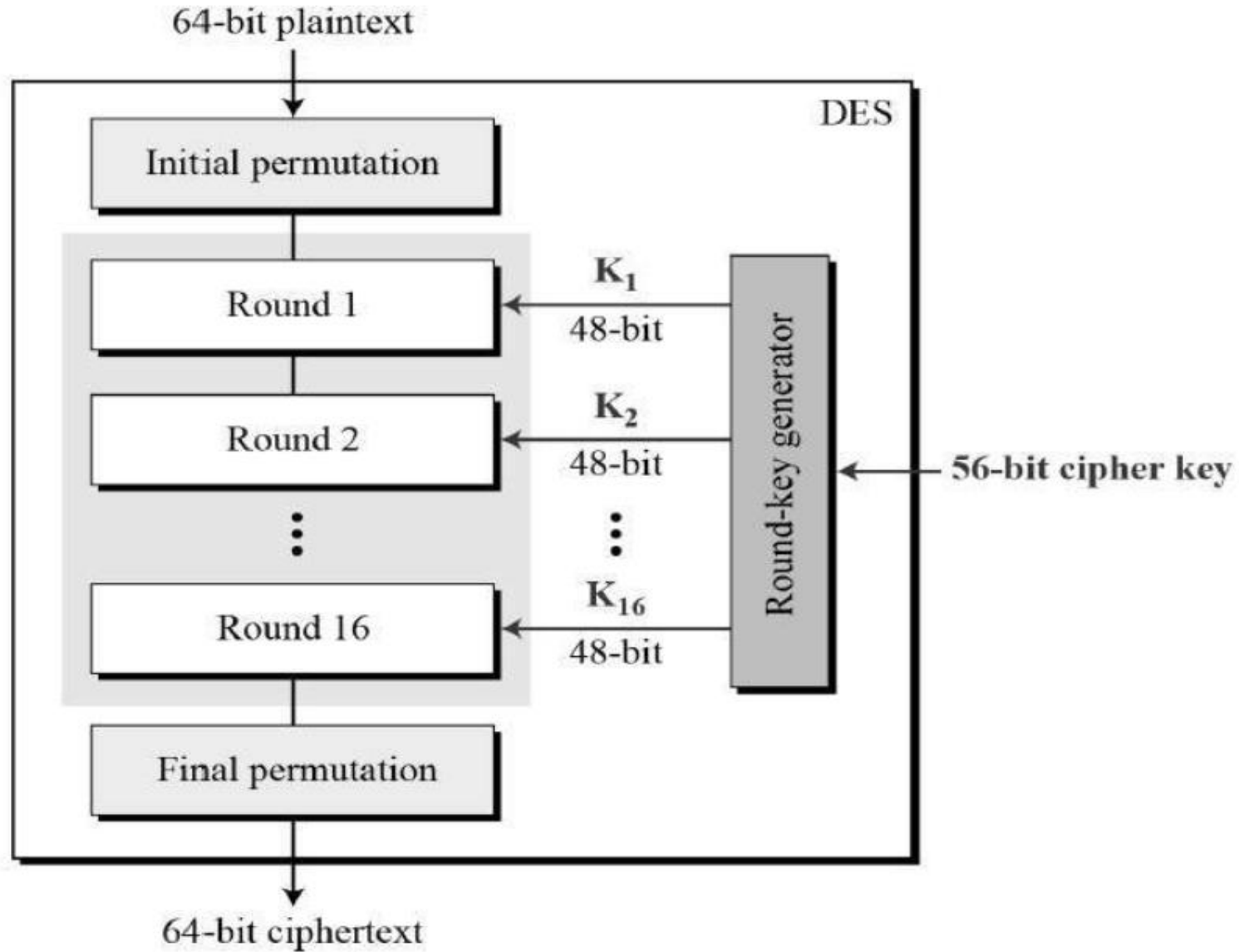
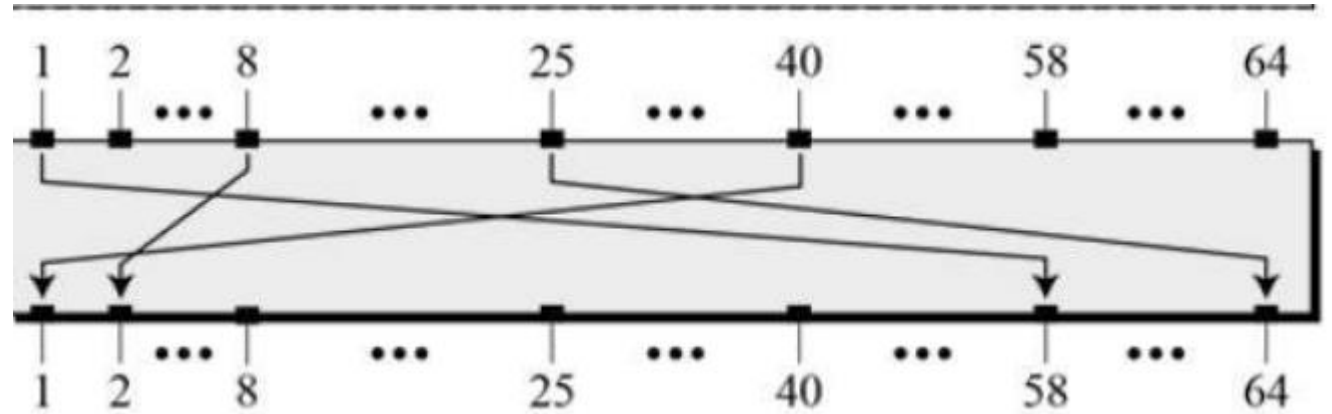


Data Encryption Standard



Initial permutation

1	2	3	4	5	6	7	8
9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24
25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48
49	50	51	52	53	54	55	56
57	58	59	60	61	62	63	64



Initial Permutation							
58	50	42	34	26	18	10	02
60	52	44	36	28	20	12	04
62	54	46	38	30	22	14	06
64	56	48	40	32	24	16	08
57	49	41	33	25	17	09	01
59	51	43	35	27	19	11	03
61	53	45	37	29	21	13	05
63	55	47	39	31	23	15	07

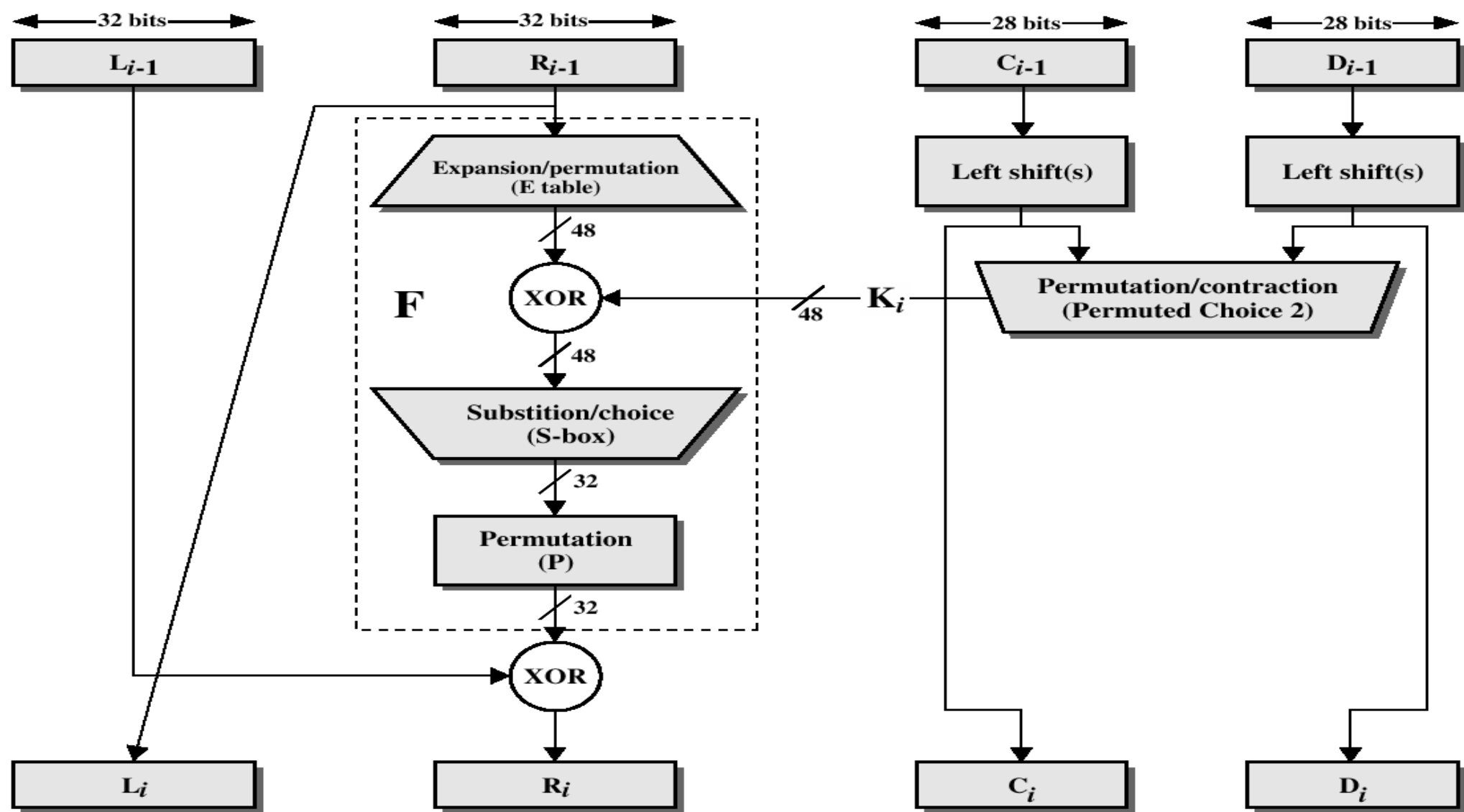
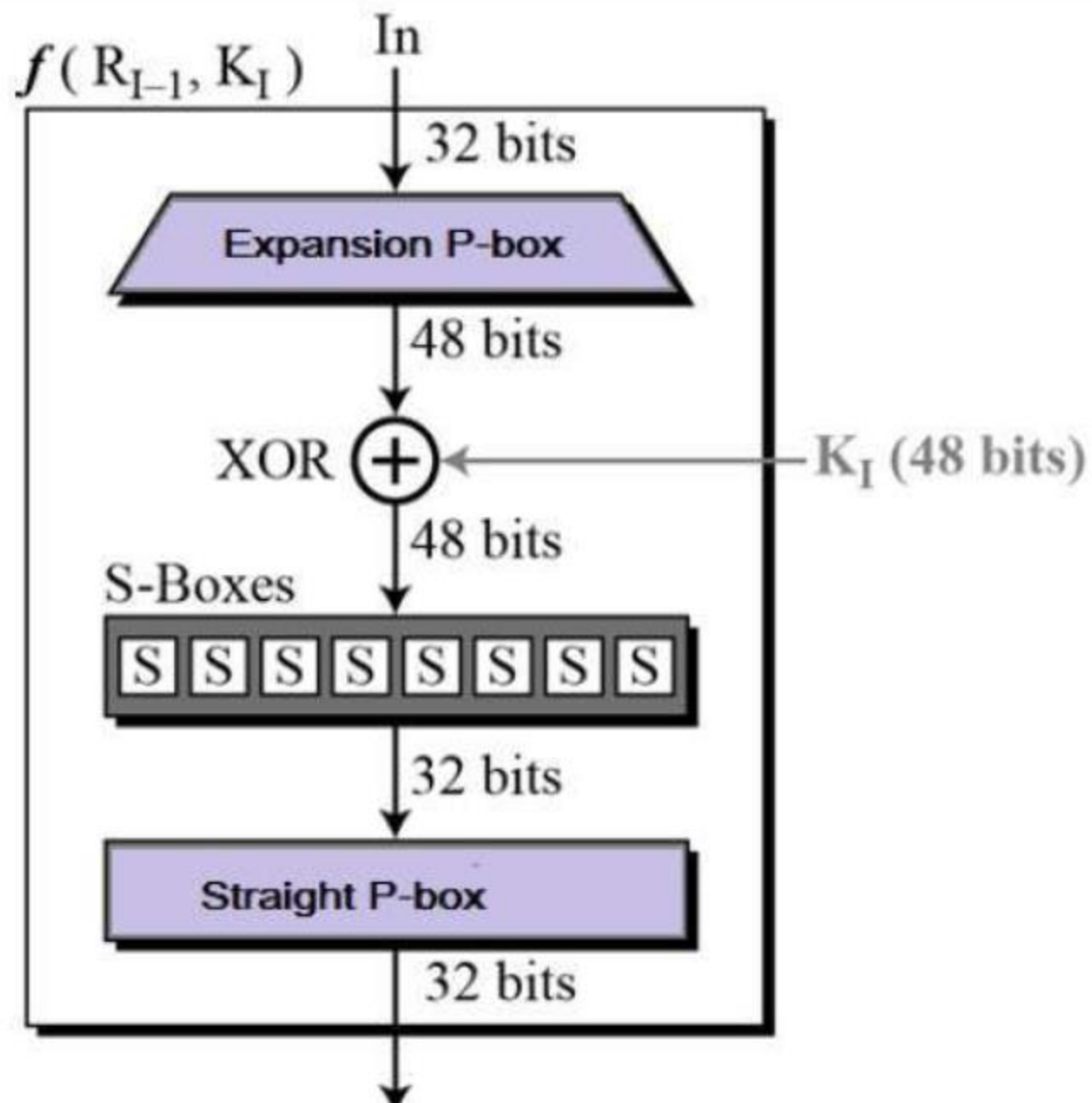
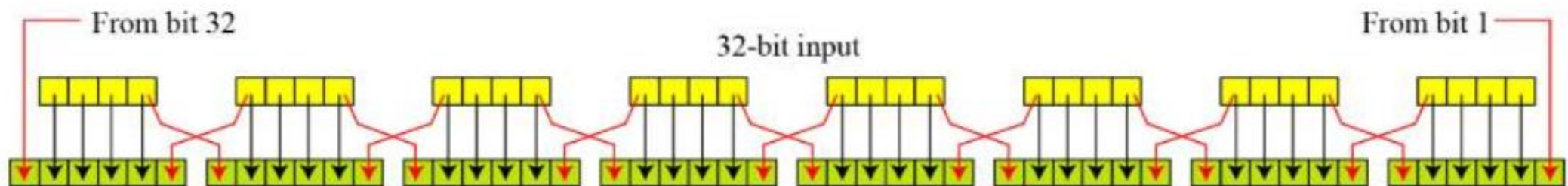
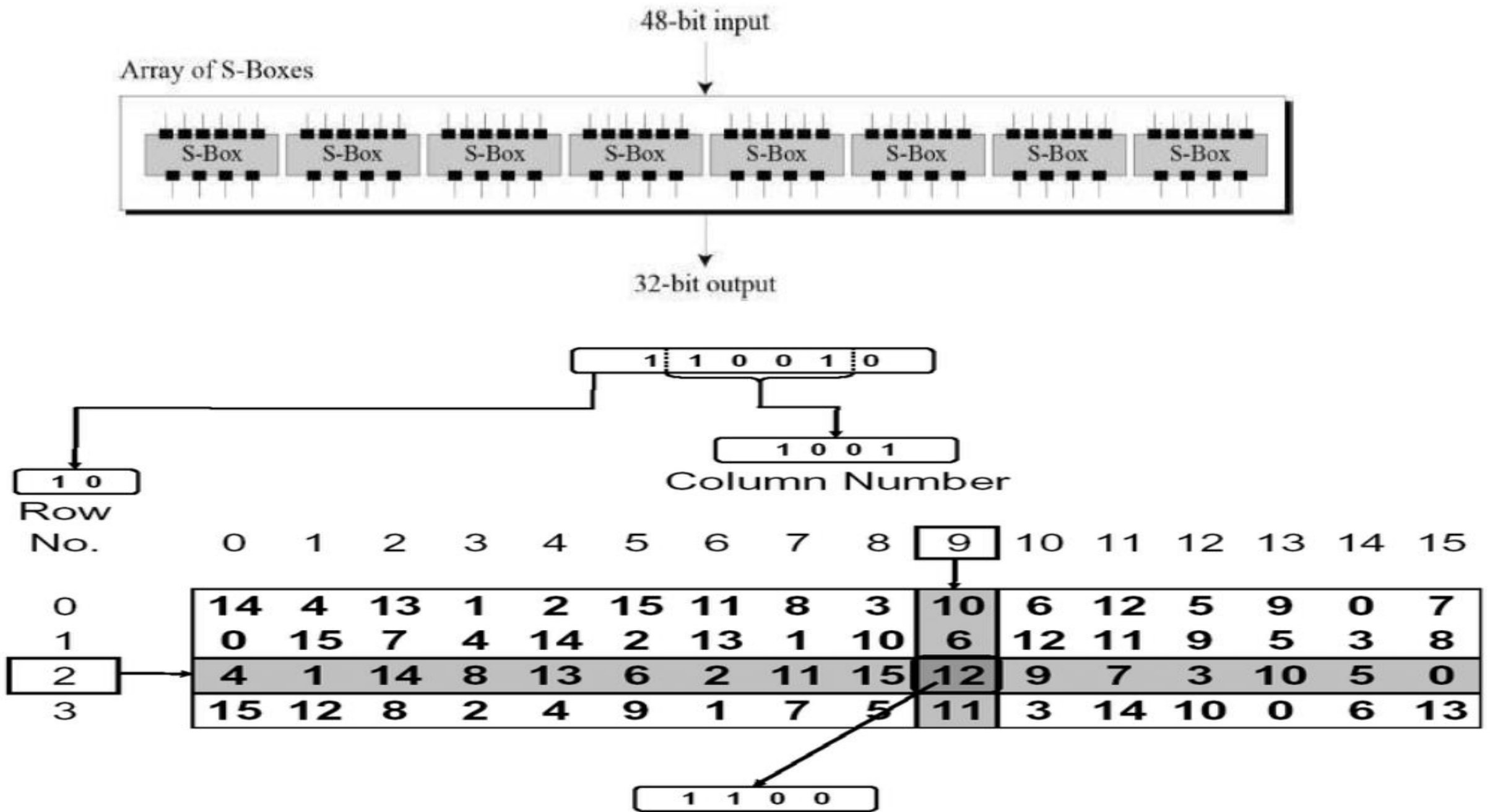


Figure 2.4 Single Round of DES Algorithm



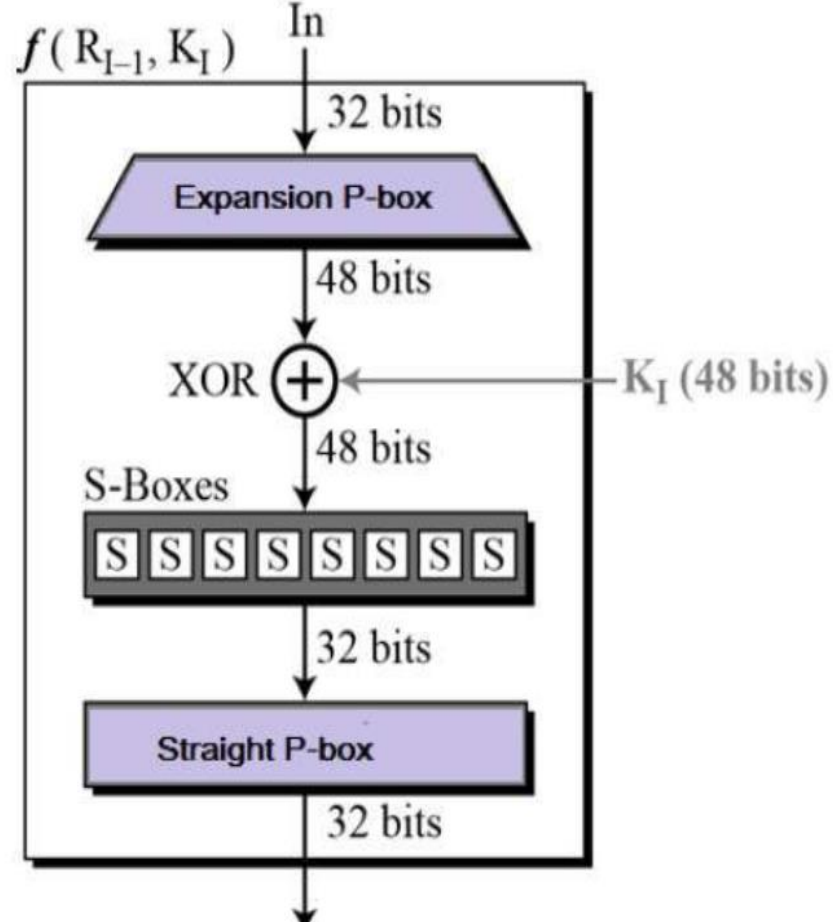
Expansion permutation





Straight Permutation

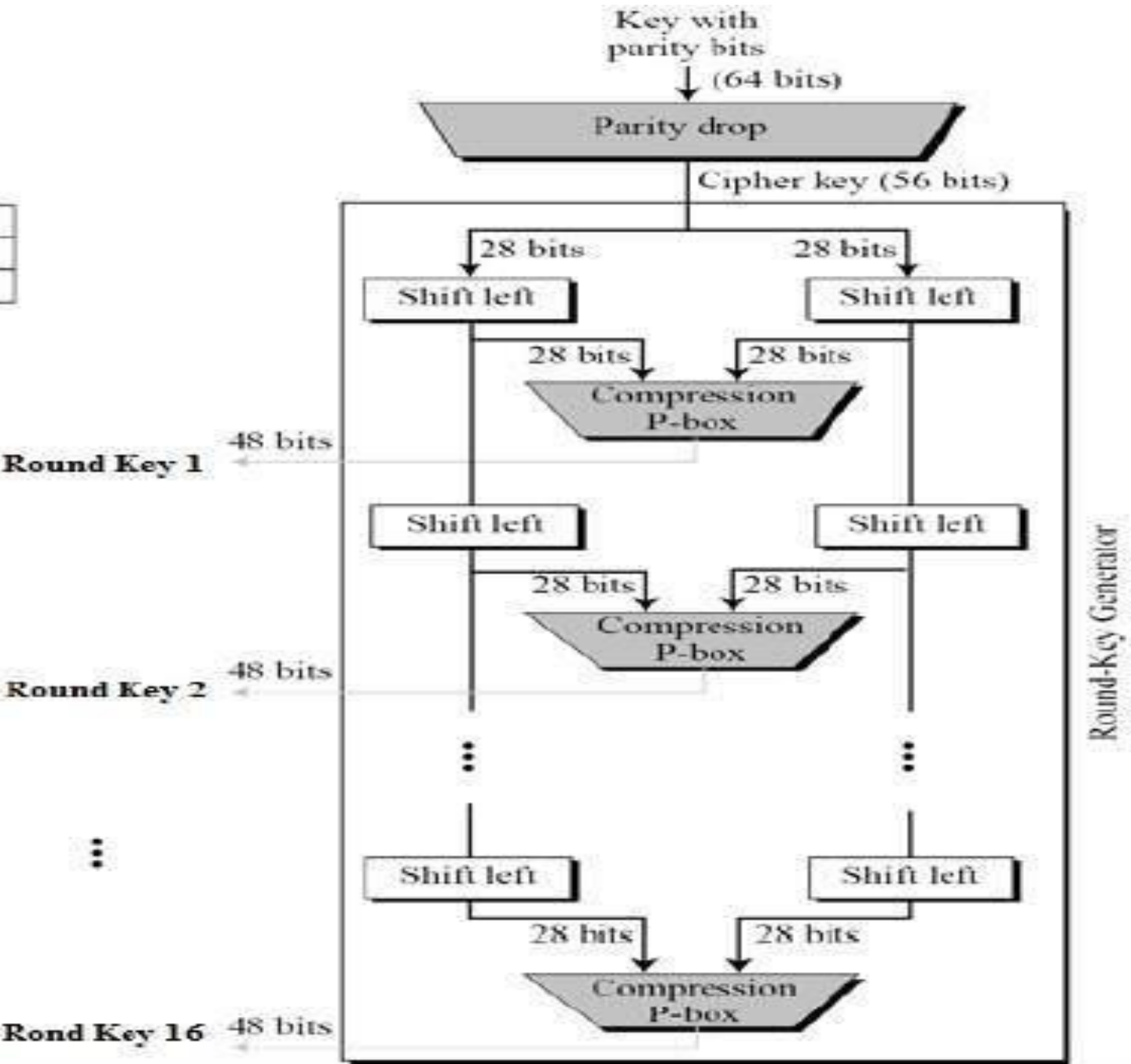
16	07	20	21	29	12	28	17
01	15	23	26	05	18	31	10
02	08	24	14	32	27	03	09
19	13	30	06	22	11	04	25



Key Generation

Shifting

Rounds	Shift
1, 2, 9, 16	one bit
Others	two bits



DES Analysis

The DES satisfies both the desired properties of block cipher. These two properties make cipher very strong.

Avalanche effect – A small change in plaintext results in a very great change in the ciphertext.

Completeness – Each bit of ciphertext depends on many bits of plaintext.

Double DES Encryption

Given a plaintext P and two encryption keys K_1 and K_2 , a cipher text can be generated as,

$$C = E(K_2, E(K_1, P))$$

Decryption requires that the keys be applied in reverse order,

$$P = D(K_1, D(K_2, C))$$

