Référence d'Équipe

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1 Dynamic Programming

1.1 convex hull trick

```
struct line {
 long long m, b;
 line (long long a, long long c) : m(a), b(c) {}
 long long eval(long long x) {
   return m * x + b;
long double inter(line a, line b) {
 long double den = a.m - b.m;
 long double num = b.b - a.b;
 return num / den;
/**
* \min m_i * x_j + b_i, for all i.
      x_j \le x_{j+1}
      m_i >= m_{j+1}
struct ordered_cht {
 vector<line> ch;
 int idx; // id of last "best" in query
  ordered_cht() {
   idx = 0:
  void insert_line(long long m, long long b) {
   line cur(m, b);
   // new line's slope is less than all the previous
```

```
while (ch.size() > 1 &&
       (inter(cur, ch[ch.size() - 2]) >= inter(cur,
      ch[ch.size() - 1]))) {
        // f(x) is better in interval [inter(ch.back(),
      cur), inf)
        ch.pop_back();
    ch.push_back(cur);
  long long eval(long long x) { // minimum
    // current x is greater than all the previous x,
    // if that is not the case we can make binary search.
    idx = min<int>(idx, ch.size() - 1):
    while (idx + 1 < (int)ch.size() && ch[idx + 1].eval(x)
      <= ch[idx].eval(x))
      idx++:
    return ch[idx].eval(x);
};
// Dynammic convex hull trick
typedef long long int64;
typedef long double float128;
const int64 is_query = -(1LL<<62), inf = 1e18;</pre>
struct Line {
  int64 m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {</pre>
    if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    if (!s) return 0;
    int64 x = rhs.m;
    return b - s \rightarrow b < (s \rightarrow m - m) * x;
};
struct HullDynamic : public multiset<Line> { // will
     maintain upper hull for maximum
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
    return (float128)(x->b-y->b)*(z->m-y->m) >=
      (float128)(y->b - z->b)*(y->m - x->m);
  void insert_line(int64 m, int64 b) {
    auto y = insert({ m, b });
```

1.2 divide and conquer

1.3 dp on trees

```
/**
 * for any node, save the total answer and the answer of
     every children.
 * for the query(node, pi) the answer is ans[node] -
     partial[node][pi]
 * cases:
      - all children missing
      - no child is missing
     - missing child is current pi
void add_edge(int u, int v) {
 int id_u_v = g[u].size();
 int id_v_u = g[v].size();
 g[u].emplace_back(v, id_v_u); // id of the parent in the
     child's list (g[v][id] -> u)
 g[v].emplace_back(u, id_u_v); // id of the parent in the
     child's list (g[u][id] -> v)
```

2 Geometry

2.1 all

```
double INF = 1e100;
double EPS = 1e-12;
struct PT {
 double x, y;
 PT() {}
 PT(double x, double y) : x(x), y(y) {}
 PT(const PT \&p) : x(p.x), y(p.y) {}
 PT operator + (const PT &p) const { return PT(x+p.x,
    v+p.v); }
 PT operator - (const PT &p) const { return PT(x-p.x,
    y-p.y); }
                              const { return PT(x*c.
 PT operator * (double c)
    y*c ); }
 PT operator / (double c)
                              const { return PT(x/c,
    y/c ); }
double dot(PT p, PT q)
                          { return p.x*q.x+p.v*q.v; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
 return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.v,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
 double r = dot(b-a,b-a);
 if (fabs(r) < EPS) return a;</pre>
 r = dot(c-a, b-a)/r;
 if (r < 0) return a;</pre>
 if (r > 1) return b;
 return a + (b-a)*r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane
double DistancePointPlane(double x, double y, double z,
   double a, double b, double c, double d) {
 return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel
     or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;</pre>
```

```
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
    && fabs(cross(a-b, a-c)) < EPS
    && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
        dist2(b, c) < EPS || dist2(b, d) < EPS) return
    if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b,
     d-b) > 0)
     return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b=(a+b)/2;
  c=(a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c,
     c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon
     (by William
// Randolph Franklin); returns 1 for strictly interior
     points, 0 for
// strictly exterior points, and 0 or 1 for the remaining
     points.
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++){</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
          p[j].y \le q.y \&\& q.y \le p[i].y) \&\&
        q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y)
     /(p[j].y - p[i].y))
      c = !c;
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()],
     q), q) < EPS)
     return true;
  return false;
```

```
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double
     r) {
  vector<PT> ret;
  b = b-a;
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with
     radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r,
     double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R || d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly
// polygon, assuming that the coordinates are listed in a
     clockwise or
// counterclockwise fashion. Note that the centroid is
     often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0:
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++){</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[i])*(p[i].x*p[i].y - p[i].x*p[i].y);
  return c / scale;
```

2.2 center 2 points + radius

```
vector<point> find_center(point a, point b, long double r)
    {
    point d = (a - b) * 0.5;
    if (d.dot(d) > r * r) {
        return vector<point> ();
    }
    point e = b + d;
    long double fac = sqrt(r * r - d.dot(d));
    vector<point> ans;
    point x = point(-d.y, d.x);
    long double 1 = sqrt(x.dot(x));
    x = x * (fac / l);
    ans.push_back(e + x);
    x = point(d.y, -d.x);
    x = x * (fac / l);
    ans.push_back(e + x);
    return ans;
}
```

2.3 closest pair

```
struct point {
  double x, y;
  int id;
  point() {}
 point (double a, double b) : x(a), y(b) {}
double dist(const point &o, const point &p) {
 double a = p.x - o.x, b = p.y - o.y;
 return sqrt(a * a + b * b);
double cp(vector<point> &p, vector<point> &x,
     vector<point> &y) {
  if (p.size() < 4) {</pre>
   double best = 1e100;
   for (int i = 0; i < p.size(); ++i)</pre>
      for (int j = i + 1; j < p.size(); ++j)</pre>
        best = min(best, dist(p[i], p[j]));
    return best:
```

```
int ls = (p.size() + 1) >> 1;
  double 1 = (p[ls - 1].x + p[ls].x) * 0.5;
  vector<point> xl(ls), xr(p.size() - ls);
  unordered_set<int> left;
  for (int i = 0; i < ls; ++i) {</pre>
    xl[i] = x[i];
   left.insert(x[i].id);
  for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
  vector<point> y1, yr;
  vector<point> pl, pr;
  yl.reserve(ls); yr.reserve(p.size() - ls);
  pl.reserve(ls); pr.reserve(p.size() - ls);
  for (int i = 0; i < p.size(); ++i) {</pre>
    if (left.count(y[i].id)) yl.push_back(y[i]);
    else yr.push_back(y[i]);
    if (left.count(p[i].id)) pl.push_back(p[i]);
    else pr.push_back(p[i]);
  double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
  for (int i = 0; i < p.size(); ++i) {</pre>
    if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
  for (int i = 0; i < yp.size(); ++i) {</pre>
    for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
     d = min(d, dist(yp[i], yp[j]));
  }
double closest_pair(vector<point> &p) {
  vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const point
     &b) {
    return a.x < b.x;</pre>
  vector<point> y(p.begin(), p.end());
  sort(y.begin(), y.end(), [](const point &a, const point
     &b) {
    return a.y < b.y;</pre>
  });
  return cp(p, x, y);
```

2.4 convex hull

```
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
   T x, y;
```

```
PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return</pre>
     make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
  bool operator == (const PT &rhs) const { return
     make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c)
     + cross(c.a): }
#ifdef REMOVE REDUNDANT
bool between(const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x)
     <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2],
     up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2],
     dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--)
     pts.push_back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i]))
     dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back():
    dn.pop_back();
 pts = dn;
#endif
```

2.5 rotating calipers

```
typedef long double gtype;
const gtype pi = M_PI;
typedef complex<gtype> point;
typedef complex<gtype> point;
#define x real()
#define y imag()
#define polar(r, t) polar((gtype) (r), (t))
// vector
```

```
#define rot(v, t) ((v) * polar(1, t))
#define crs(a, b) ( (conj(a) * (b)).y )
#define dot(a, b) ( (conj(a) * (b)).x )
#define pntLinDist(a, b, p) ( abs(crs((b)-(a), (p)-(a)) /
     abs((b)-(a))) )
bool cmp_point(point const& p1, point const& p2) {
   return p1.x == p2.x ? (p1.y < p2.y) : (p1.x < p2.x);
// O(n) - rotating calipers (works on a ccw closed convex
gtype rotatingCalipers(vector<point> &ps) {
   int aI = 0, bI = 0;
   for (size_t i = 1; i < ps.size(); ++i)</pre>
        aI = (ps[i].y < ps[aI].y ? i : aI), bI = (ps[i].y
    > ps[bI].y ? i : bI);
    gtype minWidth = ps[bI].y - ps[aI].y, aAng, bAng;
   point aV = point(1, 0), bV = point(-1, 0);
   for (gtype ang = 0; ang < pi; ang += min(aAng, bAng)) {</pre>
        aAng = acos(dot(ps[aI + 1] - ps[aI], aV)
            / abs(aV) / abs(ps[aI + 1] - ps[aI]));
       bAng = acos(dot(ps[bI + 1] - ps[bI], bV)
            / abs(bV) / abs(ps[bI + 1] - ps[bI]));
        aV = rot(aV, min(aAng, bAng)), bV = rot(bV,
     min(aAng, bAng));
        if (aAng < bAng)</pre>
            minWidth = min(minWidth, pntLinDist(ps[aI],
     ps[aI] + aV, ps[bI]))
            , aI = (aI + 1) \% (ps.size() - 1);
        else
            minWidth = min(minWidth, pntLinDist(ps[bI],
     ps[bI] + bV, ps[aI]))
            , bI = (bI + 1) \% (ps.size() - 1);
   return minWidth;
```

2.6 split convex polygon

```
typedef long double Double;
typedef vector <Point> Polygon;
// This is not standard intersection because it returns
// when the intersection point is exactly the t=1 endpoint
// the segment. This is OK for this algorithm but not for
     general
// use.
bool segment_line_intersection(Double x0, Double y0,
   Double x1, Double y1, Double x2, Double y2,
   Double x3, Double y3, Double &x, Double &y){
   Double t0 = (y3-y2)*(x0-x2) - (x3-x2)*(y0-y2);
   Double t1 = (x1-x0)*(y2-y0) - (y1-y0)*(x2-x0);
   Double det =(y1-y0)*(x3-x2) - (y3-y2)*(x1-x0);
   if (fabs(det) < EPS){ //Paralelas</pre>
       return false;
   }else{
        t0 /= det:
        t1 /= det:
        if (cmp(0, t0) \le 0 \text{ and } cmp(t0, 1) \le 0)
```

```
x = x0 + t0 * (x1-x0);
            y = y0 + t0 * (y1-y0);
            return true;
        return false;
// Returns the polygons that result of cutting the CONVEX
// polygon p by the infinite line that passes through (x0,
     y0)
// and (x1, y1).
// The returned value has either 1 element if this line
// doesn't cut the polygon at all (or barely touches it)
// or 2 elements if the line does split the polygon.
vector<Polygon> split(const Polygon &p, Double x0, Double
                      Double x1, Double y1) {
    int hits = 0, side = 0;
    Double x, y;
    vector<Polygon> ans(2);
    for (int i = 0; i < p.size(); ++i) {</pre>
        int j = (i + 1) % p.size();
        if (segment_line_intersection(p[i].x, p[i].y,
            p[j].x, p[j].y, x0, y0, x1, y1, x, y)) {
            ans[side].push_back(p[i]);
            if (cmp(p[i].x, x) != 0 or cmp(p[i].y, y) !=
     0) {
                ans[side].push_back(Point(x, y));
            side ^= 1;
            ans[side].push_back(Point(x, y));
            ans[side].push_back(p[i]);
    return hits < 2 ? vector<Polygon>(1, p) : ans;
```

2.7 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

 UPJV

3 Graphs

3.1 Dijkstra

#include<stdio.h>

```
#include<stdlib.h>
#define MAX 100
#define INF 1e9
typedef struct {
    int V; // Nombre de sommets
    int adj[MAX][MAX]; // Matrice d'adjacence (pour les
} Graph;
void initGraph(Graph *g, int V) {
    g \rightarrow V = V;
    for(int i = 0; i < V; i++) {</pre>
        for(int j = 0; j < V; j++) {</pre>
            g->adj[i][j] = 0; // ou INF si on prfre
     indiquer qu'il n'y a pas de lien
   }
void addEdge(Graph *g, int src, int dest, int weight) {
    g->adj[src][dest] = weight:
    // g->adj[dest][src] = weight; // Si le graphe est non
     dirig
void dijkstra(Graph *g, int src) {
    int dist[MAX];
    int visited[MAX] = {0};
    for(int i = 0; i < g->V; i++) {
        dist[i] = INF;
    dist[src] = 0;
    for(int i = 0; i < g->V - 1; i++) {
        int u = -1;
        // Trouver le sommet avec la distance minimale,
     parmi les sommets non traits.
        for(int j = 0; j < g->V; j++) {
            if(!visited[j] && (u == -1 || dist[j] <</pre>
     dist[u])) {
                u = j;
        }
        visited[u] = 1;
        for(int v = 0; v < g->V; v++) {
            if(!visited[v] && g->adj[u][v] && dist[u] +
     g->adj[u][v] < dist[v]) {</pre>
                dist[v] = dist[u] + g->adj[u][v];
       }
```

```
// Affichage des distances
for(int i = 0; i < g->V; i++) {
    printf("Distance du sommet %d au sommet %d =
    %d\n", src, i, dist[i]);
}
```

3.2 Recherche en Largeur

```
#include<stdio.h>
#include<stdlib.h>
#define MAX 100
typedef struct {
   int V; // Nombre de sommets
   int adj[MAX][MAX]; // Matrice d'adjacence
void initGraph(Graph *g, int V) {
   g \rightarrow V = V;
   for(int i = 0; i < V; i++) {</pre>
        for(int j = 0; j < V; j++) {</pre>
            g->adj[i][j] = 0;
void addEdge(Graph *g, int src, int dest) {
   g->adj[src][dest] = 1;
   g->adj[dest][src] = 1; // Si le graphe est non dirig
void BFS(Graph *g, int start) {
   int visited[MAX] = {0};
   int queue[MAX], front = -1, rear = -1;
   void enqueue(int v) {
        if(rear == MAX-1) return;
        if(front == -1) front = 0;
        queue[++rear] = v;
   int dequeue() {
        if(front == -1) return -1;
        int v = queue[front];
        if(front == rear) front = rear = -1;
        else front++;
        return v;
   printf("%d ", start);
   visited[start] = 1;
   enqueue(start);
   while(front != -1) {
        int curr = dequeue();
        for(int i = 0; i < g->V; i++) {
            if(g->adj[curr][i] == 1 && !visited[i]) {
               printf("%d ", i);
                visited[i] = 1;
```

```
enqueue(i);
       }
    }
int main() {
    Graph g;
    initGraph(&g, 6); // Cration d'un graphe avec 5
     sommets (0,1,2,3,4)
    addEdge(&g, 0, 1);
    addEdge(&g, 0, 2);
    addEdge(&g, 1, 3);
    addEdge(&g, 5, 5);
    addEdge(&g, 3, 4);
    printf("BFS partir du sommet 0: ");
    BFS(&g, 3);
    return 0;
}
```

3.3 Rechernche en profondeur

```
#include <stdbool.h>
#include <stdio.h>
#define MAX_NODES 1000
bool visited[MAX_NODES];
int graph[MAX_NODES] [MAX_NODES];
void dfs(int node, int n) {
    visited[node] = true;
    printf("Visited node: %d\n", node);
    for (int i = 0; i < n; i++) {</pre>
        if (graph[node][i] && !visited[i]) {
            dfs(i, n);
   }
//Exemple implementation
#include <stdbool.h>
#include <stdio.h>
#define MAX_SIZE 1000
char map[MAX_SIZE][MAX_SIZE];
bool visited[MAX_SIZE][MAX_SIZE];
int n, m;
void dfs(int x, int y) {
    if (x < 0 || x >= n || y < 0 || y >= m) return; //
     Vrifie les limites
    if (map[x][y] == '#' || visited[x][y]) return; //
     Vrifie les murs et les zones visites
```

```
visited[x][y] = true;
    dfs(x + 1, y);
    dfs(x - 1, y);
    dfs(x, y + 1);
    dfs(x, y - 1);
}
int main() {
    scanf("%d %d", &n, &m);
    for (int i = 0; i < n; i++) {</pre>
        for (int j = 0; j < m; j++) {
            scanf(" %c", &map[i][j]);
            visited[i][j] = false;
        }
    }
    int rooms = 0;
    for (int i = 0; i < n; i++) {</pre>
        for (int j = 0; j < m; j++) {
            if (!visited[i][j] && map[i][j] == '.') {
                dfs(i, j);
                rooms++;
            }
        }
    }
    printf("%d\n", rooms);
    return 0:
}
```

3.4 euler formula

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

```
f + v = e + c + 1
```

4 Math

4.1 FFT

```
typedef long double T;
const T pi = acos(-1);
struct cpx {
   T real, image;
    cpx(T _real, T _image) {
        real = real:
        image = _image;
    cpx() {}
};
cpx operator + (const cpx &c1, const cpx &c2) {
    return cpx(c1.real + c2.real, c1.image + c2.image);
cpx operator - (const cpx &c1, const cpx &c2) {
    return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
    return cpx(c1.real * c2.real - c1.image * c2.image ,
     c1.real *c2.image + c1.image * c2.real);
int rev(int id, int len) {
    int ret = 0:
    for (int i = 0; (1 << i) < len; i++) {
        ret <<= 1;
        if (id & (1 << i)) ret |= 1;</pre>
   }
    return ret;
}
void fft(cpx *a, int len, int dir) {
   for (int i = 0: i < len: i++) { A[rev(i, len)] = a[i]:</pre>
   for (int s = 1; (1 << s) <= len; s++) {
        int m = (1 << s):
        cpx wm = cpx(cos(dir * 2 * pi / m), sin(dir * 2 *
     pi / m ));
        for (int k = 0; k < len; k += m) {</pre>
            cpx w = cpx(1, 0);
            for (int j = 0; j < (m >> 1); j++) {
                cpx t = w * A[k + j + (m >> 1)];
                cox u = A[k + j];
                A[k + j] = u + t;
                A[k + j + (m >> 1)] = u - t;
                w = w * wm:
        }
    if (dir == -1) for (int i = 0; i < len; i++) A[i].real</pre>
     /= len, A[i].image /= len;
    for (int i = 0; i < len; i++) a[i] = A[i];</pre>
```

4.2 Proba

```
#include <stdio.h>
double Tab[600][100];
double prob(int n, int s, int a, int b) //s = 6n
{
    if(Tab[s][n] != -1)
```

```
return Tab[s][n];
    }
        if(n==0)
             if( a<=s && s<=b)
                 return 1;
             else
                 return 0;
        }else
        {
             return Tab[s][n] = 1./6. * prob(n-1, s, a, b)
      + 1./6. * prob(n-1, s-1, a, b) + 1./6. * prob(n-1, a)
      s-2, a, b) +
                     1./6. * prob(n-1, s-3, a, b) + 1./6. *
      prob(n-1, s-4, a, b) + 1./6. * prob(n-1, s-5, a, b);
}
int main()
    int a, b, n;
    scanf("%d %d %d", &n, &a, &b);
    for (int i = 0; i<=6*n; i++)</pre>
        for (int y = 0; y<=n; y++)</pre>
            Tab[i][y] = -1;
    printf("\frac{1}{n}", prob(n, 6*n, a, b));
}
```

4.3 fibonacci

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

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4.4 lucas

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:
$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

5 Structure

5.1 treeDiameter

```
#include <stdio.h>
#include <stdlib.h>
#define MAX_N 200000
// Structure to represent an edge
struct Edge {
    int to;
    struct Edge* next;
struct Edge* graph[MAX_N + 1]; // Adjacency list
     representation of the tree
int maxDepth = 0;
// Depth-First Search to find the diameter
int dfs(int node, int parent) {
    int maxDepth1 = 0, maxDepth2 = 0;
    struct Edge* edge = graph[node];
    while (edge != NULL) {
        int neighbor = edge->to;
        if (neighbor != parent) {
            int depth = dfs(neighbor, node);
            if (depth > maxDepth1) {
                maxDepth2 = maxDepth1;
                maxDepth1 = depth;
            } else if (depth > maxDepth2) {
                maxDepth2 = depth;
        edge = edge->next;
    }
    // Update the diameter
```

```
maxDepth = (maxDepth > maxDepth1 + maxDepth2) ?
     maxDepth : maxDepth1 + maxDepth2;
    // Return the maximum depth rooted at this node
    return maxDepth1 + 1;
}
int main() {
    int n:
    scanf("%d", &n);
    // Initialize the graph
    for (int i = 1; i <= n; i++) {</pre>
        graph[i] = NULL;
    // Build the tree
   for (int i = 0; i < n - 1; i++) {</pre>
        int a, b;
        scanf("%d %d", &a, &b);
        struct Edge* edge1 = (struct
     Edge*)malloc(sizeof(struct Edge));
        edge1->to = b;
        edge1->next = graph[a];
        graph[a] = edge1;
        struct Edge* edge2 = (struct
     Edge*)malloc(sizeof(struct Edge));
        edge2->to = a;
        edge2->next = graph[b];
        graph[b] = edge2;
   dfs(1, 0); // Start the DFS from node 1 as the root
   printf("%d\n", maxDepth);
    return 0;
```

6 X - Misc

6.1 equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by x = -b/2a.

$$ax + by = e$$

$$cx + dy = f$$

$$\Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

In general, given an equation Ax = b, the solution to a

variable x_i is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where A'_i is A with the *i*'th column replaced by b.

6.2 Trigonometry

$$\sin(v+w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v+w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$
$$\sin v + \sin w = 2\sin\frac{v+w}{2}\cos\frac{v-w}{2}$$
$$\cos v + \cos w = 2\cos\frac{v+w}{2}\cos\frac{v-w}{2}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$
where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

6.3 Triangles

Side lengths:
$$a,b,c$$

Semiperimeter: $p=\frac{a+b+c}{2}$
Area: $A=\sqrt{p(p-a)(p-b)(p-c)}$
Circumradius: $R=\frac{abc}{4A}$
Inradius: $r=\frac{A}{p}$

Length of median (divides triangle into two equal-area triangles): $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c} \right)^2 \right]}$$
Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$
Law of cosines: $a^2 = b^2 + c^2 - 2bc\cos\alpha$
Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

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6.4 Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F=b^2+d^2-a^2-c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

6.5 Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \arccos(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \operatorname{atan2}(y, x) \end{array}$$

6.6 Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \qquad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \qquad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

6.7 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c-1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

6.8 Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$
$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

6.9 Geometric series

$$r \neq 1$$

$$a + ar + ar^{2} + ar^{3} + \dots + ar^{n-1} = \sum_{k=0}^{n-1} ar^{k} = a\left(\frac{1-r^{n}}{1-r}\right)$$