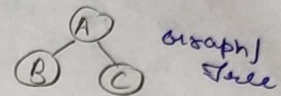


Graphs

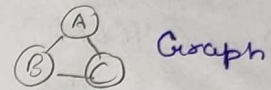
→ set of (V, E) pairs
 $V \rightarrow$ vertices set
 $E \rightarrow$ set of edges

→ vertices → Represented as circles
 also known as nodes.

→ Edges → Represented as lines
 connecting 2 vertices/nodes



AB, C → vertices
 $\{A, B\}$
 $\{A, C\}$ → edges

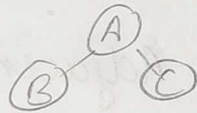


AB
 BC → edges
 AC

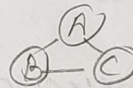
Terminology →

① Adjacent nodes → $A \rightarrow B$ $B \rightarrow C$ which connect same edge

② Degree of Node → No. of edges connected to that node

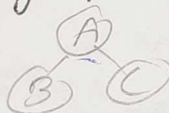


$A \rightarrow 2$
 $C \rightarrow 1$

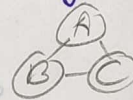


$A \rightarrow 2$
 $B \rightarrow 2$
 $C \rightarrow 2$

③ Size of graph → Total no. of edges in graph



Size = 2



Size → 3

④ path → sequence of vertices from source node to destination node. (vertices & edges not repeated)



$A \rightarrow C \rightarrow$
 $S \quad D$

$A \rightarrow B \rightarrow C$
 or
 $A \rightarrow D \rightarrow C$

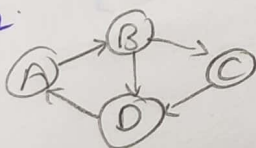
$\langle A, B, C \rangle$

⑤ Trail → a walk in which no edge is repeated

Types of Graphs →

① undirected

directed or digraph

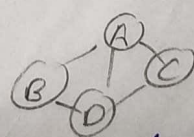


(B, C)

$(A, B) \neq (B, A)$

in 1 direction

undirected



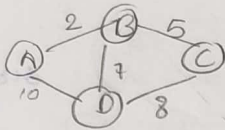
bidirectional
 $\{B, C\}$

$(A, B) = (B, A)$

$(B, C) = (C, B)$

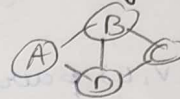
can go in any direction

② weighted

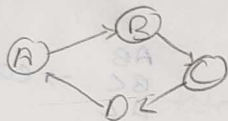


cost/weight is specified for every edge

unweighted



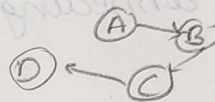
③ CYCLIC



starting node = end node

A → B → C → D → A

ACYCLIC

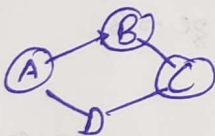


Self loop
A can have edge with itself

↓ can be used to refresh a pan

Representation →

① By using multichm. array → Adjacency Matrix



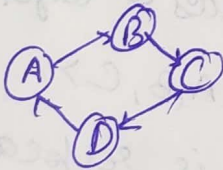
	A	B	C	D
A	0	1	0	1
B	1	0	1	0
C	0	1	0	1
D	1	0	1	0

0 → if edge not

1 → if edge yes

A to A → 0

A to B → 1

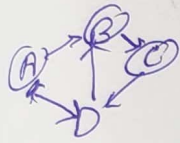


	A	B	C	D
A	0	1	0	0
B	0	0	1	0
C	0	0	0	1
D	1	0	0	0

* V are represented as → {v₁, v₂, v₃}
E as → { {v₁, v₂ }, {v₃, v₄ } }

(A, B) = (B, A)
(A, A) = (A, A)

using list →



→ A is having edge to B, D

A: [B] → [D] X
 B: [A] → [C] X
 C: [B] → [D] X
 D: [A] → [B] → [C] X

A: B, D
 B: A, C
 C: B, D
 D: A, B, C

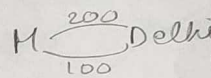
A: [B] → [D] X
 B: [C] X
 C: [D] X
 D: [B] X

A: B, D
 also rep.
 like this

Properties →

① Multiedge / Parallel edge →

like in airlines



could have
diff cost

② No. of edges →

min → 0



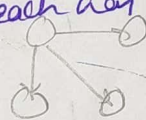
$V = \{v_1, v_2, v_3, v_4\}$

$|V| = 4$

$|E| = \emptyset$

max
in directed

→ Node x edge each way
 $= 4 \times 3$
 $= 12$



if $|V| = n$

then $0 \leq |E| \leq n(n-1)$ (directed)

$0 \leq |E| \leq \frac{n(n-1)}{2}$ (undirected)

Graph is called dense → too many edges
 called sparse → less edges