Developing iSCSI Storage Systems on Linux State of Linux iSCSI support

Tomonori Fujita

NTT Cyber Solutions Laboratories

tomof@acm.org
http://zaal.org/

Ottawa Linux Symposium, 2005/07/21, Ottawa, Canada

Outline

- What is iSCSI ?
- Initiator implementations
- Targets implementations
- Performance
- Remaining issues

What is iSCSI?

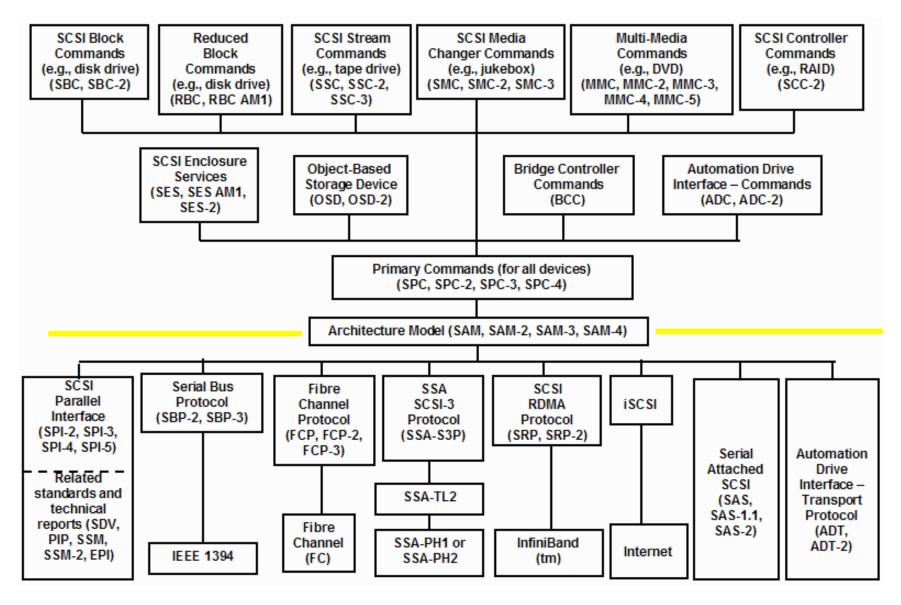
SCSI over TCP/IP

- SCSI commands are encapsulated into IP packets
- The host and storage systems are connected with Ethernet

TCP/IP is used as one of SCSI transports

- SCSI is a layered architecture
 - ▶ Device type specific command sets (disk, tape, etc)
 - Shared command sets (for all device types)
 - > Transport protocols
- Various transports are supported
 - ▶ Parallel SCSI (SPI) Parallel cable
 - ▶ Fibre Channel (FC) Fibre Channel
 - ▶ Serial Attached SCSI (SAS) Serial cable
 - iSCSI TCP/IP

SCSI Standard Architecture



taken from http://t10.org/

Why is iSCSI?

New networked storage technology - Storage Area Network (SAN)

- Today the dominant networked storage technology is Fibre Channel (SAN)
 - Specialized expensive hardware
- iSCSI can build inexpensive SAN (IP-SAN)

Why not NFS?

- Some applications does not like NFS
 - Database
- iSCSI works with NFS nicely (it does not replace NFS)
 - ▶ A large NFS server can have lots of disk drives by using SAN

iSCSI initiator implementations in Linux

Initiator

It issues SCSI commands to request services (hosts)

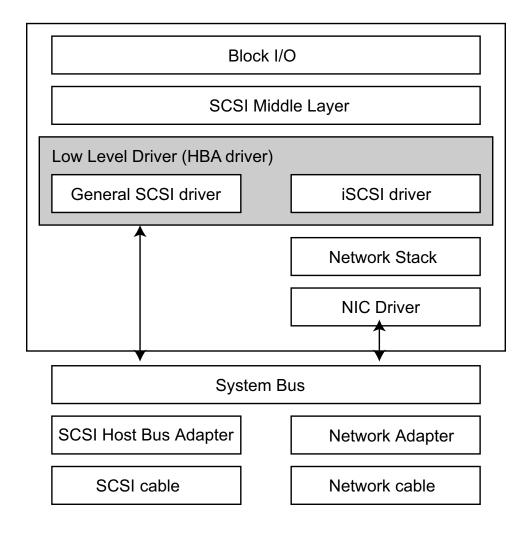
We have several implementations under GPL license

- UNH
- linux-iscsi (sfnet)
- core-iscsi
- open-iscsi

The common design of initiators in the Linux kernel

Implemented as a SCSI Host Bus Adapter (HBA) driver

- General LLDs communicate with its HBA using DMA.
- iSCSI initiators communicate with the network stack.



What does iSCSI disk look like?

No difference between iSCSI and general disk

- Three DAS (direct attached storage) disk drives (SEAGATE ST37345LC)
- One iSCSI disk drive (IET VERTUAL-DISK)

```
Debian GNU/Linux comes with ABSOLUTELY NO WARRANTY, to the extent
permitted by applicable law.
 You have mail.
Last login: Tue Jul 19 09:53:09 2005 from rouen.hil.ntt.co.jp
lales:~$ cat /proc/scsi/scsi
lAttached devices:
Host: scsi0 Channel: 00 Id: 00 Lun: 00
   Vendor: SEAGATE Model: ST373453LC
                                           Rev: DX10
                                           ANSI SCSI revision: 03
  Type: Direct-Access
Host: scsi0 Channel: 00 Id: 01 Lun: 00
   Vendor: SEAGATE Model: ST373453LC
                                           Rev: DX10
                                           ANSI SCSI revision: 03
   Type: Direct-Access
Host: scsi0 Channel: 00 Id: 02 Lun: 00
   Vendor: SEAGATE Model: ST373453LC
                                           Rev: DX10
         Direct-Access
                                           ANSI SCSI revision: 03
Host: scsi0 Channel: 00 Id: 06 Lun: 00
                 Model: 1x6 SCSI BP
   Vendor: PE/PV
                                           Rev: 1.0
                                           ANSI SCSI revision: 02
  Type: Processor
Host: scsi2 Channel: 00 Id: 00 Lun: 00
   Vendor: IET
                  Model: VIRTUAL-DISK
                                           Rev: 0
                                           ANSI SCSI revision: 04
   Type:
         Direct-Access
√ales:~$ [
```

UNH initiator - Initiators in Linux (1)

Things to know

- Implemented by University of New Hampshire people
- Maintained by HP and University of New Hampshire people
- The first iSCSI initiator implementation in Linux
- Aimed for a reference implementation

Advantages

 Most of features in the iSCSI RFC are supported (though some of them are not so useful)

Disadvantages

- Not tested intensively in production environments
- Configuration scheme is not handy
- Not designed properly from Linux kernel perspective
 - ▶ It doesn't use SCSI Transport attributes feature in the Linux kernel
 - Everything in kernel space
 - ▶ Wrong coding style

Linux-iscsi (sfnet) - Initiators in Linux (2)

Things to know

- Implemented by Cisco
- Maintained mainly by Cisco and major storage vendor people
- It WAS considered as the standard iSCSI initiator implementation in Linux
- Some distributions include this driver by default
- Note that linux-iscsi 5.x is not sfnet (it's open-iscsi)
- 165 mailing list subscribers

Advantages

- Tested intensively in production environments
- Supported by lots of commercial iSCSI storage systems
- Probably the most robust
- Decent configuration scheme

Disadvantages

- Not more development though still maintained
 - ▶ It cannot be compiled with 2.6.11 or the newer versions

open-iscsi - Initiators in Linux (3)

Things to know

- Implemented and maintained by Dmitry Yusupov and Alex Aizman
- It will be merged to the mainline kernel (hopefully in the near future)
- It's also known as linux-iscsi version 5.X
 - ▶ linux-iscsi version 4.X and 3.X refer to sfnet
- 111 mailing list subscribers

Advantages

- Designed properly from Linux kernel perspective
 - ▶ A large portion in user space
- The most active development
 - ▶ Many iSCSI developers are focusing on this implementation

Disadvantages

- It needs more tests in real (production) environments
- Some important features are not implemented yet

core-iscsi - Initiators in Linux (4)

Things to know

- Implemented and maintained by PyX Technologies, Inc
- It had been distributed with PyX Technologies target systems (Not GPL)

Advantages

- Tested intensively with PyX Technologies target systems
- Decent configuration scheme

Disadvantages

- Not active development
 - ▶ It cannot be compiled with 2.6.12-rc1 or the newer versions
- Not designed properly from Linux kernel perspective
 - ▶ It doesn't use SCSI Transport attributes feature
 - Everything in kernel space

Summary - Initiators in Linux

- Linux iSCSI initiators are ready for most production environments
- Try open-iscsi in your environment now because it will be merged into the mainline kernel
- sfnet would be a safe choice if you cannot use the latest kernel

How to implement iSCSI storage systems (target)

Target

It provides services to initiators (storage systems)

Commercial solutions

- Commercial iSCSI storage system use optimized operating systems and specialized hardware
 - ▶ iSCSI specialized NIC, NVRAM, etc

Open source solutions

- Linux and commodity hardware
 - ▶ PC and Gigabit NIC
- Linux provides many features useful for iSCSI target systems
 - ▶ Software RAID
 - \triangleright LVM
- You can implement whatever you want
 - Management interfaces
 - ▶ Access control

iSCSI target implementations in Linux

We have several implementations under GPL license

- UNH
- Ardis
- iSCSI Enterprise target software (aka IET)

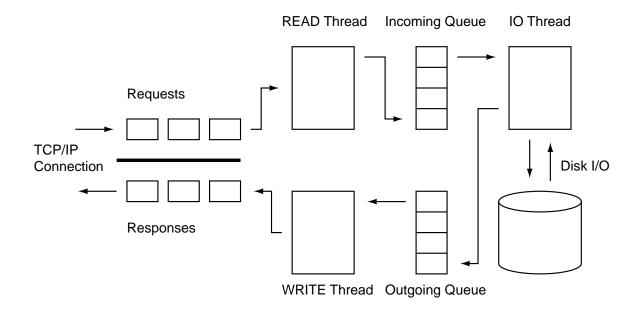
Features

- Virtual devices
 - ▶ Exporting files (including block device files) as disk to initiators
 - Exporting iso image files as cdrom drives to initiators
 - ▶ Exporting files (including block device files) as tape drives to initiators
- Real devices
 - Exporting real devices like tape and cdrom drives

The common design of targets in the Linux kernel

All the implementations run in Linux space

- They need the full control of page cache
 - ▶ Intelligent caching like commercial storage systems
- They interact with hardware like NVRAM



UNH target - targets in Linux

Things to know

- Implemented by University of New Hampshire people
- Maintained by HP and University of New Hampshire people
- Probably the first iSCSI target implementation in Linux
- Aimed to debug UNH initiator implementation

Advantages

Most of features in the iSCSI RFC are supported

Disadvantages

- Unnecessarily complicated design
 - Designed for multiple transport protocols at the beginning
- Configuration scheme is not handy
 - ▶ No dynamic changes to configurations
 - Very complicated

Ardis target - targets in Linux

Things to know

- Implemented by Ardis technologies
- It keeps unchanged for already one and a half years

Advantages

Decent design

Disadvantages

- Probably dead
- Support only 2.4 kernels
- Need a kernel patch
- Configuration scheme is not handy
 - ▶ No dynamic changes to configurations
- Buggy

iSCSI Enterprise Target (IET) - targets in Linux

Things to know

- Started as a fork from Ardis
- Maintained by me
- The most part of Ardis code has gone (especially in kernel space)
- 203 mailing list subscribers

Advantages

- Decent configuration scheme
 - ▶ Very simple
 - ▶ All configurations are dynamically changeable
 - Access control based on initiator address and target name patterns
- Tested heavily
 - Some brave people use it in production use

Disadvantages

Some important features are not implemented yet

The detailed designs of initiators

Things to know

- open-iscsi integrates iSCSI processing into the network stack
 - ▶ A bit complicated but effective code

UNH, sfnet, core-iscsi

- tx thread writes to sockets
- rx thread read from sockets
- queuecommand function links SCSI commands to an internal link
- tx thread unlinks SCSI commands, builds iSCSI PDUs, and send them
- rx thread calls cmnd->done

open-iscsi

- queuecommand function builds iSCSI PDUs and send them
- sock->data_ready calls cmnd->done

The performance comparison of initiators

Things to know

- There is no performance difference due to the bottleneck of target's disk in most cases
- All initiators are fast enough for 1 Gigabit Ethernet and normal slow target's disk
- You can see the difference if you use the combination of 10 Gigabit Ethernet and a really fast target like RAM disk

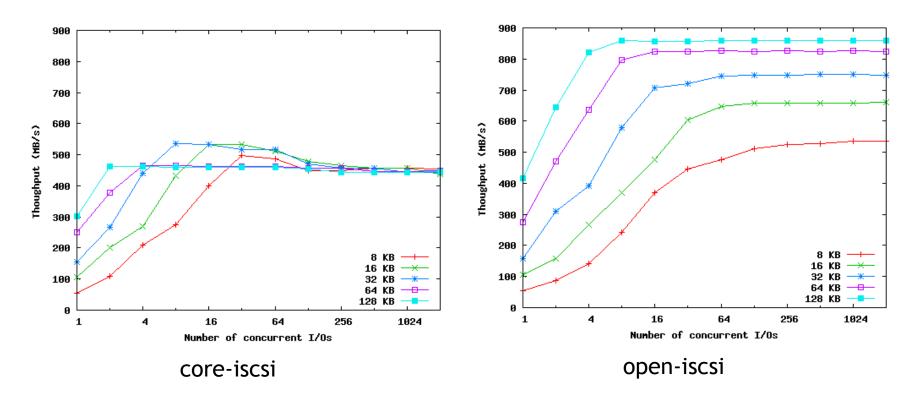
core-iscsi vs. open-iscsi

- Both claim that it's a high-performance iSCSI initiator
- core-iscsi 1.6.2.0-rc1 / open-iscsi 0.3rc6-363

Experiment environment

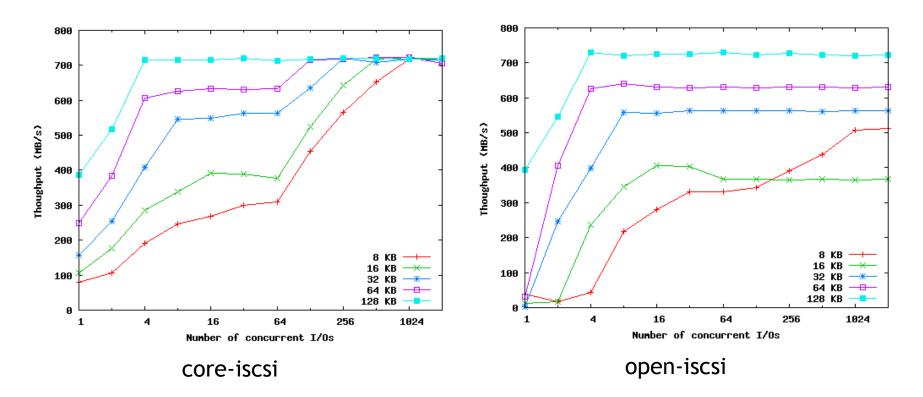
- Two Opteron (2.4 GHz) boxes are connected directly
- Chelsio T110 10GbE NIC (full TOE NIC, all TCP functionality is implemented on NIC chip)
- Benchmark software runs in kernel space (uses bio functions)
- IET runs in nullio mode (dumps received data and sends meaningless data)

Sequential read



 open-iscsi is much faster (due to the iSCSI processing integrated into the network stack)

Sequential write



- Both can achieve 700 MB/s
- open-iscsi performs terribly with single concurrent I/O (since the combination of open-iscsi and Chelsio network stack leads to some bugs)

Remaining issues with initiators (1)

System locks in out-of-memory situations

Symptoms

- Suppose that the system is in an OOM situation due to lots of dirty pages
- ▶ The system tries to write dirty pages to iSCSI disk to reclaim memory
- ▶ The network stack fails to allocate memory since the system is already in an OOM situation
- ▶ The system cannot write dirty pages thus it cannot reclaim memory

Problem

- You cannot preallocate some memory in the network stack for an OOM situations
- Even in an OOM situations, you must send and receive several packets to write one write request at least

Status

▶ The developers cannot do nothing (New features in the network stack are necessary)

Remaining issues with initiators (2)

Unnecessarily I/O errors happens under heavy load

Symptoms

- ▶ iSCSI initiators have some time limit for some operations
- Suppose that the target sends a ping thus the initiator must respond immediately
- ▶ The initiator cannot send since the system is under heavy load
- ▶ The target judges the initiator dead and closes the connection

Problem

- ▶ iSCSI initiator needs some real-time stuff (scheduling and network processing)
- The issue gets complicated especially with open-iscsi because the most parts are in user-space

Status

- > open-iscsi use mlock to avoid page out
- ▶ New feature in the network stack are necessary again

Remaining issues with targets (1)

Unnecessary I/O error happens under heavy load

- Symptoms
 - ▶ iSCSI targets have some time limit for some operations
 - Suppose that the initiator sends a request thus the target must respond immediately
 - ▶ The target cannot send since the system is under heavy load
 - ▶ The initiator judges the target dead and aborts the request

Problem

▶ iSCSI target needs some real-time stuff (scheduling and network processing)

Status

▶ New features in the network stack are necessary again

Remaining issues with targets (2)

Write back vs. write through (disk drive cache algorithm)

- Symptoms
 - ► The target uses a file with the vfs interface to export it as disk to initiators
 - ▶ With writeback, the target doesn't need to sync h dirty page
 - ▶ The amount of dirty page cache increases too much (several hundreds mega bytes) compared with normal disk cache (several mega bytes)
- Problem
 - ▶ The target cannot control the amount of dirty page cache
- Status
 - ▶ IET supports only write through