

# Algorithm for Multiplication

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First, lets look at decimal multiplication:

$$\begin{array}{r} 256 \\ \times 387 \\ \hline 1792 \\ 2048 \\ 768 \\ \hline 99072 \end{array}$$

*Final result* → 99072

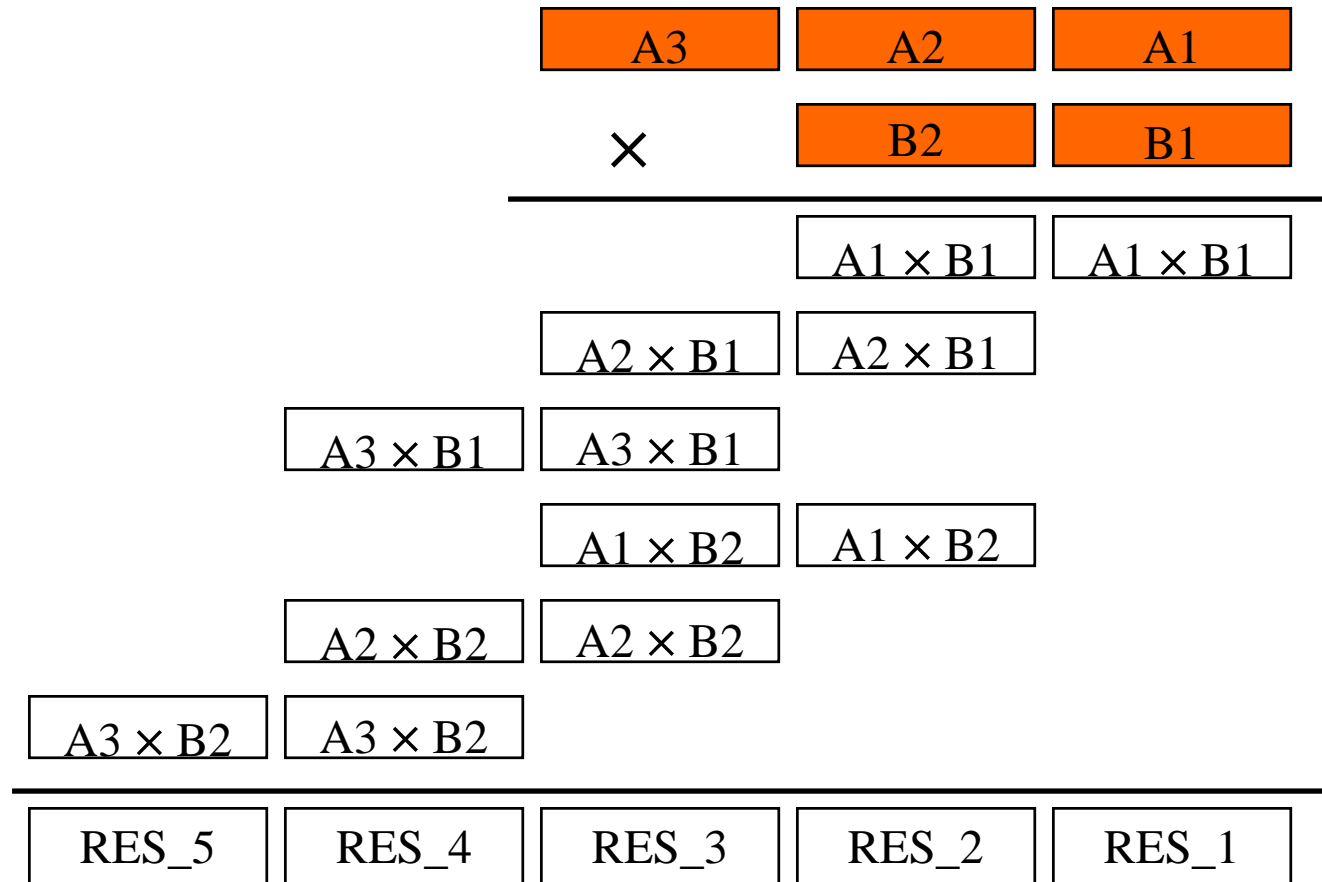
7 × 256 is first intermediate result → 1792

8 × 256 shifted and added to intermediate result → 2048

3 × 256 shifted and added to intermediate result → 768

Now, lets look at an example for binary multiplication:

Instead of each bit, it is possible to multiply bytes:



**Example:** Write a program to multiply \$5678 by \$1234 and store the result at memory locations \$1100 to \$1103. Use the method that we just illustrated in previous slide.

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LDA	#78	;load $M_L$ into A
LDB	#34	;load $N_L$ into B
MUL		;multiply $M_L$ by $N_L$
STD	\$1102	;store the partial product $M_L N_L$ at 02 & 03
LDA	#56	;load $M_H$ into A
LDB	#12	;load $N_H$ into B
MUL		;multiply $M_H$ by $N_H$
STD	\$1100	;store the partial product $M_H N_H$ at 00 & 01
LDA	#56	; load $M_H$ into A
LDB	#34	; load $N_L$ into B
MUL		;generate partial product $M_H N_L$

\* The following two instructions add  $M_H N_L$  to memory locations at \$1101 and \$1102

ADD	\$1101
STD	\$1101

## *Program continued...*

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\* The following three instructions add the C flag to memory location at \$1100

LDAA    \$1100

ADCA    #0

STAA    \$1100

LDAA    #\$78                    ;load  $M_L$  into A

LDAB    #\$12                    ;load  $N_H$  into B

MUL                            ;generate the partial product  $M_L N_H$

\* The following two instructions add  $M_L N_H$  to memory locations at \$1101 and \$1102

ADDD    \$1101

STD      \$1101

\* The following three instructions add the C flag to memory location at \$1100

LDAA    \$1100

ADCA    #0

STAA    \$1100

END

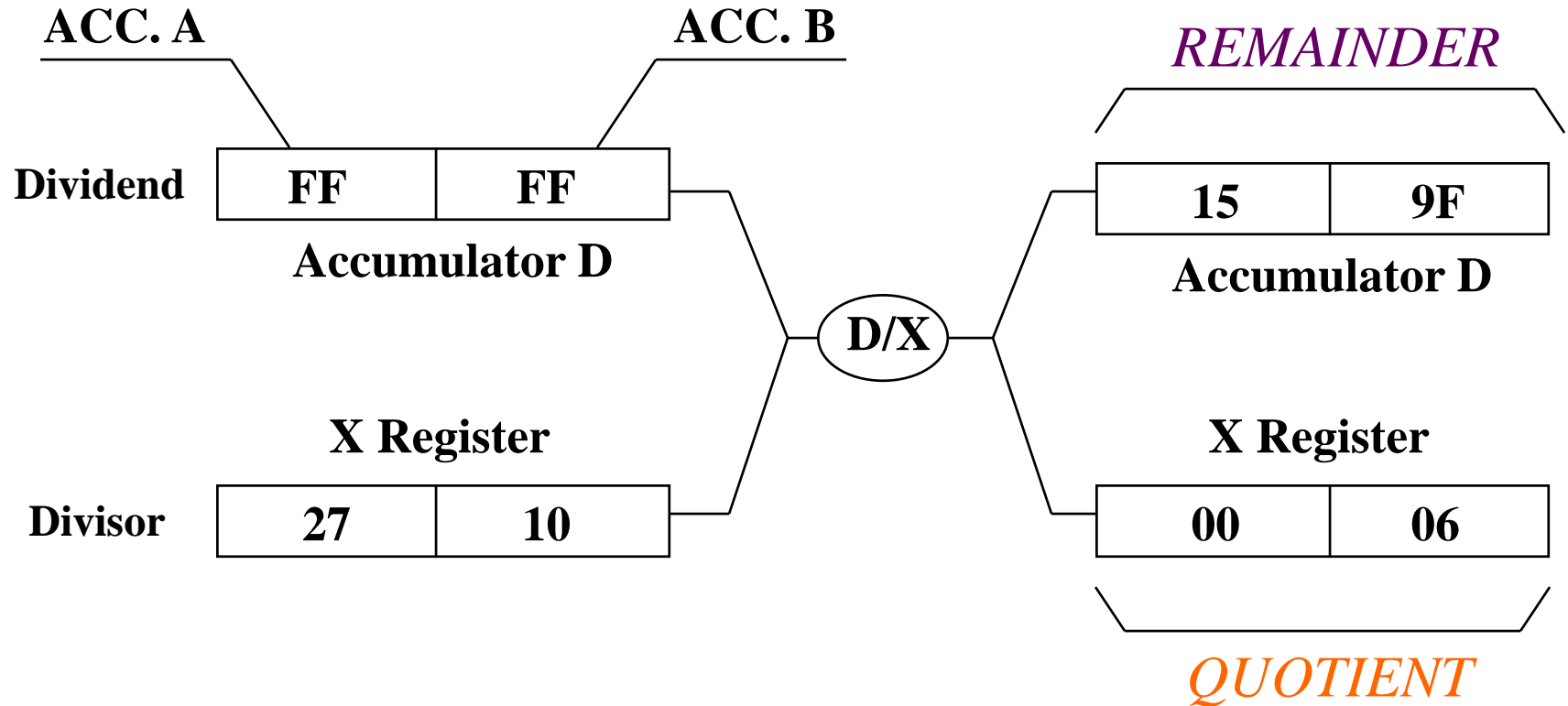
## Example For IDIV:

Program instructions:

LDD # \$FFFF

LDX # \$2710

IDIV



## Example:

Write a program to convert the 16-bit binary number stored at location called 'Bin' to decimal format and store the result in 5 consecutive memory locations starting at address 'Dec'. Each BCD digit is stored in one byte.

	ORG	\$1000	
	LDY	#Dec	; set pointer at most significant digit
	LDD	Bin	; get the 16-bit binary number
	LDX	#10	;
	IDIV		; compute the least significant digit
	STAB	4,Y	; save the least significant digit in place
	XGDX		; place the quotient in D to get the next digit
	LDX	#10	
	IDIV		; compute the next significant digit
	STAB	3,Y	; save the next significant digit in place
	XGDX		; place the quotient in D to get the next digit
	LDX	#10	
	IDIV		; compute the next significant digit
	STAB	2,Y	; save the next significant digit in place
	XGDX		; place the quotient in D to get the next digit
	LDX	#10	
	IDIV		; compute the next significant digit
	STAB	1,Y	; save the next significant digit in place
	XGDX		; swap the most significant digit to B
	STAB	0,Y	; save the most significant BCD digit
	SWI		
Bin	FDB	\$F5AC	; \$F5AC = 62892
Dec	RMB	5	
	END		

Try to implement previous code segment using loop:

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	<b>ORG</b>	<b>\$1000</b>	
	<b>LDY</b>	<b>#Dec+4</b>	<b>; set pointer at least significant digit</b>
	<b>LDX</b>	<b>Bin</b>	<b>; get the 16-bit binary number</b>
<b>Loop</b>	<b>XGDX</b>		<b>; put the binary number or quotient in D</b>
	<b>LDX</b>	<b>#10</b>	<b>;</b>
	<b>IDIV</b>		<b>; compute the next significant digit start from least</b>
	<b>STAB</b>	<b>1,Y-</b>	<b>; save next significant digit in place &amp; adjust pointer</b>
	<b>CPY</b>	<b>#Dec</b>	<b>; as long as within the range, continue Loop</b>
	<b>BHI</b>	<b>Loop</b>	
	<b>XGDX</b>		<b>; get the MSD into acc. B</b>
	<b>STAB</b>	<b>0,Y</b>	<b>; save it in place</b>
	<b>SWI</b>		
<b>Bin</b>	<b>FDB</b>	<b>\$F5AC</b>	
<b>Dec</b>	<b>RMB</b>	<b>5</b>	
	<b>END</b>		

## Branch Instructions:

Unary Branches		
Mnemonic	Function	Operation
BRA	Branch Always	1 = 1
BRN	Branch Never	1 = 0
Simple Branches		
Mnemonic	Function	Operation
BCC	Branch if Carry Clear	C = 0
BCS	Branch if Carry Set	C = 1
BEQ	Branch if Equal	Z = 1
BMI	Branch if Minus	N = 1
BNE	Branch if not Equal	Z = 0
BPL	Branch if Plus	N = 0
BVC	Branch if Overflow Clear	V = 0
BVS	Branch if Overflow Set	V = 1



## Branch Instructions *continued* ...

<b>Unsigned Branches</b>		
Mnemonic	Function	Operation
BHI	Branch if Higher	$C + Z = 0$
BHS	Branch if Higher or Same	$C = 0$
BLO	Branch if Lower	$C = 1$
BLS	Branch if Lower or Same	$C + Z = 1$
<b>Signed Branches</b>		
Mnemonic	Function	Operation
BGE	Branch if Greater than or Equal	$N \oplus V = 0$
BGT	Branch if Greater than	$Z + (N \oplus V) = 0$
BLE	Branch if Less than or Equal	$Z + (N \oplus V) = 1$
BLT	Branch if Less than	$N \oplus V = 1$

- These are all short branch instructions. One-byte relative offset.
- There is another set of long branch instructions. Mnemonic-wise the only difference is that they start with letter **L**. They all have two-byte relative offset. E.g. **LBEQ** → Long Branch if Equal

## Compare and Test Instructions:

<b>Compare Instructions</b>		
Mnemonic	Function	Operation
CBA	Compare A to B	$(A) - (B)$
CMPA	Compare A to Memory	$(A) - (M)$
CMPB	Compare B to Memory	$(B) - (M)$
CPD	Compare D to Memory	$(D) - (M:M+1)$
CPS	Compare SP to Memory	$(SP) - (M:M+1)$
CPX	Compare X to Memory	$(X) - (M:M+1)$
CPY	Compare Y to Memory	$(Y) - (M:M+1)$
<b>Test Instructions</b>		
Mnemonic	Function	Operation
TST	Test Memory for zero or minus	$(M) - \$00$
TSTA	Test A for zero or minus	$(A) - \$00$
TSTB	Test B for zero or minus	$(B) - \$00$

# Loop Primitive Instructions

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- A lot of the program loops are implemented by incrementing or decrementing a loop count.
- The branch is taken when either the loop count is equal to zero or not equal to zero, depending on the application.
- The HCS12 provides a set of loop primitive instructions for implementing this type of looping mechanism.
- These instructions test a counter value in a register or accumulator (A, B, D, X, Y, or SP) for zero or nonzero value as a branch condition.
- There are predecrement, preincrement, and test only versions of these instructions.
- The branch range is one byte relative from the instruction immediately following the loop primitive instruction.
- The syntax for these instructions is as follow:

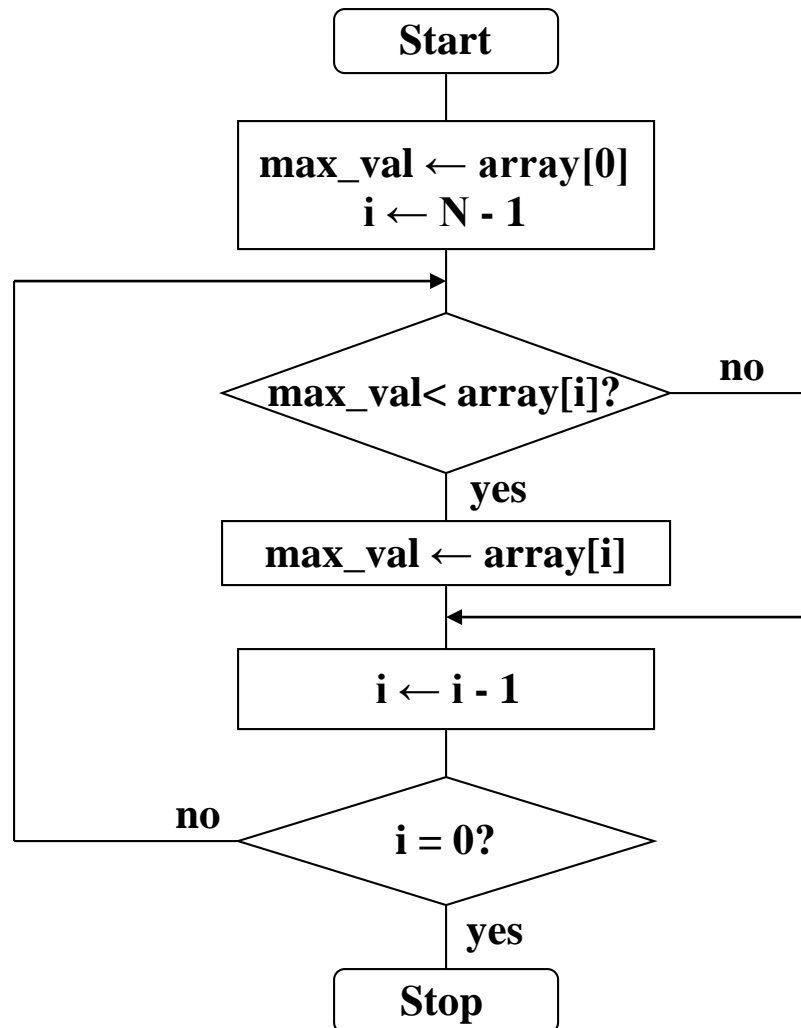
DBEQ        X,loop

## Loop Primitive Instructions

Mnemonic	Function	Operation
DBEQ	Decrement counter and branch if = 0 (counter = A, B, D, X, Y, or SP)	counter $\leftarrow$ (counter) - 1 If (counter) = 0' then branch else continue to next instruction
DBNE	Decrement counter and branch if $\neq$ 0 (counter = A, B, D, X, Y, or SP)	counter $\leftarrow$ (counter) - 1 If (counter) $\neq$ 0' then branch else continue to next instruction
IBEQ	Increment counter and branch if = 0 (counter = A, B, D, X, Y, or SP)	counter $\leftarrow$ (counter) + 1 If (counter) = 0' then branch else continue to next instruction
IBNE	Increment counter and branch if $\neq$ 0 (counter = A, B, D, X, Y, or SP)	counter $\leftarrow$ (counter) + 1 If (counter) $\neq$ 0' then branch else continue to next instruction
TBEQ	Test counter and branch if = 0 (counter = A, B, D, X, Y, or SP)	If (counter) = 0' then branch else continue to next instruction
TBNE	Test counter and branch if $\neq$ 0 (counter = A, B, D, X, Y, or SP)	If (counter) $\neq$ 0' then branch else continue to next instruction

### **Example:**

Write a program to find the maximum element from an array of N 8-bit elements using the primitive loop instruction.



	ORG	\$1000	
	LDAA	array	; set array[0] as temporary array max
	STAA	max_val	;
	LDX	#array+N-1	; start from the end of array
	LDAB	#N-1	; compute the least significant digit
loop	LDAA	max_val	;
	CMPA	0,X	; compare max_val with array[i]
	BHS	chk_end	; no update if max_val is larger
	MOVB	0,X,max_val	; update max_val
chk_end	DEX		; move the array pointer
	DBNE	B,loop	; decrement loop count, branch if not zero yet
	LDD	max	; push the print parameter into stack
	PSHD		;
	LDD	#msg	; store starting address of msg in acc. D
	LDX	printf	; call printf subroutine
	JSR	0,X	;
	LEAS	2,SP	; balance the stack
	SWI		; return to D-Bug12 monitor
max	DB	0	
max_val	RMB	1	
array	DB	1,3,5,6,19,41,53,28,13,42,76,14,20,54,64,74,29,33,41,45	
msg	FCC	'The maximum element of the vector is %u'	
	DB	\$0D,\$0A,0	
printf	EQU	\$EE88	
N	EQU	20	
	END		

## Decrementing and Incrementing Instructions:

<b>Decrement Instructions</b>		
Mnemonic	Function	Operation
DEC	Decrement memory by 1	$M \leftarrow [M] - \$01$
DECA	Decrement A by 1	$A \leftarrow [A] - \$01$
DECB	Decrement B by 1	$B \leftarrow [B] - \$01$
DES	Decrement SP by 1	$SP \leftarrow [SP] - \$01$
DEX	Decrement X by 1	$X \leftarrow [X] - \$01$
DEY	Decrement Y by 1	$Y \leftarrow [Y] - \$01$
<b>Increment Instructions</b>		
Mnemonic	Function	Operation
INC	Increment memory by 1	$M \leftarrow [M] + \$01$
INCA	Increment A by 1	$A \leftarrow [A] + \$01$
INCB	Increment B by 1	$B \leftarrow [B] + \$01$
INS	Increment SP by 1	$SP \leftarrow [SP] + \$01$
INX	Increment X by 1	$X \leftarrow [X] + \$01$
INY	Increment Y by 1	$Y \leftarrow [Y] + \$01$

# Branch Instructions Based on Bit Condition

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- In certain applications, one needs to make branch decisions based on the value of a few bits.
- The HCS12 provides two special conditional branch instructions for this purpose.
- The syntax of these two branch instructions are

<b>BRCLR</b>	<b>opr,msk,rel</b>
<b>BRSET</b>	<b>opr,msk,rel</b>

where

**opr** – specifies the memory location to be checked and can be specified using direct, extended, and all index addressing modes

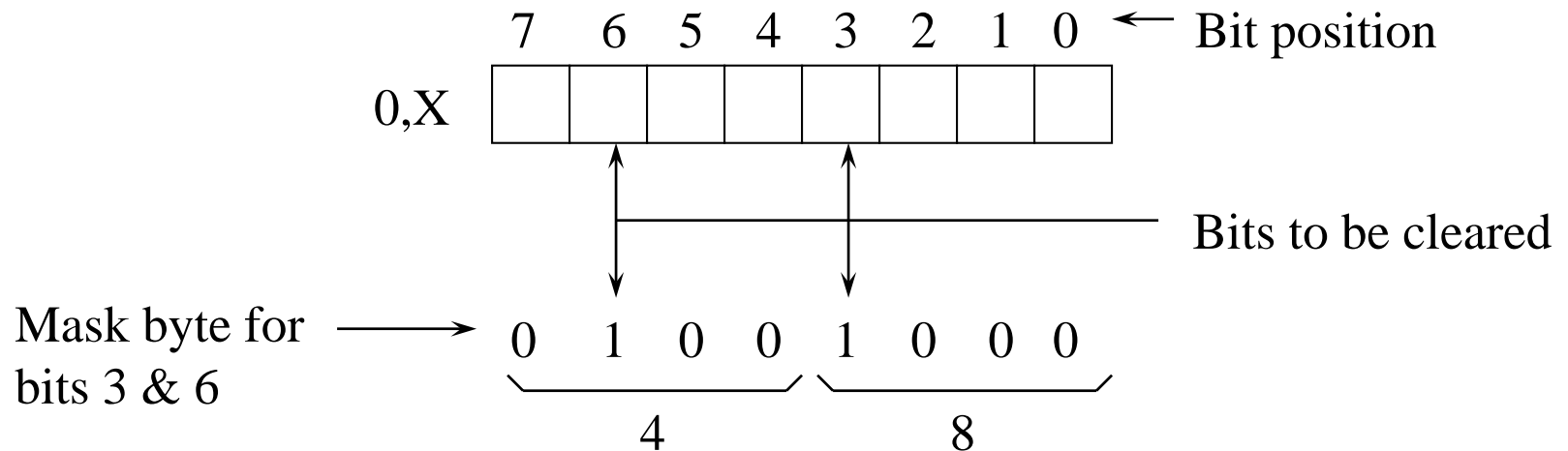
**msk** – is an 8-bit mask that specifies the bits of the memory location to be checked. The bits to be checked correspond to those bit positions that are 1s in the mask.

**rel** – is the branch offset and it is specified in 8-bit relative mode.



## Example:

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Loop    LDAA    3,Y
        ⋮
        BRCLR   0,X $48 Loop
```



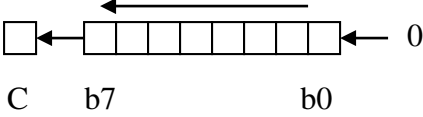
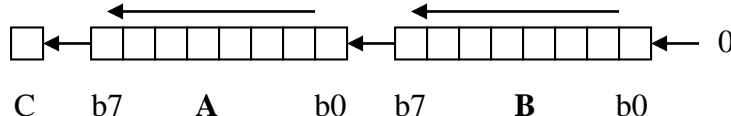
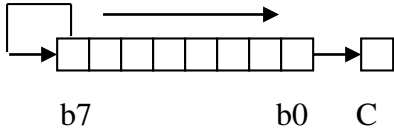
**Example:** Write a program to count the number of elements that are divisible by four in an array of N 8-bit numbers.

	ORG	\$1000	
	CLR	total	; initialize total to zero
	LDX	#array	; use X reg as array pointer
	LDAB	#N	; use acc. B as a loop counter
loop	BRCLR	0,X,\$03,yes	; if number divisible by 4
	BRA	chkend	; continue
yes	INC	total	; add one to the total
chkend	INX		; move the array pointer
	DBNE	B,loop	; decrement loop count, branch if not zero yet
	LDAB	total	; push the print parameter into stack
	CLRA		;
	PSHD		;
	LDD	#msg	; store starting address of msg in acc. D
	LDX	printf	; call printf subroutine
	JSR	0,X	;
	LEAS	2,SP	; balance the stack
	SWI		; return to D-Bug12 monitor
total	RMB	1	
array	DB	2,3,4,8,12,13,19,24,33,32,20,18,53,52,80,82,90,94,100,102	
msg	FCC	'The number of elements that are divisible by 4 are %u'	
	DB	\$0D,\$0A,0	
printf	EQU	\$EE88	
N	EQU	20	
	END		

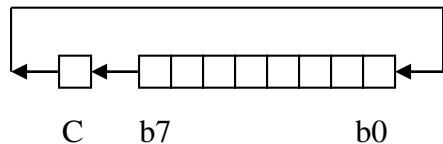
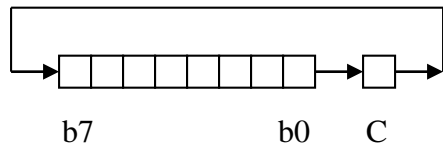
## Shift Instructions:

<b>Logical Shift Instructions</b>		
Mnemonic	Function	Operation
LSL <opr> LSLA LSLB	Logical shift left memory Logical shift left A Logical shift left B	
LSLD	Logical shift left D	
LSR <opr> LSRA LSRB	Logical shift right memory Logical shift right A Logical shift right B	
LSRD	Logical shift right D	

# Arithmetic Shift Instructions

Mnemonic	Function	Operation
ASL <opr> ASLA ASLB	Arithmetic shift left memory Arithmetic shift left A Arithmetic shift left B	
ASLD	Arithmetic shift left D	
ASR <opr> ASRA ASRB	Arithmetic shift right memory Logical shift right A Logical shift right B	

# Rotate Instructions

ROL <opr> ROLA ROLB	Rotate left memory thru carry Rotate left A through carry Rotate left B through carry	
ROR <opr> RORA RORB	Rotate right memory thru carry Rotate right A through carry Rotate right B through carry	

**Example:** Write a program to count the number of 0s contained in memory locations \$1000 ~ \$1001 and save the result at memory location \$1005.

	<b>ORG</b>	<b>\$1000</b>	
	<b>DW</b>	<b>\$2355</b>	
<b>zero_cnt</b>	<b>RMB</b>	<b>1</b>	
<b>lp_cnt</b>	<b>RMB</b>	<b>1</b>	
	<b>ORG</b>	<b>\$1010</b>	
	<b>CLR</b>	<b>zero_cnt</b>	<b>; initialize the zero count to 0</b>
	<b>MOVB</b>	<b>#16,lp_cnt</b>	<b>; initialize lp_cnt to 16</b>
	<b>LDD</b>	<b>\$1000</b>	<b>; place the 16-bit number in acc. D</b>
<b>loop</b>	<b>LSRD</b>		<b>;</b>
	<b>BCS</b>	<b>chkend</b>	<b>; branch if the least sig. bit is 1</b>
<b>yes</b>	<b>INC</b>	<b>zero_cnt</b>	<b>; add one to the total</b>
<b>chkend</b>	<b>DEC</b>	<b>lp_cnt</b>	<b>;</b>
	<b>BNE</b>	<b>loop</b>	<b>; have we checked all 16 bit yet?</b>
	<b>MOVB</b>	<b>zero_cnt,\$1005</b>	<b>; store the result in location \$1005</b>
	<b>SWI</b>		<b>; return to D-Bug12 monitor</b>
	<b>END</b>		

**Example:** Write a subroutine called mult10 such that every time the subroutine is called, it will multiply the content of accumulator D by 10 using shift command.

	ORG	\$1000	
	⋮		
	LDD	#20	
	JSR	mult10	
	⋮		
Mult10	LSLD		; multiply D by 2
	STD	temp	; save DX2 in temp
	LSLD		; multiply D by 2 again
	LSLD		; create D X 8
	ADDD	temp	; sum DX2 and DX8 to get DX10
	RTS		;
Temp	RMB	2	
	END		

## Boolean Logic Instructions

Mnemonic	Function	Operation
ANDA <opr>	AND A with memory	$A \leftarrow (A) \cdot (M)$
ANDB <opr>	AND B with memory	$B \leftarrow (B) \cdot (M)$
ANDCC<opr>	AND CCR with memory	$CCR \leftarrow (CCR) \cdot (M)$
EORA <opr>	Exclusive OR A with memory	$A \leftarrow (A) \oplus (M)$
EORB <opr>	Exclusive OR B with memory	$B \leftarrow (B) \oplus (M)$
ORA <opr>	OR A with memory	$A \leftarrow (A) + (M)$
ORB <opr>	OR B with memory	$B \leftarrow (B) + (M)$
ORCC <opr>	OR CCR with memory	$CCR \leftarrow (CCR) + (M)$
CLC	Clear C bit in CCR	$C \leftarrow 0$
CLI	Clear I bit in CCR	$I \leftarrow 0$
CLV	Clear V bit in CCR	$V \leftarrow 0$
COM <opr>	One's complement memory	$M \leftarrow \$FF - (M)$
COMA	One's complement A	$A \leftarrow \$FF - (A)$
COMB	One's complement B	$B \leftarrow \$FF - (B)$
NEG <opr>	Two's complement memory	$M \leftarrow \$00 - (M)$
NEGA	Two's complement A	$A \leftarrow \$00 - (A)$
NEGB	Two's complement B	$B \leftarrow \$00 - (B)$

## Bit Test and Bit Manipulate Instructions

Mnemonic	Function	Operation
BCLR <opr> <sup>1</sup> ,msk8	Clear bits in memory	$M \leftarrow (M) \cdot (\overline{mm})$
BITA <opr> <sup>2</sup>	Bit test A	$(A) \cdot (M)$
BITB <opr> <sup>2</sup>	Bit test B	$(B) \cdot (M)$
BSET <opr> <sup>1</sup> ,msk8	Set bits in memory	$M \leftarrow (M) + (mm)$

Note: 1. <opr> can be specified using direct, extended, and indexed (exclude indirect) addressing modes.

2. <opr> can be specified using all except relative addressing modes for BITA and BITB.

3. msk8 is an 8-bit value.

**Example:** Write a set of instructions to set bits 2, 5, and 7 in the memory location \$2500 and clear bits 1, 3, 6, and 7 in memory location \$2510.

BSET    \$2500,\$A4

BCLR    \$2510,\$CA



### **Example:**

Write a program to determine the largest of the twenty 8-bit numbers stored in consecutive memory locations starting at address 'Nums'. Save the largest number at memory location called 'Max'. All the numbers are positive.

	<b>NAM</b>	<b>EXAMPLE</b>
	<b>ORG</b>	<b>\$1000</b>
<b>Nums</b>	<b>FCB</b>	<b>250,28,38,168,251,222,38,55,183,232</b>
	<b>FCB</b>	<b>12,88,198,209,246,77,253,9,133,200</b>
<b>Max</b>	<b>RMB</b>	<b>1</b>
	<b>ORG</b>	<b>\$1050</b>
	<b>SWI</b>	
	<b>END</b>	

## **Example:**

Write a program to determine the largest of the twenty 8-bit numbers stored in consecutive memory locations starting at address 'Nums'. Save the largest number at memory location called 'Max'. All the numbers are positive.

	<b>NAM</b>	<b>EXAMPLE</b>	
	<b>ORG</b>	<b>\$C000</b>	
<b>Nums</b>	<b>FCB</b>	<b>250,28,38,168,251,222,38,55,183,232</b>	
	<b>FCB</b>	<b>12,88,198,209,246,77,253,9,133,200</b>	
<b>Max</b>	<b>RMB</b>	<b>1</b>	
	<b>ORG</b>	<b>\$C050</b>	
	<b>LDAB</b>	<b>#19</b>	<b>; set the counter</b>
	<b>LDX</b>	<b>#Nums</b>	<b>; set the pointer to first number</b>
	<b>LDAA</b>	<b>0,X</b>	<b>; put the first number in Acc. A</b>
<b>LOOP</b>	<b>CMPA</b>	<b>1,X</b>	<b>; check the next number with A, if A is bigger or</b>
	<b>BHS</b>	<b>Next</b>	<b>; equal go to next number, otherwise replace the</b>
	<b>LDAA</b>	<b>1,X</b>	<b>; number with the content of A.</b>
<b>Next</b>	<b>INX</b>		
	<b>DECB</b>		
	<b>BNE</b>	<b>LOOP</b>	<b>; repeat the process till all numbers are checked</b>
	<b>STAA</b>	<b>Max</b>	<b>; put the largest number in memory location Max</b>
	<b>SWI</b>		
	<b>END</b>		

# Finding the Square Root

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Consider the following series :

$$\sum_{i=0}^{n-1} i = \frac{n(n-1)}{2}$$

e.g. if  $n = 5$ ,

$$1 + 2 + 3 + 4 = \frac{5(4)}{2} = 10$$

for  $n = 8$ ,

$$1 + 2 + 3 + 4 + 5 + 6 + 7 = \frac{8(7)}{2} = 28$$

### Finding the Square Root *continued* ...

Lets rewrite the series as follow :

$$\sum_{i=0}^{n-1} 2i = n^2 - n \quad \Rightarrow \quad n^2 = \sum_{i=0}^{n-1} (2i + 1)$$

Suppose, we want to compute the square root of  $q$ , and  $n$  is the integer value that is closest to true square root. One of the three following relationships is satisfied :

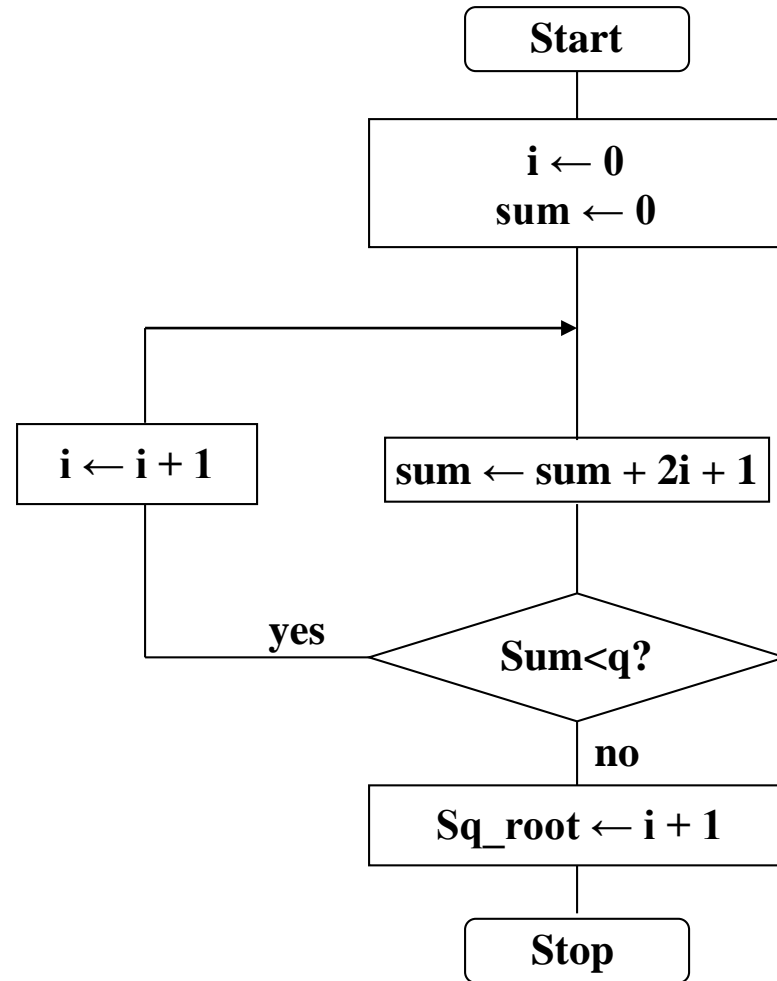
$$n^2 < q$$

$$n^2 = q$$

$$n^2 > q$$

## Finding the Square Root *continued* ...

Algorithm for  
finding the  
square root of  
integer  $q$



### Finding the Square Root *continued* ...

**Accuracy:** lets assume  $q = 820$ , this procedure creates  $n = 28$ . which gives  $n^2=784$  and is not very accurate.

**Remedy:** If  $q$  is less than two-byte, then multiply it with 10000 and find the square root of the new number. Then divide the answer by 100 and quotient is integer part of the square root and the remainder is the first two decimal digits.

$$q = 60,000 \Rightarrow q_1 = 60,000 \times 10,000 = 600,000,000$$

$$n_1 = 24496 \Rightarrow n = \frac{n_1}{100} = 244 \quad \text{and} \quad \frac{96}{100} = 0.96$$

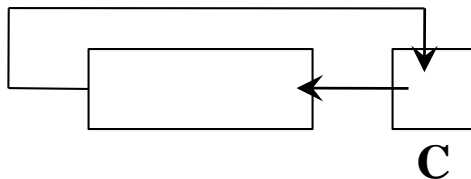
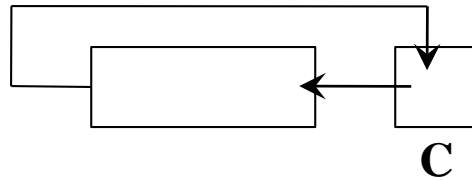
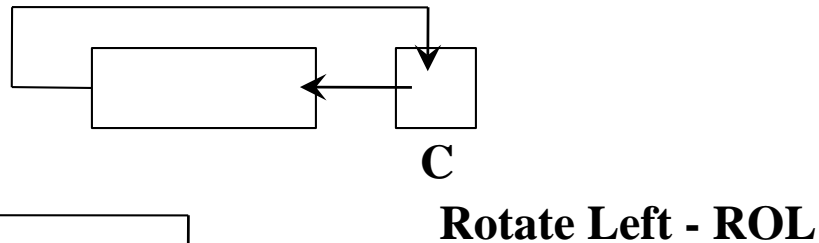
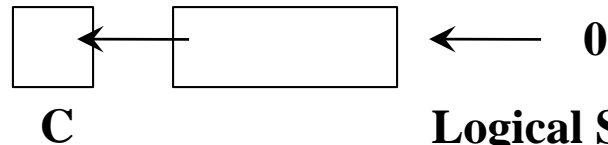
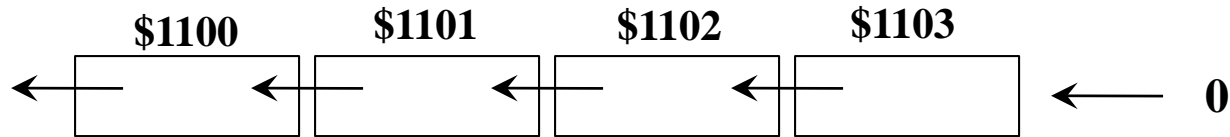
$$\text{Therefore,} \quad n = 244 \Rightarrow n^2 = 59536$$

$$\text{but if} \quad n = 244.96 \Rightarrow n^2 = 60005.4$$

## Project 2 – Finding the Square Root of a number

- Ask user to enter a number up to 4,294,967,295 ( $2^{32} - 1$ ).
- Make sure number is correct and within the range.
- Pack the number into binary (up to 4-bytes).
- If number is 2-byte or less, multiply it by 10000.
- If number is 3-bytes, multiply it by 100.
- Find square root of the number and adjust for decimal point.
- Output the number and its square root to terminal in appropriate way.
- Ask if user wants to try again.

# Shift 4-byte Number to the Left by 1-bit



**Rotate Left - ROL**

**L1**

```
LDX    #$1103
CLC
LDAB   #4
ROL    1,X-
DBNE   B,L1
```



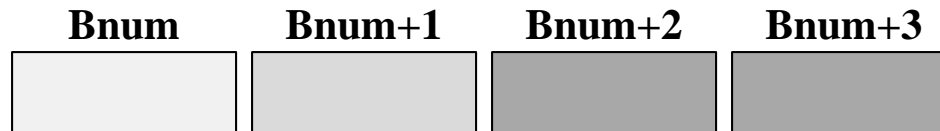
# Add two N-byte Numbers in Memory

---



	<b>LDX</b>	<b>#Num+N-1</b>	<b>; set pointer to LS Byte</b>
	<b>LDY</b>	<b>#Temp+N-1</b>	<b>; set pointer to LS Byte</b>
	<b>CLC</b>		
	<b>LDAB</b>	<b>#N</b>	<b>; set counter</b>
<b>L1</b>	<b>LDAA</b>	<b>0,X</b>	<b>; add same order bytes</b>
	<b>ADCA</b>	<b>1,Y-</b>	<b>; together w/carry and</b>
	<b>STAA</b>	<b>1,X-</b>	<b>; store it back</b>
	<b>DBNE</b>	<b>B,L1</b>	<b>; do it N-times</b>

# How to adjust 2- & 3-byte number to 4-byte



	LDD	Bnum	; get MS 2-bytes
	BEQ	two	; if 0, then mult with 10000
	CMPA	#0	; get MS-byte
	BEQ	one	; if 0, then mult by 100
	BRA	find	; otherwise, calculate square root
two	LDD	Bnum+2	; get two-byte number
	LDY	#10000	; multiply it by 10000
	EMUL		;
	STD	Bnum+2	; store the result in Bnum
	STY	Bnum	;
	BRA	find	; calculate square root
one	LDD	Bnum+2	; get 1 <sup>st</sup> 2-bytes
	LDY	#100	; multiply by 100
	EMUL		;
	STD	Bnum+2	; store 1st part of the multiplication
	LDD	Bnum	; pick up last two byte of the number
	STY	Bnum	; store 2 <sup>nd</sup> part of mult.
	LDY	#100	; multiply by 100 2 <sup>nd</sup> part
	EMUL		;
	ADDD	Bnum	; add to result last 2-byte of 1 <sup>st</sup> product
	STD	Bnum	; store the last 2-byte of the result
find			; calculate square root

# How to Remember and Adjust for Decimal

---

Use a flag to differentiate between 3 cases:

1. Flag = 0 → no decimal
2. Flag = 1 → one decimal digit
3. Flag = 2 → two decimal digit

How would you take care of the following result?

378.09

Be careful on the use of %u conversion on '*printf*' function!

# Quiz 1

## Converting degree Fahrenheit to Celsius

---

- ☐ Ask user to enter a temperature in Fahrenheit from 32 up to 255.
- ☐ Read in the temperature.
- ☐ Pack the numbers into binary (1 byte).
- ☐ Convert it to degree Celsius.
- ☐ Output the temperature to the terminal as follow:

XXX °F      YYY °C

- ☐ Ask if user wants to try again.

$$^{\circ}C = \frac{^{\circ}F - 32}{1.8} \Rightarrow ^{\circ}C = \frac{(^{\circ}F - 32) \times 10}{18}$$

$$^{\circ}C = \frac{^{\circ}F \times 5 - 160}{9}$$

LDAB	F	; get Fahrenheit value
LDAA	#5	; multiply By 5
MUL		;
SUBD	#160	; subtract 160
LDX	#9	; divide by 9
IDIV		;
PSHX		; push Celsius value into stack
CLRA		; push Fahrenheit value into stack
LDAB	F	;
PSHD		;
LDD	#outmsg	; output the result
LDX	printf	;
JSR	0,X	;
LEAS	4,SP	; balance the stack
LDD	#repmsg	; ask for repeat
LDX	printf	;
JSR	0,X	;
LDX	getchar	; check answer
JSR	0,X	;
CMPB	#'y'	; act accordingly
BEQ	start	;
SWI		

# How to Take Care of Negative Numbers

---

	<b>LDAB</b>	<b>F</b>	<b>; get Fahrenheit value</b>
	<b>LDAA</b>	<b>#5</b>	<b>; multiply By 5</b>
	<b>MUL</b>		<b>;</b>
	<b>SUBD</b>	<b>#160</b>	<b>; subtract 160</b>
	<b>BCC</b>	<b>cont</b>	<b>; if number isn't negative continue</b>
	<b>COMA</b>		<b>; complement the result</b>
	<b>COMB</b>		<b>;</b>
	<b>ADDD</b>	<b>#1</b>	<b>;</b>
	<b>MOVB</b>	<b>#\$2D,sign</b>	<b>; incorporate the negative sign</b>
<b>cont</b>	<b>LDX</b>	<b>#9</b>	<b>; divide by 9</b>
	<b>IDIV</b>		<b>;</b>
	<b>PSHX</b>		<b>; push Celsius value into stack</b>
	<b>CLRA</b>		<b>; push Fahrenheit value into stack</b>
	<b>LDAB</b>	<b>F</b>	<b>;</b>
	<b>PSHD</b>		<b>;</b>
	<b>LDD</b>	<b>#outmsg</b>	<b>; output the result</b>
	<b>LDX</b>	<b>printf</b>	<b>;</b>
	<b>JSR</b>	<b>0,X</b>	<b>;</b>
	<b>LEAS</b>	<b>4,SP</b>	<b>; balance the stack</b>

## How to Take Care of Negative Numbers *cont'd...*

	<b>LDD</b>	<b>#repmsg</b>	<b>; ask for repeat</b>
	<b>LDX</b>	<b>printf</b>	<b>;</b>
	<b>JSR</b>	<b>0,X</b>	<b>;</b>
	<b>LDX</b>	<b>getchar</b>	<b>; check answer</b>
	<b>JSR</b>	<b>0,X</b>	<b>;</b>
	<b>CMPB</b>	<b>#'y'</b>	<b>; act accordingly</b>
	<b>BEQ</b>	<b>start</b>	<b>;</b>
	<b>SWI</b>		
<b>outmsg</b>	<b>DB</b>	<b>CR,LF</b>	
	<b>FCC</b>	<b>' %u F is equal to '</b>	
<b>sign</b>	<b>DB</b>	<b>'+'</b>	
	<b>FCC</b>	<b>'%u C'</b>	
	<b>DB</b>	<b>CR,LF,0</b>	
<b>repmsg</b>	<b>FCC</b>	<b>'Would you like to enter another value (y/n)?'</b>	
	<b>DB</b>	<b>CR,LF,0</b>	
<b>CR</b>	<b>EQU</b>	<b>\$0D</b>	
<b>LF</b>	<b>EQU</b>	<b>\$0A</b>	
<b>Printf</b>	<b>EQU</b>	<b>\$EE88</b>	

# How to Add Decimal Point

$$^{\circ}C = \frac{^{\circ}F - 32}{1.8} \Rightarrow ^{\circ}C = \frac{(^{\circ}F - 32) \times 50}{9} = \frac{^{\circ}F \times 50 - 1600}{9}$$

$$\frac{^{\circ}C}{10} = ^{\circ}C(\text{integer}) \text{ and remainder is decimal number}$$

	<b>LDAB</b>	<b>F</b>	<b>; get Fahrenheit value</b>
	<b>LDAA</b>	<b>#50</b>	<b>; multiply By 5</b>
	<b>MUL</b>		<b>;</b>
	<b>SUBD</b>	<b>#1600</b>	<b>; subtract 160</b>
	<b>BCC</b>	<b>cont</b>	<b>; if number isn't negative continue</b>
	<b>COMA</b>		<b>; complement the result</b>
	<b>COMB</b>		<b>;</b>
	<b>ADDD</b>	<b>#1</b>	<b>;</b>
<b>cont</b>	<b>MOVB</b>	<b>#'-',sign</b>	<b>; incorporate the negative sign</b>
	<b>LDX</b>	<b>#9</b>	<b>; divide by 9</b>
	<b>IDIV</b>		<b>;</b>
	<b>XGDX</b>		<b>; separate integer and decimal</b>
	<b>LDX</b>	<b>#10</b>	<b>;</b>
	<b>IDIV</b>		<b>;</b>
	<b>PSHD</b>		<b>; push decimal point of Celsius</b>
	<b>PSHX</b>		<b>; push Celsius value into stack</b>



## How to Add Decimal Point *cont'd...*

```
CLRA                ; push Fahrenheit value into stack
LDAB    F           ;
PSHD              ;
LDD    #outmsg      ; output the result
LDX    printf       ;
JSR    0,X          ;
LEAS   6,SP         ; balance the stack
LDD    #repmsg      ; ask for repeat
LDX    printf       ;
JSR    0,X          ;
LDX    getchar      ; check answer
JSR    0,X          ;
CMPB   #'y'         ; act accordingly
BEQ    start        ;
SWI
```

```
outmsg DB    CR,LF
        FCC   '    %u.0 F is equal to '
sign   DB    '+'
        FCC   '%u.%u C'
        DB    CR,LF,0
```

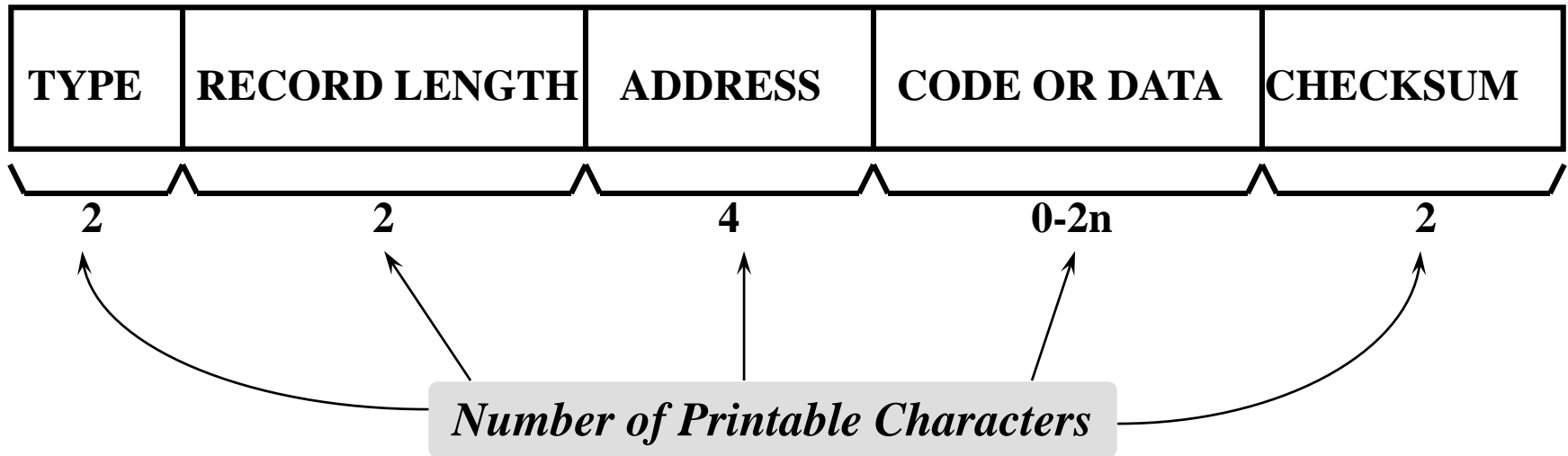
# S-Record

---

- Motorola uses character strings called S-records to encode programs and data for serial transmission between computers.
- Each string begins with upper case S and is divided into five fields.
- There are ten possible types of s-records (S0-S9).
- Only three s-record types are used with the HCS12 microcontroller (S0, S1, and S9).
- An s-record encodes the object code into a printable ASCII format.
- The last byte of each s-record is a *checksum*.
- This value is used for error checking to ensure that the instruction has been transmitted error free.

# Five Fields of S-record

---



**S0-record:** An S0-record is always the first character string and is referred to as the *header record*. The address field may be all zeros in S0 record.

### **Example:**

S00A00006368342E4F5554D0

S0	Indicates the record is a header record.
0A	The number of bytes (in Hex) to follow in this string.
0000	Address field is all zeros.
63	} ASCII representation for ch4.OUT
68	
34	
2E	
4F	
55	
54	
D0	checksum.

---

**Checksum:** The checksum is determined first by adding all the bytes in the record length, address, and code/data fields, then obtaining the one's complement of the least significant byte of the sum.

$$0A + 00 + 00 + 63 + 68 + 34 + 2E + 4F + 55 + 54 = 22F$$

Least Significant Byte of Sum = \$2F

1's Complement of \$2F = \$D0 which is the *Checksum* value.

**S1-record:** S1-records contain the instructions and data to be operated on by the microcontroller.

---

**Example:**

**S119C400BDFFCD8131270E8132270A813327068134270226EB3FBF**

S1	Indicates the record contains code/data.
19	There are 25 bytes of binary data to follow.
C400	Starting address where the data is stored.
BD FF CD	Machine code for first instruction, JSR INCHAR.
81 31	Machine code for CMPA #\$31.
. . . .	Machine code for other instructions in the program.
3F	Machine code for SWI.
BF	Checksum.

---

$$19 + C4 + 00 + \dots + 26 + EB + 3F = 840 \quad 1' \text{ Comp. of } \$40 = \$BF$$

**S9-record:** The S9-record is the termination string. The address field may also be all zeroes and there is no code/data field.

---

*Example:*

S9030000FC

S9	Indicates a termination record.
03	There are three bytes to follow.
0000	Address field is all zeroes.
FC	Checksum.

---

*Calculating Checksum:*

$$03 + 00 + 00 = 03$$

$$1' \text{Complement of } \$03 = \$FC$$

# Sample of S-records

```
S015000046696C653A2074656D706C63642E61736D0AAE
S11310001610BD16102216104B868016108FCE11A6
S11310104A16111586C016108FCE115916111506D1
S11310201003CC1121FEEE881500FEEE8415007B22
S11310301159C030860A127B1120FEEE8415007B04
S1131040115AC030FB11207B11203DF61120865A25
S113105012CE00051810B7C5C30140CE000A1810FF
S1131060CB307B1165B7C5CE000A1810CB307B118D
S113107063B7C5CE000A1810CB307B1162B7C5C167
S1131080002706CB307B11613D180B2011613D36E2
S11310904D32014C320284F044448A025A32A7A7EA
S11310A0A74D320232840F48484C32028A025A3227
S11310B0A7A7A74D3202CD00F00326FD3D180BFF74
S11310C000331610E0862816108F860C16108F86B3
S11310D00616108F860116108FCD27100326FD3DAE
S11310E0CDEA600326FD3D364C32014C320284F0D9
S11310F044448A035A32A7A7A74D320232840F48C8
S1131100484C32028A035A32A7A7A74D3202CD00B7
S1131110F00326FD3DA63027061610E70611153DFF
S11311210D20456E746572206465677265652043A0
S1131131656C63697573206265747765656E2030CB
S113114130202D2035303A0D005468652054656DEA
S10B1151702E206973203A009E
S109115B20DF43202D20DB
S10411642E58
S107116620DF46003C
S9030000FC
```

# How to Divide Two 32-bit Numbers

$$EDIV \rightarrow \frac{Y:D}{X} \Rightarrow \begin{cases} \text{Quotient} & \rightarrow X \\ \text{Remainder} & \rightarrow D \end{cases}$$

$$2^{32} - 1 = 4,294,967,295$$

	3,589,204,795	<div>0</div>	counter
Carry Clear	- 1,000,000,000		
	<hr/> 2,589,204,795	<div>1</div>	
Carry Clear	- 1,000,000,000		
	<hr/> 1,589,204,795	<div>2</div>	
Carry Clear	- 1,000,000,000		
	<hr/> 589,204,795	<div>3</div>	Most significant digit
Carry Set	- 1,000,000,000		

dnum	
dnum+9	



## How to Divide Two 32-bit Numbers *cont'd...*

589,204,795	<b>0</b>	
- 100,000,000		1,000,000,000
<hr/>		
489,204,795	<b>1</b>	100,000,000
- 100,000,000		10,000,000
<hr/>		
389,204,795	<b>2</b>	1,000,000
- 100,000,000		100,000
<hr/>		
289,204,795	<b>3</b>	10,000
- 100,000,000		1,000
<hr/>		
189,204,795	<b>4</b>	100
- 100,000,000		10
<hr/>		
89,204,795	<b>5</b>	
- 100,000,000		
<hr/>		
CARRY SET		

2<sup>nd</sup> most significant digit

Continue the process till you get to the unit digit – this is called  
*Successive Subtraction Method*

## Project 3 – Multiplying two 10-digit numbers

- ❑ Ask user to enter two numbers up to 4,294,967,295.
- ❑ Make sure numbers are correct and within the range.
- ❑ Pack the numbers into binary (up to 4-bytes).
- ❑ Multiply the numbers to get a binary number up to 8-bytes.
- ❑ Generate 8-byte long power of tens and store them into memory.
- ❑ Use Successive Subtraction method to convert 8-byte binary result into 20-digit decimal number.
- ❑ Output the result as follow:

$$85000 \times 125000 = 10625000000$$

- ❑ Ask if user wants to try again.

*Bonus:* Try to output the result as follow:

$$85,000 \times 125,000 = 10,625,000,000$$

no extra spaces are displayed.