



## THE UNIVERSITY OF THE WEST INDIES ST. AUGUSTINE

## **EXAMINATIONS OF** December 2007

Code and Name of Course:

COMP2000 – Data Structures

Date and Time: Monday 10th December 2007

A·M.

Duration: 2 Hours

INSTRUCTIONS TO CANDIDATES: This paper has 3 pages and 3 questions

## Answer all questions

- The integer elements of an almost complete binary tree are stored in an array A[1..n], 1. with the root in location 1. Write a function to rearrange the elements of A so that the smallest integer is in location 1, and the smallest integer of each subtree is also in the root of that subtree. [6]
  - (b) Draw a non-degenerate binary tree of 5 nodes such that the pre-order and level-order traversals produce identical results. [3]
  - The following are the inorder and postorder traversals of the nodes of a binary tree:

Inorder:

GDPKENFATL

Postorder:

GPDKFNTALE

Draw the tree.

[5]

- (d) Write a function which, given a pointer to the root of a binary tree, returns the sum of the levels of the nodes in the tree. Assume the root is at level 0. [3]
- (e) In a binary tree of n nodes, an 'external' node is attached to each null pointer. If I is the sum of the levels of the (internal) nodes of the tree and E is the sum of the levels of the external nodes, prove that E - I = 2n. [3]
- Write a structured, non-recursive algorithm to traverse a binary tree in post-order. The only pointer fields in each node are left and right. You may assume the existence of [5] the usual functions for manipulating a stack.

- 2. (a) Integers are inserted in an integer hash table list[1] to list[n] using "open addressing with double hashing". Assume that the function h1 produces the initial hash location and the function h2 produces the increment. An available location has the value Empty and a deleted location has the value Deleted.
  - Write a function to search for a given value key. If found, the function returns the location containing key. If not found, the function inserts key in the *first* deleted location encountered (if any) in searching for key, or an Empty location, and returns the location in which key was inserted. You may assume that list contains room for a new integer.
  - (b) In another hashing application, keys which hash to the same location are held on a linked list sorted in ascending order with the hashtable location containing a pointer to the first item on the list. Each item in the linked list consists of an integer key and a pointer to the next element in the list. Use the names key and next for these fields. Storage for a linked list item is allocated as needed. Assume that the hash table is of size 53.

Write programming code, including all relevant declarations, to

- (i) initialize the hash table; [2]
- (ii) search for the number m. If found, return the pointer to the node containing m. If not found, insert m in its appropriate position. [6]
- (c) An n x n matrix A is used to store the points obtained in cricket matches among n teams. A team gets 3 points for a win, 1 point for a tie and 0 points for a loss. A[i, i] is set to 3 if team i beats team j; it is set to 1 if the match is tied and it is set to 0 if team i loses to team j. In order to conserve storage, the values in the (strictly) lower triangle of A are stored in an array B[1..m] in row order.
  - (i) What is the value of m? [1]
  - (ii) Write a function score(i, i) which, by accessing B, returns the value of A[i, j].

    If i or j is invalid, the function returns -1.

    [5]
- (d) A function is given an integer array A and two subscripts m and n. The function must rearrange the elements A[m] to A[n] and return a subscript d such that all elements to the left of d are less than or equal to A[d] and all elements to the right of d are greater than A[d].

  [5]

3. (a) A directed graph **G(V, E)** is represented using *adjacency lists*. The graph has *n* nodes stored in an array **G[1..n]**. Each node, **G[j]**, has the following fields:

char id; //the name of the node; a single character e.g. A, F, T int colour;

int parent; //an array subscript indicating the location of a parent node GEdgePtr first; //pointer to the first edge from the node; null, if none

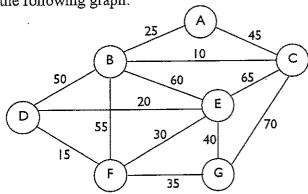
Each edge node has the following fields:

int child; //a subscript in G; G[child] is the child node int weight;

GEdgePtr next; //pointer to the next edge; null, if none

- (i) Write a function which, given **G** and **n**, outputs the graph, listing the nodes in the order in which they appear in **G**. Each node is followed by the name and weight of the edges leaving it. [4]
- (ii) Write a function which, given G, n and a node name S, performs a depth-first traversal of G starting at S. Output the name of a node as it is visited and set the parent fields so that a depth-first path can be determined. [8]
- (iii) Assuming that a depth-first traversal has already been done as in (ii), write a function which, given **G**, **n** and a node name **D**, prints the names of the nodes, in the order visited, on the depth-first path from the source node to **D**. [5]
- (b) The *indegree* of a node is the number of edges leading into it. Assume that the field **inDegree** is added to a graph node defined in part (a), above. Write a function which, given **G** and **n**, correctly sets the **inDegree** field of all the nodes. [3]

(c) Given the following graph:



- (i) Give the depth-first and breadth-first traversals starting at A. Edges of a node are processed in alphabetical order. [2]
- (ii) Draw the minimum spanning tree obtained using Kruskal's algorithm. [3]