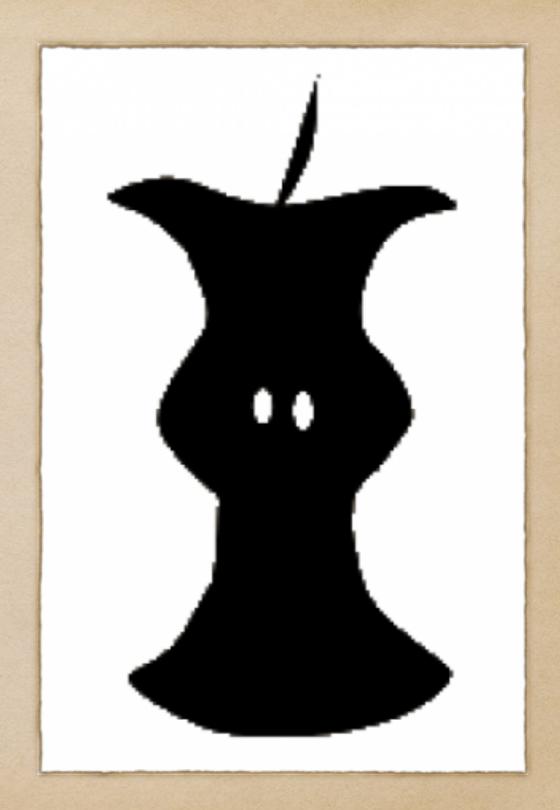


CAUTION

THIS IS SPARTA

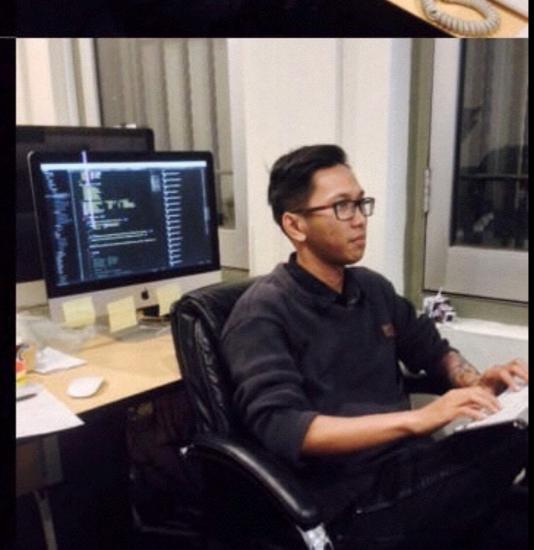
Алексей Охрименко

Twitter: @Ai_boy



{}
BACKEND
DEVELOPER

</>
FRONTEND
DEVELOPER







ПРОСТИ, ЧЕБУРАШКА ты родился в Спарте





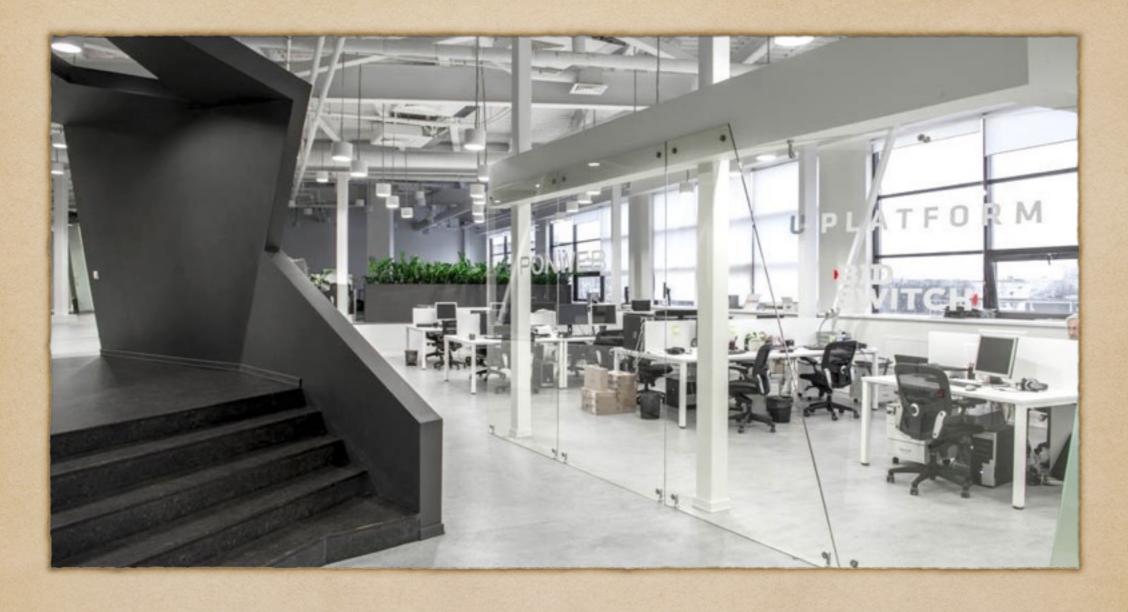




Lionidus & HIGHLOAD 2015



Tumyp infosystems jet



IPONWEB

RTB



BH [070BH?

```
<input type="email">
<input type="email" multiple>
```

Live Demo

" "@[IPv6:2001:db8::1]

user@gmail.com

Отправить



Firefox 4+





Safari 5+



Safari Mobile iOS 3.1+



Chrome 10+





Opera 10.6+





1E 10+



value = <u>e-mail address</u>

A single e-mail address.

Value: Any string that matches the following [ABNF] production:

```
1*( atext / "." ) "@" ldh-str 1*( "." ldh-str )
```

...where atext is as defined in [RFC 5322], and Idh-str is as defined in [RFC 1034].

That is, any string which matches the following regular expression:

```
/^[a-zA-Z0-9.!#$%&'*+/=?^_`{|}~-]+@[a-zA-Z0-9-]+(?:\.[a-zA-Z0-9-]+)*$/
```

Examples:

foo-bar.baz@example.com

RFC 822, 2822, 5322

- https://tools.ietf.org/html/rfc822 (1982 год)
- https://tools.ietf.org/html/rfc2822 (2001 год)
- https://tools.ietf.org/html/rfc5322 (2008 год)

(?:(?:\r\n)?[\t])*(?:(?:(?:[^()<>ê,;;\\".\[\] \000-\031]+(?:(?:(?:\r\n)?[\t])+|\z|(?=[\["()<\e,;:\\".\[\]]))|"(?:[^\"\r\\]|\\.|(?:(?:\r\n)?[\t]))*"(?:(?: \r\n)?[\t])*)(?:\.(?:(?:\r\n)?[\t])*(?:[^()<>0,;:\\".\[\] \000=\031]+(?:(?:(?:\r\n)?[\t])+\\&|(?=[\["()\<\e,;:\\".\[\]]))|"(?:[^\"\\]|\\.|(?:(?:\r\n)?[\t]))*"(?:(?:\r\n)?[\t])*))*ë(?:(?:\r\n)?[\t])*(?:["()<>ê,;:\\".\[\] \000-\0 31]+(?:(?:(?:\r\n)?[\t])+|\z|(?=[\["()⇔ê,;:\\".\[\]]))|\[([^\[\]\r\\]|\\.)*\](?:(?:\r\n)?[\t])*)(?:\.(?:(?:\r\n)?[\t])*(?:[^()<>0,;:\\".\[\] \000-\031]+ (?:(?:\r\n)?[\t])+\\\\z\(?=[\["()\<\;.\\\\])))\\(([^\[\]\\\))\\\)(?: (7:\r\n)?[\t])*))*|(7:[^()<>ê,;:\\".\[\]\000-\031]+(7:(7:\r\n)?[\t])+\Z [(?=[\["()<=ê,;:\\".\[\]]))|"(?:[^\"\r\\]|\\.|(?:(?:\r\n)?[\t]))*"(?:(?:\r\n) ?[\t])*)*\<(?:(?:\r\n)?[\t])*(?:@(?:[^()<>@,;:\\".\[\] \000-\031]+(?:(?:(?:\ r\n)?[\t])+\%[(?=[\["()<>0,;:\\".\[\]]))|\[([^\[\]\r\\]|\\.)*\](?:(?:\r\n)?[\t])*)(?:\.(?:\r\n)?[\t])*(?:[^()<>0,;:\\".\[\]\000-\031]+(?:(?:\r\n) ?[\t])+\\Z|(?=[\["()<0,;:\\".\[\]]))\\[([*\[\]\x\\]|\\.)*\](?:(?:\x\n)?[\t])*))*(?:,@(?:\r\n)?[\t])*(?:[^()<>@,;:\\".\[\] \000-\031]+(?:(?:\r\n)?[\t|)+|\Z|(?=[\["()<>0,;:\\".\[\]]))|\[([^\[\]\r\\]|\\.)*\](?:\r\n)?[\t])*)(?:\.(?:(?:\r\n)?[\t])*(?:[*()<>ê,;:\\".\[\] \000-\031]+(?:(?:(?:\r\n)?[\t])+\\Z|(?=[\["()<>ê,;;\\".\[\]]))\\[([^\[\]\x\\]|\\.)*\](?:(?:\x\n)?[\t])*))*) *:(?:(?:\r\n)?[\t])*)?(?:[^()<>0,;:\\".\[\] \000-\031]+(?:(?:\r\n)?[\t])+ \n)?[\t])*)(?:\.(?:(?:\r\n)?[\t])*(?:[^()<\e,;:\\".\[\] \000-\031]+(?:(?:(?: \r\n)?[\t])+\\Z|(?=[\["()\0,;:\\".\[\]]))|"(?:[^\"\r\\]|\\.|(?:(?:\r\n)?[\t]))*"(?:(?:\r\n)?[\t])*000-\031]+(?:(?:(?:\r\n)?[\t])+\Z|(?=[\["()<>0,;:\\".\[\]]))|\[([^\[\]\r\\]|\\.)*\](?:(?:\r\n)?[\t])*)(?:\.(?:(?:\r\n)?[\t])*(?:[^() \cdot \eta,;:\\".\[\] \000-\031]+(? :(?:(?:\r\n)?[\t])+\\z|(?=[\["()<0,;:\\".\[\]]))\\(([^\(\]\r\\]|\\.)*\](?:(? |\r\n)?[\t])*\>(?:(?:\r\n)?[\t])*)|(?:[^()<=,::\\".\[\]\000-\031]+(?:(? :(?:\r\n)?[\t])+\\Z|(?=[\["()<>0,::\\".\[\]]))|"(?:[^\"\r\\]|\\.|(?:(?:\r\n)? 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(?:[^()⇔ê,;:\\".\[\] \000-\031]+(?:(?:\r\n)?[\t])+\\Z|(?=[\["()⇔ê,;:\\". \[\]]))|"(?:[^\"\r\\]|\\.|(?:(?:\r\n)?[\t]))*"(?:(?:\r\n)?[\t])*)(?:\.(?:(?: \r\n)?[\t])*(?:[*()<0,;:\\".\[\] \000-\031]+(?:(?:\r\n)?[\t])+\\Z|(?=[\["()<>#,;:\\".\[\]]))|"(?:[^\"\r\\]|\\.|(?:\r\n)?[\t]))+"(?:(?:\r\n)?[\t]) *))*@(?:(?:\r\n)?[\t])*(?:[^()<>@,;:\\".\[\] \000-\031]+(?:(?:(?:\r\n)?[\t]) +|\Z|(?=[\["()<>ê,;:\\",\[\]))|\[([*\[\]\r\\]|\\.)*\](?:(?:\r\n)?[\t])*)(?:\ .(?:(?:\r\n)?[\t])*(?:[^()<>ê,;:\\".\[\] \000-\031]+(?:(?:\r\n)?[\t])+|\Z |(?=[\["()<>0,;:\\".\[\]]))|\[([^\[\]\\:)|\\.)+\](?:(?:\r\n)?[\t])+\)+\>(?:(?:\r\n)?[\t])*))*)?;\s*)

http://www.ex-parrot.com/ ~pdw/Mail-RFC822-Address.html Acknowledgements

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Regular expression Denial of Service - ReDoS

This is an Attack. To view all attacks, please see the Attack Category page.

Last revision (mm/dd/yy): 11/9/2015

Introduction

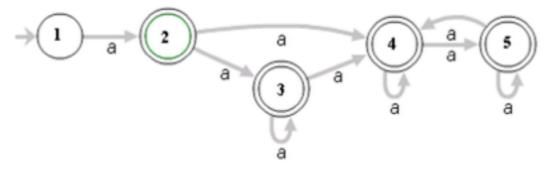
The Regular expression Denial of Service (ReDoS) is a Denial of Service attack, that exploits the fact that most Regular Expression is work very slowly (exponentially related to input size). An attacker can then cause a program using a Regular Expression to enter these

Description

The problematic Regex naïve algorithm

The Regular Expression naïve algorithm builds a Nondeterministic Finite Automaton (NFA) , which is a finite state machine where for expression to the state of th states. Then the engine starts to make transition until the end of the input. Since there may be several possible next states, a determinis paths (if needed) until a match is found (or all the paths are tried and fail).

For example, the Regex $^{4}(a+)+$$ is represented by the following NFA:



For the input aaaaX there are 16 possible paths in the above graph. But for aaaaaaaaaaaaaaX there are 65536 possible paths, and to where the naïve algorithm is problematic, because it must pass on many many paths, and then fail.

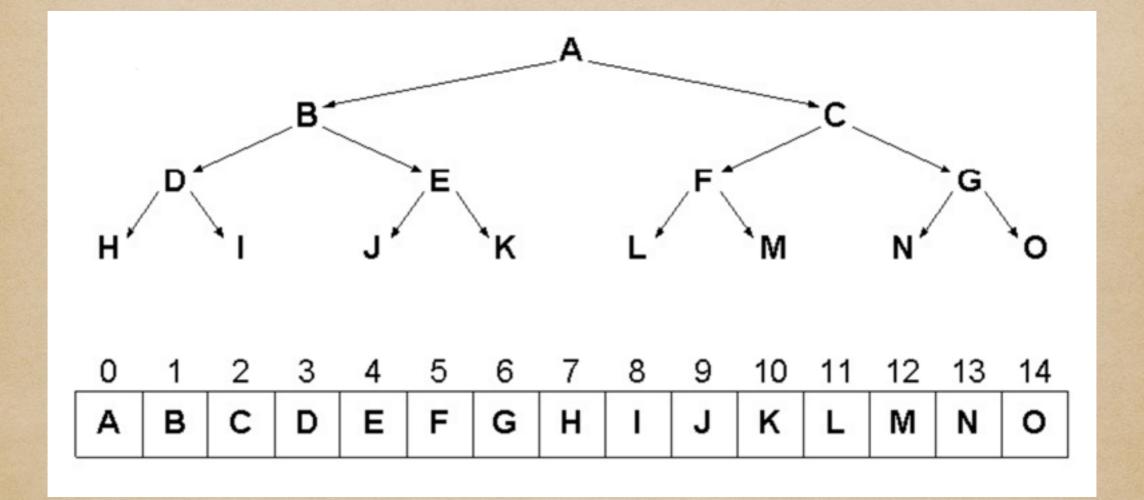
Notice, that not all algorithms are naïve, and actually Regex algorithms can be written in an efficient way. Unfortunately, most Regex en Regexes with "special additions", such as back-references that cannot be always be solved efficiently (see Patterns for non-regular la

DEMO TIME

Parsing or syntactic analysis is the process of analysing a string of symbols, either in natural language or in computer languages, conforming to the rules of a formal grammar

В основе любой сложной задачи стоит неправильно заданный вопрос.

2 3 4 5 6 7 8 9 10 11 12 13 14 F ВС Ε Н Α G K D J M Ν 0



Что нам нужно?

- Парсер
- Грамматика
- Абстрактная формальная грамматика
- Система описания синтаксиса

CODE TIME

Что дальше?

PEG.js

Parser Generator for JavaScript

Home

Online Version

Documentation

Development

PEG.js is a simple parser generator for JavaScript that produces fast parsers with excellent error reporting. You can use it to process complex data or computer languages and build transformers, interpreters, compilers and other tools easily.

Features

- Simple and expressive grammar syntax
- Integrates both lexical and syntactical analysis
- Parsers have excellent error reporting out of the box
- Based on <u>parsing expression grammar</u> formalism more powerful than traditional LL(k) and LR(k) parsers
- Usable <u>from your browser</u>, from the command line, or via JavaScript API

Try PEG.js online

— or —

npm install pegjs

— or —

bower install pegjs

— or —

Download browser version

- PEG.js minified
- PEG.js development

CANOPY

a parser compiler

JAVA JAVASCRIPT PYTHON RUBY

GRAMMAR SYNTAX
MATCHING STRINGS
CHARACTER CLASSES
OPTIONAL NODES
REPEATED NODES
SEQUENCES
LOOKAHEADS
ORDERED CHOICES
CROSS-REFERENCES
BUILDING PARSE TREES

GITHUB

Canopy, a parser compiler

Canopy is a **PEG** parser compiler. It lets you describe the grammar of the language you're trying to parse using a simple, terse syntax, and it generates a parser for the language from this definition.

You can install the command-line tools through npm:

\$ npm install -g canopy

Canopy can generate parsers in the following languages:

- Java
- JavaScript
- Python
- Ruby

Compilers Principles, Techniques, & Tools Second Edition Alfred V. Aho Monica S. Lam Ravi Sethi Jeffrey D. Ullman



docs demos try install community

Calculator demo

This demo parses mathematical expressions and returns the answer, keeping the correct order of operations.

Enter an expression to evaluate, such as PI*4^2 + 5:

PI*4*2 + 5 equals

The grammar

This Jison grammar was used to create the parser/evaluator:

```
/* description: Parses end evaluates mathematical expressions. */
/* lexical grammar */
%lex
%%
\s+ {/* skip whitespace */}
[0-9]+("."[0-9]+)?\b {return 'NUMBER';}
```

Eslint

JSCS

CSSO

https://github.com/rdio/jsfmt

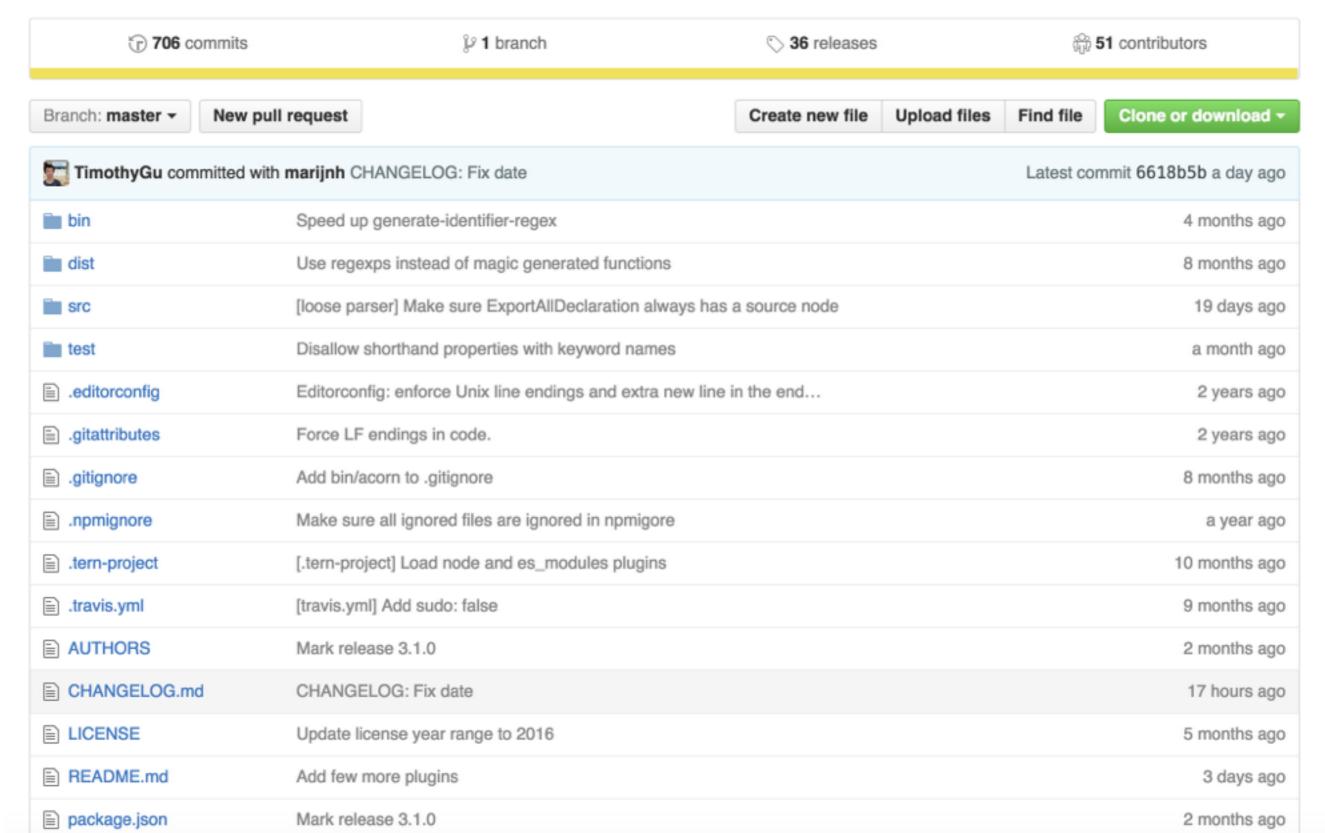
Example

```
var jsfmt = require('jsfmt');
var fs = require('fs');

var js = fs.readFileSync('each.js');

js = jsfmt.rewrite(js, "_.each(a, b) -> a.forEach(b)");
```





ECMAScript parsing infrastructure for multipurpose analysis

Esprima is a high performance, standard-compliant ECMAScript parser written in ECMAScript (also popularly known as JavaScript).

Features

Esprima

- Full support for ECMAScript 6 (ECMA-262)
- Sensible syntax tree format as standardized by EStree project
- Optional tracking of syntax node location (index-based and line-column)
- Heavily tested (~1200 tests with full code coverage)

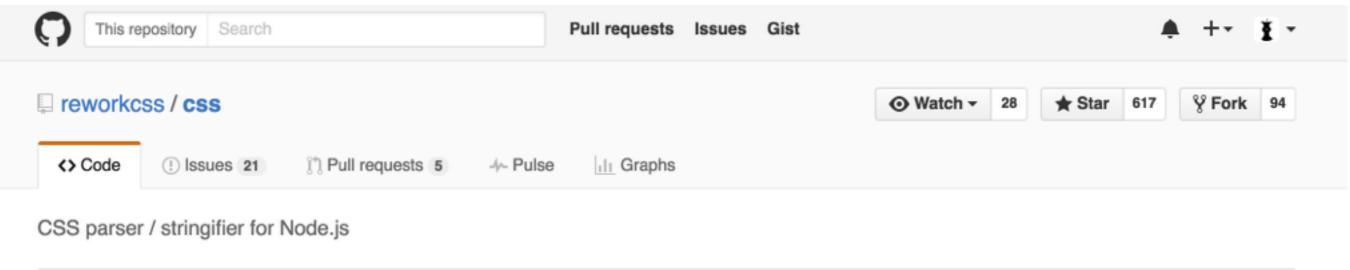
Esprima serves as an important **building block** for some JavaScript language tools, from code instrumentation to editor autocompletion.

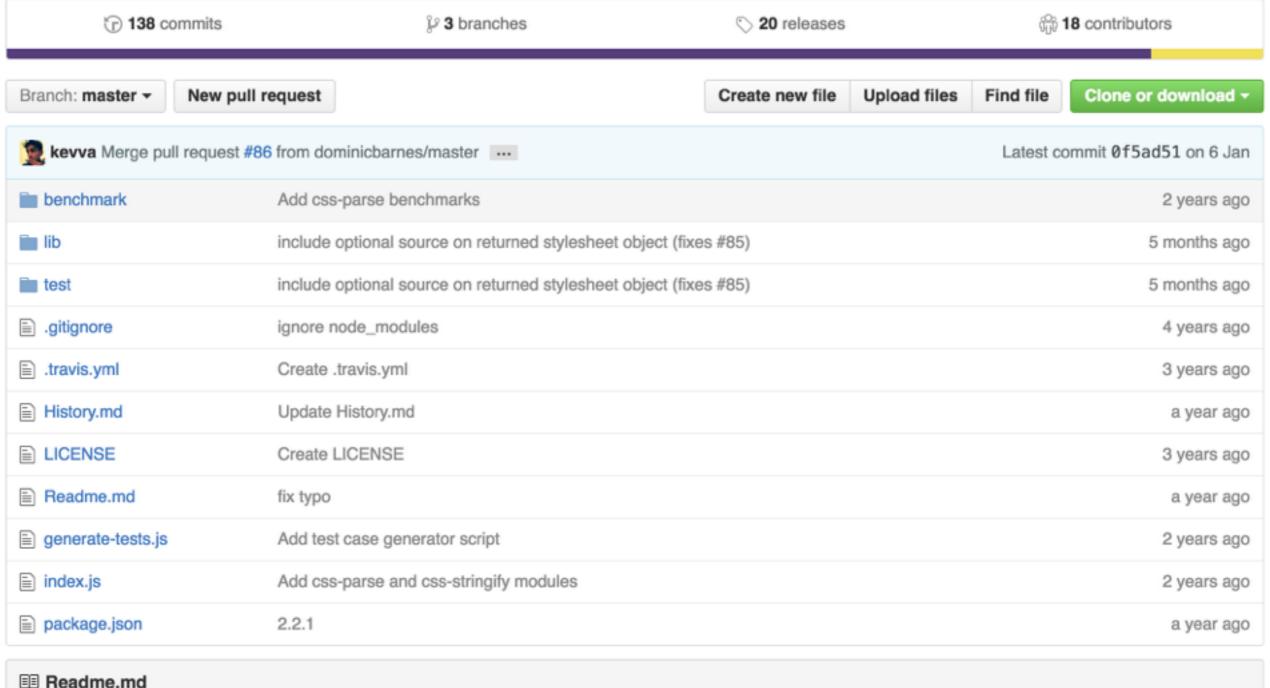
```
var capitalDb = {
2
       Indonesia: 'Jakarta',
3
       Germany: 'Berlin',
       Norway: 'Oslo'
4
5
   };
6
   // Property completion: "capitalDb." and press Ctrl+Space.
8
   capitalDb.
9
              Germany : String
              Indonesia : String
              Norway : String
              hasOwnProperty(property) : Boolean
              isPrototypeOf(object) : Boolean
              propertyIsEnumerable(property) : Boolean
              toLocaleString() : String
              toString() : String
```

Once the full syntax tree is obtained, various static code analysis can be applied to give an insight to the code: syntax visualization, code validation, editing autocomplete (with type inferencing) and many others.

Regenerating the code from the syntax tree permits a few different types of **code transformation**, from a simple rewriting (with specific formatting) to a more complicated minification.

Esprima runs on many popular web browsers, as well as other ECMAScript platforms such as Rhino, Nashorn, and Node.js. It is distributed under the BSD license.







Q&A?

http://bit.ly/1U1Fpa8

Алексей Охрименко



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