

KRUSKAL'S MST

COP4533 PROJECT

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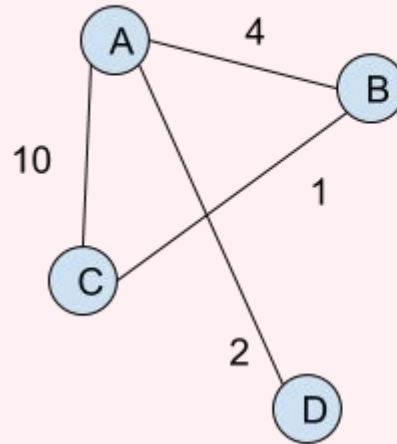
ALGORITHM OVERVIEW

Background

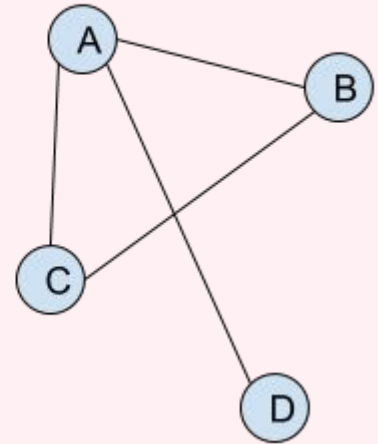
Graph: A structure comprising vertices (nodes) connected by edges (lines).

Weighted Graph: A graph where each edge has an associated weight (e.g., cost, distance, time).

Weighted Graph



Unweighted Graph

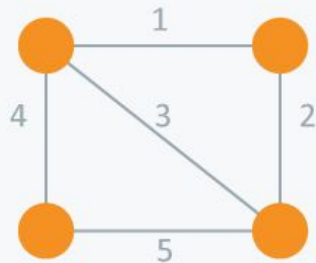


ALGORITHM OVERVIEW

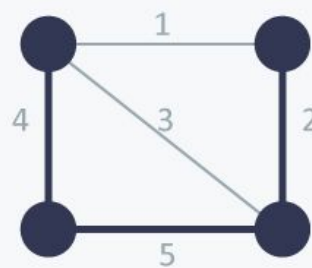
Background

Spanning Tree: A subgraph that connects all vertices with the minimum number of edges and no cycles.

Minimum Spanning Tree: A spanning tree with the lowest possible total edge weight.

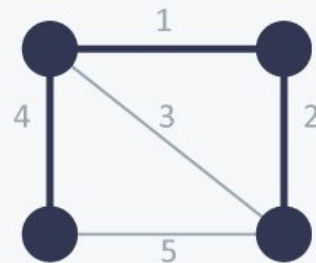


Undirected
Graph



Spanning
Tree

Cost = $11 (= 4 + 5 + 2)$



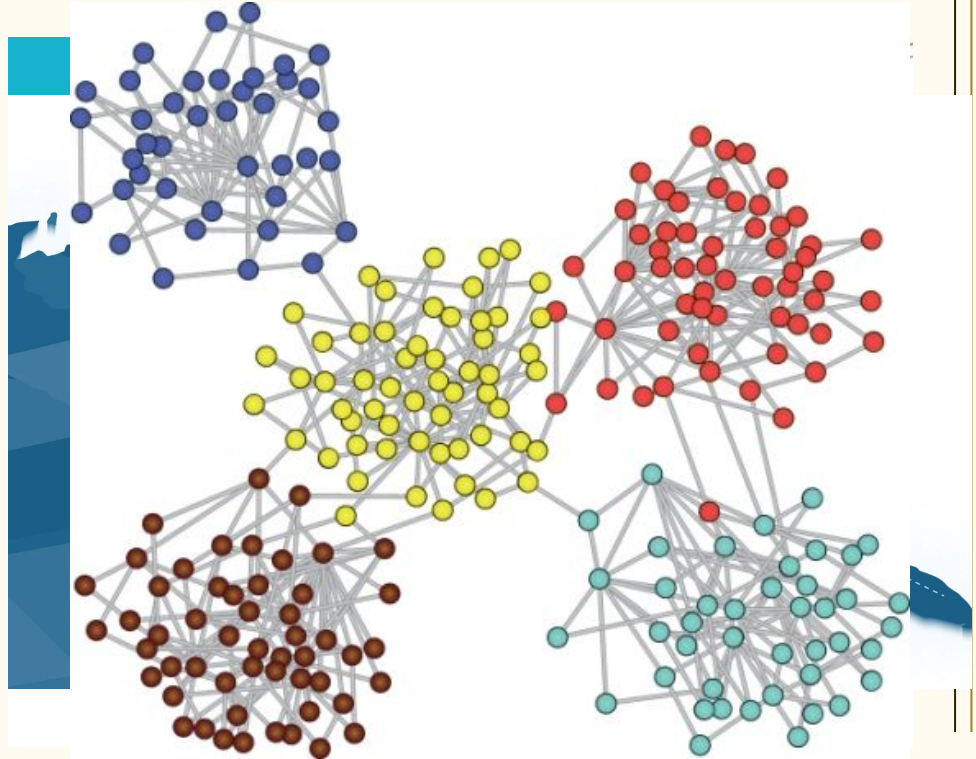
Minimum Spanning
Tree

Cost = $7 (= 4 + 1 + 2)$

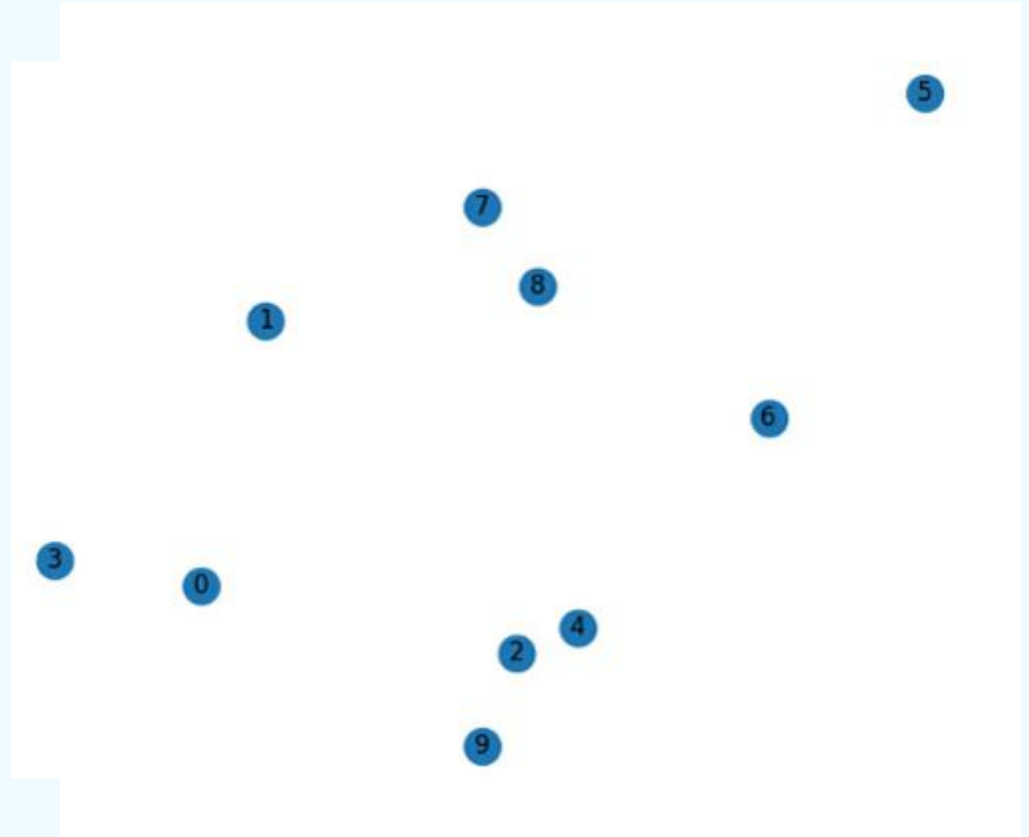
↑ This is what Kruskal's
Algorithm finds.

Purpose & Real World Examples

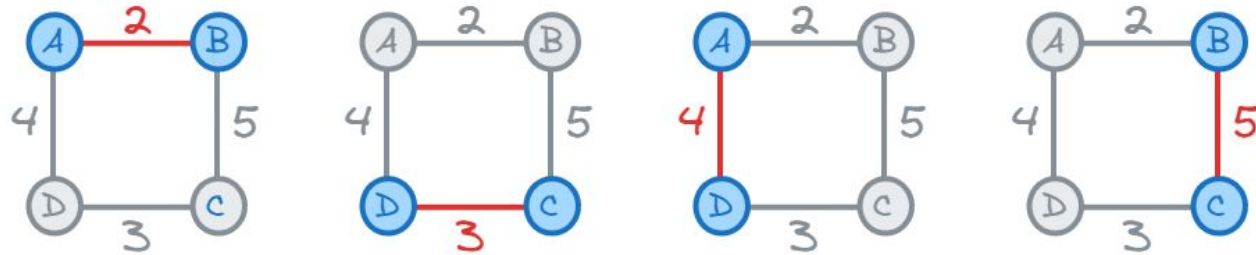
Find the cheapest way to connect all nodes in a network without creating any loops.



How Kruskal's Algorithm works

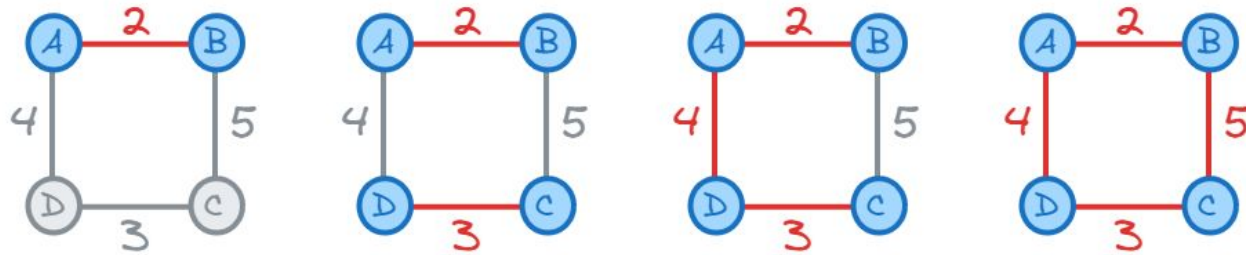


1. Sort Edges by Weight



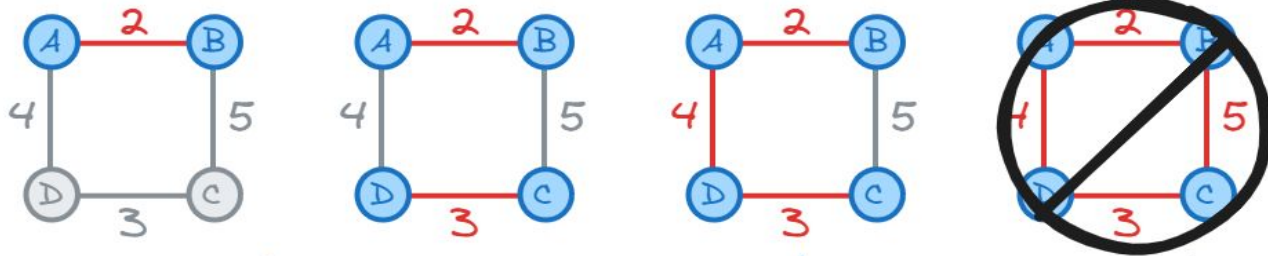
We'll create an array which will hold an array of each edge in our graph, then sort the array on edge weight.

2. Add Edges While Avoiding Cycles



We then construct the MST using the edges, checking to ensure we don't create any cycles (loops).

3. Terminate When MST is Complete



When the edges we construct are 1 less than the number of vertices, we have completed the MST!

Implementation Details

Code Review

Challenges Faced and How We Addressed Them:

- Lauren had to learn how Kruskal's MST worked again
- Lydia had all the DSU code written in Java already, but she had to translate it into C++ for the sake of this implementation

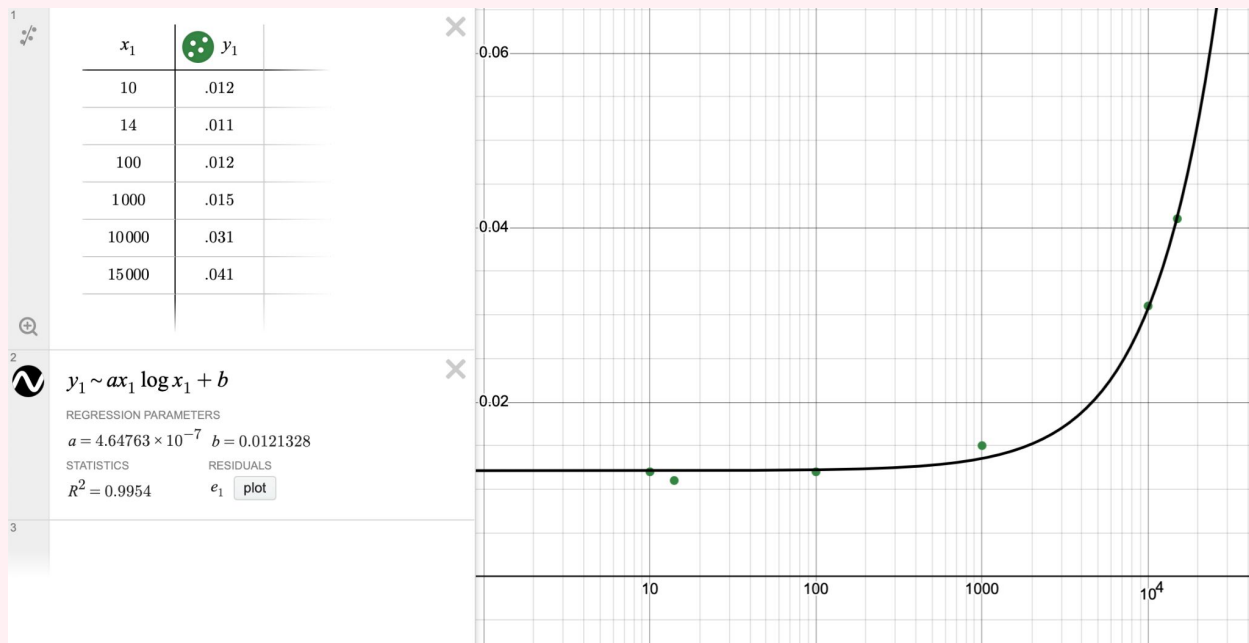


PERFORMANCE ANALYSIS

RUNNING TIME

Background

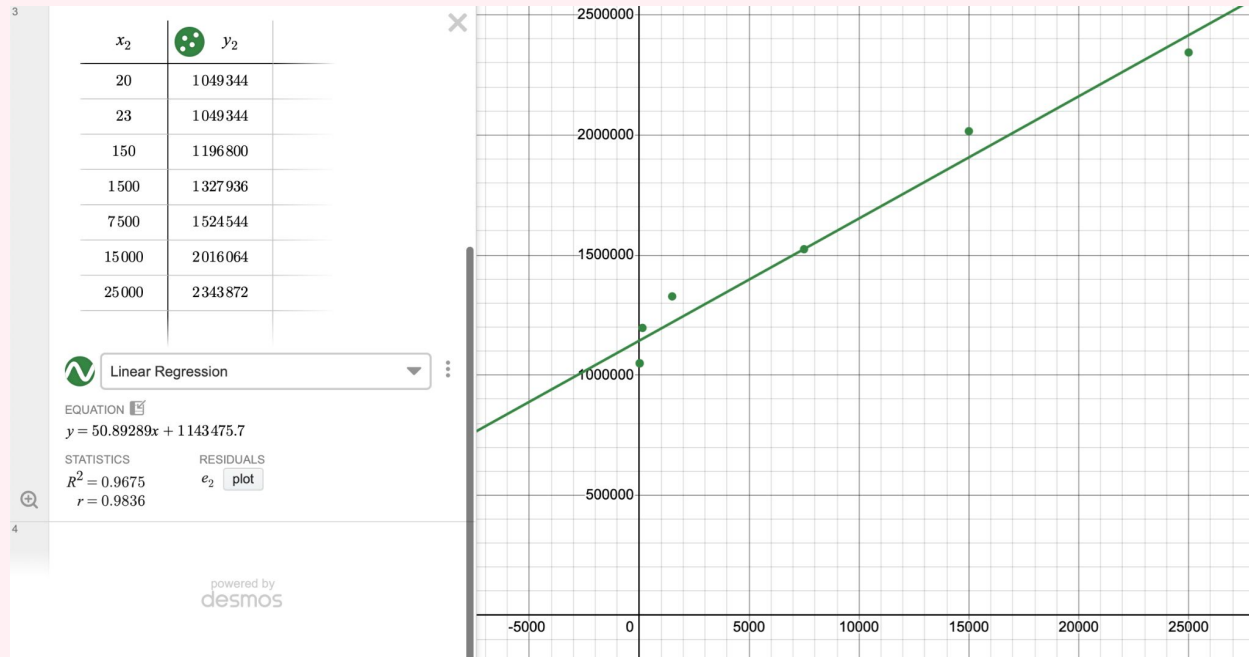
Kruskal's Algorithm runs in $O(E \log E)$ time, where E is the number of edges. Sorting the edges dominates the cost. The Disjoint Set Union (DSU) operations run in nearly constant time with path compression and union by rank.



SPACE COMPLEXITY

Memory Usage

The space complexity is $O(V + E)$ because we store all the edges and use arrays in DSU of size V .



REFLECTION

Strengths, Limitations, Improvements

1

Strengths

Efficient for sparse graphs. Uses sorting and DSU to build MST quickly with low time complexity.

2

Limitations

Slower on dense graphs due to edge sorting. Memory grows with number of edges.

3

Improvements

Added path compression & union by rank to speed up DSU. File I/O used for flexible testing.

**THANK
YOU!**