

# SF3580

## HW 3

Anna Broms & Fredrik Fryklund

2018/11/29

### Task 2

### Task 4

### Task 5

(a)

The matrix  $A$  is diagonalizable with the eigendecomposition  $A = QDQ^{-1}$ , where  $D$  is a diagonal matrix. For such structures it holds that  $\sin(A) = Q \sin(D)Q^{-1}$ . Thus we can validate the result for the Schur-Parlett method, which is

$$\sin(A) = \sin\left(\begin{bmatrix} 1 & 4 & 4 \\ 3 & -1 & 3 \\ -1 & 4 & 4 \end{bmatrix}\right) \approx \begin{bmatrix} 0.846192 & 0.0655435 & -0.187806 \\ 0.33476 & 0.385017 & -0.141244 \\ -0.190921 & 0.192478 & 0.848269 \end{bmatrix}. \quad (1)$$

which in norm differs  $4.28e - 16$  from  $Q \sin(D)Q^{-1}$ .

(b) & (c)

It is clear from Figure 1 that the number of flops required for Schur-Parlett is not discernibly affected by  $N$ , at least for  $N \in 10, 50, 100, 150, 200, 250, 300$ . This is not surprising, as often the most computationally demanding part of the Schur-Parlett method, is in performing the Schur decomposition, which scales like  $O(n^3)$ . Once obtained, the function  $f$  is only applied to the diagonal elements, which are scalars.

For the naive approach the number of flops is proportional to  $N$ . A matrix multiplication is of  $O(n^3)$ , thus performing  $N$  matrix gives  $O(Nn^3)$ , which we read from Figure 1. The black line corresponds to the line  $0.08 + 0.07 N$ .

### Task 6

(a)

The matrix

$$A = \begin{bmatrix} \pi & 1 \\ 0 & \pi + \varepsilon \end{bmatrix}, \quad (2)$$

with  $\varepsilon > 0$ , has two eigenvalues:  $\lambda_1 = \pi$  and  $\lambda_2 = \pi + \varepsilon$ . Let  $A = X \operatorname{diag}(J_1, J_2) X^{-1}$ , be the Jordan canonical form, with  $J_1 = \lambda_1$  and  $J_2 = \lambda_2$ .

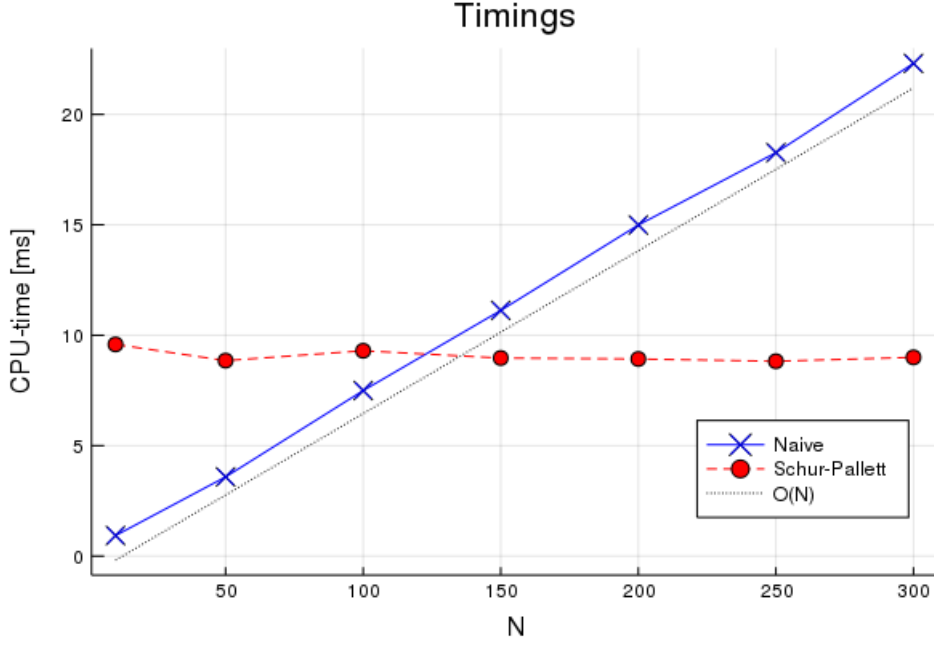


Figure 1: Task 5, (b) & (c): CPU-time in milliseconds, as a function of  $N$ .

The Jordan canonical form definition gives that

$$p(A) = X \text{diag}(p(J_1), p(J_2)) X^{-1}. \quad (3)$$

A simple consequence is

$$g(A) = X \text{diag}(g(\lambda_1), g(\lambda_2)) X^{-1} = X \text{diag}(p(\lambda_1), p(\lambda_2)) X^{-1} = p(A), \quad (4)$$

since the polynomial  $p$  interpolates the function  $g$  in the eigenvalues of  $A$ , i.e.  $p(\lambda_1) = g(\lambda_1)$  and  $p(\lambda_2) = g(\lambda_2)$ .

Two points defines a unique polynomial of order 1, thus we may choose a  $p$  in  $\mathbb{P}^1$  and write  $p(z) = \alpha + \beta z$ . The unknown coefficients are obtained by solving

$$\begin{cases} \alpha + \beta \lambda_1 = g(\lambda_1) \\ \alpha + \beta \lambda_2 = g(\lambda_2) \end{cases} \Leftrightarrow \begin{cases} \alpha = \frac{g(\lambda_1)\lambda_2 - g(\lambda_2)\lambda_1}{\lambda_2 - \lambda_1} \\ \beta = \frac{g(\lambda_2) - g(\lambda_1)}{\lambda_2 - \lambda_1} \end{cases} \quad (5)$$

**(b)**

Given  $g := \exp$  we have from (a) and (b) that

$$p(A) = \alpha I + \beta A = \frac{\exp(\pi)(\pi + \varepsilon) - \exp(\pi + \varepsilon)\pi}{\varepsilon} I + \frac{\exp(\pi + \varepsilon) - \exp(\pi)}{\varepsilon} A \quad (6)$$

$$= \frac{\exp(\pi)}{\varepsilon} (\varepsilon I + (1 - \exp(\varepsilon))(\pi I - A)) \quad (7)$$

**(c)**

It is known that the Jordan decomposition is unstable for non-symmetric matrices, as the eigenvalues may lie close to each other. For the given matrix  $A$  this can be tuned artificially by setting  $\varepsilon = |\lambda_1 - \lambda_2|$ . As we see in Figure (to be added, some layout to fix). However, the function `exp` in Julia is analogous to `expm` in Matlab, (we compared the results for  $\varepsilon = 1e-1, \dots, 1e-10$ ). Using Matlabs `expm(A)` as reference value, the error from the exact result in (b) also increases as  $\varepsilon$ . This is most likely due to cancellation.

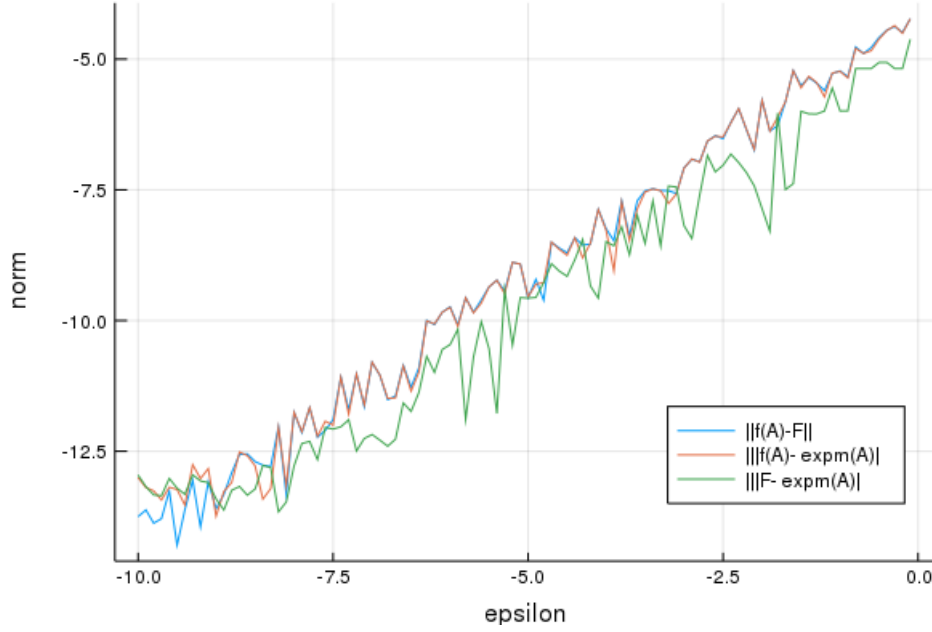


Figure 2: Task 5, (b) & (c): CPU-time in milliseconds, as a function of  $N$ .

## Task 7

(a)

Consider the function  $f(z, t) = e^{tz}$ . We want to investigate the matrix valued function  $f(A, t) = e^{tA}$ . Let  $\mu \in \mathbb{C}$  be an expansion point. Then,

$$f(A, t) = \sum_{i=0}^{\infty} \frac{f^{(i)}(\mu)}{i!} (A - \mu I)^i = \sum_{i=0}^{\infty} \frac{t^i e^{t\mu}}{i!} (A - \mu I)^i. \quad (8)$$

If  $A \in \mathbb{C}^{n \times n}$ , then  $f : \mathbb{C}^{n \times n} \rightarrow \mathbb{C}^{n \times n}$ . Now, compute the derivative of  $f(A, t)$  with respect to time:

$$\begin{aligned} \frac{d}{dt} e^{tA} &= \frac{d}{dt} \sum_{i=0}^{\infty} \frac{t^i e^{t\mu}}{i!} (A - \mu I)^i = \\ &= \sum_{i=0}^{\infty} \frac{d}{dt} \left( \frac{t^i e^{t\mu}}{i!} (A - \mu I)^i \right) = \\ &= \sum_{i=0}^{\infty} \frac{d}{dt} \left( \frac{t^i e^{t\mu}}{i!} \right) (A - \mu I)^i = \\ &= \sum_{i=0}^{\infty} \left( \frac{it^{i-1} e^{t\mu} + t^i \mu e^{t\mu}}{i!} \right) (A - \mu I)^i. \end{aligned} \quad (9)$$

The last expression can be identified as  $g(A)$ , where  $g(z) = z e^{tz}$  as the expression

$$(it^{i-1} e^{t\mu} + t^i \mu e^{t\mu}) \quad (10)$$

is the  $i$ th derivative of the product  $z \cdot e^{tz}$ , which can be seen using the general Leibniz rule  $((f_1 f_2)^{(n)}) = \sum_{k=0}^n \binom{n}{k} f_1^{(n-k)}(x) f_2^{(k)}(x)$ . Thus, we can conclude that  $\frac{d}{dt} e^{tA} = A e^{tA}$ . The matrix function  $e^{tA} A$  has the same Taylor expansion expression as  $A e^{tA}$ . Thus,  $\frac{d}{dt} e^{tA} = A e^{tA} = e^{tA} A$ , which is what we wanted to show.

(d)

Task: Let  $C_k = [C_{k-1}, A]$ , with  $C_0 = B$ . We want to show that  $\|C_k\| \leq 2^k \|A\|^k \|B\|$ .

The proof is done by induction. For  $k = 0$  we have that  $\|C_0\| = \|B\| \leq 2^0 \|A\|^0 \|B\|$ . Now, assume that  $\|C_k\| \leq 2^k \|A\|^k \|B\|$ . We want to show that  $\|C_{k+1}\| \leq 2^{k+1} \|A\|^{k+1} \|B\|$ :

$$\begin{aligned} \|C_{k+1}\| &= \|C_k A - A C_k\| = \|C_k A + (-A C_k)\| \leq \|C_k A\| + \|-A C_k\| = \|C_k A\| + \|A C_k\| \leq \|C_k\| \|A\| + \|A\| \|C_k\| = \\ &= 2 \|A\| \|C_k\| = 2^{k+1} \|A\|^{k+1} \|B\|, \end{aligned} \quad (11)$$

which is what we wanted to show.

(e)

Suppose  $\|A\| < \frac{1}{2}$  and  $t \leq 1$ . Let  $G_N(t)$  be the truncation of  $G(t)$ , where

$$G(t) = \sum_{k=0}^{\infty} \frac{t^k}{k!} C_k. \quad (12)$$

Then,

$$\begin{aligned} \|G_N(t) - G(t)\| &= \left\| \sum_{k=N+1}^{\infty} \frac{t^k}{k!} C_k \right\| \leq \sum_{k=N+1}^{\infty} \left( \frac{t^k}{k!} \right) \|C_k\| \leq \\ &\leq \sum_{k=N+1}^{\infty} \left( \frac{t^k}{k!} \right) 2^k \|A\|^k \|B\| \leq \sum_{k=N+1}^{\infty} \left( \frac{t^k}{k!} \right) 2^k \left( \frac{1}{2} \right)^2 \|B\| \leq \frac{\|B\|}{(N+1)!} \sum_{k=N+1}^{\infty} t^k = \frac{\|B\|}{(N+1)!} \sum_{k=0}^{\infty} t^{k+(N+1)} = \\ &= \frac{\|B\| t^{N+1}}{(N+1)!} \cdot \frac{1}{1-t}. \end{aligned} \quad (13)$$