

# Chapter 4

# Digital Transmission

Jirasak Sittigorn

Data Communications

Department of Computer Engineering, KMITL

# Digital Transmission

- DIGITAL-TO-DIGITAL CONVERSION
  - Line Coding
  - Block Coding
  - Scrambling
- ANALOG-TO-DIGITAL CONVERSION
  - Pulse Code Modulation (PCM)
  - Delta Modulation (DM)



# DIGITAL-TO-DIGITAL CONVERSION

- LINE CODING
- LINE CODING SCHEMES
- BLOCK CODING
- SCRAMBLING

# Line Coding

- the process of converting digital data to digital signals
  - converts a sequence of bits to a digital signal

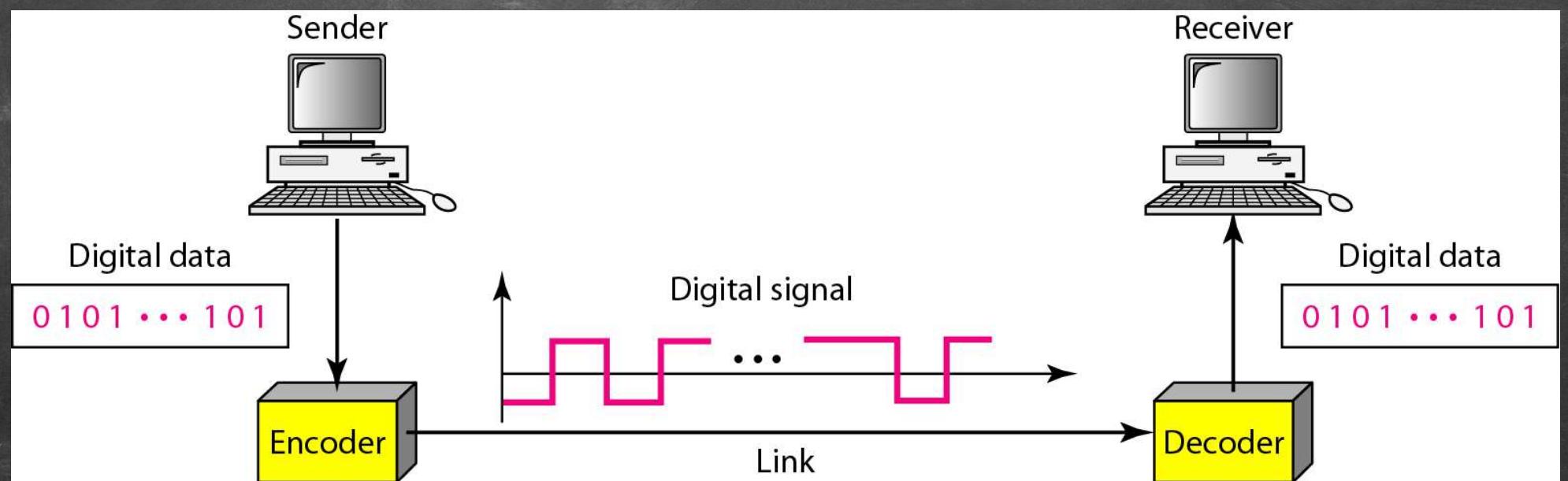
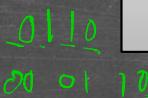


Figure 4.1 Line coding and decoding

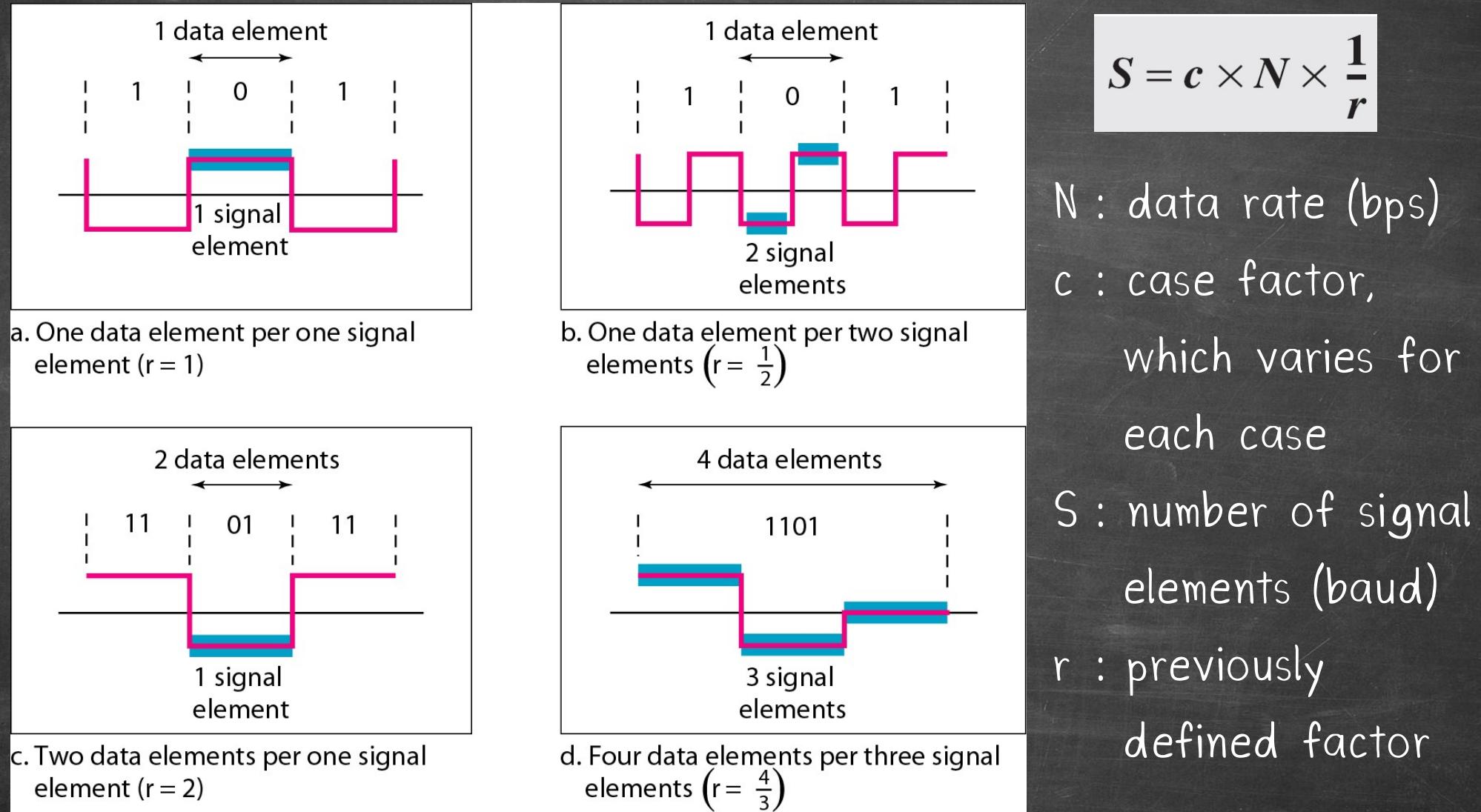
# Line Coding : Characteristics



- Signal Element are the carriers
- Data Element are being carried  
  
a signal element carries data elements
- Ratio ( $r$ ) is the number of data elements carried by each signal element  

- Data Rate (bit rate) defines the number of data elements (bits) sent in 1s : bits per second (bps)
- Signal Rate (pulse rate, modulation rate or baud rate) is the number of signal elements sent in 1s : baud

# Line Coding : Characteristics



$$S = c \times N \times \frac{1}{r}$$

N : data rate (bps)

c : case factor,  
which varies for  
each case

S : number of signal  
elements (baud)

r : previously  
defined factor

Figure 4.2 Signal element versus data element

## Example 4.1

$$S = c \times N \times \frac{1}{r}$$

- A signal is carrying data in which one data element is encoded as one signal element ( $r = 1$ ). If the bit rate is 100 kbps, what is the average value of the baud rate if  $c$  is between 0 and 1?
- Solution
  - We assume that the average value of  $c$  is  $1/2$ . The baud rate is then

$$S = c \times N \times \frac{1}{r} = \frac{1}{2} \times 100,000 \times \frac{1}{1} = 50,000 = 50 \text{ kbaud}$$

# Line Coding : Characteristics

- Bandwidth : Although the actual bandwidth of a digital signal is infinite (Chapter 3), the effective bandwidth is finite

## Example 4.2

$$S = c \times N \times \frac{1}{r}$$

- The maximum data rate of a channel (see Chapter 3) is  $N_{\max} = 2 \times B \times \log_2 L$  (defined by the Nyquist formula). Does this agree with the previous formula for  $N_{\max}$ ?
- Solution
  - A signal with  $L$  levels actually can carry  $\log_2 L$  bits per level. If each level corresponds to one signal element and we assume the average case ( $c = 1/2$ ), then we have

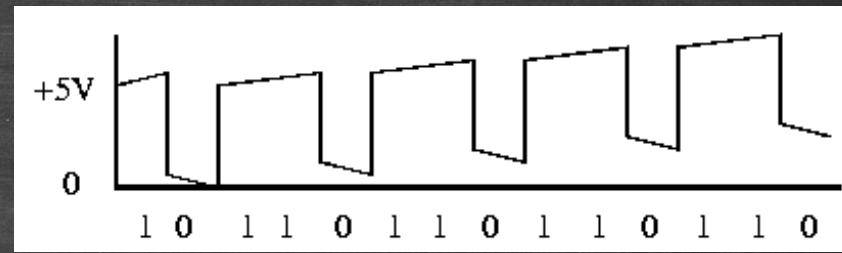
$$N_{\max} = \frac{1}{c} \times B \times r = 2 \times B \times \log_2 L$$

# Line Coding : Characteristics

- **Baseline Wandering**
  - Baseline is average of the received signal power [0, 1]
  - A long string of 0 or 1 can cause a drift in the baseline
- **DC Components**
  - voltage level in a digital signal is constant : very low frequencies
  - Problems : cannot pass low frequencies or a system that uses electrical coupling (via a transformer)

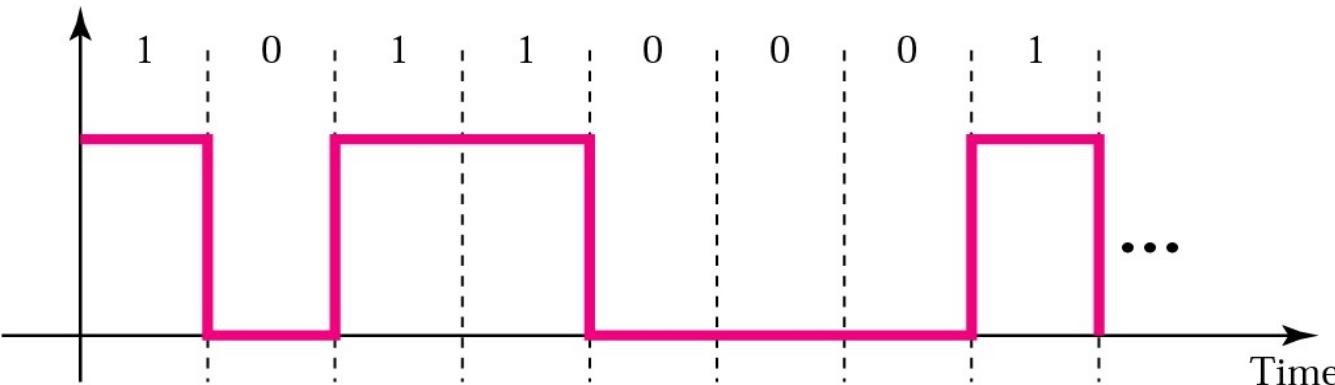
Wire A                          Wire B

Stray Capacitance  $\frac{1}{T}$   $v(t) = 1 - e^{-\frac{t}{RC}}$



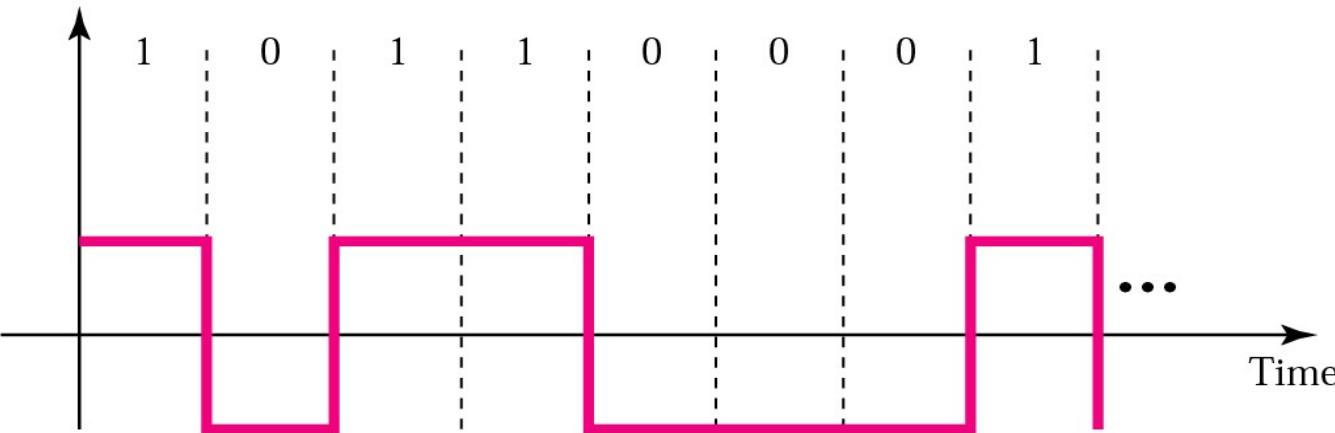
# Line Coding : Characteristics

Amplitude



a. A signal with dc component

Amplitude



b. A signal without dc component

# Line Coding : Characteristics

- Self-synchronization

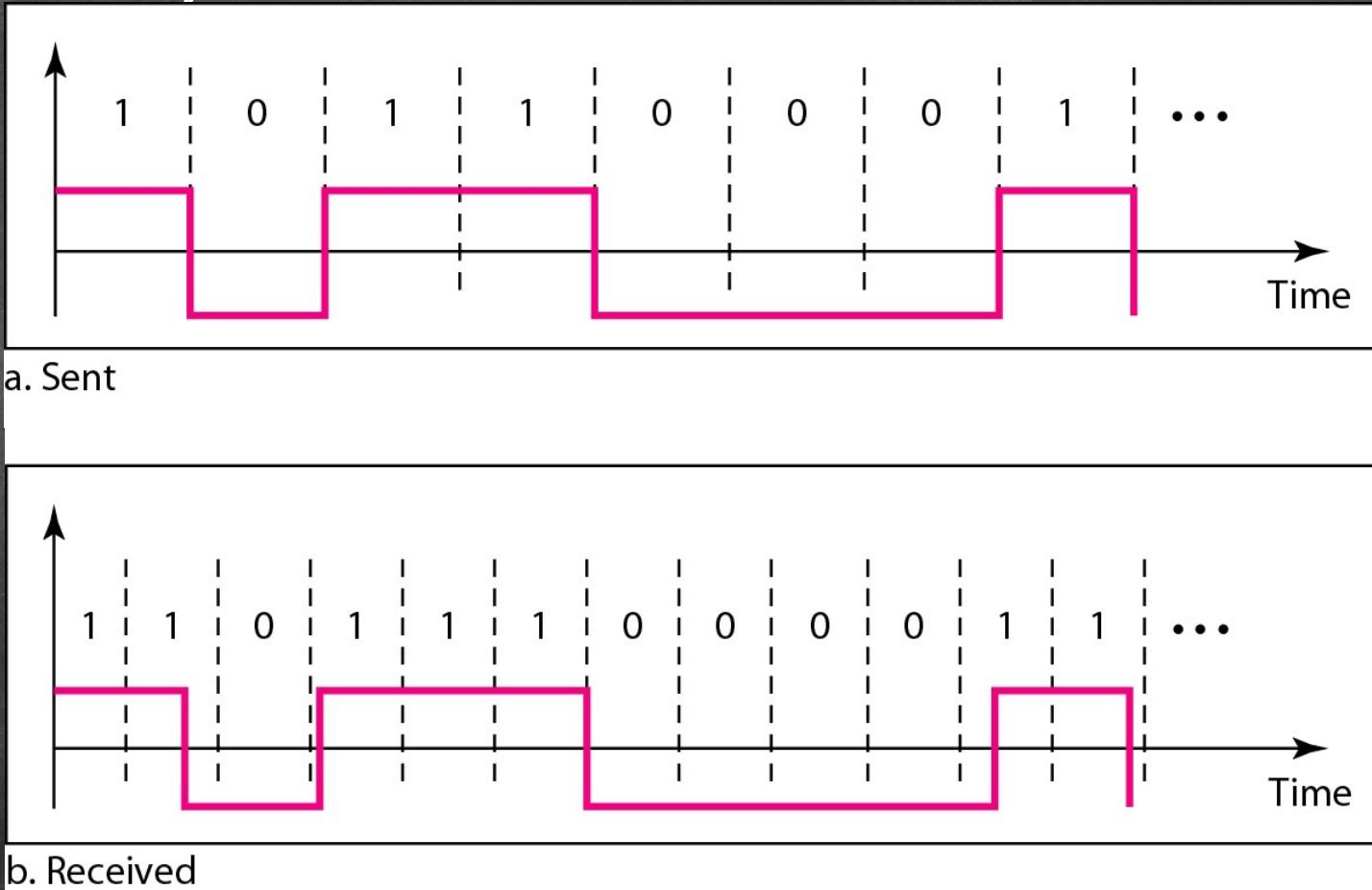


Figure 4.3 Effect of lack of synchronization

## Example 4.3

- In a digital transmission, the receiver clock is 0.1 percent faster than the sender clock. How many extra bits per second does the receiver receive if the data rate is 1 kbps? How many if the data rate is 1 Mbps?
  - Solution
    - At 1 kbps, the receiver receives 1001 bps instead of 1000 bps.
    - At 1 Mbps, the receiver receives 1,001,000 bps instead of 1,000,000 bps.

1000 bits sent

1001 bits received

1 extra bps

1,000,000 bits sent

1,001,000 bits received

1000 extra bps

# Line Coding Schemes

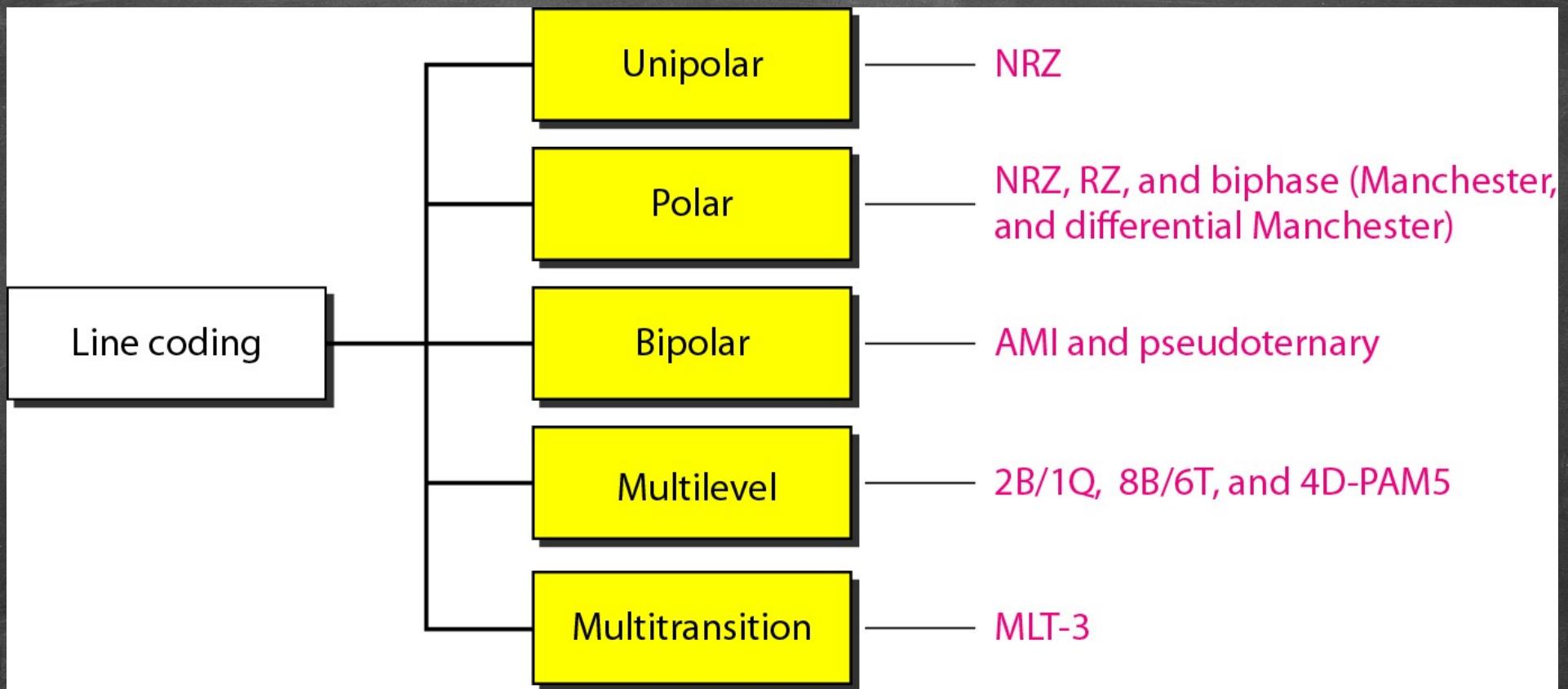


Figure 4.4 Line coding schemes

# Line Coding Schemes

- **Unipolar Scheme** : all the signal levels are on one side of the time axis, either above or below.
  - NRZ (Non-Return-to-Zero)
    - bit 1 = positive voltage (*maybe or negative*) / bit 0 = zero voltage

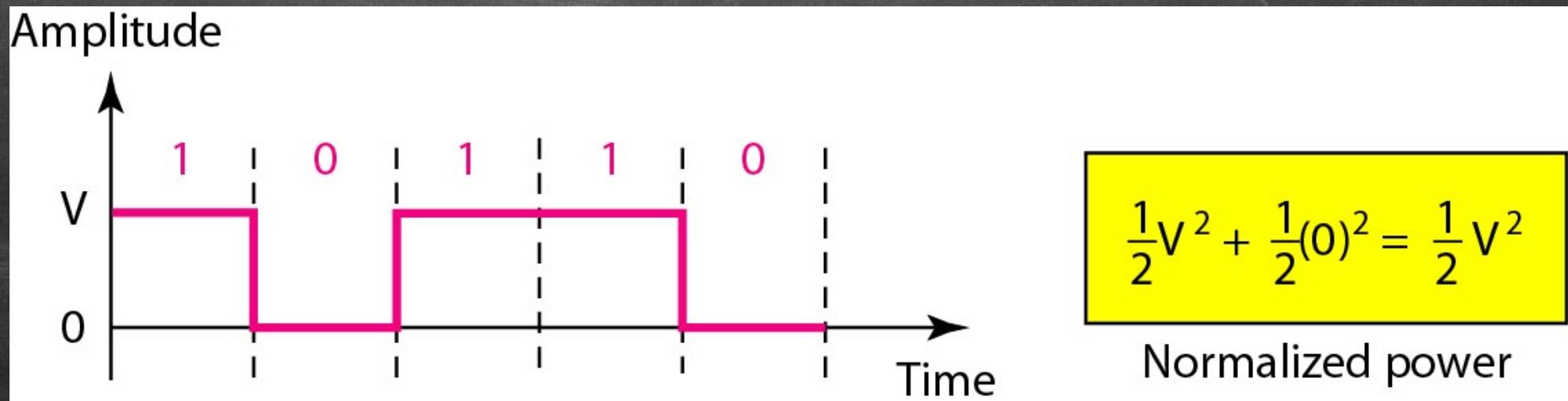
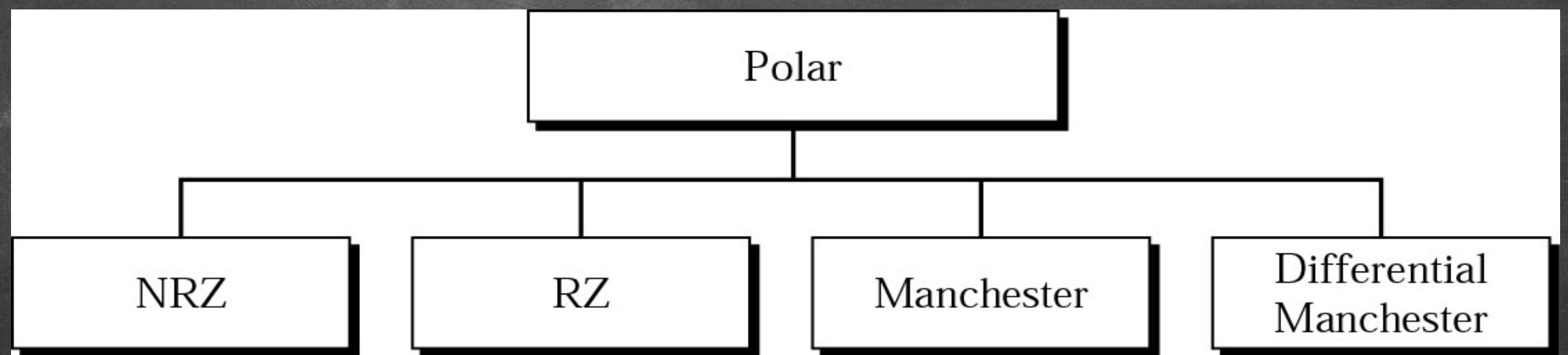


Figure 4.5 Unipolar NRZ scheme

B. A. Forouzan, Data Communications and Networking, 4th edition, McGRAW-HILL

# Line Coding Schemes

- Polar Schemes : the voltages are on the both sides of the time axis



# Line Coding Schemes

- Polar Schemes (Cont.)
  - Non-Return-to-Zero (NRZ)
  - NRZ-I (NRZ-Invert)

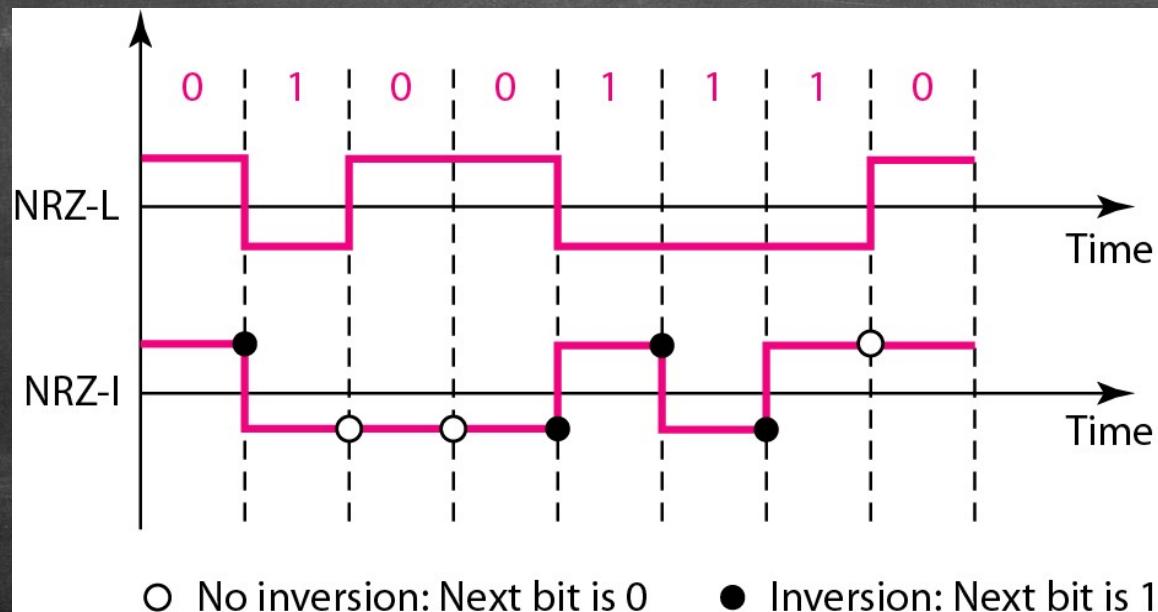
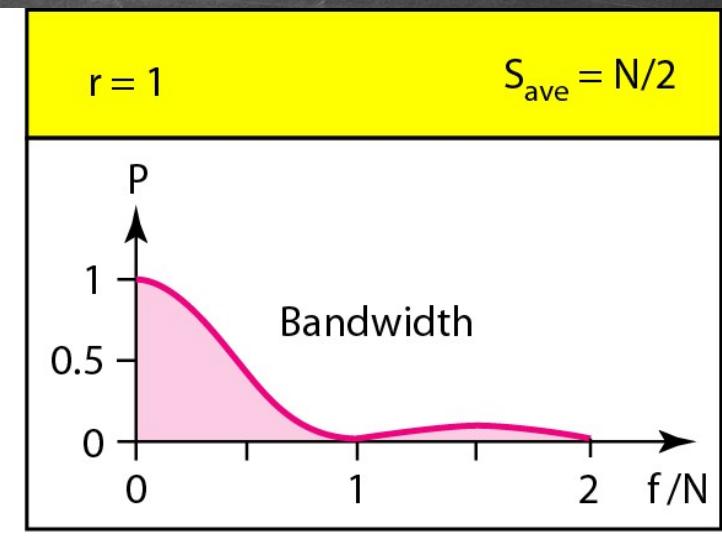


Figure 4.6 Polar NRZ-L and NRZ-I schemes



# Line Coding Schemes

- In NRZ-L the level of the voltage determines the value of the bit.
- In NRZ-I the inversion or the lack of inversion determines the value of the bit.
- NRZ-L and NRZ-I both have an average signal rate of  $N/2$  Bd.
- NRZ-L and NRZ-I both have a DC component problem.

## Example 4.4

- A system is using NRZ-I to transfer 10-Mbps data.  
What are the average signal rate and minimum bandwidth?
- Solution
  - The average signal rate is  $S = N/2 = 500$  baud. The minimum bandwidth for this average baud rate is  $B_{\min} = S = 500$  kHz.

# Line Coding Schemes

- Polar Schemes (Cont.)
  - Return to Zero (RZ)
  - A good encoded digital signal must contain a provision for synchronization.

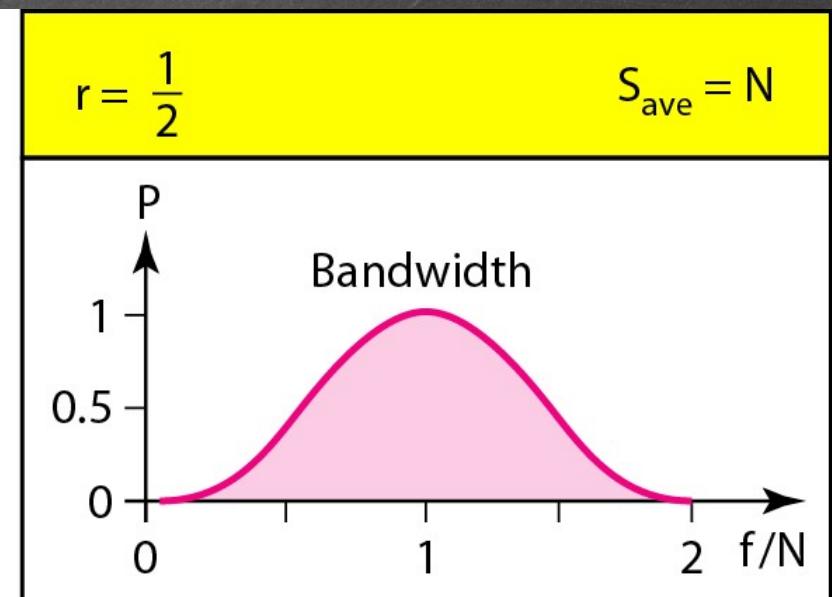
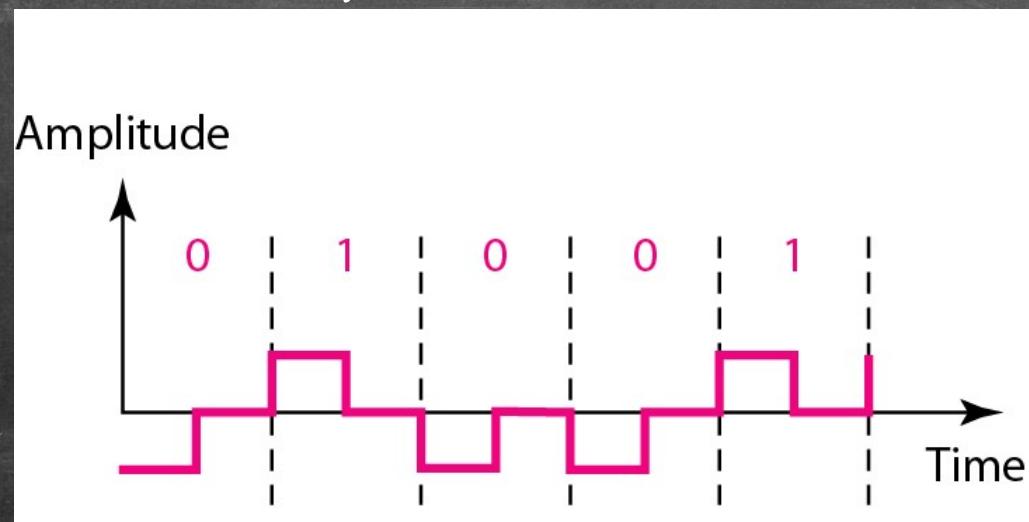


Figure 4.7 Polar RZ scheme

# Line Coding Schemes

- Polar Schemes (Cont.)
  - Biphasic : Manchester & Differential Manchester
    - the transition at the middle of the bit is used for synchronization.
    - The minimum bandwidth is 2 times that of NRZ.
    - In bipolar encoding, we use three levels: positive, zero, and negative.

# Line Coding Schemes

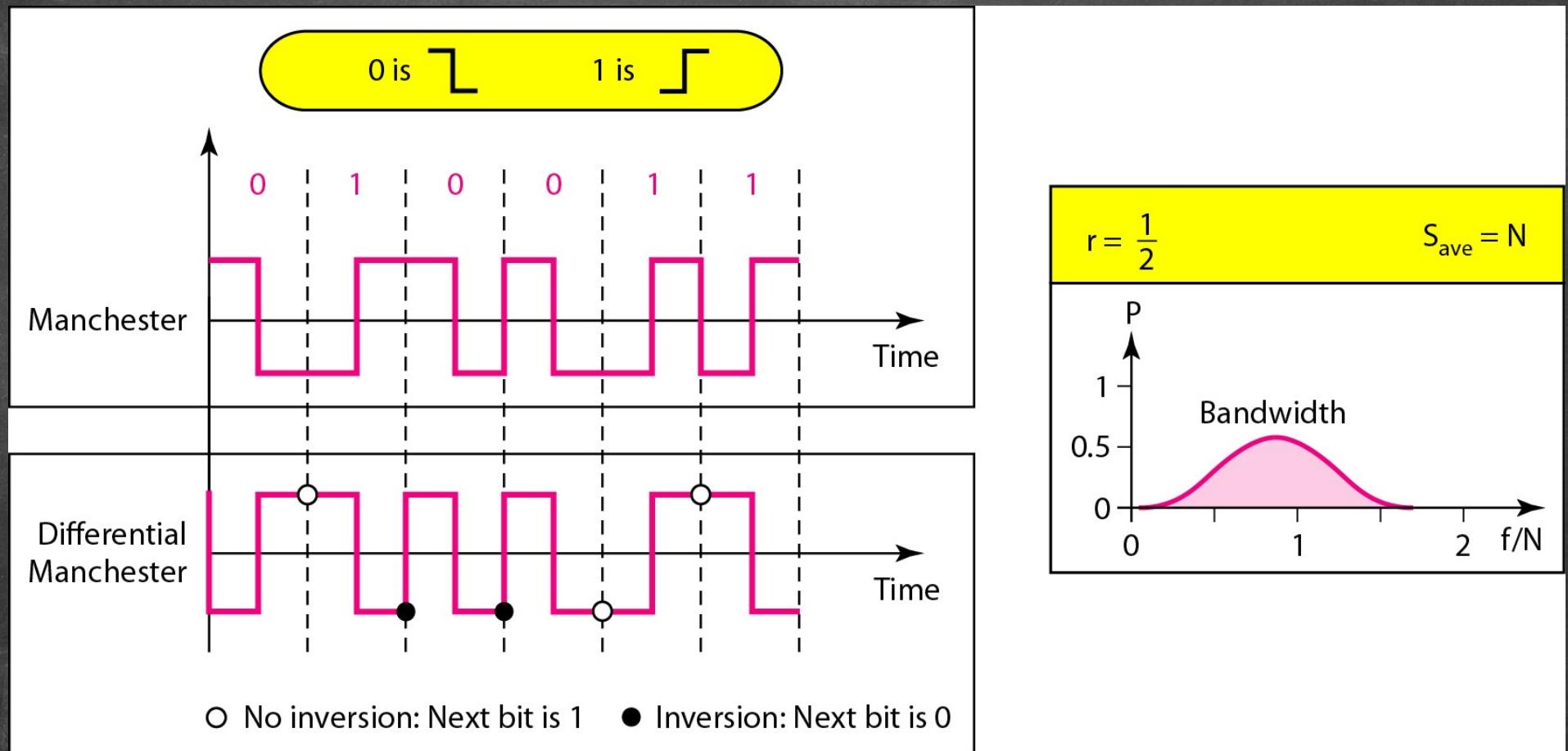


Figure 4.8 Polar biphasic: Manchester and differential Manchester schemes

# Line Coding Schemes

- Applications of Polar Schemes
  - NRZ encoding:
    - RS232 based protocols
  - Manchester encoding:
    - Ethernet networks
    - Hard drive
  - Differential Manchester encoding:
    - token-ring networks
  - NRZ-Inverted encoding:
    - Fiber Distributed Data Interface (FDDI)

# Line Coding Schemes

- **Bipolar Schemes** : bipolar encoding (multilevel binary) / three voltage levels : positive, negative, and zero / voltage level for one data element is at zero
  - AMI (alternate mark inversion)
  - Pseudoternary

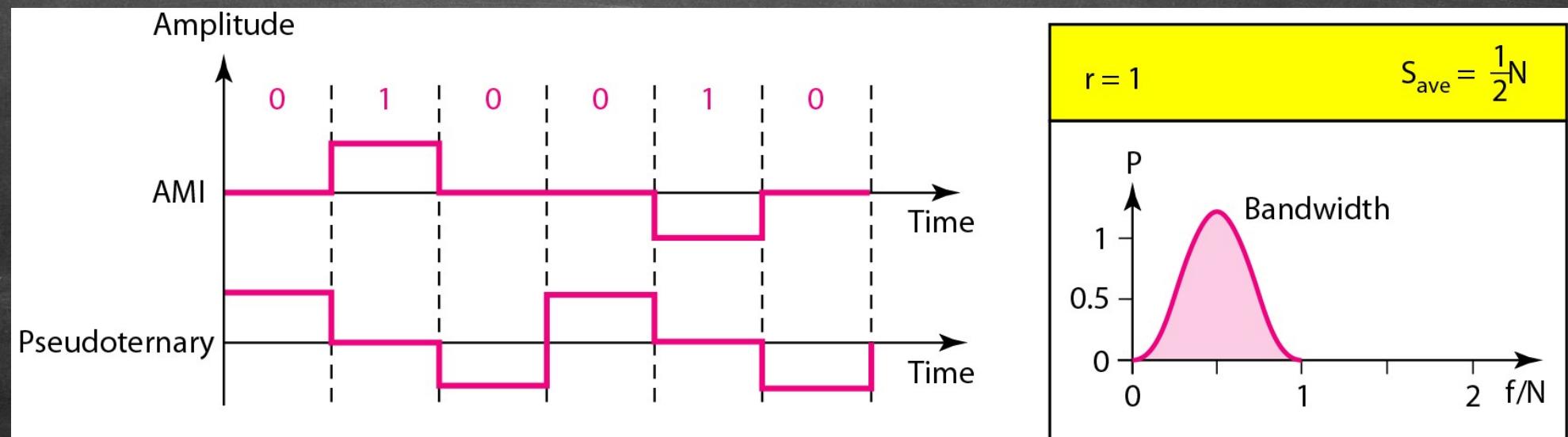


Figure 4.9 Bipolar schemes: AMI and pseudoternary

B. A. Forouzan, Data Communications and Networking, 4th edition, McGRAW-HILL

# Line Coding Schemes

- Multilevel Schemes : increase the number of bits per baud by encoding a pattern of  $m$  data elements into a pattern of  $n$  signal elements
  - data elements : combination of  $2^m$  data patterns
  - signal elements : combinations of  $L_n$  signal patterns
- $2^m = L_n$  : then each data pattern is encoded into one signal pattern
- $2^m < L_n$  : data patterns occupy only a subset of signal patterns
- $2^m > L_n$  because some of the data patterns cannot be encoded

# Line Coding Schemes

- Multilevel Schemes (Cont.)

- $mBnL$  : a pattern of  $m$  data elements is encoded as a pattern of  $n$  signal elements in which  $2^m \leq L_n$

- $m$  : length of the binary pattern
    - $B$  : binary data
    - $n$  : length of the signal pattern
    - $L$  : number of levels in the signaling

# Line Coding Schemes

- Multilevel Schemes (Cont.)
  - 2B1Q (two binary, one quaternary)
    - uses data patterns of size 2 and encodes the 2-bit patterns as one signal element belonging to a four-level signal ( $m=2$ ,  $n=1$ , and  $L=4$ )
    - used in DSL (Digital Subscriber Line)

# Line Coding Schemes

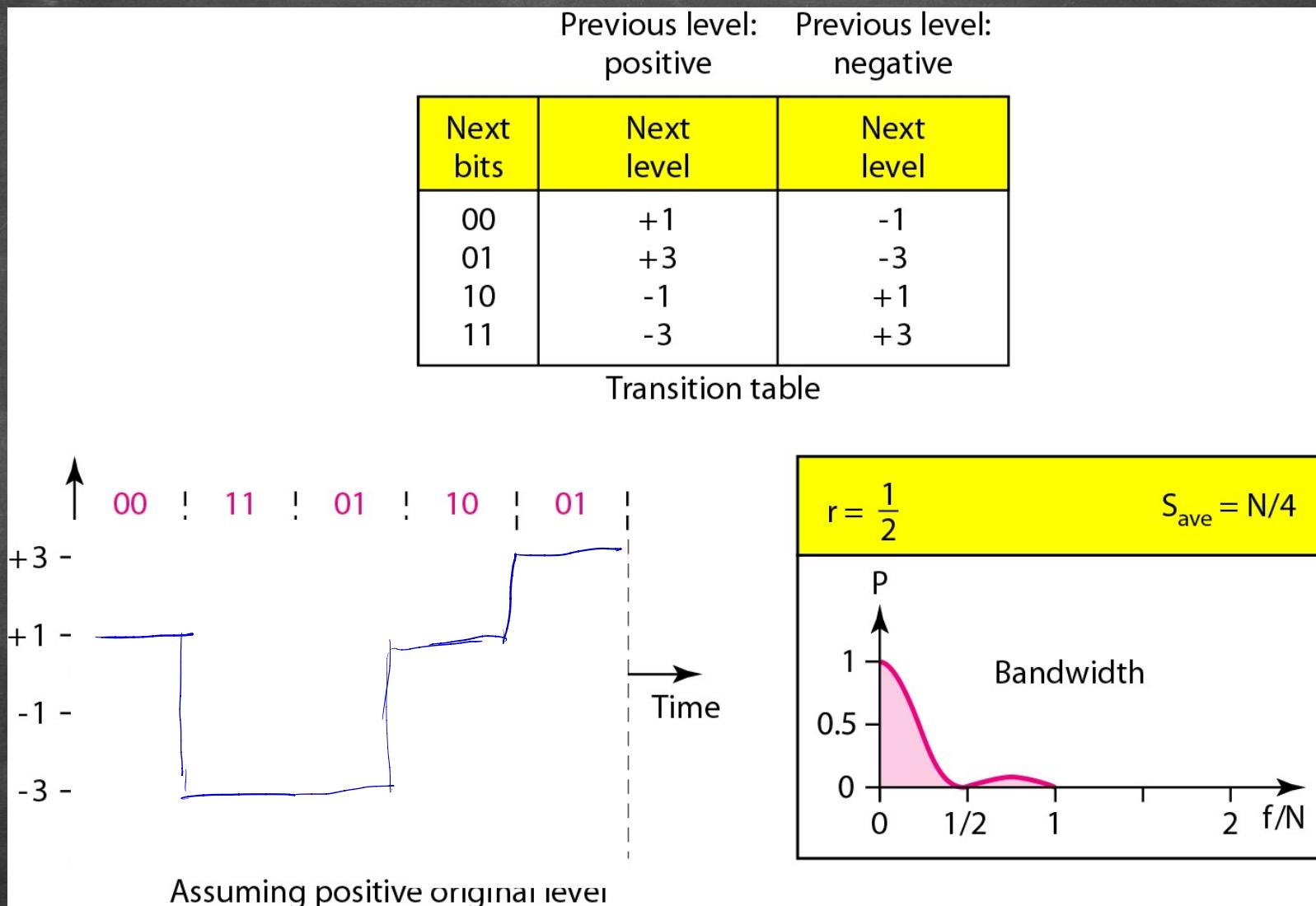


Figure 4.10 Multilevel: 2B1Q scheme

# Line Coding Schemes

- Multilevel Schemes (Cont.)

- 8B6T (eight binary, six ternary)

- pattern of 8 bits as a pattern of 6 signal elements, where the signal has three levels (ternary)
      - $2^8=256$  different data patterns
      - $3^6=478$  different signal patterns
    - is used with 100BASE-4T cable

APPENDIX F 8B6T Code							
The symbols in a collection of 8B6T code patterns. The 8-bit data sequence is transmitted in 6 pulses. The three levels of the waveform represent three symbols. Each symbol consists of two consecutive bits, which is equivalent to one byte of information.							
A							
B							
C							
D							
E							
F							
G							
H							
I							
J							
K							
L							
M							
N							
O							
P							

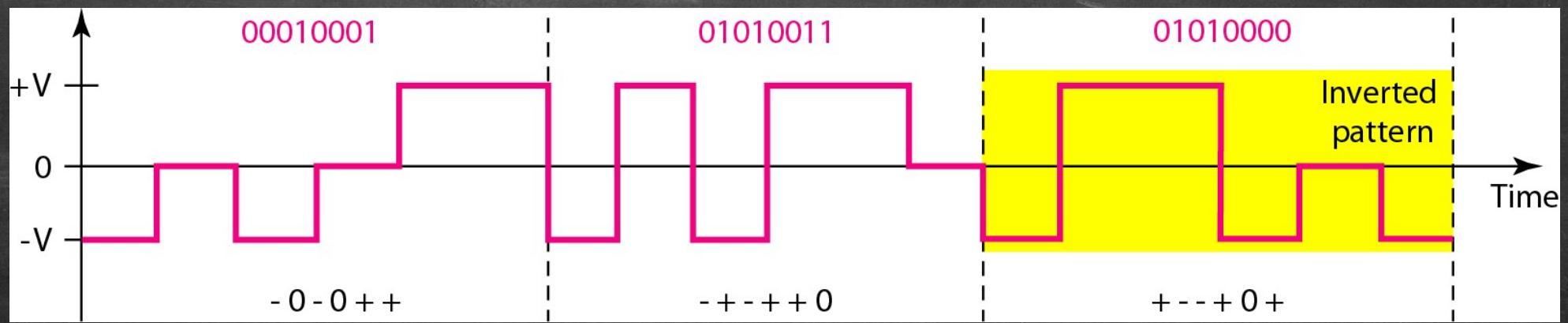


Figure 4.11 Multilevel: 8B6T scheme

B. A. Forouzan, Data Communications and Networking, 4th edition, McGRAW-HILL

# Line Coding Schemes

- Multilevel Schemes (Cont.)
  - 4D-PAMS (four dimensional five-level pulse amplitude modulation)
    - pattern of 8 bits as a pattern of 6 signal elements, where the signal has three levels (ternary)
      - $2^8$ -256 different data patterns
      - $3^6$ -478 different signal patterns
    - is used with 100BASE-4T cable

# Line Coding Schemes

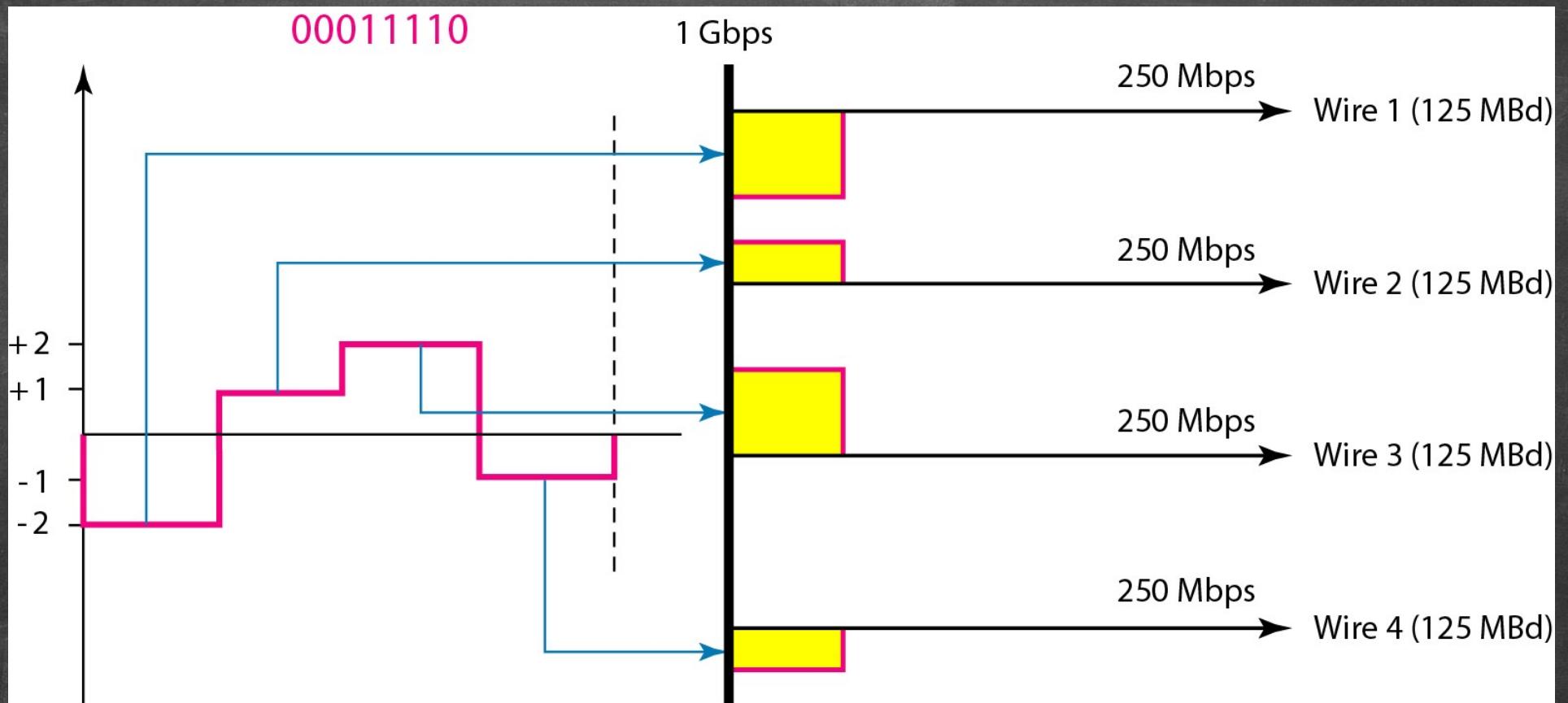


Figure 4.12 Multilevel: 4D-PAM5 scheme

# Line Coding Schemes

- Multiline Transmission : three level (MLT-3) scheme uses three levels (+v, 0, and - V) and three transition rules to move between the levels
  - If the next bit is 0, there is no transition.
  - If the next bit is 1 and the current level is not 0, the next level is 0.
  - If the next bit is 1 and the current level is 0, the next level is the opposite of the last nonzero level.
- First introduced by Cisco System for FDDI (Fiber Distributed Data Interface: Token Ring)
- Used in 100Base-TX (100 Mbit/s Ethernet)

# Line Coding Schemes

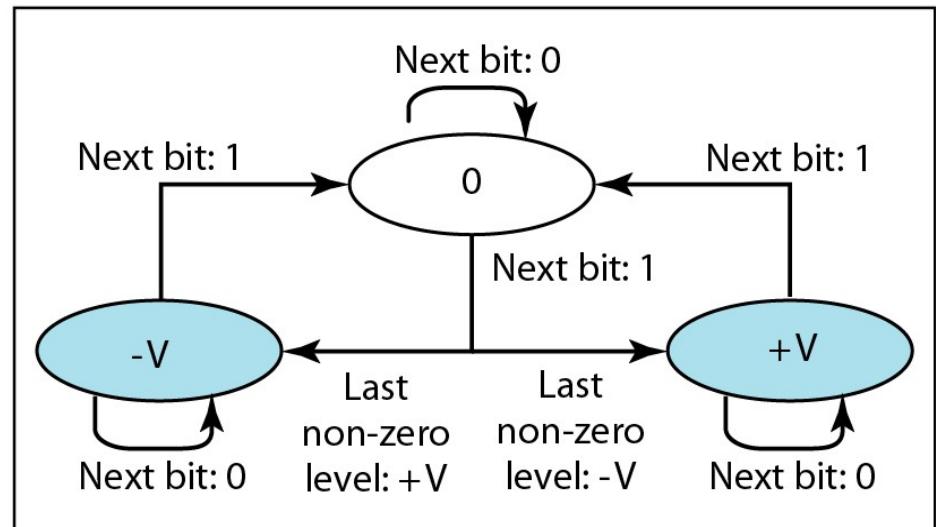
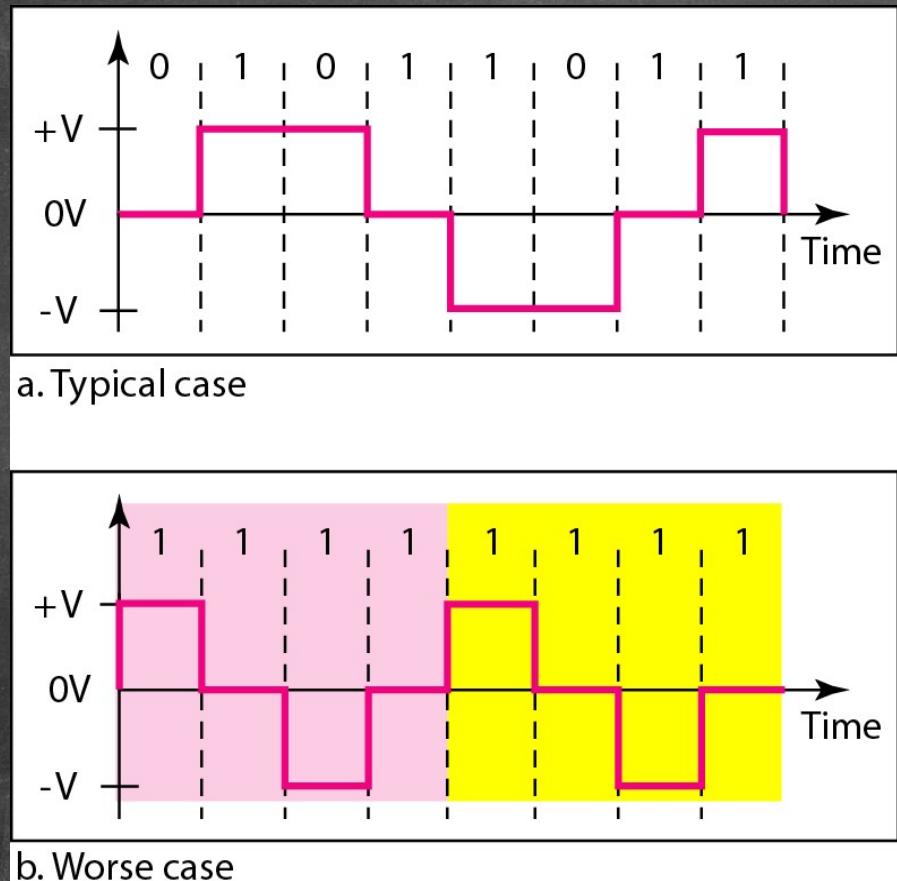


Figure 4.13 Multitransition: MLT-3 scheme

# Line Coding Schemes

<i>Category</i>	<i>Scheme</i>	<i>Bandwidth (average)</i>	<i>Characteristics</i>
Unipolar	NRZ	$B = N/2$	Costly, no self-synchronization if long 0s or 1s, DC
Unipolar	NRZ-L	$B = N/2$	No self-synchronization if long 0s or 1s, DC
	NRZ-I	$B = N/2$	No self-synchronization for long 0s, DC
	Biphase	$B = N$	Self-synchronization, no DC, high bandwidth
Bipolar	AMI	$B = N/2$	No self-synchronization for long 0s, DC
Multilevel	2B1Q	$B = N/4$	No self-synchronization for long same double bits
	8B6T	$B = 3N/4$	Self-synchronization, no DC
	4D-PAM5	$B = N/8$	Self-synchronization, no DC
Multiline	MLT-3	$B = N/3$	No self-synchronization for long 0s

Table 4.1 Summary of line coding schemes

# Block coding

- Block coding is normally referred to as  $mB/nB$  coding; it replaces each  $m$ -bit group with an  $n$ -bit group.

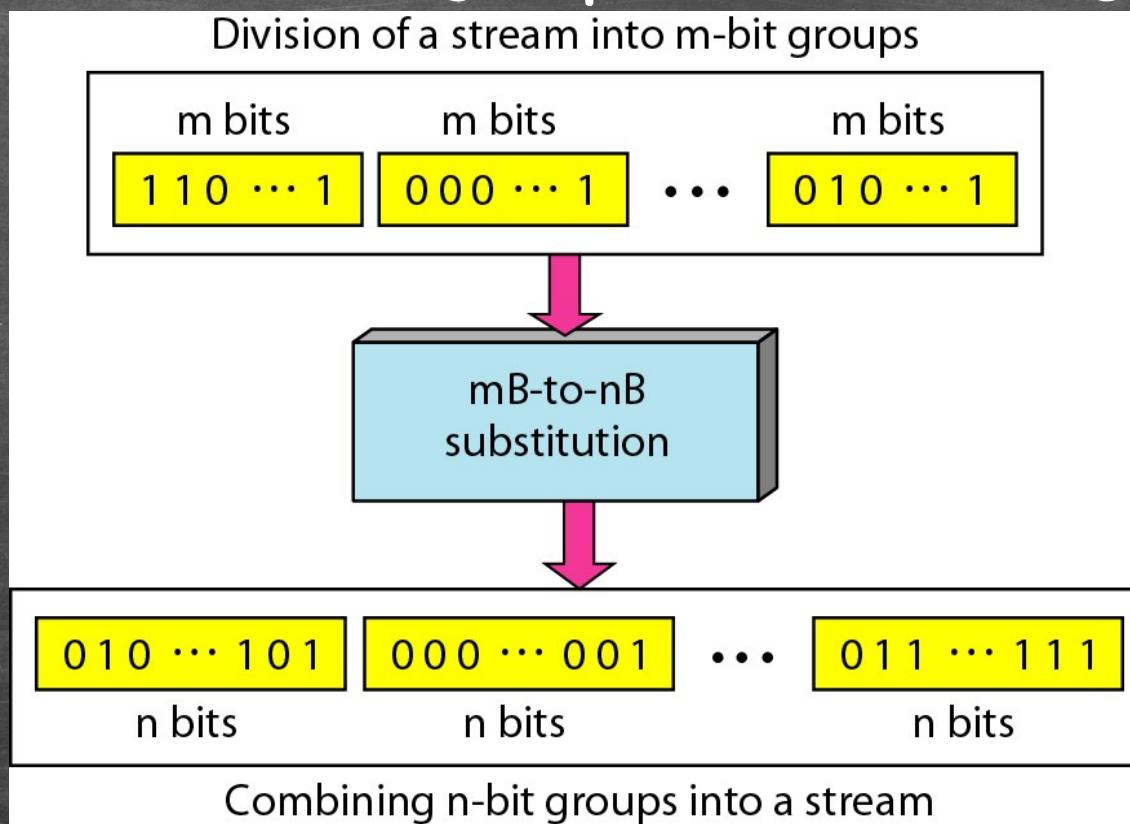


Figure 4.14 Block coding concept

# Block coding

- 4B/5B
  - NRZ-I has a good signal rate : one-half that of the biphase, but it has a synchronization problem

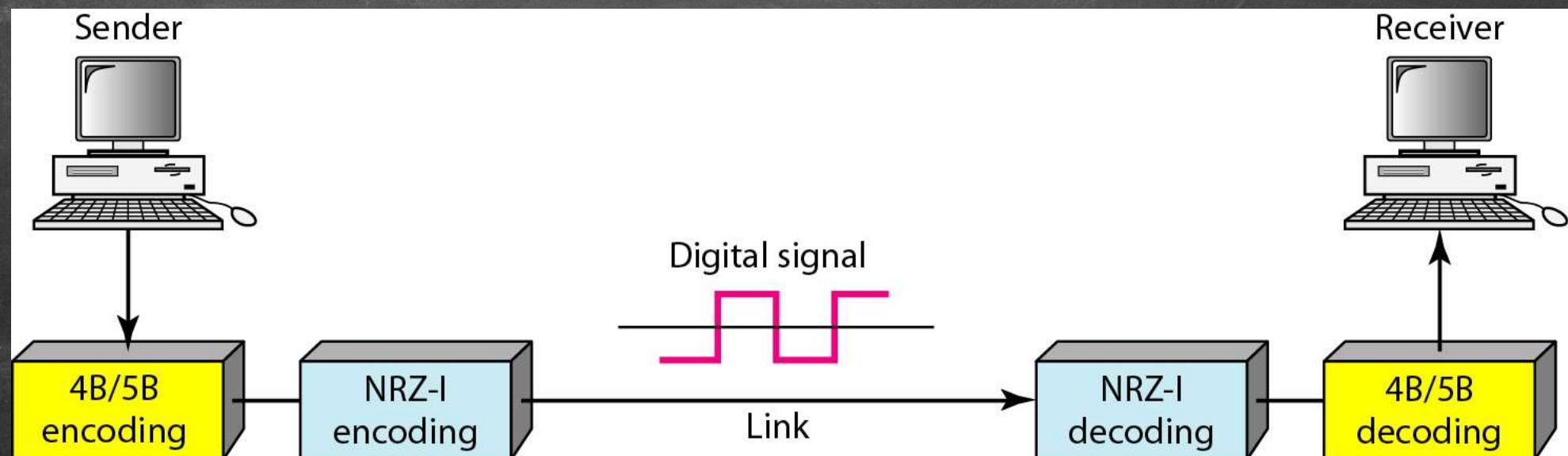


Figure 4.15 Using block coding 4B/5B with NRZ-I line coding scheme

# Table 4.2 4B/5B mapping codes

<i>Data Sequence</i>	<i>Encoded Sequence</i>	<i>Control Sequence</i>	<i>Encoded Sequence</i>
0000	11110	Q (Quiet)	00000
0001	01001	I (Idle)	11111
0010	10100	H (Halt)	00100
0011	10101	J (Start delimiter)	11000
0100	01010	K (Start delimiter)	10001
0101	01011	T (End delimiter)	01101
0110	01110	S (Set)	11001
0111	01111	R (Reset)	00111
1000	10010		
1001	10011		
1010	10110		
1011	10111		
1100	11010		
1101	11011		
1110	11100		
1111	11101		

# Block coding

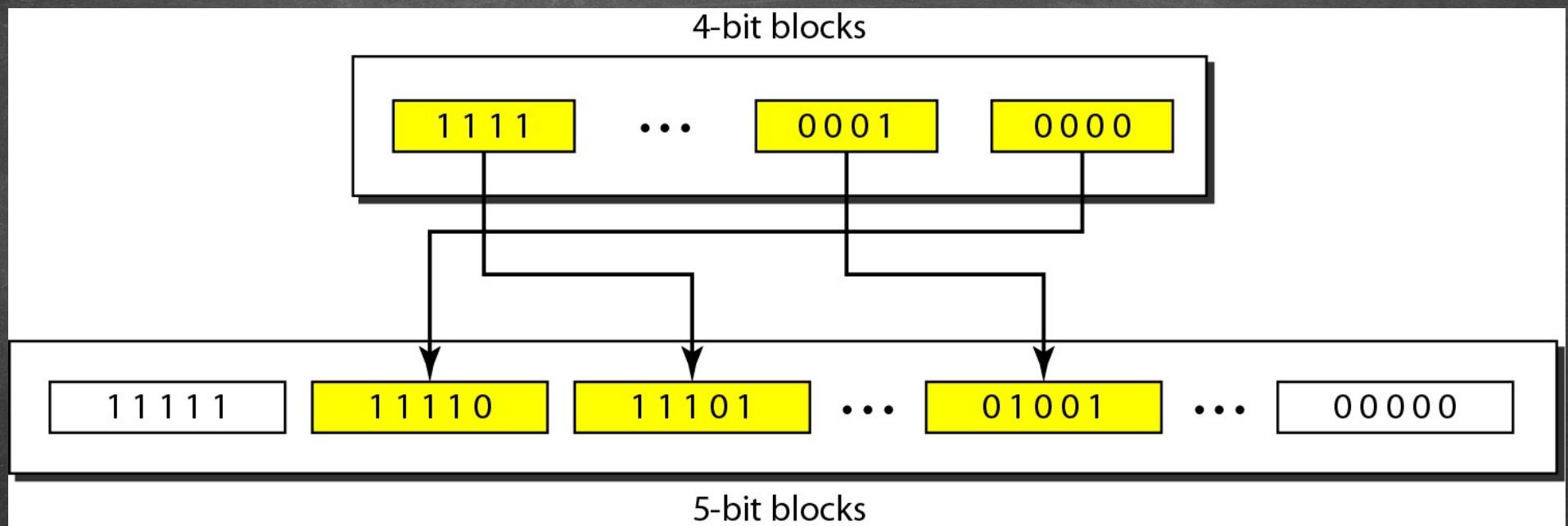


Figure 4.16 Substitution in 4B/5B block coding

## Example 4.5

- We need to send data at a 1-Mbps rate. What is the minimum required bandwidth, using a combination of 4B/5B and NRZ-I or Manchester coding?
- Solution
  - First 4B/5B block coding increases the bit rate to 1.25 Mbps. The minimum bandwidth using NRZ-I is  $N/2$  or 625 kHz. The Manchester scheme needs a minimum bandwidth of 1 MHz. The first choice needs a lower bandwidth, but has a DC component problem; the second choice needs a higher bandwidth, but does not have a DC component problem.

# Block coding

- 8B/10B

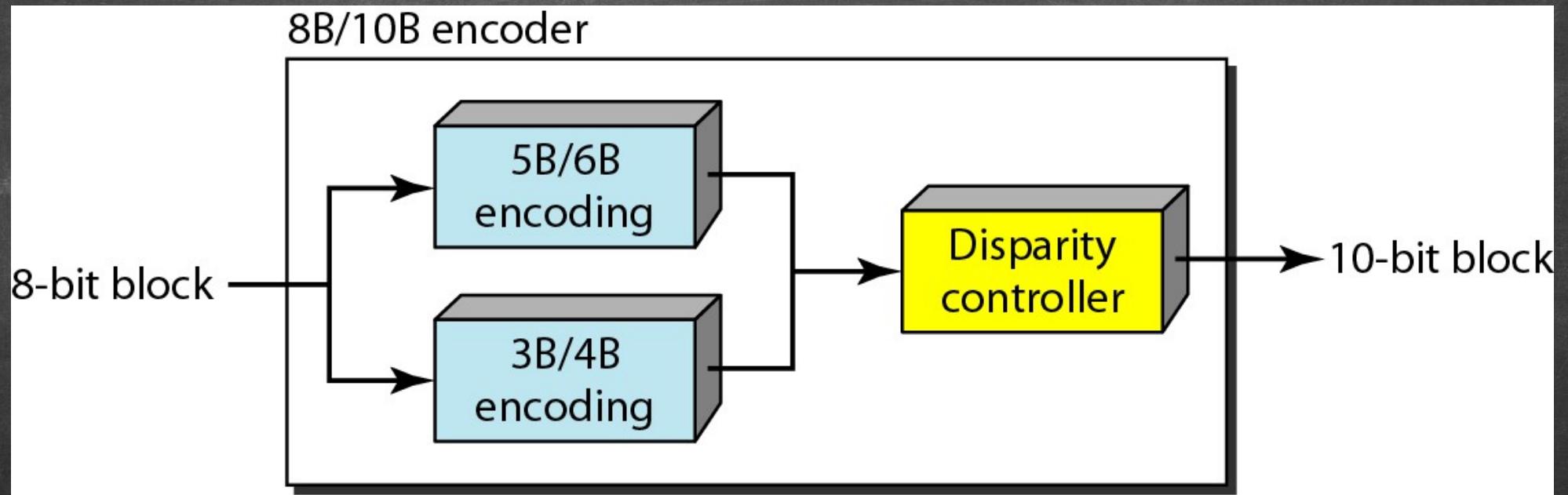


Figure 4.17 8B/10B block encoding

# Scrambling

- Biphase schemes => LAN
  - Not suitable for long-distance communication (wide bandwidth requirement)
- Block coding and NRZ
  - Not suitable for long-distance encoding either (DC component)
- Bipolar AMI encoding
  - narrow bandwidth and does not create a DC component
  - void a long sequence of 0s

# Scrambling

උරුම සංස්කීරණය

- B8ZS substitutes eight consecutive zeros with 000VB0VB.
- HDB3 substitutes four consecutive zeros with 000V or B00V depending on the number of nonzero pulses after the last substitution.

උරුම සංස්කීරණය

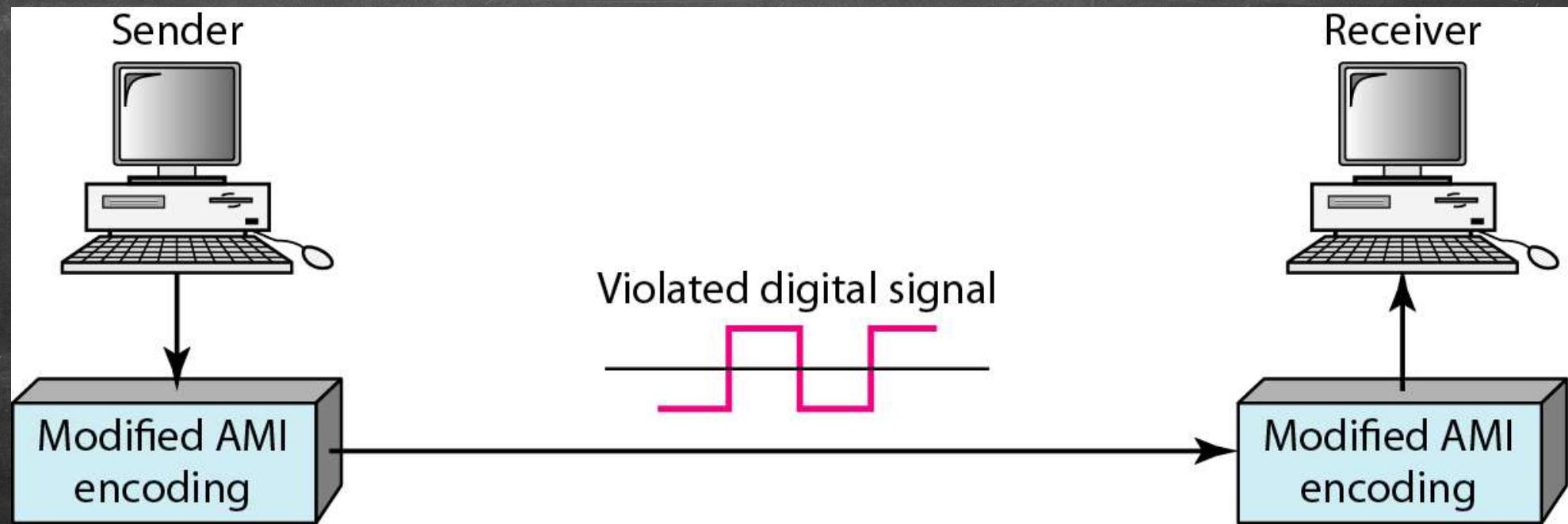
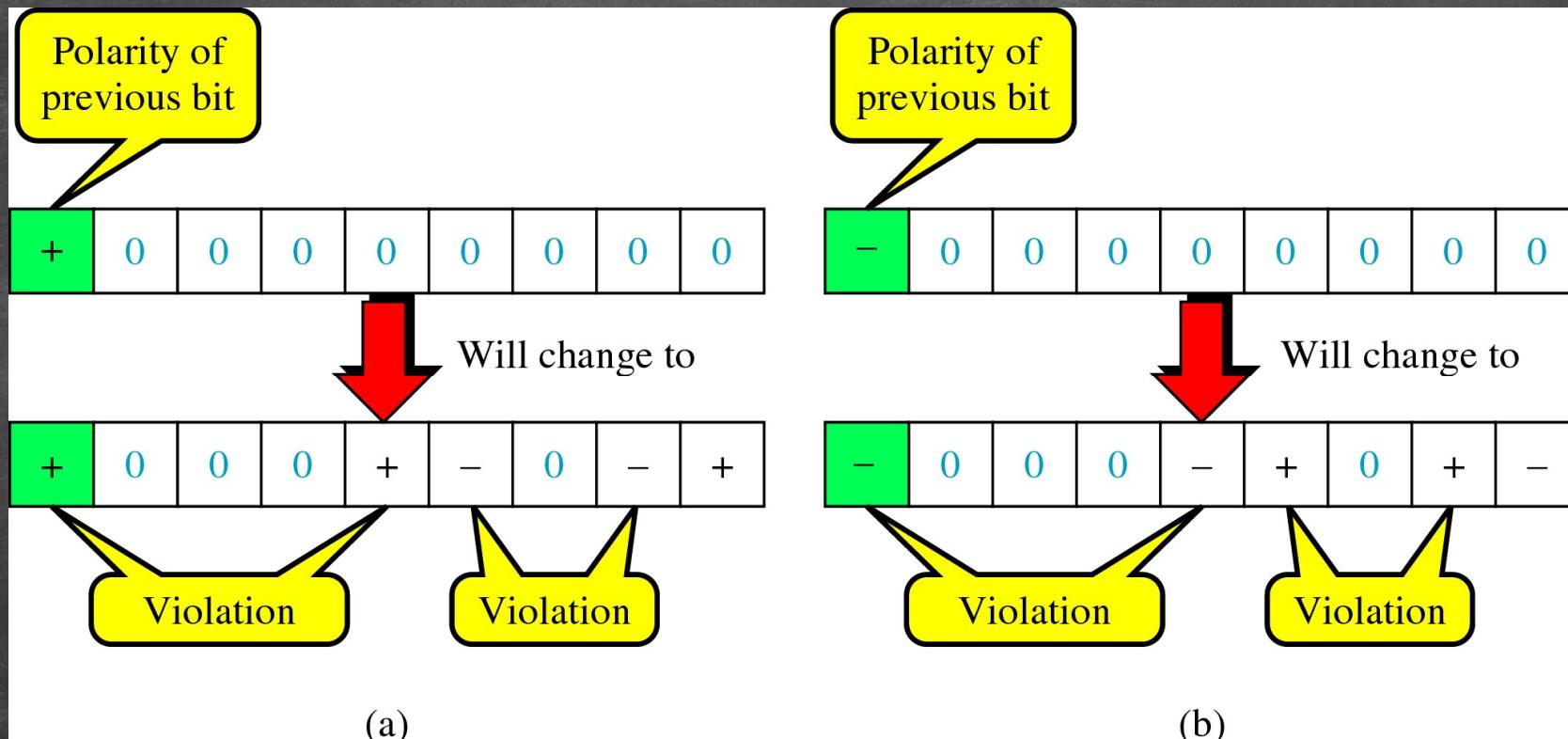


Figure 4.18 AMI used with scrambling

B. A. Forouzan, Data Communications and Networking, 4th edition, McGRAW-HILL

# Scrambling

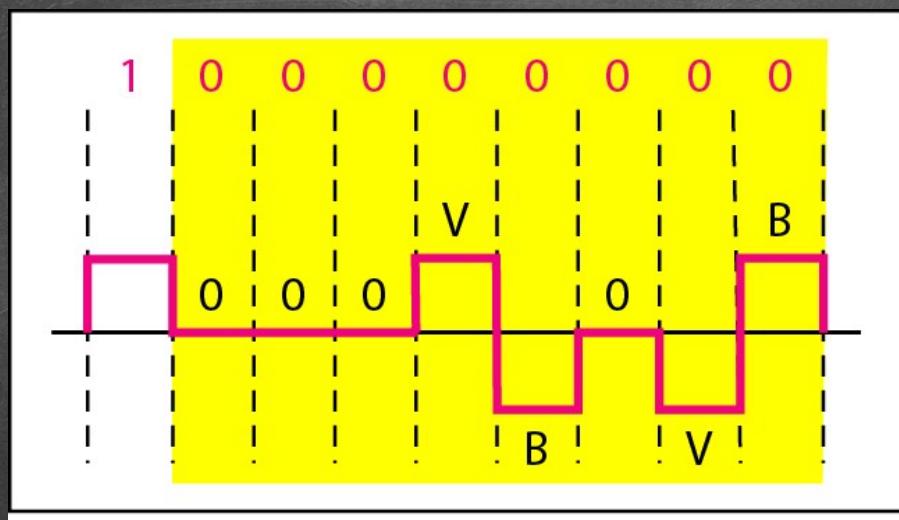
- B8ZS : Bipolar with S-zero substitution



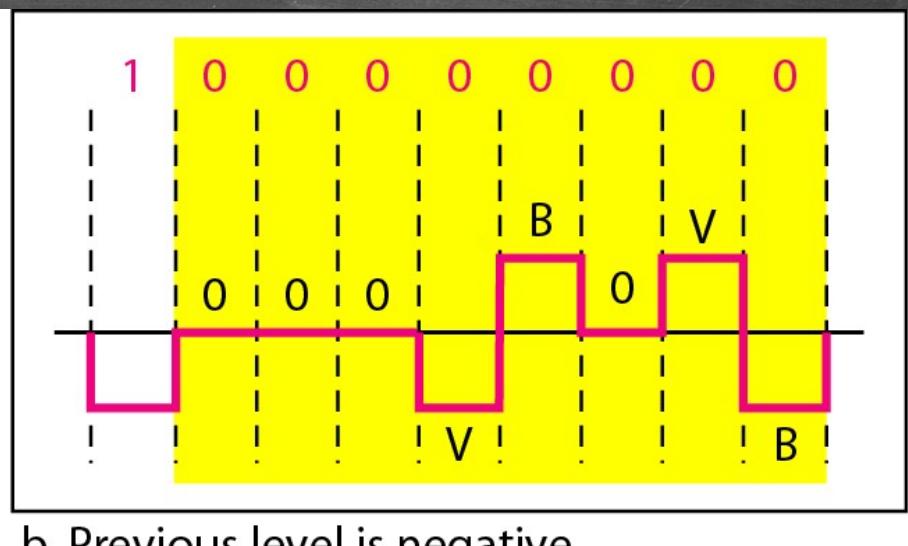
0 0 0 V B 0 V B

V = Violation, B = Bipolar

# Scrambling



a. Previous level is positive.

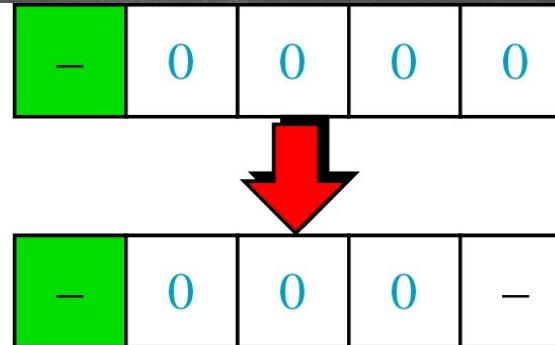
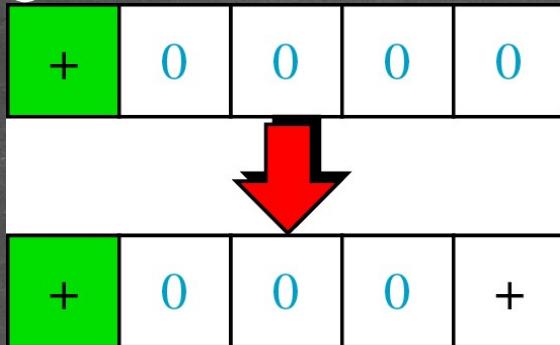


b. Previous level is negative.

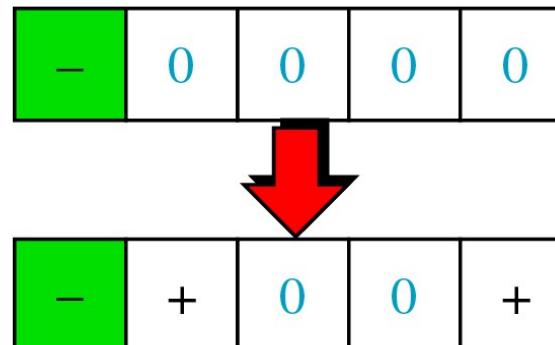
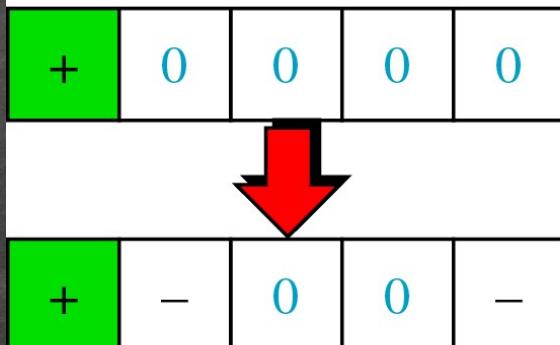
Figure 4.19 Two cases of B8ZS scrambling technique

# Scrambling

- HDB3



(a) If the number of 1s since the last substitution is odd



(b) If the number of 1s since the last substitution is even

↑ ປົກ  
V

0 0 0 V

or

B 0 0 V

↑ ປົກ

# Scrambling

000V or B00V

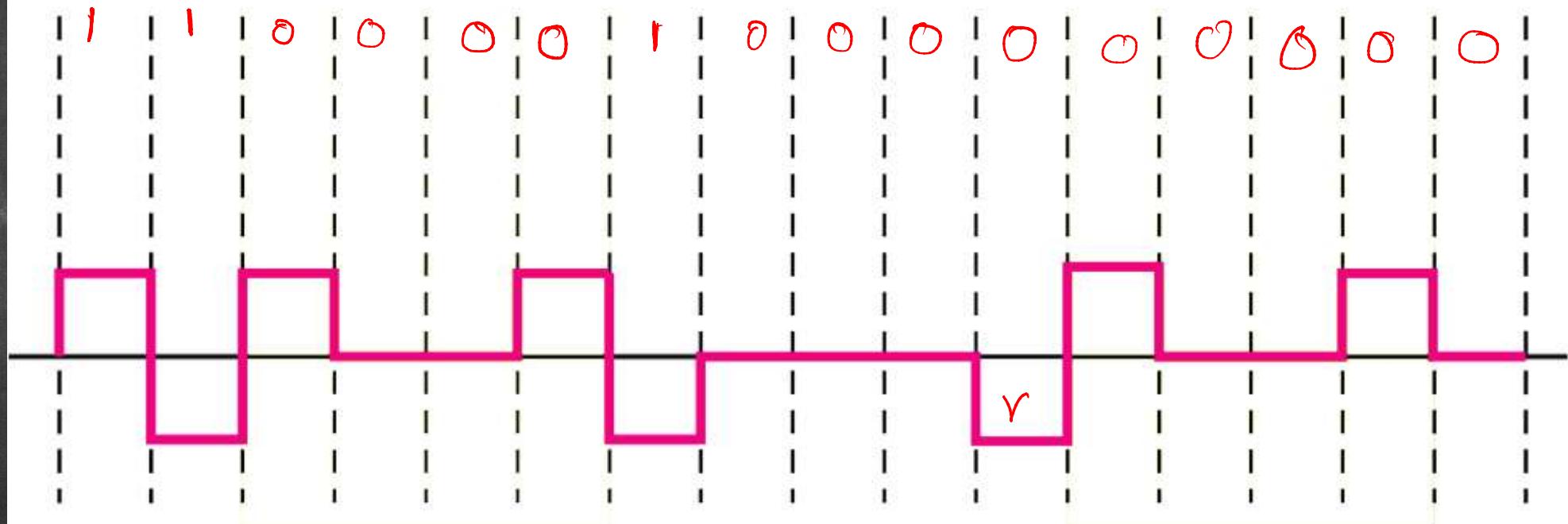


Figure 4.20 Different situations in HDB3 scrambling technique