

## UNIT-3

### ALL PAIR SHORTEST PATH

- ❖ Let  $G = \langle N, A \rangle$  be a directed graph 'N' is a set of nodes and 'A' is the set of edges.
- ❖ Each edge has an associated non-negative length.
- ❖ We want to calculate the length of the shortest path between each pair of nodes.
- ❖ Suppose the nodes of G are numbered from 1 to n, so  $N = \{1, 2, \dots, N\}$ , and suppose G matrix L gives the length of each edge, with  $L(i, j) = 0$  for  $i = 1, 2, \dots, n, L(i, j) = \infty$  for all i & j, and  $L(i, j) = \infty$ , if the edge (i, j) does not exist.
- ❖ The principle of optimality applies: if k is the node on the shortest path from i to j then the part of the path from i to k and the part from k to j must also be optimal, that is shorter.
- ❖ First, create a cost adjacency matrix for the given graph.
- ❖ Copy the above matrix-to-matrix D, which will give the direct distance between nodes.
- ❖ We have to perform N iteration after iteration k. the matrix D will give you the distance between nodes with only (1, 2, ..., k) as intermediate nodes.
- ❖ At the iteration k, we have to check for each pair of nodes (i, j) whether or not there exists a path from i to j passing through node k.

#### COST ADJACENCY MATRIX:

$$D_0 = L = \begin{vmatrix} 0 & 5 & \infty & \infty \\ 50 & 0 & 15 & 5 \\ 30 & \infty & 0 & 15 \\ 15 & \infty & 5 & 0 \end{vmatrix}$$

$$\begin{vmatrix} 1 & 7 & 5 & \infty & \infty \\ 2 & 7 & \infty & \infty & 2 \\ 3 & \infty & 3 & \infty & \infty \\ 4 & 4 & \infty & 1 & \infty \end{vmatrix} \quad \begin{vmatrix} 11 & 12 & - & - \\ 21 & - & - & 24 \\ - & 32 & - & - \\ 41 & - & 43 & - \end{vmatrix}$$

**vertex 1:**

$$\left| \begin{array}{cccc} 7 & 5 & \infty & \infty \\ 7 & \mathbf{12} & \infty & 2 \\ \infty & 3 & \infty & \infty \\ 4 & \mathbf{9} & 1 & \infty \end{array} \right| \left| \begin{array}{cccc} 11 & 12 & - & - \\ 21 & \mathbf{212} & - & 24 \\ - & 32 & - & - \\ 41 & \mathbf{412} & 43 & - \end{array} \right|$$

vertex 2:

$$\left| \begin{array}{cccc} 7 & 5 & \infty & \mathbf{7} \\ 7 & 12 & \infty & 2 \\ \mathbf{10} & 3 & \infty & \mathbf{5} \\ 4 & 9 & 1 & \mathbf{11} \end{array} \right| \left| \begin{array}{cccc} 11 & 12 & - & \mathbf{124} \\ 21 & 212 & - & 24 \\ \mathbf{321} & 32 & - & \mathbf{324} \\ 41 & 412 & 43 & \mathbf{4124} \end{array} \right|$$

vertex 3:

$$\left| \begin{array}{cccc} 7 & 5 & \infty & 7 \\ 7 & 12 & \infty & 2 \\ 10 & 3 & \infty & 5 \\ 4 & \mathbf{4} & 1 & \mathbf{6} \end{array} \right| \left| \begin{array}{cccc} 11 & 12 & - & 124 \\ 21 & 212 & - & 24 \\ 321 & 32 & - & 324 \\ 41 & \mathbf{432} & 43 & \mathbf{4324} \end{array} \right|$$

vertex 4:

$$\left| \begin{array}{cccc} 7 & 5 & \mathbf{8} & 7 \\ \mathbf{6} & \mathbf{6} & \mathbf{3} & 2 \\ \mathbf{9} & 3 & \mathbf{6} & 5 \\ 4 & 4 & 1 & 6 \end{array} \right| \left| \begin{array}{cccc} 11 & 12 & \mathbf{1243} & 124 \\ \mathbf{241} & \mathbf{2432} & \mathbf{243} & 24 \\ \mathbf{3241} & 32 & \mathbf{3243} & 324 \\ 41 & 432 & 43 & 4324 \end{array} \right|$$

❖ At 0<sup>th</sup> iteration it nil give you the direct distances between any 2 nodes

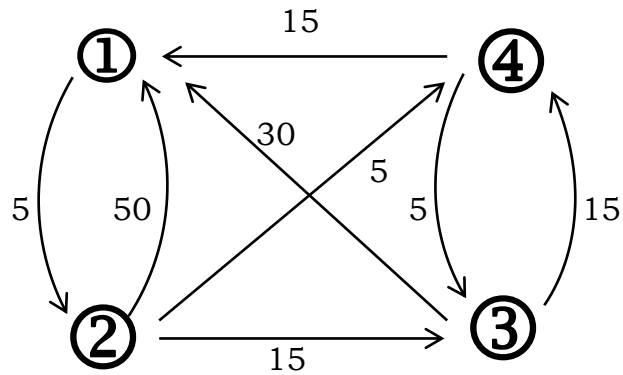
$$D_0 = \left| \begin{array}{cccc} 0 & 5 & \infty & \infty \\ 50 & 0 & 15 & 5 \\ 30 & \infty & 0 & 15 \\ 15 & \infty & 5 & 0 \end{array} \right|$$

❖ At 1<sup>st</sup> iteration we have to check the each pair(i,j) whether there is a path through node 1. if so we have to check whether it is minimum than the previous value and if I is so than the distance through 1 is the value of d1(i,j). at the same time we have to solve the intermediate node in the matrix position p(i,j).

$$D_1 = \left| \begin{array}{cccc} 0 & 5 & \infty & \infty \\ 50 & 0 & 15 & 5 \\ 30 & \mathbf{35} & 0 & 15 \\ 15 & 5 & 5 & 0 \end{array} \right| \quad \begin{array}{l} p[3,2]= 1 \\ p[4,2]= 1 \end{array}$$

2

15 **20** 5 0



**Fig: floyd's algorithm and work**

- ❖ likewise we have to find the value for N iteration (ie) for N nodes.

$$D2 = \begin{vmatrix} 0 & 5 & \mathbf{20} & \mathbf{10} \\ 50 & 0 & 15 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{vmatrix} \quad \begin{matrix} P[1,3] = 2 \\ P[1,4] = 2 \end{matrix}$$

$$D3 = \begin{vmatrix} 0 & 5 & 20 & 10 \\ \mathbf{45} & 0 & 15 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{vmatrix} \quad P[2,1]=3$$

$$D4 = \begin{vmatrix} 0 & 5 & \mathbf{15} & 10 \\ 20 & 0 & \mathbf{10} & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{vmatrix} \quad \begin{matrix} P[1,3]=4 \\ P[2,3]=4 \end{matrix}$$

- ❖ D4 will give the shortest distance between any pair of nodes.
- ❖ If you want the exact path then we have to refer the matrix p. The matrix will be,

$$\begin{vmatrix} 0 & 0 & 4 & 2 \\ 3 & 0 & 4 & 0 \\ & & 3 & \end{vmatrix} \quad 0 \longrightarrow \text{direct path}$$

$$P = \begin{matrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{matrix}$$

- ❖ Since,  $p[1,3]=4$ , the shortest path from 1 to 3 passes through 4.
- ❖ Looking now at  $p[1,4]$  &  $p[4,3]$  we discover that between 1 & 4, we have to go to node 2 but that from 4 to 3 we proceed directly.
- ❖ Finally we see the trips from 1 to 2, & from 2 to 4, are also direct.
- ❖ The shortest path from 1 to 3 is 1,2,4,3.

### ALGORITHM :

Function Floyd ( $L[1..r, 1..r]$ ): array  $[1..n, 1..n]$   
array  $D[1..n, 1..n]$

```
D = L
For k = 1 to n do
For i = 1 to n do
For j = 1 to n do
D [ i , j ] = min (D[ i , j ], D[ i , k ] + D[ k , j ])
Return D
```

### ANALYSIS:

This algorithm takes a time of  $\theta(n^3)$

Additional Resources which are useful to refer:

<https://www.youtube.com/watch?v=4OQeCuLYj-4>

<https://www.youtube.com/watch?v=4NQ3HnhYNfQ>

[https://www.youtube.com/watch?v=nV\\_wOZnhbog](https://www.youtube.com/watch?v=nV_wOZnhbog) [important ]

<https://www.youtube.com/watch?v=DzfmJoFq1pc>

<https://www.youtube.com/watch?v=oNI0rf2P9gE>