Project 1 - FYS3150 Computational Physics

Fredrik Østrem (fredost)

September 19, 2016

Contents

| 1 | Introduction | 2 |
|--------------|---|---|
| 2 | Methods 2.1 Approximate solution to $-u''(x) = f(x) \dots \dots \dots \dots$ 2.2 General algorithm for solving tridiagonal matrix equations $\dots \dots$ | |
| 3 | 3 Results | |
| 4 | Conclusions | |
| \mathbf{A} | Appendix | 6 |

1 Introduction

In this project, we solve the one-dimensional Possion equation

$$-u''(x) = f(x) \tag{1}$$

in the case where u(0) = u(1) = 0 and $f(x) = 100e^{-10x}$. We shall set up a matrix equation

$$A\mathbf{v} = \mathbf{b} \tag{2}$$

and generate efficient algorithms for solving this equation to get a numerical solution to the differential equaion. We shall also compare these algorithms to LU decomposition, and see how much more efficient the specialized algorithms are.

2 Methods

2.1 Approximate solution to -u''(x) = f(x)

We consider a differential equation on the form -u''(x) = f(x) with $x \in (0,1)$ and u(0) = u(1) = 0, and where f is a known function.

We consider the functions u and f at a sequence of equally spaces points x_0, x_1, \ldots, x_n in the interval [0,1], so that $x_i = ih$ where $h = \frac{1}{n+1}$. We let $v_i = \tilde{u}(x_i) \approx u(x_i)$ be the value of our approximate solution at x_i , and let $b_i = h^2 f(x_i)$. Since u(0) = u(1), we have the boundary condition $v_0 = v_{n+1} = 0$.

We can approximate u''(x) by using the Taylor's expansion of u(x), which yields:

$$u''(x_i) \approx \frac{u(x_i + h) + u(x_i - h) - 2u(x_i)}{h^2} \approx \frac{v_{i+1} + v_{i-1} - 2v_i}{h^2}$$
(3)

which, when substituted into our differential equation gives

$$2v_i - v_{i+1} - v_{i-1} = b_i \tag{4}$$

Since this is a linear equation with v_{i-1}, v_i, v_{i+1} as unknowns, we can write this as

$$\underbrace{\begin{pmatrix} 0 & 0 & \cdots & -1 & 2 & -1 & \cdots & 0 & 0 \end{pmatrix}}_{\mathbf{a}_{i}} \begin{pmatrix} v_{0} \\ v_{1} \\ \vdots \\ v_{i-1} \\ v_{i} \\ v_{i+1} \\ \vdots \\ v_{n-1} \\ v_{n} \end{pmatrix}}_{= b_{i} \tag{5}$$

We can do this for every index i, and \mathbf{a}_i is shifted a single step to the right compared to \mathbf{a}_{i-1} . Therefore, when we look at all values of $i = 1, \ldots, n$, we get the tridiagonal

matrix A:

$$A = \begin{pmatrix} \mathbf{a}_0 \\ \vdots \\ \mathbf{a}_{i-1} \\ \mathbf{a}_i \\ \mathbf{a}_{i+1} \\ \vdots \\ \mathbf{a}_n \end{pmatrix} = \begin{pmatrix} 2 & -1 & 0 & \cdots & 0 \\ -1 & 2 & -1 & 0 & \ddots \\ 0 & -1 & 2 & \ddots & 0 & \vdots \\ \vdots & 0 & \ddots & 2 & -1 & 0 \\ & \ddots & 0 & -1 & 2 & -1 \\ 0 & & \cdots & 0 & -1 & 2 \end{pmatrix}$$
(6)

such that $A\mathbf{v} = \mathbf{b}$.

2.2 General algorithm for solving tridiagonal matrix equations

In general, we can write an $n \times n$ tridiagonal matrix A as

$$A = \begin{pmatrix} b_1 & c_1 & 0 & \dots & \dots & \dots \\ a_1 & b_2 & c_2 & \dots & \dots & \dots \\ & a_2 & b_3 & c_3 & \dots & \dots \\ & & & \ddots & \dots & \dots & \dots \\ & & & & a_{n-2} & b_{n-1} & c_{n-1} \\ & & & & a_{n-1} & b_n \end{pmatrix}$$
 (7)

We can solve the matrix equation $A\mathbf{v} = \mathbf{b}$ in three steps:

1. Eliminate the lower diagonal $a_1, a_2, \ldots, a_{n-1}$ through forward substitution.

•
$$a'_{i-1} = 0$$
; $b'_i = b_i - \frac{a_{i-1}}{b'_{i-1}} \cdot c_{i-1}$; $f'_i = f_i - \frac{a_{i-1}}{b'_{i-1}} \cdot f'_{i-1}$

2. Eliminate the upper diagonal $c_1, c_2, \ldots, c_{n-1}$ through backward substitution.

•
$$c_i'' = 0$$
; $f_i'' = f_i' - \frac{c_i}{b_{i+1}'} \cdot f_{i+1}''$

3. Divide each row i by b_i to get only 1 elements along the main diagonal.

•
$$v_i = \frac{f_i''}{b_i'}$$

When doing these steps, we transform A into the identity matrix, and transform \mathbf{b} into \mathbf{v} . (This algorithm has been implemented in the file tridiagonal.hh.) We can see the results of the approximation in figure 1.

From the implementation, we can see that the first step uses 5 floating-point operations for every row except the first, the second step uses 5 for every row except the last, and the third step uses 2 for every row; in total, the number of floating-point operations is 5(N-1) + 5(N-1) + 2N = 12N - 10.

3 Results

We run the algorithm from section 2.2 from different values of n. In figure 1 we see that for small values of N, the curve keeps the right shape but is quite a bit off from the analytic solution; however, for N = 1000 it is almost indistinguishable from the analytic solution. In table 2, where we have listed the maximum value of

$$\varepsilon_i = \log_{10} \left| \frac{v_i - u_i}{u_i} \right| \tag{8}$$

for each N, as a function of $\log_{10}(h)$ where h is the step length, we see that the relative error becomes decreases exponentially as h decreases.

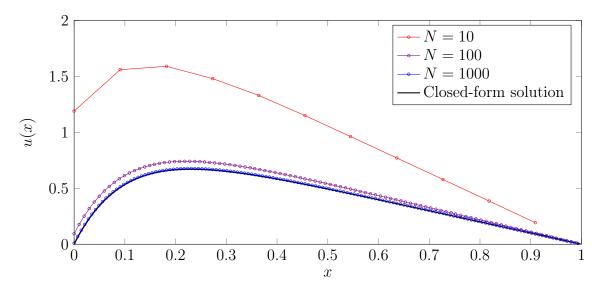


Figure 1: Plot of approximate solutions based on the method described in section 2.2

| $\log_{10}(h)$ | $arepsilon_{max}$ |
|----------------|----------------------|
| -1 | $3.17 \cdot 10^{-1}$ |
| -2 | $3.85 \cdot 10^{-2}$ |
| -3 | $3.95 \cdot 10^{-3}$ |
| -4 | $3.96 \cdot 10^{-4}$ |
| -5 | $3.96\cdot10^{-5}$ |
| -6 | $4.10 \cdot 10^{-6}$ |
| -7 | $1.34\cdot10^{-6}$ |

Table 2: ε_{max} as a function of $\log_{10}(h)$.

Table 1

If we compare the runtime of our tridiagonal algorithm to one using LU decomposition of the matrix (the implementation given in the course's code files), we see that

| N | Time (tridiagonal) / ms | Time (LU decomposition) / ms |
|-------|-------------------------|------------------------------|
| 10 | 0.001 | 0.008 |
| 100 | 0.004 | 1.925 |
| 1,000 | 0.050 | 1839.400 |

Table 3: Running times for our specialized algorithm, and a LU decomposition algorithm.

the tridiagonal algorithm is very fast, even for large N, while the LU decomposition runs at about 1s for N = 1000:

There's a big difference here because our simple algorithm for tridiagonal matrices only needs on the order of N float-point operations, so it run in less than a millisecond, even for very large matrices. However, LU decomposition is much more general, and needs on the order of N^3 different floating-point operations for each matrix. If N is very large, the number of operations grows very fast.

4 Conclusions

From the results we got, we can conclude that numerically solving a differential equation by using a matrix equation can work well, but we a small enough step length (N = 10 is way off, for instance) and a sufficiently efficient algorithm for solving the matrix equation (LU decomposition is too slow for large N).

A Appendix

All files used in this project can be found at https://github.com/frxstrem/fys3150/tree/master/project1. The following code files are used:

- tridiagonal.hh
- oppg_b.cc
- oppg_d.cc
- oppg_e.cc