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Davis Putnam Logemann Loveland (DPLL)

$$C_1 \wedge C_2 \wedge \dots \wedge C_k$$

$SAT(\varphi, PA)$ returns $(\text{status}, \text{assignment})$

↓

sat unsat $x_1=0, x_2=1, \dots$

0. If $\varphi = \top$ return (sat, PA)
If $\varphi = \perp$ return (unsat, PA)

1. If C_i is a unit clause (clause with single literal l)

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return SAT( $\varphi[l=1]$ ,  $PA \cup \{l=1\}$ )
```

Unit-propagation

simplify Q after setting $l=1$

2. If a literal l doesn't appear negated in any clause

```
return SAT( $\varphi[l=1]$ , PAU $\{l=1\}$ )
```

pure

literal

elimination

```

3.  $x := \text{choose\_variable}(\varphi)$ 
    $v := \text{choose\_value}() \dots \in \{0, 1\}$ 
   if  $\text{SAT}(\varphi[x=v], PA \cup \{x=v\}).\text{status} = \text{sat}$ 
       return  $(\text{sat}, PA \cup \{x=v\})$ 

```

decision

```

else if SAT( $\varphi[x=1-v]$ ,  $PA \cup \{x=1-v\}$ ).status = sat
return (sat,  $PA \cup \{x=1-v\}$ )

```

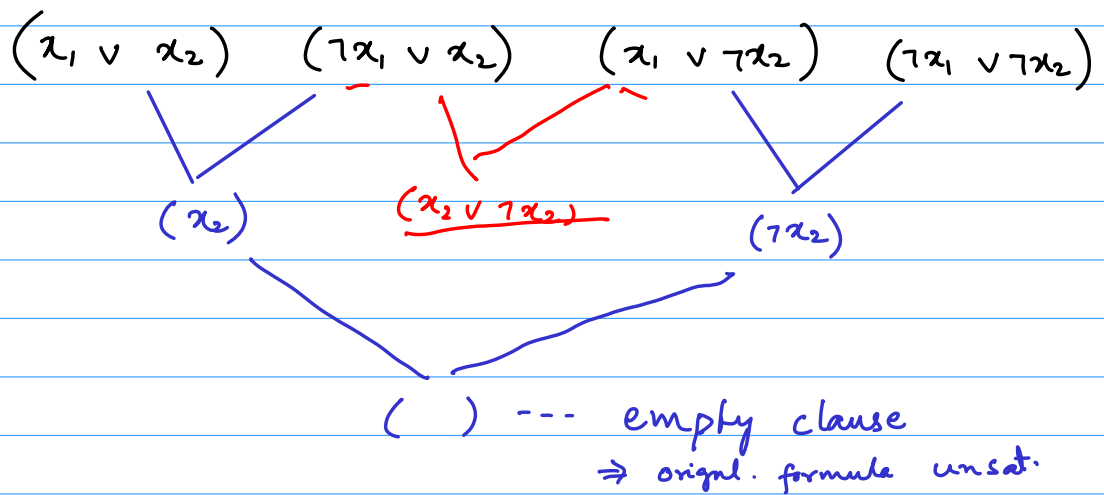
```
else return (unsat, PA)
```

back track

Illustrating how a Horn formula may force DPLL to backtrack
if we start with 1 as the first choice of value for a decision variable

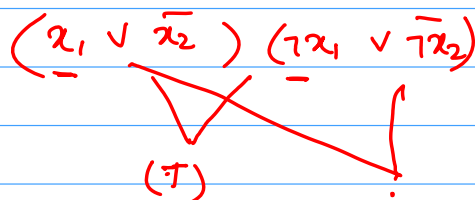
		<u>$p_1 = 1$</u>	<u>$p_1 = 0$</u>
$p_1 \rightarrow p_2$	$(\neg p_1 \vee p_2)$	p_2	1
$p_2 \rightarrow p_1$	$(\neg p_2 \vee p_1)$	1	$\neg p_2$
$p_1 \wedge p_2 \rightarrow \perp$	$(\neg p_1 \vee \neg p_2)$	$\neg p_2$	1

Resolution.
$$\frac{(l_1 \vee l_2 \vee l_3 \vee \underline{x}) \quad (l_4 \vee l_5 \vee l_6 \vee l_7 \vee \underline{\neg x})}{\text{Resolvent} \rightarrow (l_1 \vee l_2 \vee l_3 \vee l_4 \vee l_5 \vee l_6 \vee l_7)}$$



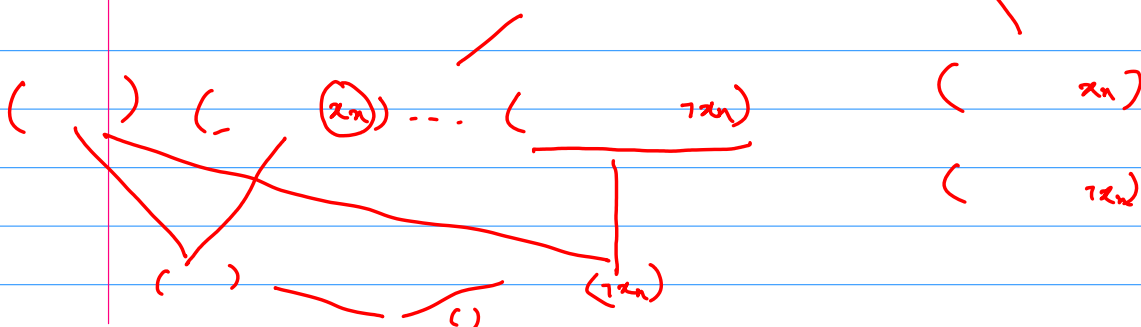
n prop vars.

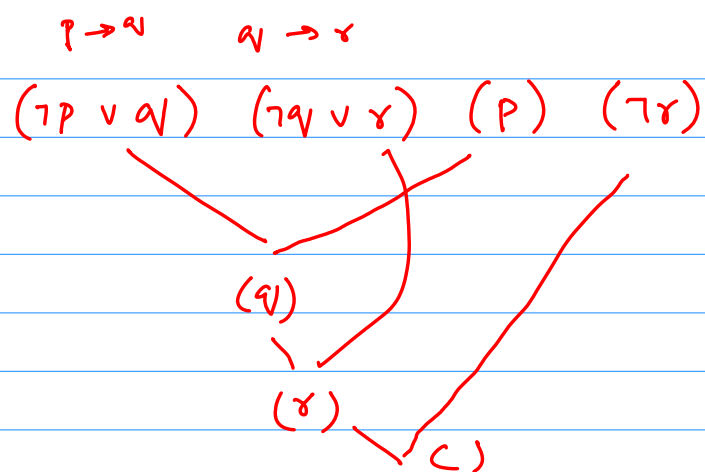
3^n clauses



Completeness of res. for unsatisfiability.

UNSAT formula. $(x_1) \wedge (\neg x_1)$ $()$ --- CNF on n vars





$$L(\underbrace{\varphi(x_1, \dots, x_n)}_{\cdot}) \subseteq \{0,1\}^n$$

$$\{0,1\} \times \{0,1\} \times \dots \times \{0,1\}$$

$$L(x_1 \oplus x_2 \oplus \dots \oplus x_n) \quad 0011010$$

$$=$$