

Proof of Chebyshev's theorem

$P_{n+1}(x)$ is a monic polynomial of degree $n+1$.
 For any $x \in [-1, 1]$, assume $|P_{n+1}(x)| < \frac{1}{2^n}$ (*)

Note that $T_{n+1}(x)$ alternates between -1 and 1 a total of $(n+2)$ times

* At these $(n+2)$ points, $P_{n+1} - \frac{T_{n+1}}{2^n}$ is alternatively positive and negative. Why? 2^n

* We note that at one of these points z

$$\begin{array}{ccc}
 P_{n+1}(z) - \frac{1}{2^n} T_{n+1}(z) & & \\
 \swarrow & & \searrow \\
 = P_{n+1}(z) - \frac{1}{2^n} & & P_{n+1}(z) + \frac{1}{2^n} \\
 < 0 & & > 0
 \end{array}$$

using intermediate value theorem, $P_{n+1} - \frac{1}{2^n} T_{n+1}$ has at least $n+1$ zeros.

* However, $P_{n+1} - \frac{1}{2^n} T_{n+1}$ is at most degree $n+1$. This (since it is monic) for $P_{n+1} \neq T_{n+1}$ leads to a contradiction.

\Rightarrow What if $P_{n+1} - \frac{1}{2^n} T_{n+1} = 0$. Then $P_{n+1} = \frac{1}{2^n} T_{n+1}$

In that case, $|P_{n+1}| = \left| \frac{1}{2^n} T_{n+1}(x) \right| = \frac{1}{2^n}$. This contradicts definition of P_{n+1} (see *)

Note that T_{n+1} is a polynomial with $n+1$ roots $x_0 \dots x_n$.
 Therefore, $T_{n+1}(x) = c(x-x_0)(x-x_1) \dots (x-x_n)$

Let's require $T_{n+1}(x) = 1 \Rightarrow c = 1$.
 Leading coefficient of

How about the stability?

$$\Lambda_n(x) < \frac{2}{\pi} \left(\log(n+1) + \gamma + \log \frac{8}{\pi} \right) + \frac{\pi}{72(n+1)^2}$$

$\gamma \equiv$ Euler constant

This is much less sensitive to perturbation errors compared to Lagrange interpolation.

* Chebyshev nodes are optimal when f is a smooth explicit function

* However there are many functions in applications that don't have smoothness or we only have access to discrete data.

Piecewise linear interpolation

x_0, x_1, \dots, x_n

$I_i = [x_i, \dots, x_{i+1}]$

Given: $(x_i, f(x_i)) \quad (x_{i+1}, f(x_{i+1}))$

Linear interpolant: $\pi_i f(x) \equiv f(x_i) + \frac{f(x_{i+1}) - f(x_i)}{(x_{i+1} - x_i)} (x - x_i) \quad x \in I_i$

Main Result If $f \in C^2(I)$, where $I = [x_0, x_n]$ then

$$\max_{x \in I} |f(x) - \pi_i f(x)| \leq \frac{H^2}{8} \max_{x \in I} |f''(x)|$$

$H \equiv$ Maximum length of intervals I_i

$\pi_i f(x) \rightarrow f(x)$ as $H \rightarrow 0$

proof Recall proof of Lagrange interpolation

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi_x) \prod_{i=0}^n (x - x_i)$$

set $n=1$

$$f(x) - p(x) = \frac{1}{2!} f''(\xi_x) \prod_{i=0}^1 (x - x_i)$$

$$|f(x) - p(x)| \leq \max_{x \in I} |f''(x)| \frac{H^2}{8}$$

Applications: Fitting time-series data • fplot in MATLAB

$$(x_0, y_0) \quad (x_1, y_1) \quad (x_2, y_2) \quad \dots \quad (x_n, y_n)$$

$$P_n(x) = a_0 + a_1 x + \dots + a_n x^n$$

$$\begin{aligned} P_n(x_0) &= y_0 \\ P_n(x_1) &= y_1 \\ &\vdots \\ P_n(x_n) &= y_n \end{aligned} \Rightarrow \underbrace{\begin{bmatrix} 1 & x_0 & x_0^2 & \dots & x_0^{n-1} & x_0^n \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & x_n & x_n^2 & \dots & x_n^{n-1} & x_n^n \end{bmatrix}}_{V = \text{Vandermonde matrix}} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_n \end{bmatrix}$$

Exercise : show that V is invertible (assuming distinct points)

condition number of matrix

$$\begin{aligned} \frac{\|A(x + \Delta x) - Ax\|}{\|Ax\|} &= \frac{\|A \Delta x\|}{\|Ax\|} \\ &\leq \frac{\|A\| \|\Delta x\|}{\|Ax\|} \\ &\leq \left(\frac{\|A\| \|\Delta x\|}{\|Ax\|} \right) \frac{\|\Delta x\|}{\|\Delta x\|} = \kappa(x) \frac{\|\Delta x\|}{\|\Delta x\|} \end{aligned}$$

Assume A is invertible.

$$\begin{aligned} \|\Delta x\| &= \|A^{-1} Ax\| = \|A^{-1}\| \|Ax\| \\ \frac{\|A\| \|\Delta x\|}{\|Ax\|} &\leq \frac{\|A\| \|A^{-1}\| \|Ax\|}{\|Ax\|} = \|A\| \|A^{-1}\| = \kappa(A) \end{aligned}$$

It turns out that V is highly ill-conditioned as a function of n