

```
#include <stdlib.h>
#include <string.h>
#include <ctype.h>
```

```
#define MAXPAROLA 30
#define MAXRIGA 80
```

```
int main(int argc, char *argv[])
{
    int freq[MAXPAROLA]; /* vettore di contatori
delle frequenze delle lunghezze delle parole */
    char riga[MAXRIGA];
    int i, inizio, lunghezza;
    FILE *f;
```

```
for(i=0; i<MAXPAROLA; i++)
    freq[i]=0;
```

```
if(argc != 2)
```

```
{
    fprintf(stderr, "ERRORE, serve un parametro con il nome del file\n");
    exit(1);
}
```

```
f = fopen(argv[1], "r");
if(f==NULL)
```

```
{
    fprintf(stderr, "ERRORE, impossibile aprire il file %s\n", argv[1]);
    exit(1);
}
```

```
while( fgets( riga, MAXRIGA, f ) != NULL )
```



Graphs

Single Source Shortest Paths for DAGs

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Shortest paths on weighted DAGs

- ❖ For a DAG the SSSPs problem can be solved with a simplified algorithm
- ❖ Shortest paths are always well defined even if there are negative-weight edges
 - This is because, obviously, negative-weight cycles cannot exist in a DAG

Shortest paths on weighted DAGs

❖ As there are no cycles it is enough to

➤ Topologically sort the DAG

- Impose a linear order on the vertices

Perform a DFS computing
end-processing times
Order vertices using the end-
processing times

➤ Relax all vertices following the sorted order given
by the topological sort

- In other words, it suffices to make just one pass
over the vertices in the topological sorted order
- As we process a vertex, we relax each edge that
leaves the vertex

SSSP for DAGs

Pseudo-code

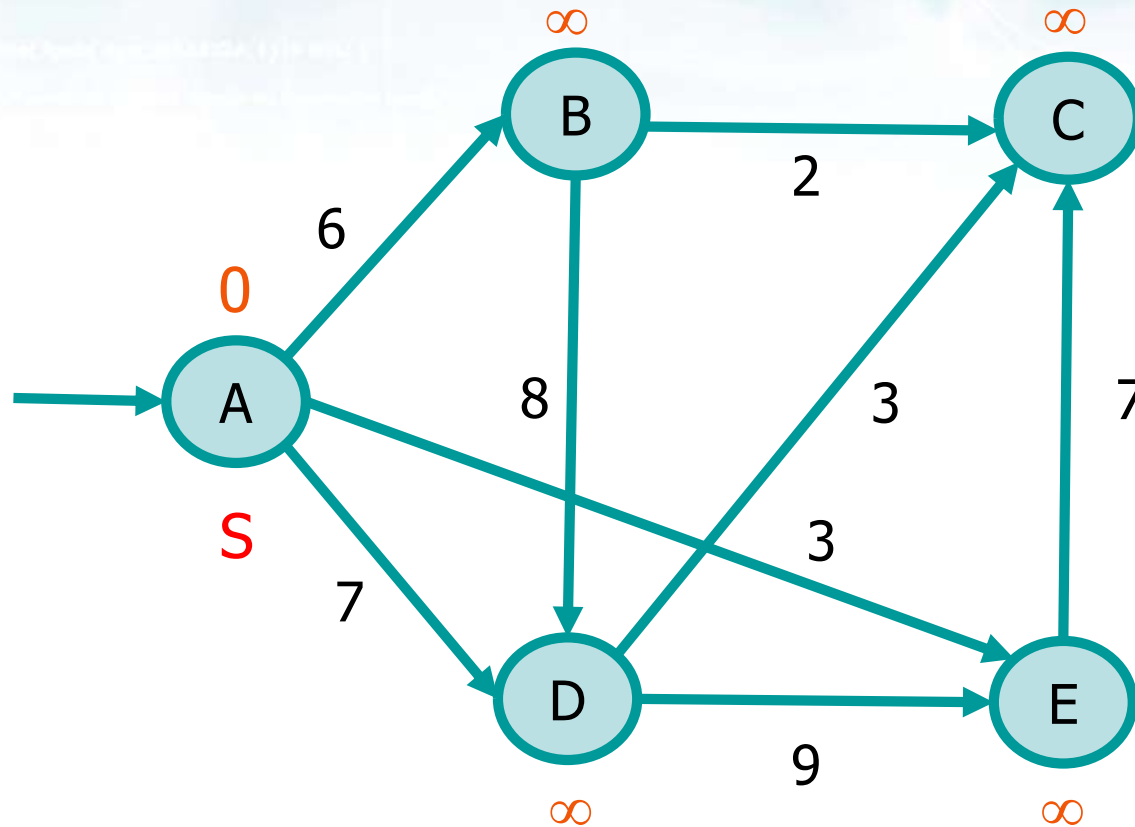
```

sssp_for_DAGs (G, w, s)
    topological sort the vertices of G
    initialize_single_source (G, s)

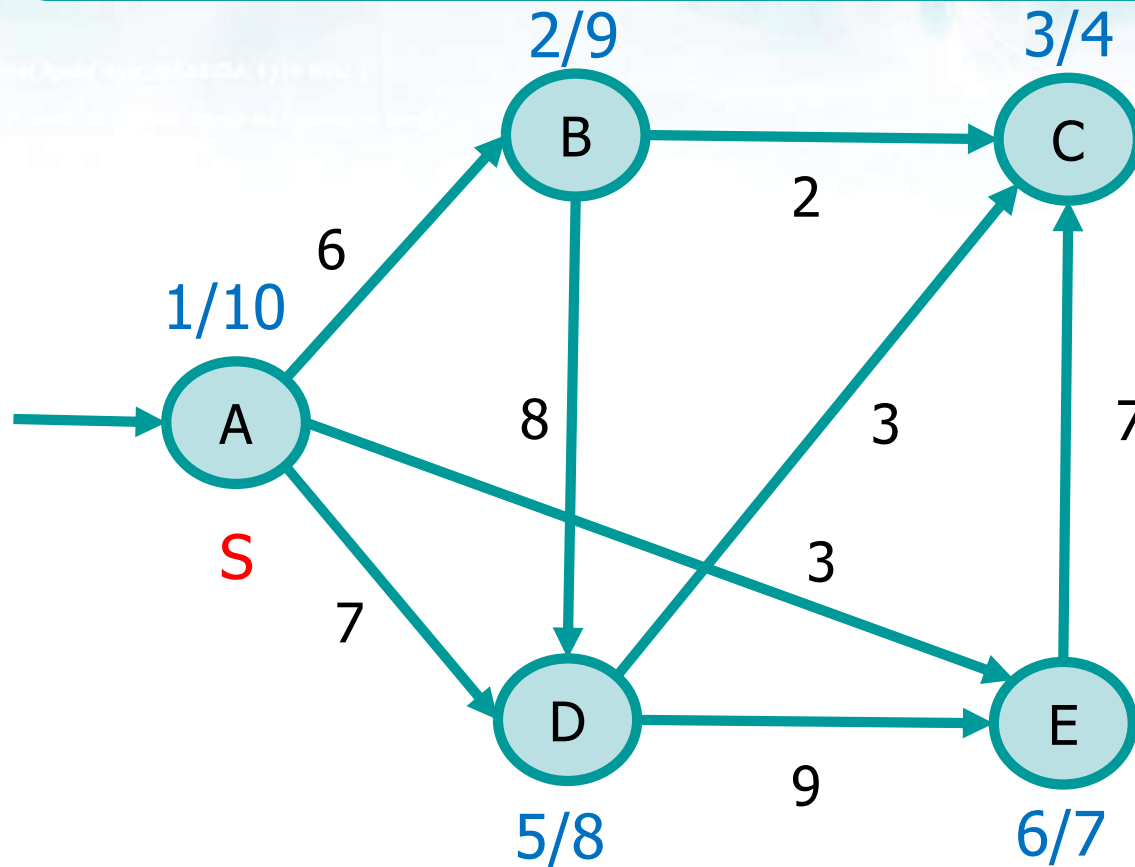
    for each vertex u ∈ V
        for each vertex v ∈ adjacency list of u
            relax (u, v, w)
    
```

Taken in topologically
sorted order

Example



Solution



Relaxation order

(A, B)

(A, D)

(A, E)

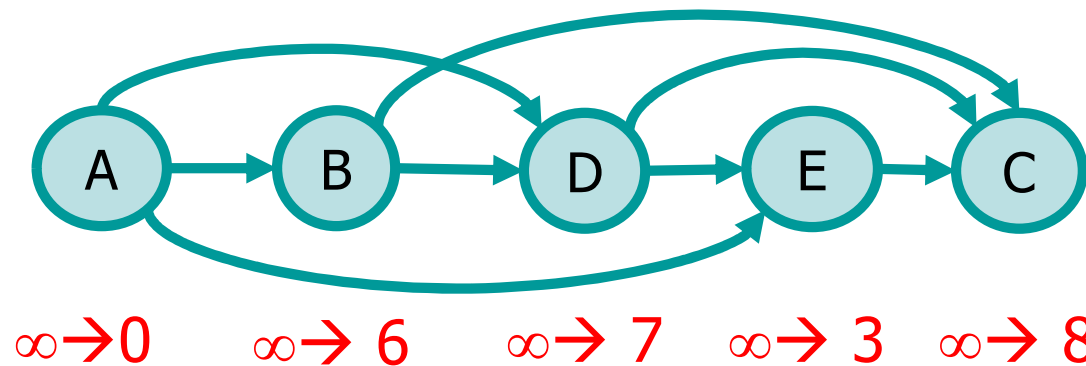
(B, D)

(B, C)

(D, E)

(D, C)

(E, C)



Complexity

Pseudo-code

```
sssp_for_DAGs (G, w, s)  
  topological sort the vertices of G  
  initialize_single_source (G, s)
```

```
  for each vertex  $u \in V$   
    for each vertex  $v \in$  adjacency list of  $u$   
      relax ( $u, v, w$ )
```

$\Theta(|V| + |E|)$

$\Theta(|V|)$

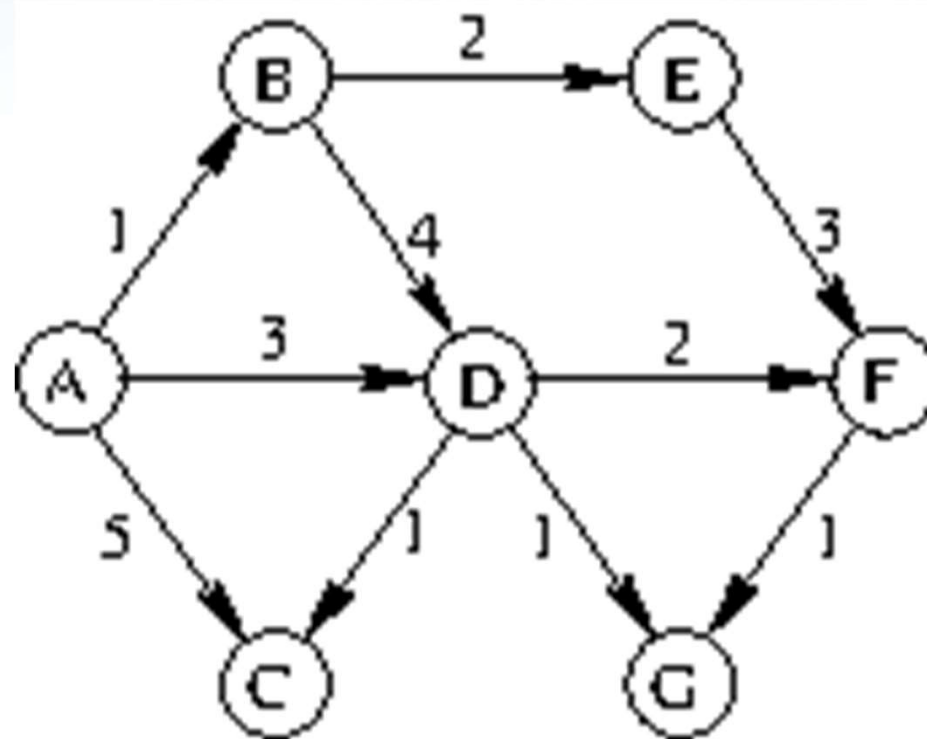
Executed E times
altogether

$\Theta(1) \rightarrow \Theta(|E|)$

Taken in topological
sorted order

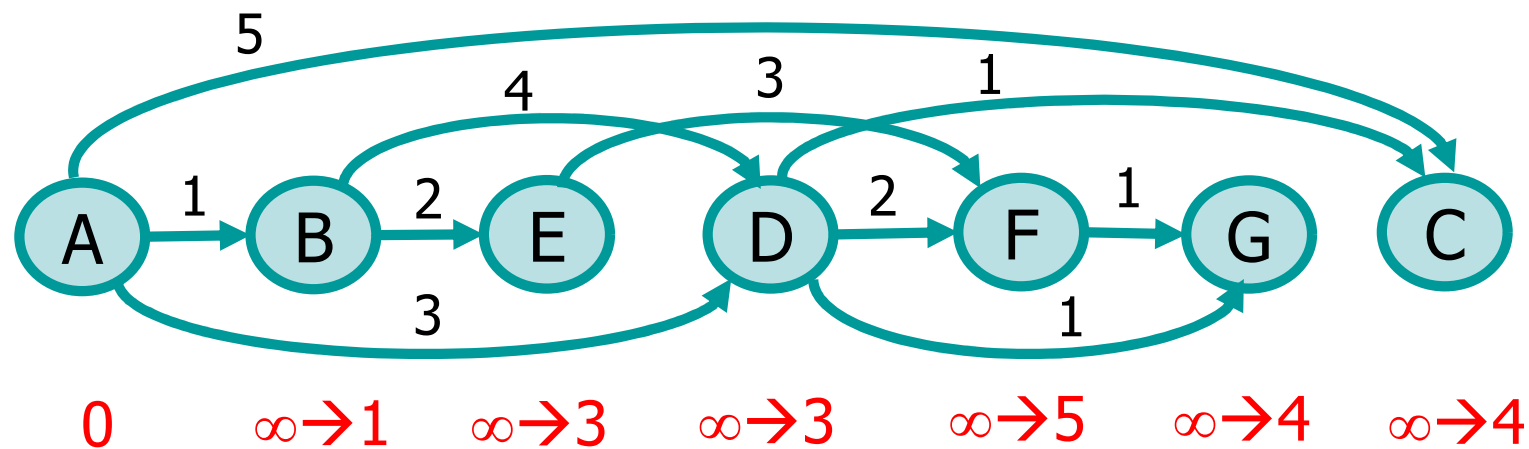
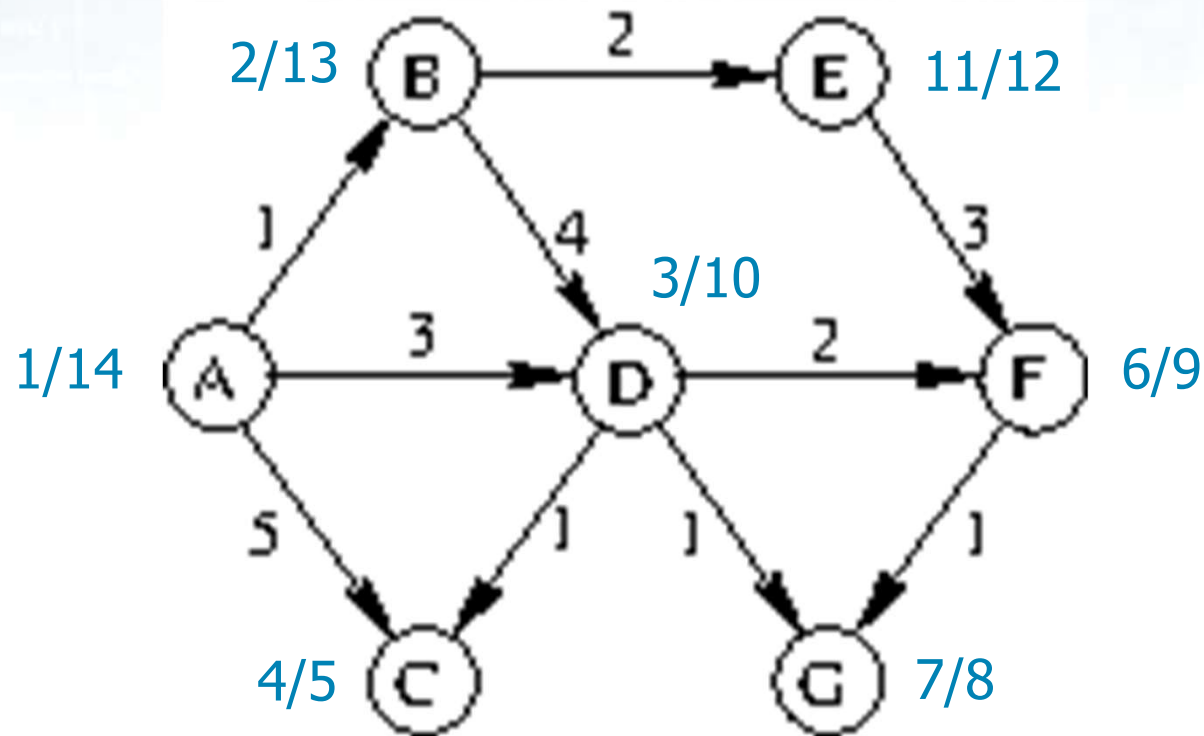
Overall running time complexity
 $T(n) = \Theta(|V| + |E|)$

Exercise

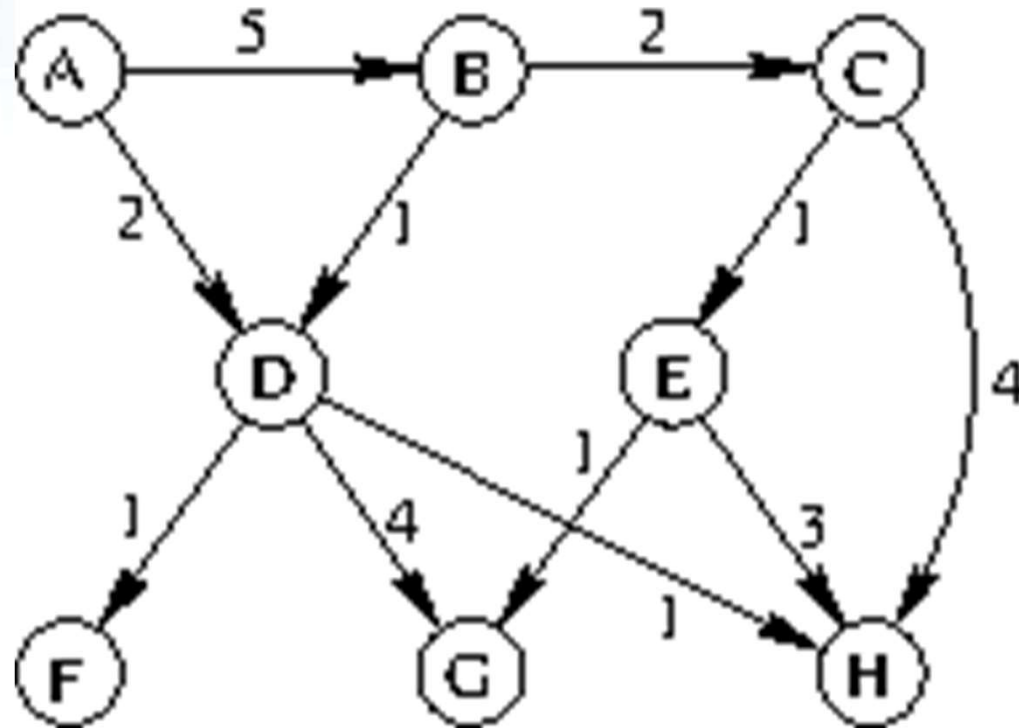


- ❖ Given the following graph find all shortest paths starting from vertex A

Solution

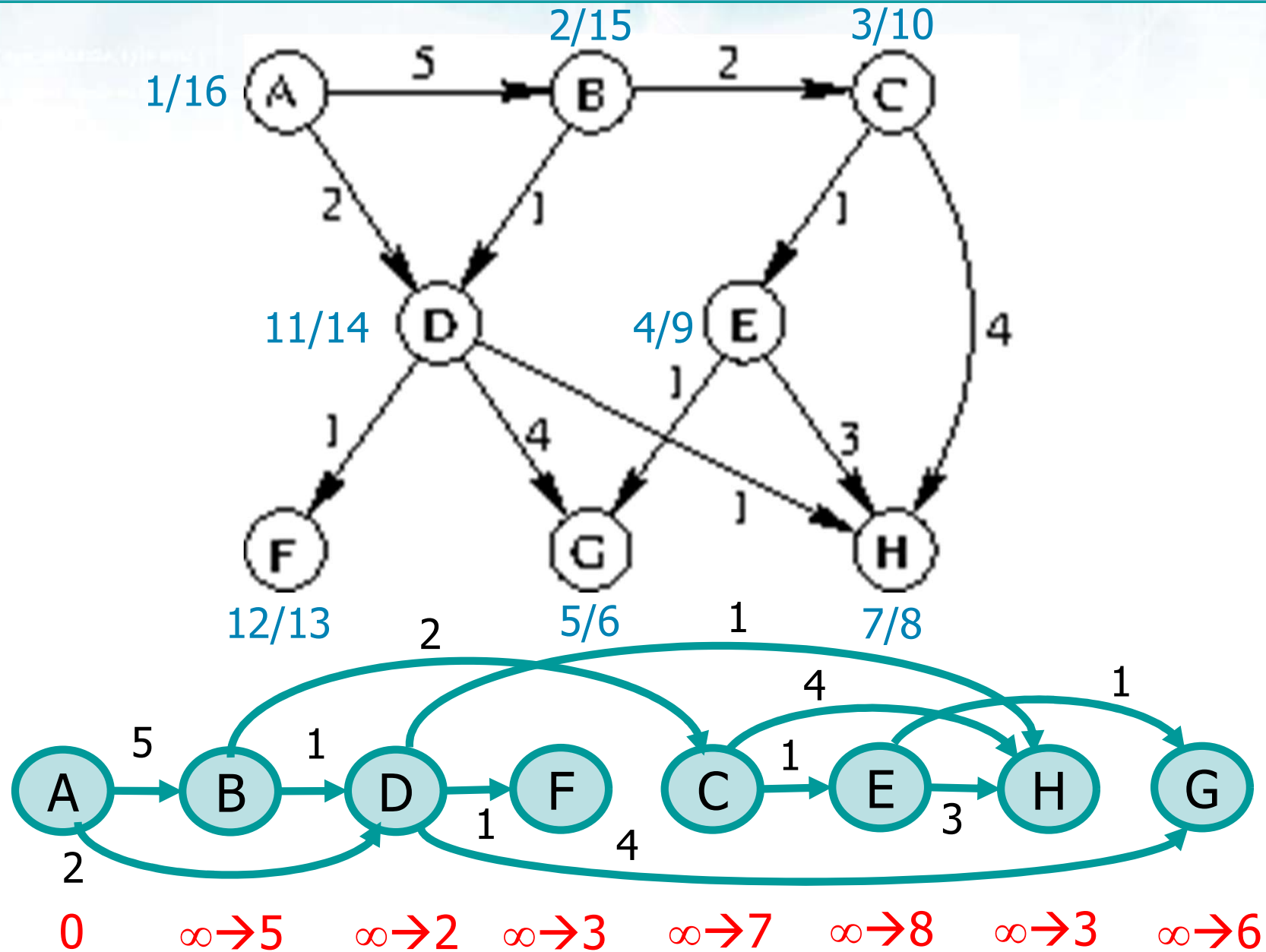


Exercise



- ❖ Given the following graph find all shortest paths starting from vertex A

Solution

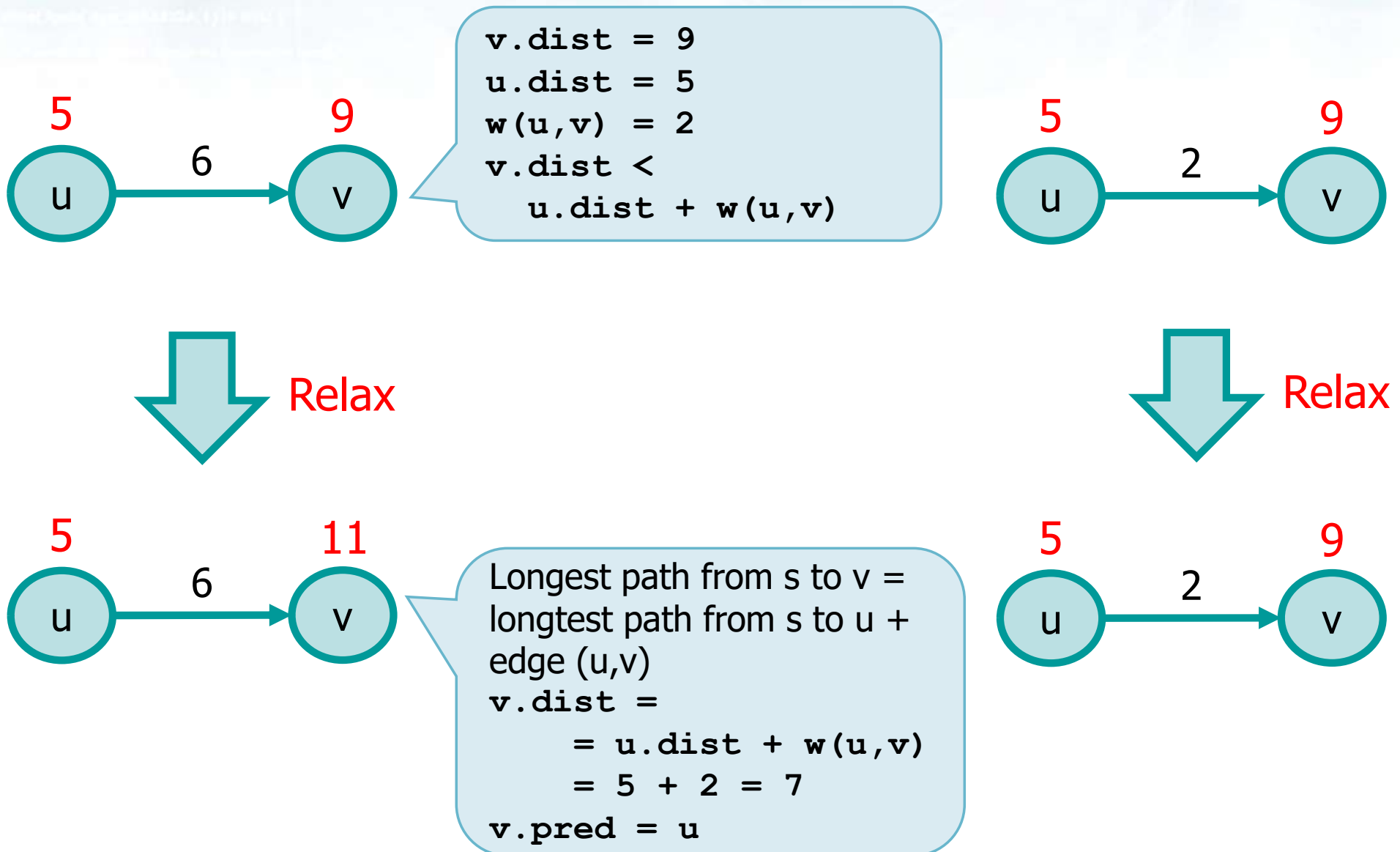


Longest path on weighted DAG

- ❖ Problem intractable on generic weighted graph
- ❖ As on a DAG there are no cycles, the problem become computationally feasible
 - Topologically sort the DAG
 - For all ordered vertices
 - Apply the "inverse" relaxation rule starting from that vertex

```
inverse_relax (u, v, w) {
    if (v.dist < u.dist + w(v,u)) {
        v.dist = u.dist + w(v,u)
        v.pred = u
    }
}
```

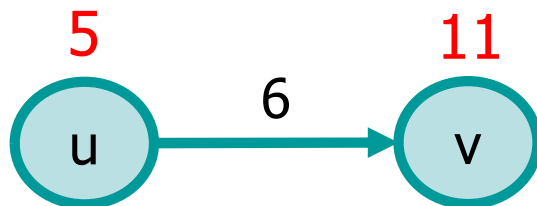
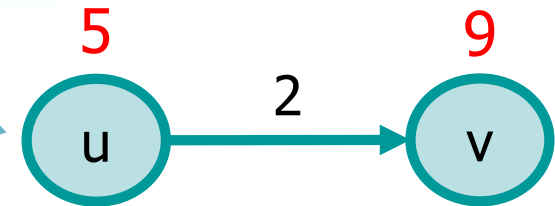
Example



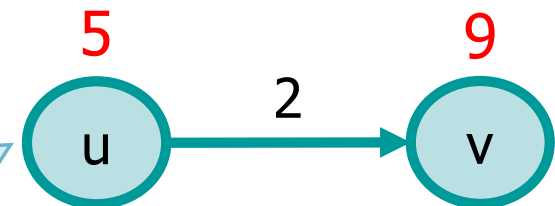
Example



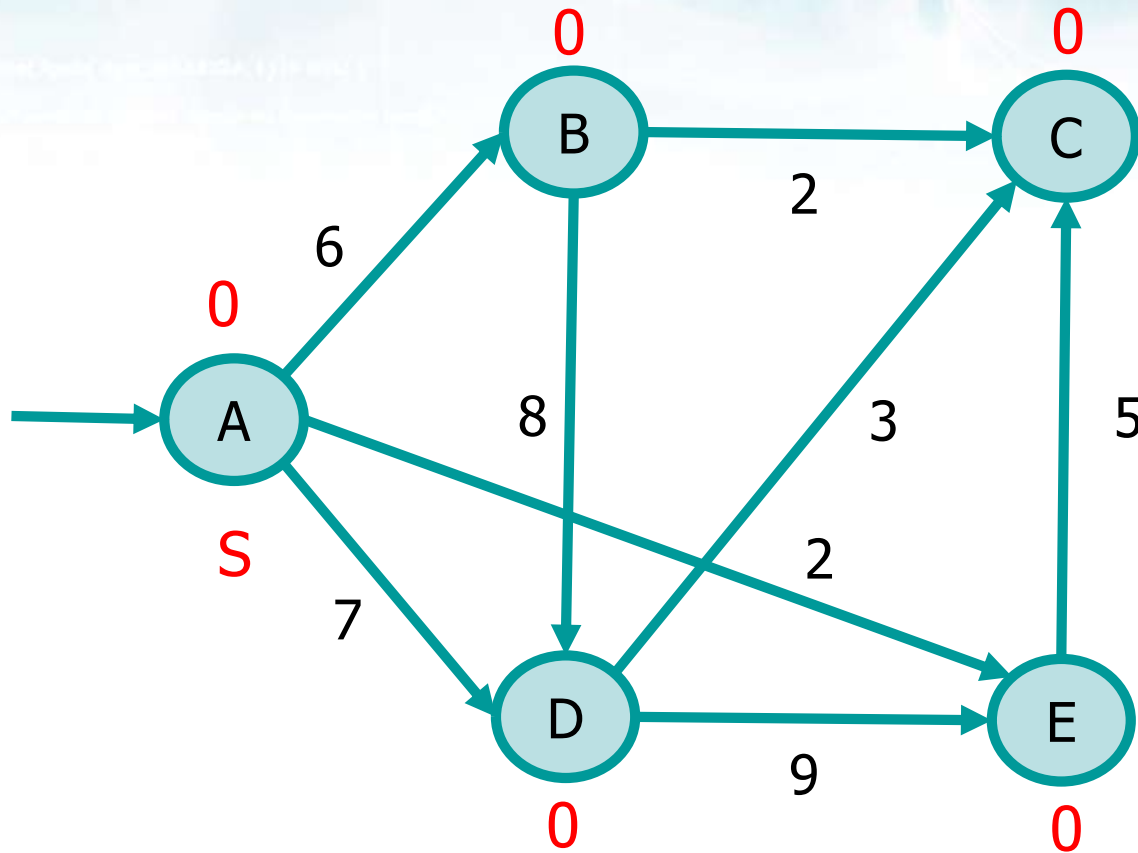
`v.dist = 9`
`u.dist = 5`
`w(u,v) = 2`
`v.dist >`
`u.dist + w(u,v)`



Relaxation has no effect
`v.dist = unchanged`
`= 6`
`v.pred = unchanged`

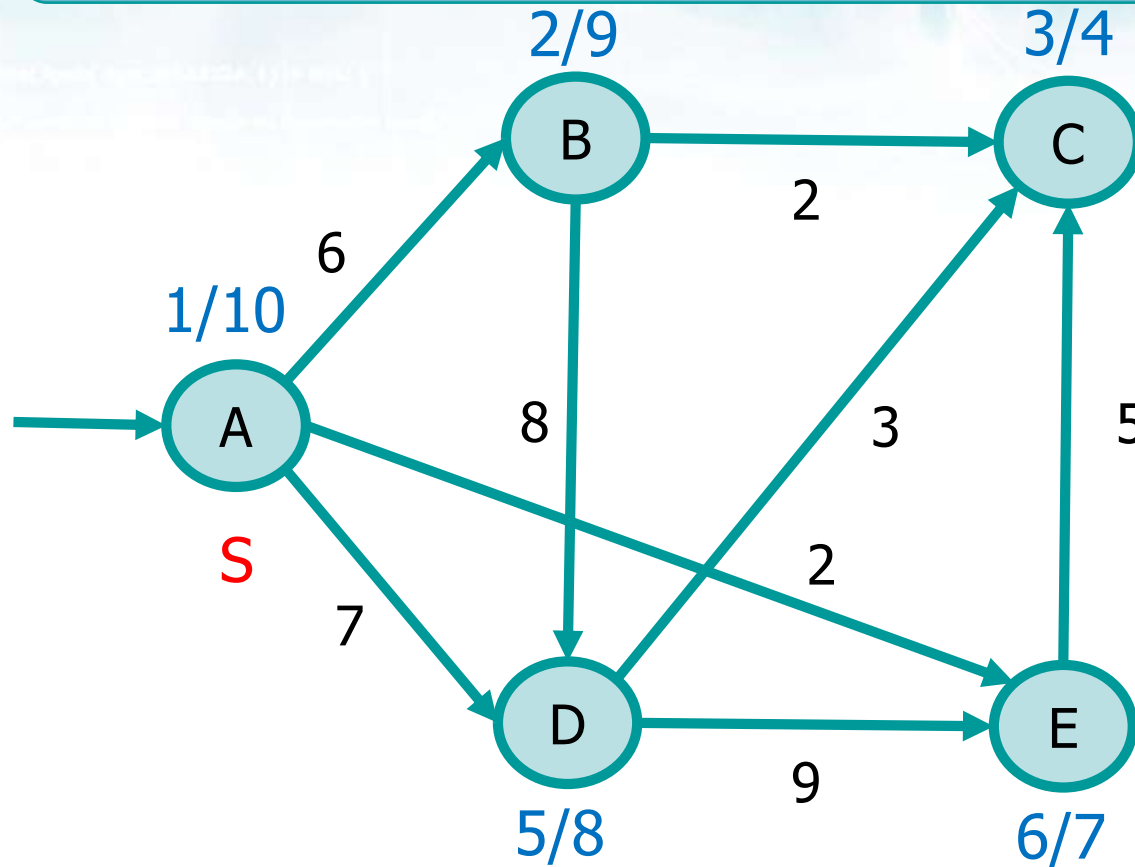


Example



The initial estimate is equal to zero for all vertices

Solution



Relaxation order

(A, B)

(A, D)

(A, E)

(B, D)

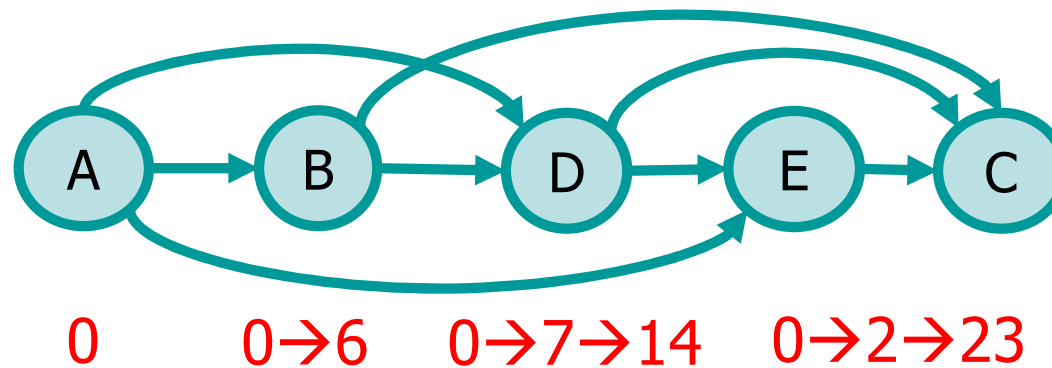
(B, C)

(D, E)

(D, C)

(E, C)

Relaxation order



0→8→17→28

0

0→6

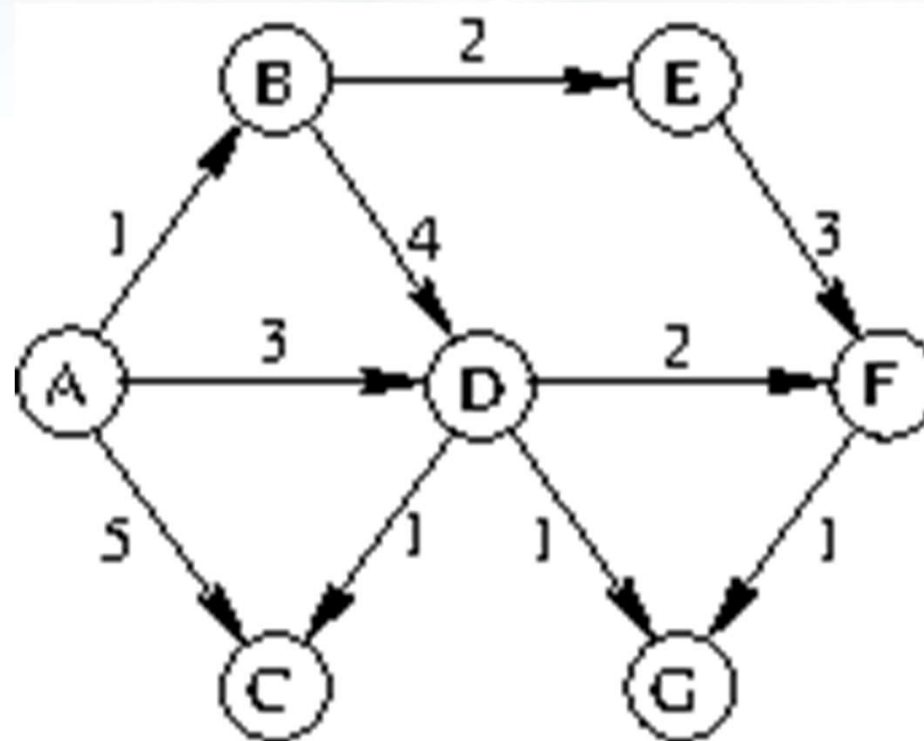
0→7→14

0→2→23

Complexity

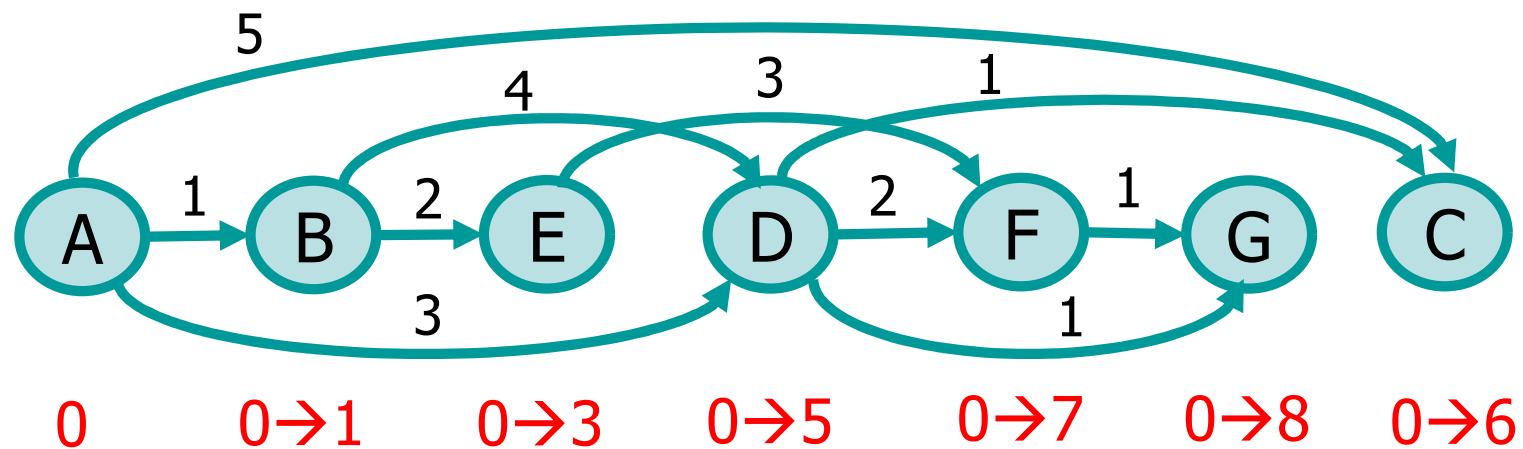
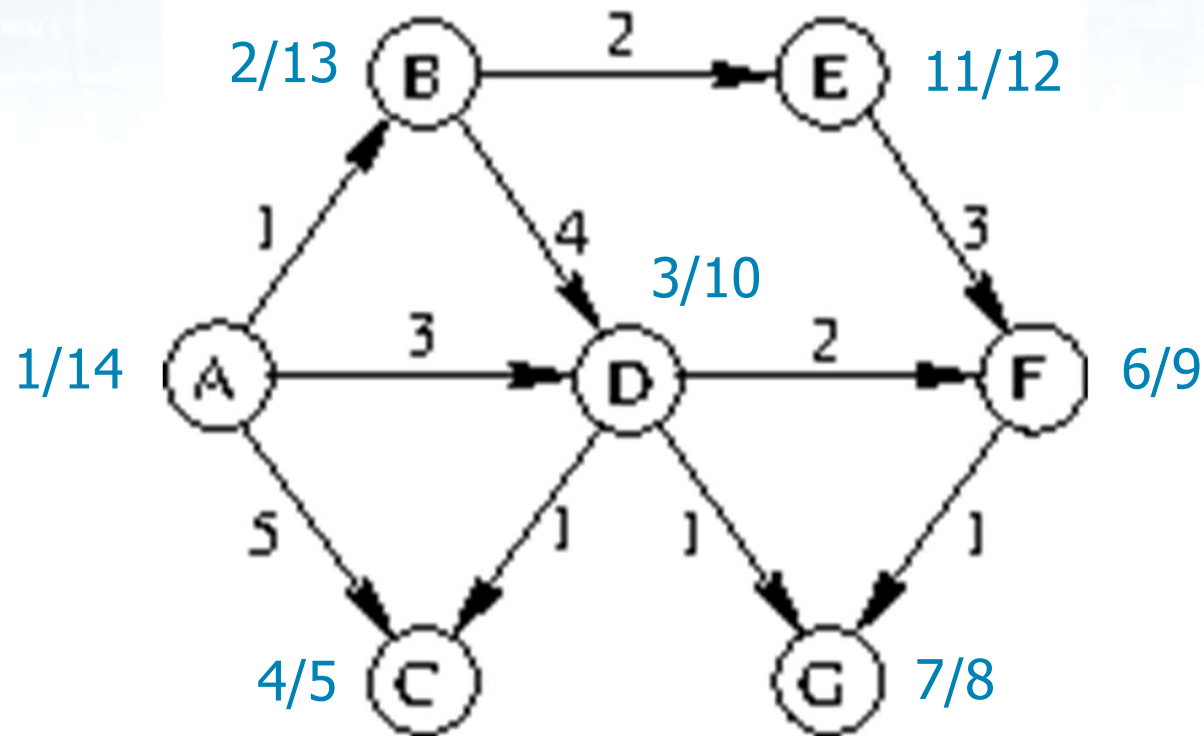
- ❖ As the algorithm for the shortest paths for DAGs

Exercise

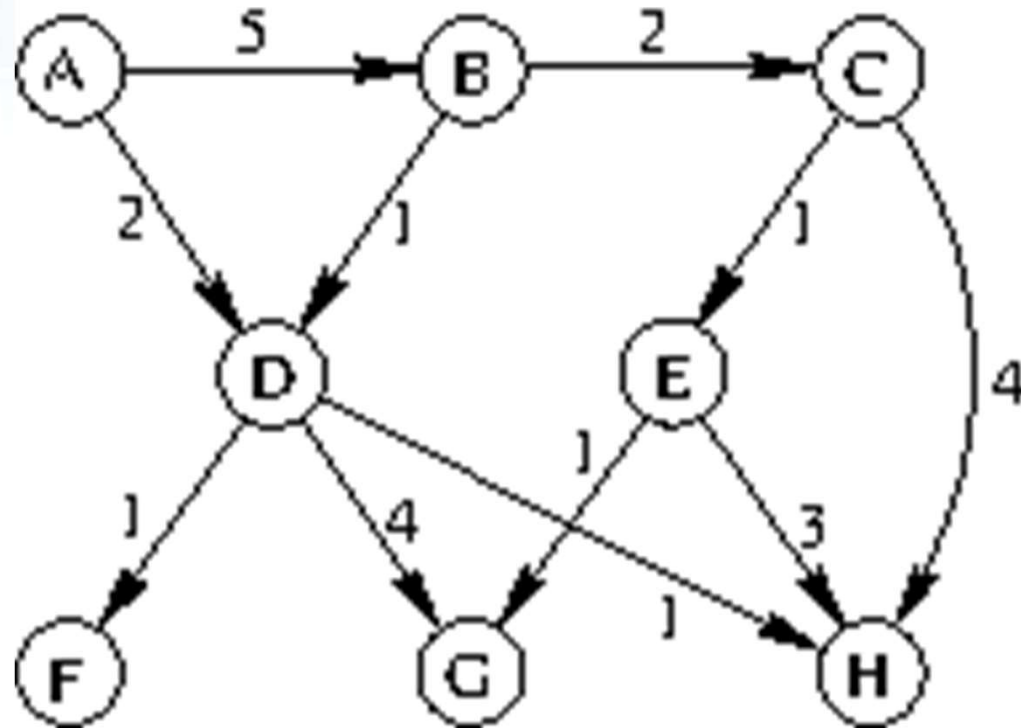


- ❖ Given the following graph find all longest paths starting from vertex A

Solution



Exercise



- ❖ Given the following graph find all longest paths starting from vertex A

Solution

