Optimal Demand Shut-offs of AC Microgrid using AO-SBQP Method

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Abstract: Microgrids are increasingly being utilized to improve the resilience and operational flexibility of power grids, and act as a backup power source during grid outages. However, it necessitates that the microgrid itself could provide power to the critical loads. This paper presents an algorithm named alternating optimization based sequential boolean quadratic programming tailored for solving optimal demand shut-offs problems arising in microgrids. Moreover, we establish local superlinear convergence of the proposed approximate Boolean quadratic programming method over nonconvex problems. In the end, the performance of the proposed method is illustrated on the modified IEEE 30-bus case study.

Keywords: Demand Shut-offs, Nonconvex, Sequential Boolean Quadratic Programming

1. INTRODUCTION

Microgrids are small-scale power grids that can operate autonomously or in collaboration with other small power grids, which are increasingly being utilized to increase the resilience and operational flexibility of power grids. They act as a backup power supply in the case of grid outages caused by devastating disasters. Abbey et al. (2014) summarized the measures taken by Sendai region in Japan to cope with the shortage of circuit supply caused by the nuclear power accident after the 2011 East Japan Earthquake. Similarly, Panora et al. (2014) describes a successful case of rapid restoration of local microgrid integrated energy systems after infrastructure damage caused by Superstorm Sandy in Manhattan Island, 2012. However, this necessitates that each isolated microgrid itself be resistant and formulate its own optimal power supply strategy with the shortage of energy supply, which is still a challenging problem. Basically, a potential solution is to abstract the above power grid optimization problem in a mixed Boolean nonlinear programming fashion (MBNLP).

To the best of our knowledge, classical optimization in power grid consists optimal power flow (OPF, Frank et al. (2012)), optimal reactive power dispatch (Zhu (2015)), power system state and parameter estimation (Monticelli (1999); Du et al. (2019)). Recently, Du et al. (2020, 2022) proposes the method of optimal experimental design (OED) in order to extract more relative information for assisting the admittance estimation process. Moreover,

Recently, a quadratic programming with linear complementary constraint (LCQP) problem is well studied by a series of literatures (Hall et al. (2021, 2022)). By sequentially solving the QP problem with the corresponding linearlized penalty term, the complementary constraint can be reached with finite steps. Inspired by the above literatures, alternating optimization based sequential boolean quadratic programming (AO-SBQP) (Zhu and Du (2022)) is proposed to set up a bridge between LCQP and MBNLP, leading the solution of optimization problems with Boolean variables no longer rely on B&B

Du et al. (2021) offers an adaptive method for balancing OED and the OPF cost. These mentioned power grid optimization problems are smooth and can be solved directly with interior point method, while notable recent researches optimize the discrete decision variables at the same time. Rhodes et al. (2020) and Kody et al. (2022) modeled the direct current (DC) optimal power shut-in problems in a mixed integer linear program (MILP) framework and solved them with Gurobi. However, only few literature can be found for nonconvex MBNLP in the area of alternating current (AC) power systems. On the other hand, from the algorithmic level, the solver is based on branch and bound (B&B) method (Morrison et al. (2016)), which needs to establish a tree storage structure to explore each integer variable with low efficiency. Luo et al. (2010) solves Boolean optimization in a semi-definite relaxation (SDR) fashion, however, with matrix variables. Solving nonconvex MBNLP accurately and efficiently remains an open problem in our view.

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method with tree structure searching or SDR with matrix variables.

In this paper, the idea of AO-SBQP method is inherited, and the algorithm is modified in the specific optimal distribution of limited supply scenario in microgrid. Unlike Zhu and Du (2022), our considered model is still a nonconvex problem except the Boolean variable constraints. Moreover, a local convergence analysis for the approximate BQP with constraints is introduced.

The rest of this paper is organized as follows: Section 2 reviews the basics concepts of power grid. Sections 3 proposes the optimal distribution of limited supply model. Section 4 describes the AO-SBQP algorithm in detail. And the numerical result is shown in Section 5.

Notation: For $a \in \mathbb{R}^n$ and $C \subseteq \{1, ..., n\}$, $[a_i]_{i \in C} \in \mathbb{R}^{|C|}$ collects all components of a whose index i is in C. Similarly, for $A \in \mathbb{R}^{n \times l}$ and $S \subseteq \{1, ..., n\} \times \{1, ..., l\}$, $[A_{i,j}]_{(i,j) \in S} \in \mathbb{R}^{|S|}$ denotes the concatenation of $A_{i,j}$ for all $(i,j) \in S$. $\mathfrak{i} = \sqrt{-1}$ denotes the imaginary unit, such that $\operatorname{Re}(z) + \mathfrak{i} \cdot \operatorname{Im}(z) = z \in \mathbb{C}$, and \hat{a} denotes the estimated value of a. Moreover, $\mathbf{1}$ represents a vector with all entries being one.

2. AC POWER GRID MODEL

Consider a power grid defined by the triple $(\mathcal{N}, \mathcal{L}, (G + iB))$, where $\mathcal{N} = \{1, 2, \dots N\}$ represents the set of buses, $\mathcal{L} \subseteq \mathcal{N} \times \mathcal{N}$ denotes transmission lines and $(G + iB) \in \mathbb{C}^{N \times N}$ is the complex and potentially sparse Laplacian admittance matrix

$$(G_{k,l} + \mathfrak{i}B_{k,l}) \doteq \begin{cases} \sum\limits_{i \neq k} (g_{k,i} + \mathfrak{i}\,b_{k,i}) & \text{if } k = l, \\ -(g_{k,l} + \mathfrak{i}\,b_{k,l}) & \text{if } k \neq l. \end{cases}$$

Here, $g_{k,l}$ and $b_{k,l}$ are the conductances and susceptances of the transmission line $(k,l) \in \mathcal{L}$, which aims to connect the buses. Note that $(G_{k,l} + \mathrm{i} B_{k,l}) = 0$ if $(k,l) \notin \mathcal{L}$. The set $\mathcal{G} \subseteq \mathcal{N}$ collects all nodes equipped with generators and $\mathcal{D} \subseteq \mathcal{N}$ collects all nodes with power demands. Figure 1 shows a 5-bus system with $\mathcal{N} = \{1, \ldots, 5\}, \mathcal{G} = \{1, 3, 4, 5\}, \mathcal{D} = \{2, 3, 4\}.$

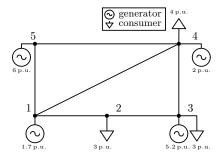


Fig. 1. Modified IEEE 5-bus system from Li and Bo (2010) with 4 generators and 3 consumers.

Let v_k denote the voltage amplitude at the k-th node and θ_k the corresponding voltage angle, $\theta_{k,l}$ denotes the angle difference between node k and l. Throughout this paper, we assume that the voltage magnitude and the voltage angle at the first node (the slack node) are fixed, $\theta_1=0$ and $v_1=\mathrm{const}$. We refer to Du et al. (2021) and

Du et al. (2022) for further discussion. The state variables of the system is defined as

$$x \doteq [v_k, \theta_k]_{k \in \mathcal{N}}^{\top} \in \mathbb{R}^{2|\mathcal{N}|}.$$

Moreover, we have active and reactive power generation of generators p_k^g and q_k^g for all $k \in \mathcal{G}$, and $d_k = (p_k^d, q_k^d)^{\top}$ denote the active and reactive power demand at demand nodes $k \in \mathcal{D}$. We consider active and reactive power supply of all generators

$$u \doteq \left[p_k^g, q_k^g \right]_{k \in \mathcal{G}}^{\top} \in \mathbb{R}^{2|\mathcal{G}|}$$

as the system input variables.

The active and reactive power flow over the transmission line $(k, l) \in \mathcal{L}$ is given by

$$\Pi_{k,l}(x) \doteq v_k^2 \begin{bmatrix} g_{k,l} \\ -b_{k,l} \end{bmatrix} - v_k v_l \begin{bmatrix} g_{k,l}, & b_{k,l} \\ -b_{k,l}, & g_{k,l} \end{bmatrix} \begin{bmatrix} \cos(\theta_{k,l}) \\ \sin(\theta_{k,l}) \end{bmatrix}.$$

The total power outflow from node $k \in \mathcal{N}$ is given by

$$P_{k}(x) \doteq v_{k}^{2} \sum_{l \in \mathcal{N}_{k}} \begin{bmatrix} g_{k,l} \\ -b_{k,l} \end{bmatrix}$$
$$-v_{k} \sum_{l \in \mathcal{N}_{k}} v_{l} \begin{bmatrix} g_{k,l}, & b_{k,l} \\ -b_{k,l}, & g_{k,l} \end{bmatrix} \begin{bmatrix} \cos(\theta_{k,l}) \\ \sin(\theta_{k,l}) \end{bmatrix}$$
$$\doteq \sum_{l \in \mathcal{N}_{k}} \Pi_{k,l}(x),$$

where

$$\mathcal{N}_k \doteq \{l \in \mathcal{N} \mid (k, l) \in \mathcal{L} \}$$

denotes the set of neighbors of node $k \in \mathcal{N}$.

3. OPTIMAL DISTRIBUTION OF LIMITED SUPPLY POWER NETWORK

In the limited power supply scenario, potentially,

$$\sum_{k \in \mathcal{G}} \overline{p_k^g} < \sum_{k \in \mathcal{D}} p_k^d,$$

here, $\overline{p_k^g}$ denotes the upper bound of active generator power input of bus k. This leads to a fact that not all the power demand will enjoy stable energy supply. Thus we introduce $y \in \mathbb{R}^{|\mathcal{D}|}$ with $y_k \in \{0,1\}$ as an auxiliary switch variable. Then the power supply at a given bus can be expressed as

$$S_{k}(u,y) \doteq \begin{cases} \left[p_{k}^{g}, q_{k}^{g}\right]^{\top} - y_{k}^{2} \left[p_{k}^{d}, q_{k}^{d}\right]^{\top}, & k \in \mathcal{D}, \mathcal{G} \\ \left[p_{k}^{g}, q_{k}^{g}\right]^{\top}, & k \notin \mathcal{D}, k \in \mathcal{G} \\ -y_{k}^{2} \left[p_{k}^{d}, q_{k}^{d}\right]^{\top}, & k \in \mathcal{D}, k \notin \mathcal{G} \end{cases}$$

Thus, the power flow equations can be written in the form

$$P(x) = S(u, y) , (1)$$

where

$$P(x) \doteq \left[P_1(x)^\top, \dots, P_N(x)^\top \right]^\top,$$

$$S(u, y) \doteq \left[S_1(u, y)^\top, \dots, S_N(u, y)^\top \right]^\top.$$

Notice that $\dim(P) = \dim(x) = 2|\mathcal{N}|$.

Optimal power demand shut-offs (weighted maximum load delivery, Rhodes et al. (2020)) can be formulated in the

following way for a given positive rank $r_k \in \mathbb{R}^+$ of each power demand.

$$\max_{x,u,y} f(x,u,y) = \sum_{k \in \mathcal{D}} y_k r_k p_k^d$$
 (2a)

s.t.
$$P(x) - S(u, y) = 0$$
 (2b)
 $\underline{x} \le x \le \overline{x}$ (2c)

$$x \le x \le \overline{x} \tag{2c}$$

$$\underline{u} \le u \le \overline{u}$$

$$y_k \in \{0, 1\}.$$
(2d)
(2e)

$$y_k \in \{0, 1\}.$$
 (2e)

Here we are trying to provide stable power supply for the most important consumers with limited resource. Unlike other classical optimization problems over power grid (Zhu (2015), Frank et al. (2012)), the Boolean variable y leads Equation (2) into a non-smooth non-convex MBNLP form. This leads it difficult for the mainstream algorithms such as interior point method to be applied directly. However, it can be reformulate as a mathematical programs with complementarity constraints problem (MPCC) (Hall et al. (2021), Hall (2021)).

We reformulate Problem (2) in the following form:

$$\max_{x,u,y} E(x,u,y) \tag{3a}$$

s.t.
$$C(x, u, y) \le 0$$
 (3b)
 $0 \le y_k \perp (1 - y_k) \ge 0$, (3c)

$$0 \le y_k \perp (1 - y_k) \ge 0,$$
 (3c)

$$\begin{split} E(x,u,y) &\doteq \sum_{k \in \mathcal{D}} y_k r_k p_k^g - \\ & \left(y_k r_k v_k \sum_{l \in \mathcal{N}} v_l (G_{k,l} \cos(\theta_{k,l}) + B_{k,l} \sin(\theta_{k,l})) \right), \end{split}$$

$$C(x, u, y) \doteq \begin{bmatrix} P(x) - S(u, y) \\ S(u, y) - P(x) \\ \frac{x - x}{x - \overline{x}} \\ \frac{u - u}{u - \overline{u}} \end{bmatrix}$$
(4)

and

$$0 \le y_k \perp (1 - y_k) \ge 0 \Leftrightarrow \begin{cases} 0 \le y_k & \text{(5a)} \\ 0 \le 1 - y_k & \text{(5b)} \\ 0 = y_k \cdot (1 - y_k). & \text{(5c)} \end{cases}$$

Note that the problem is based on AC static model at a specific time point without dynamic. The renewables and storage for microgrid is considered as the part of power supply in the model so considering these modules separately is not essential. Apart from the MPCC constraints, Problem (3) is still a nonlinear programming (NLP) problem which is relatively not easy to handle. In the next section, we will introduce a new method called Alternating Optimization Based Sequential Boolean Quadratic Programming Method (AO-SBQP) (Zhu and Du (2022)) that can deal with it.

4. ALTERNATING OPTIMIZATION BASED SEQUENTIAL BOOLEAN QUADRATIC PROGRAMMING

This section reviews the basic AO-SBQP structure (Zhu and Du (2022)) which solves Problem (3) in a sequential way by optimizing the continuous and Boolean variables into different steps.

4.1 Approximation of BQP

The Lagrangian function of Equation (3) (Bertsekas (1997)) without (5) is

$$\mathcal{L}_0(x, u, y, \lambda_0) \doteq E(x, u, y) - \lambda_0^{\top} C(x, u, y).$$

According to the standard penalty reformulation,

$$\phi(y) \doteq y^{\top} (\mathbf{1} - y)$$

is considered as the bi-linear complementarity penalty function relates to Equation (5c). A second-order Taylor expansion of $\mathcal{L}_0(y,\lambda_0)$ (Hall et al. (2021)) on penalty with respect to y is

$$v(y) \doteq \frac{1}{2} y^{\top} Q y + (g - \rho \nabla \phi(\tilde{y}))^{\top} y.$$
 (6)

Here $0 \succ Q \doteq \nabla_y^2 \mathcal{L}_0(\tilde{x}, \tilde{u}, \tilde{y}, \tilde{\lambda}_0) \in \mathbb{R}^{|\mathcal{D}| \times |\mathcal{D}|}, g =$ $\nabla_y E(\tilde{x}, \tilde{u}, \tilde{y}) \in \mathbb{R}^{|\mathcal{D}|}, \ \rho > 0, \ \text{and} \ (\tilde{x}, \tilde{u}, \tilde{y}, \tilde{\lambda}) \ \text{represents the value of} \ (x, u, y, \lambda) \ \text{from the last NLP iteration}.$

For each iteration, the following simplified QP (7) needs to be solved, here

$$\max_{y} \ \upsilon(y - \tilde{y}) \tag{7a}$$
s.t. $b + A \cdot (y - \tilde{y}) \ge 0$ $|\lambda$ (7b)
$$y \ge 0 \qquad |\mu$$
 (7c)
$$\mathbf{1} - y \ge 0 \qquad |\gamma$$
 (7d)

s.t.
$$b + A \cdot (y - \tilde{y}) \ge 0$$
 $|\lambda|$ (7b)

$$y \ge 0 \qquad \qquad |\mu \qquad (7c)$$

$$1 - y \ge 0 \qquad |\gamma \qquad (7d)$$

with $A = \nabla_y C(\tilde{x}, \tilde{u}, \tilde{y}) \in \mathbb{R}^{(8|\mathcal{N}|+4|\mathcal{G}|) \times |\mathcal{D}|}$, $b = C(\tilde{x}, \tilde{u}, \tilde{y}) \in \mathbb{R}^{(8|\mathcal{N}|+4|\mathcal{G}|)}$. As discussed in Ralph* and Wright (2004), penalty parameter ρ can be modulated to meet the complementartiy satisfaction.

4.2 Local Convergence Analysis for Approximate BQP

In this subsection, we will show a convergence analysis of the simplified QP (7). For representational convenience, we introduce $\Delta y = y - \tilde{y}$ as the primal step.

The merit function

$$\psi(y) \doteq \frac{1}{2} y^{\top} Q y + g^{\top} y - \rho \phi(y)$$
 (8)

represents the outer loop objective function. It pointed out that merit function $\psi(y)$ at y^k iteration is non-increasing towards Equation (7) for the local convexity of v(y) (Hall et al., 2021, Section III) and the property

$$\nabla \psi(y^k)^{\top} \Delta y = \nabla \psi(y^k)^{\top} \Delta y.$$

However no convergence rate is discussed.

To ensure any local solution is a regular stationary point of Equation (7), two assumptions are introduced below.

Assumption 1. Linear Independence Constraint Qualification Condition (LICQ)

The matrix A has full row rank in the optimal value of y^* thus the gradients of active inequality and equality constraints are linearly independent. We refer to (Hall, 2021, Chapter 2) for further discussion.

Assumption 2. Second Order Sufficient Condition (SOSC) We assume the Hessian matrix Q is negative semi-definite in a local neighborhood of y^* . This statement is well-known in Newton-type algorithms. Suppose Q is not negative semi-definite in the current iteration, then set $Q \leftarrow Q - \rho I$ (Bliek1ú et al. (2014)).

Theorem 1. Let Assumption 1 and 2 of Equation (7) be applicable, then the iteration of BQP can converge to the local saddle point $\phi(y^*)$ with super-linear convergence rate by using suitable line search step size α and penalty parameter ρ .

Proof. The Lagrangian function of Problem (7) shows as

$$\mathcal{L}(y,\lambda,\mu,\gamma) \doteq \frac{1}{2} \Delta y^{\top} Q \Delta y + (g - \rho \nabla \phi(\tilde{y}))^{\top} \Delta y + \lambda^{\top} (b + A \Delta y) + \mu^{\top} y + \gamma^{\top} (\mathbf{1} - y).$$
(9)

Assume $[\lambda_{\rm act}^{\intercal}, \mu_{\rm act}^{\intercal}, \gamma_{\rm act}^{\intercal}]^{\intercal}$ collects the dual variables of the active inequalities of Problem (7), $[(b+A\Delta y)_{\rm act}^{\intercal}, y_{\rm act}^{\intercal}, (1-y)_{\rm act}^{\intercal}]^{\intercal}$ collects the active constraints. The KKT (Karush-Kuhn-Tucker) system of Equation (9) can be summarized as

$$\begin{bmatrix} Q & A_{\mathrm{act}}^\top & I_{\mathrm{act}}^\mu & -I_{\mathrm{act}}^\gamma \\ A_{\mathrm{act}} & 0 & 0 & 0 \\ I_{\mathrm{act}}^\mu & 0 & 0 & 0 \\ -I_{\mathrm{act}}^\gamma & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} \Delta y \\ \lambda_{\mathrm{act}} \\ \mu_{\mathrm{act}} \\ \gamma_{\mathrm{act}} \end{bmatrix} = \begin{bmatrix} -g + \rho \nabla \phi(\tilde{y}) \\ -b \\ 0 \\ 0 \end{bmatrix}.$$

Thus the optimal updating primal variable can be expressed as

$$\Delta y^*(\rho) = -Q^{-1} \left(g - \rho \nabla \phi(\tilde{y}) + (A^{\top} \lambda)_{\text{act}} + \mu_{\text{act}} - \gamma_{\text{act}} \right).$$

Notice the residual norm of $\phi(y)$ can be expressed as

$$\left\| \frac{y^{\top} \nabla \phi(\tilde{y})}{\tilde{y}^{\top} \nabla \phi(\tilde{y})} \right\| = \left\| \frac{(\tilde{y} + \alpha \Delta y^{*}(\rho))^{\top} \nabla \phi(\tilde{y})}{\tilde{y}^{\top} \nabla \phi(\tilde{y})} \right\|$$

$$= \left\| 1 + \alpha \frac{\Delta y^{*}(\rho)^{\top} \nabla \phi(\tilde{y})}{\tilde{y}^{\top} \nabla \phi(\tilde{y})} \right\|$$
(10)

with α denotes the step size. As long as

$$\alpha = -\frac{\tilde{y}^{\top} \nabla \phi(\tilde{y})}{\Delta y^{*}(\rho)^{\top} \nabla \phi(\tilde{y})},$$

Equation (10) can provide a local super-linear convergence (Nocedal and Wright (2006)) of the penalty function $\phi(y)$.

Note that this is an extension result of (Hall et al., 2021, Section III-C) by considering the related active inequality. We give the corresponding local convergence result since the penalty parameter ρ can also be turned.

Integrated the structure of relaxed BQP, Algorithm 1 summarizes the full AO-SBQP steps for solving problem (3). The main idea is to use the alternate optimization method to solve the continuous and Boolean variables respectively.

AO1 Derived from Frank et al. (2012), this variant of Optimal Power Flow aims to maximize the linear electric power guarantee objective with system input u, state x, fixed \tilde{y} and the power grid physical constraints (4).

AO2 The Hessian Q and gradient g are evaluated jointly with $(\tilde{x}, \tilde{u}, \tilde{y}, \tilde{\lambda})$. In order to search for a global maximizer of Equation (7) without the complementarity constraints, parameter ρ of Equation (6) is set as 0 in Step (2a). The

Algorithm 1 AO-SBQP Method

Input: initial guess \tilde{y} , initial \bar{x}, \bar{u} , a termination tolerance $\epsilon > 0$, an initial factor $\rho > 0$ and update rate $\beta > 1$.

Repeat:

- (1) Linear Optimal Power Flow (AO1): Solve an optimization problem consists of (3a) and (3b) with given \tilde{y} . Then output optimal power system decision variables \tilde{x}, \tilde{u} . ¹
- (2) Sequential BQP (AO2):
 - (a) Globally Search: Solve QP consists of (7) with zero penalty parameter. Output optimal switch variable \hat{y} .
 - (b) Update Penalty Function Approximate:

$$\begin{aligned} \phi(y) &\approx \phi(\hat{y}) + (y - \hat{y})^{\top} \nabla \phi(\hat{y}) \\ &= (\phi(\hat{y}) - \hat{y}^{\top} \nabla \phi(\hat{y})) + y^{\top} \nabla \phi(\hat{y}). \end{aligned}$$

- (c) Locally Search: Maximize the penalty QP (7).
- (d) Line Search and Inner Termination Criterion:
 α = StepLength(ŷ, ỹ, ρ);
 ŷ ≈ ŷ + α(ỹ ŷ).
 Check if |φ(ŷ)| ≤ ε, if not, go to Step (2e); if yes, ỹ ∉
 - Check if $|\phi(\hat{y})| \le \epsilon$, if not, go to Step (2e); if yes, $\tilde{y} \leftarrow \hat{y}$ and go to Step (3).
- (e) Penalty Parameter Update:
 - $\rho = \beta \cdot \rho$ and return Step (2a)
- (3) Outer Termination Criterion: Check if

$$\|((\tilde{x}, \tilde{u})|\tilde{y}) - (\bar{x}, \bar{u})\| \le \epsilon,$$

if not, go back to Step (1) and set $(\bar{x}, \bar{u}) = ((\tilde{x}, \tilde{u})|\tilde{y}), \, \bar{y} = \tilde{y};$ if yes, output the result.

Output: $(x^*, u^*, y^*) \leftarrow (\tilde{x}, \tilde{u}, \tilde{y}).$

later steps aiming to asymptotically meet Equation (5) by increasing ρ in Step (2e). Note that, both Step (2a) and Step (2c) are simple convex QP which can be solved by any stable QP solver. We refer (Zhu and Du, 2022, Section III) as a reference for more details.

Remark 1. (Relaxation of AO2) For the robustness of switching from AO2 to AO1, (7b) can be arbitrarily replaced by the following Inequality (11) in Step (2a) and Step (2c) which named as $mixed\ AO2$,

$$\begin{cases}
\sum_{k \in \mathcal{G}} p_k^g - \sum_{k \in \mathcal{D}} y_k p_k^d \ge 0 \\
\sum_{k \in \mathcal{G}} \overline{q_k^g} - \sum_{k \in \mathcal{D}} y_k q_k^d \ge 0 \\
\sum_{k \in \mathcal{D}} y_k q_k^d - \sum_{k \in \mathcal{G}} \underline{q_k^g} \ge 0.
\end{cases}$$
(11)

Or even, in some cases, the entire AO2 process can be replaced by solving the *relaxed AO2* (12) module below,

$$\max_{y} \sum_{k \in \mathcal{D}} y_k^2 r_k p_k^d \quad \text{s.t. (11), (3c)}.$$
 (12)

5. NUMERICAL RESULT

In this section, we illustrate the numerical result of AO-SBQP method drawing upon the modified 30-bus power network. The power sources of microgrid are considered into the system.

5.1 Data and Implementation

The problem data is obtained from MATPOWER dataset Zimmerman et al. (2011) and the implementation of Algorithm (1) relays on Casadi-v3.5.5 with IPOPT (Andersson et al. (2019)). Though MATPOWER repository is for

¹ This step can be solved by any NLP solver.

transmission networks which are high voltage networks, the mathematical model of microgrid is same.

In the modified 30-bus case, $\mathcal{G} = \{1, 2, 13, 22, 23, 27\}$. To create a demand-to-power mismatch scenario, we increase the active and reactive power demands of buses with loadby 2.5 p.u. (per unit) and 0.7 p.u. respectively. The lower and upper bounds of reactive power inputs are reduced to half of the previous ones while upper bound of active ones are reduced to 70%. In addition, r_k 's are randomly set into five levels from 1 to 5 for each demand and the criterion (i.e. complementarity tolerance) is set as 10^{-6} . Problem (3) consists 60 status, 12 system inputs and 30 switch variables when all buses contain consumers.

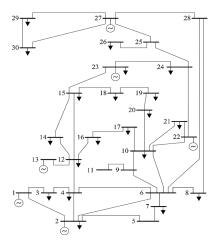


Fig. 2. Modified IEEE 30-bus system from Christie (2000).

Since we did not set multi demands for each bus, this indicates $|\mathcal{D}| = |\mathcal{N}|$, and (7b) is going to be an overdetermined system. Therefore our implementation focus on the other three relax versions of AO2 mentioned in Remark 1.

5.2 Numerical Comparison

In this section, we show the comparison of Algorithm (1) with different variations of AO2. Note that all the variations consist inequality constraints (11), (7c), (7d), and the only difference are the objectives, a) Mixed: (7a), b) Relaxed I: $\sum_{k \in \mathcal{D}} y_k^2 r_k p_k^d - \rho \phi(y)$, c) Relaxed II: $\sum_{k \in \mathcal{D}} y_k^2 r_k p_k^d - \rho \nabla \phi(\tilde{y})^{\top} y$.

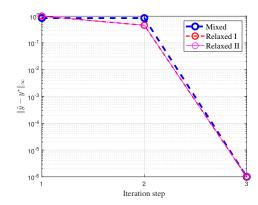


Fig. 3. Convergence of y with three variations of AO2.

Figure (3) and Figure (4) shows the convergence of the switch variable y and the complementarity satisfaction $\phi(y)$ respectively. It can be seen that even if the dimension of y is 30, all the variations of SBQP can converge to the given complementarity tolerance in only a few steps but converge to different solutions. This shows completely different properties than the B&B based solvers.

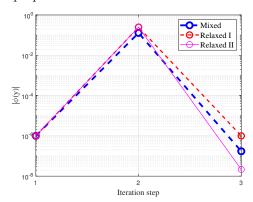


Fig. 4. Convergence of $\phi(y)$ with three variations of AO2.

Table 1. Convergence Comparison

Method	Mixed	Relaxed I	Relaxed II
$ \phi(y) $	1.7e-07	1.0e-06	2.2e-08
Time(s)	0.021	0.046	0.019
Iteration	2	2	2

Table 1 shows the comparison among complementarity satisfaction, operation time and iteration by using different variations. It can be seen that all indicators of the three methods are similar which is different from the results seen in Zhu and Du (2022). Since the three relaxation variations in this paper decrease the number of inequality constraints in AO2 that induce it easier to solve. Notice that, due to the nonlinear structure of (3a), the implementation is more complex than (2a), therefore only AO1 benefits from the reformulation of (3a).

Table 2. Performance Comparison

Method	Mixed	Relaxed I	Relaxed II	
Objective	5.421	2.428	5.877	
$\sum_{k \in \mathcal{D}} y_k p_k^d \text{ (p.u.)}$	1.567	0.825	1.592	
$\sum_{k \in \mathcal{D}} y_k q_k^d \text{ (p.u.)}$	0.702	0.347	0.709	

Table 2 shows the comparison of final performance comparison by using different methods. As can be seen, different variations converge to different local optimum. At least in this case, Relaxed II gets a bit better performance than the other two and Relaxed I is a bit conservative. Note that, both Table 1 and Table 2 show the benefits of the approximate BQP method (Mixed and Relaxed II).

Table 3. Optimal Power Inputs

Bus #	1	2	13	22	23	27
pg (p.u.)	0.400	0.400	0.134	0.250	0.150	0.275
qg (p.u.)	-0.037	0.161	0.213	0.246	0.070	0.125

Table (3) shows the numerical result of optimal system inputs by using Relaxed II based Algorithm (1). The active

demands relate with bus $\{1,2,4,9,10,11,12,13,14,16,17,18,21,22,23,24,25,26,30\}$.

6. CONCLUSION

In this paper, we proposed an effective and fast convergence method named AO-SBQP to optimize microgrid demand shut-offs problems. Importantly, local convergence theory of approximate BQP has been proposed. Moreover, a numerical result on modified IEEE 30-bus case study illustrates the potential of AO-SBQP in this area. Different from B&B and SDR, AO-SBQP can achieve a feasible local optimal solution without tree storage structure or matrix variables. Future research will investigate multistage optimal demand shut-offs and time varing priority of single bus-multiple demands on larger case studies. Moreover, the rank evaluating priority of each demand can vary with time. Comparison of accuracy and computation time of B&B and SDR will also be considered.

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