

Numerical Analysis: homework 03

Due on Tuesday, March 21, 2017

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1 Introduction

For modern circuit design, the scale of circuit is always such large that human cannot handle it manually. To overcome this, we frequently use a lot of EDA tool such as **hspice** to solve these problems. However, the license of these EDA tool is usually expensive. As a result, EDA provider such as **Synopsys** benefits a lot from it every year, which means these circuit analysis technology is quite valuable.

In this project, we will implement several simple **Resistor Networks** and use **LU Decomposition** to analyze it.

1.1 Resistor Networks

Each resistor networks only consists of several resistors as shown in Figure 1.

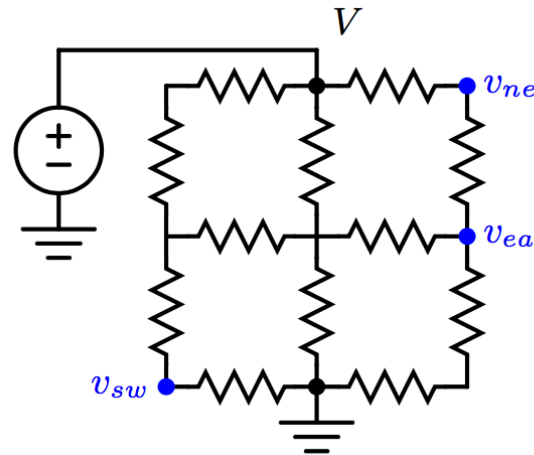


Figure 1: Simple resistor network

In this assignment, there are six networks to implement:

1. 2 resistors each side and resistance is $1K\Omega$
2. 4 resistors each side and resistance is 500Ω
3. 10 resistors each side and resistance is 200Ω
4. 20 resistors each side and resistance is 100Ω
5. 40 resistors each side and resistance is 50Ω
6. 50 resistors each side and resistance is 40Ω

For each of them, we have to show the equivalent resistance and the voltage of V_{ne} , V_{ea} and V_{sw} .

2 Implementation

2.1 Algorithm

To model the problem into A and b , I divide the network into two segments by horizontal and vertical resistor as shown in Figure 1. And then update matrix A and b with **Algorithm 1.4.1** in class material.

Algorithm 1.4.1 System Equation for a Resistor Network

For a network with N nodes in each side, create $N^2 \times N^2$ zero matrix A and N^2 zero vector b .

nodeIndex = 0

$G = resistance^{-1}$

for each $i \in \{1 \dots N\}$ **do**

for each $j \in \{1 \dots N-1\}$ **do**

$A[nodeIndex][nodeIndex] += G$

$A[nodeIndex][nodeIndex+1] -= G$

$A[nodeIndex+1][nodeIndex+1] += G$

$A[nodeIndex+1][nodeIndex] -= G$

 nodeIndex += 1

end for

nodeIndex += 1

end for

nodeIndex = 0

for each $i \in \{1 \dots N-1\}$ **do**

for each $j \in \{1 \dots N\}$ **do**

$A[nodeIndex][nodeIndex] += G$

$A[nodeIndex][nodeIndex+N] -= G$

$A[nodeIndex+N][nodeIndex+N] += G$

$A[nodeIndex+N][nodeIndex] -= G$

 nodeIndex += 1

end for

end for

for $v, i \in$ voltage and index at fixed voltage point **do**

$A[i] = 0$

$A[i][i] = 1$

$b[i] = v$

end for

$LU = \text{luDecompose}(A)$

$Y = \text{forwardSub}(LU, b)$

$X = \text{backwardSub}(LU, Y)$

X is the voltage of each node.

3 Discussion

3.1 Complexity

In all the following part, I will refer N as the square of the number of nodes at each side of networks. For the network in Figure 1, N is $3 \times 3 = 9$.

When solving the problem, the most time-consuming part is the **LU Decomposition**, which is $O(n^3)$. As a result, the system is a $O(n^3)$ problem.

3.2 Performance Evaluation

In this project, I use five numbers to indicate the efficiency of our method:

1. **Runtime**(total execution time of program)
2. **PROBLEM**(execution time of modeling problem into martrix A and vector b)
3. **LU**(time taken by LU Decomposition)
4. **FWD**(time taken by forward substitution)
5. **BCK**(time taker by backward substitution)

The detail result of each resistor networks is show in Table 1

N	9	25	121	441	1681	2601
R(each)	1000	500	200	100	50	40
R(equivalent)	1000	681.818182	376.009408	229.423008	136.042809	114.395519
V_{ne}	0.75	0.7	0.648693	0.622178	0.603088	0.598084
V_{ea}	0.5	0.5	0.5	0.5	0.5	0.5
V_{sw}	0.25	0.3	0.351307	0.377822	0.396912	0.401916
Runtime(s)	0.003	0.003	0.007	0.2	9.948	36.1
PROBLEM(s)	1.50E-05	2.10E-05	0.000199	0.001173	0.020578	0.052503
LU(s)	2.00E-06	3.30E-05	0.003537	0.194482	9.89144	35.9966
FWD(s)	1.00E-06	5.00E-06	4.40E-05	0.000489	0.008348	0.017034
BCK(s)	1.00E-06	3.00E-06	4.60E-05	0.000532	0.010656	0.019895

Table 1: Experiment result

3.3 Runtime and LU

Because LU Decomposition should be the most time-consuming part of algorithm, so after N is large enough, LU should be very close to Runtime. Figure 2 shows the result of Runtime/LU and N^3 with log-scale:

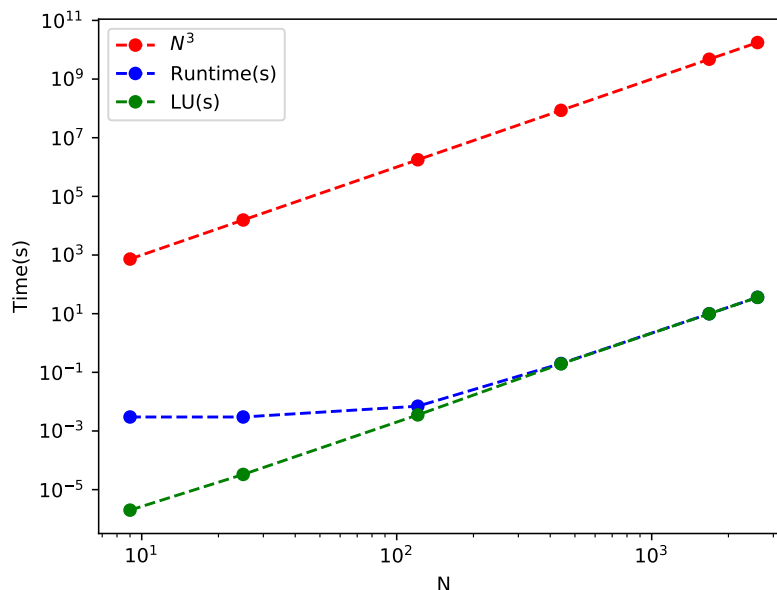


Figure 2: Runtime/LU vs N

Because when N is small, execution time will have huge error, so we should ignore former part of the line of Runtime. However, when N is large, we can see that Runtime and LU almost overlap and their slope is same as N^3 , which means the system is a $O(n^3)$ problem and satisfies my analysis in Section 3.1.

3.4 Forward/Backward Substitution

Because the complexity of Forward/Backward Substitution is $O(n^2)$ as we discuss in previous homework assignment, the slope of them should be same as N^2 in log-scale plot. Figure 3 shows the result with log-scale.

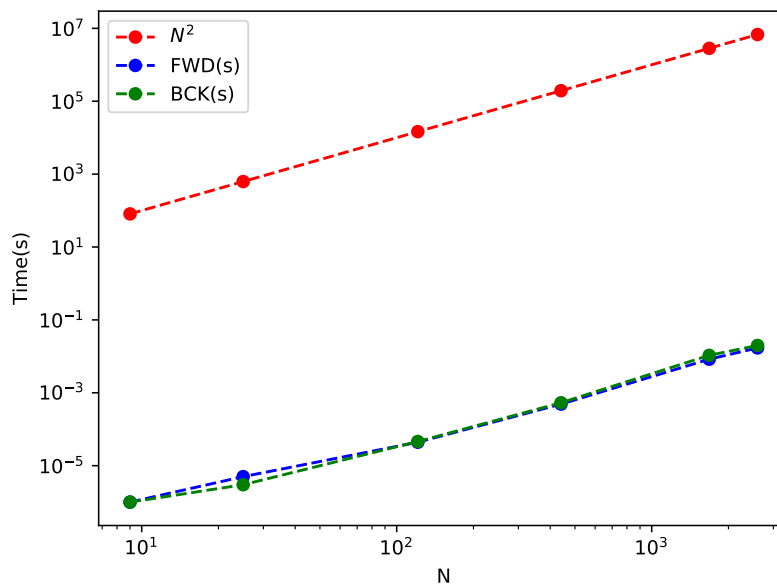


Figure 3: FWD/BCK vs N

From Figure 3, we can clearly see that the slope of FWD and BCK is same as N^2 , which means they are $O(n^2)$ problem.

3.5 Problem Modeling

When modeling the problem into matrix A , we need to create an $N \times N$ empty matrix, which is an $O(n^2)$ problem, Figure 4 show the modeling time and N^2 with log-scale:

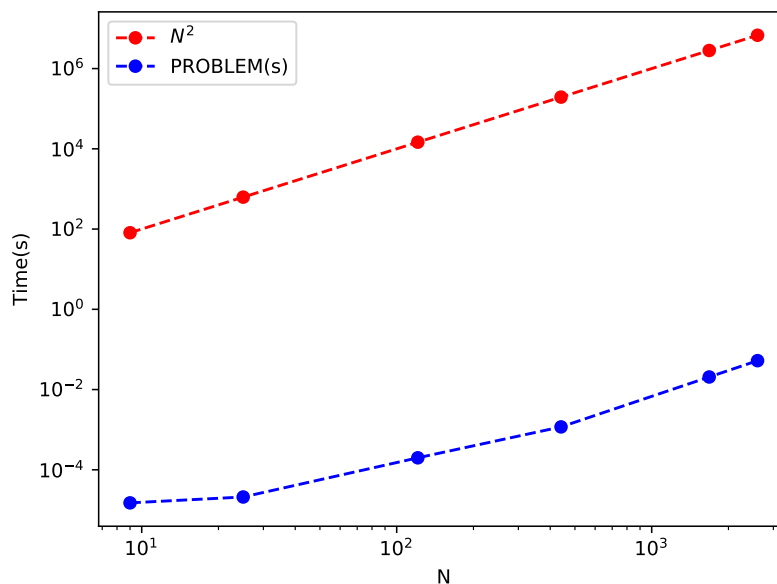


Figure 4: PROBLEM vs N

In the begin of Figure 4, we can see the unstable slope. However, after N is large enough, the slope is same as N^2 , which means this is a $O(n^2)$ problem.

4 Code Usage

To compile the code, just type `$ g++ hw03.cpp MAT.cpp VEC.cpp` in terminal.

To run the program, use `$./a.out 10`, 10 means there are 10 resistors at each side of the resistor network.