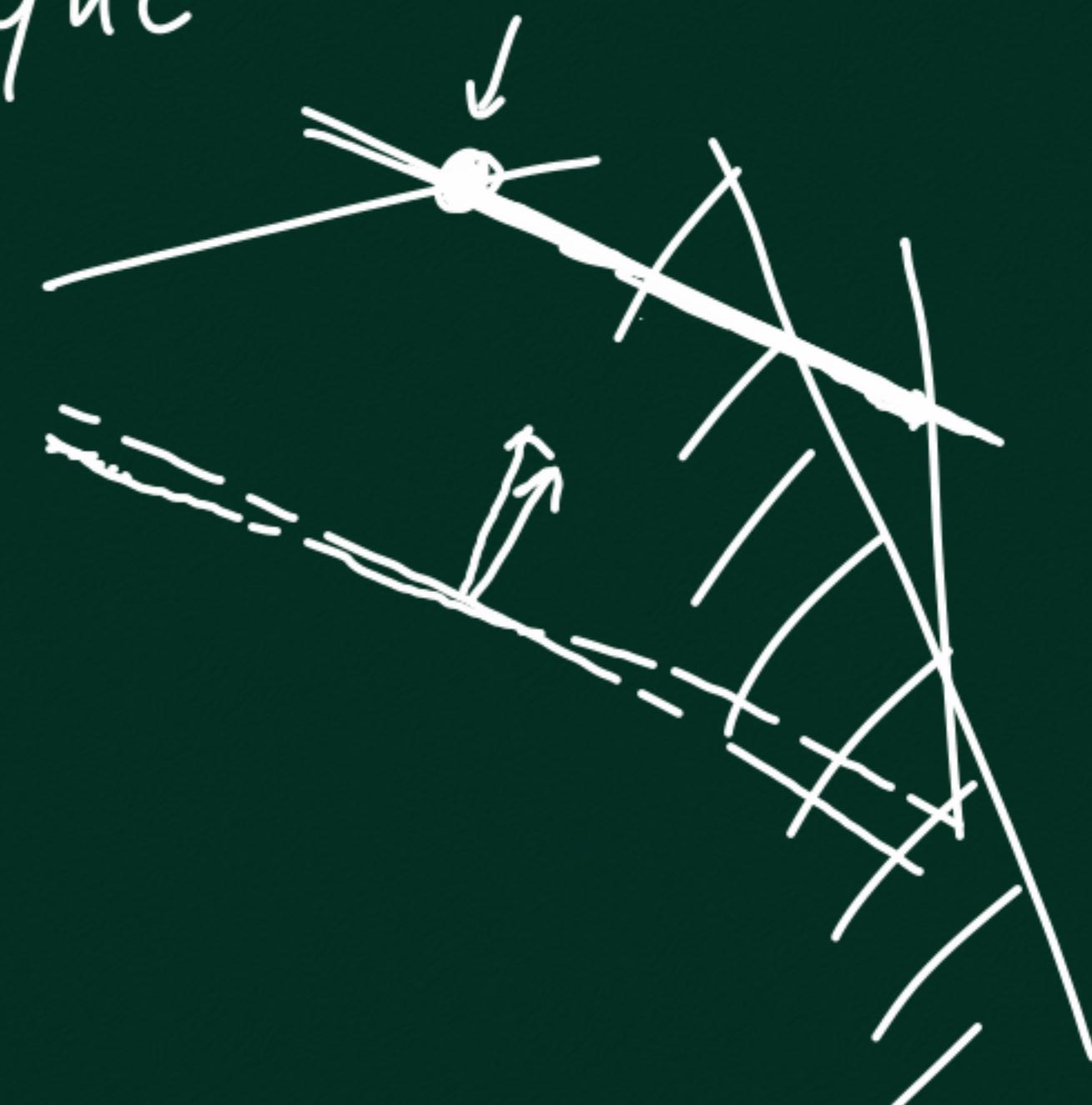
$$-L < y\text{-coord} < L$$



Assumption : optimal solution is unique.

$$C_0 \leftarrow \text{Set of constraints} \leftarrow -L \leq x \leq L$$
$$-L \leq y \leq L$$

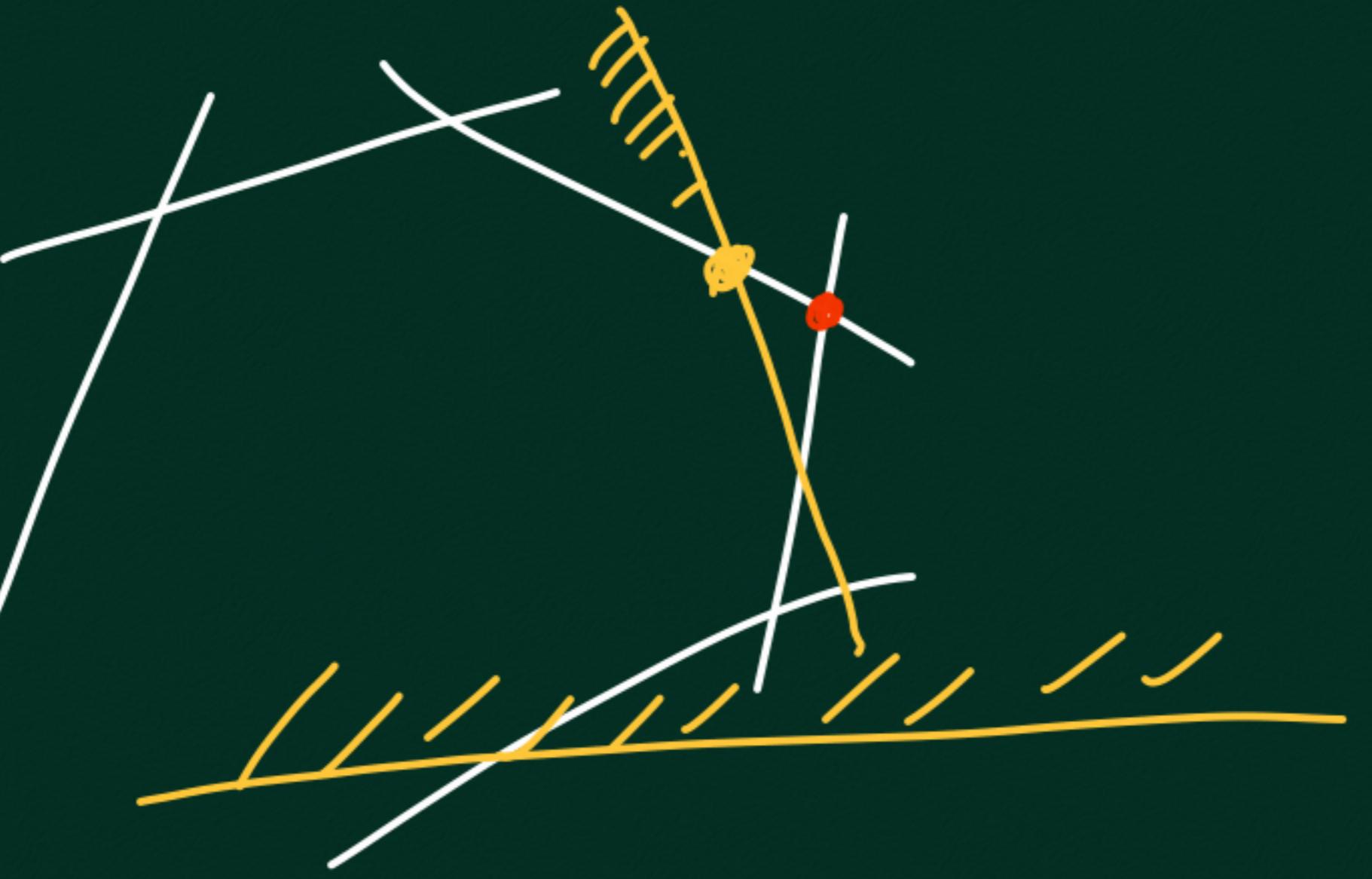
$$V_0 \leftarrow \text{Optimal point} \leftarrow (L, L)$$



Add constraints one by one and
update the optimal point

$$(C_{i-1}, V_{i-1}) \longrightarrow (C_i, V_i)$$

$$C_i = C_{i-1} \cup \{h_i\}$$



- if v_{i-1} satisfies h_i
then $v_i = v_{i-1}$

