Hw 2 1) Reflexive X=Cxo,yo) then xo +yo=xo+yo x~x

Symmetric X= Cxo,yo) y= Cx,yyi) x~y so

Xo · yo = Xi · yi → yi+xi = yo+xo =o y~x Transitive) X = Cxuya) y = Cxuya) Z= Cxzyz) X my and y NZ so:

Xo + yo = Xi + yi = Xo + yo = So

Xo + yo = Xi + yi = Xo + yo = So

Xo + yo = xi + yo and x nz

This equivalence class is set of circles centered

at (go) w/rad ws Jx + yo 2) There isn't necessarily a b s.t. and so it doesn't hold. 3 a) Reflexive:  $f(x_1) = f(x_1)$  so  $x_1 \sim x_1$ Symmetric: let  $x_1 \sim x_2$   $f(x_1) = f(x_2) - f(x_2) = f(x_1)$ so  $x_2 \sim x_1$ trons live. Let XONX, and X, NX2 50 f(xo)=f(x,) f(xi)=f(x2) 80 f(xo) = f(x2) and xo NX2. Therefore, this is an equivalence relation b) Let  $f_2: X/\sim \rightarrow Y$   $(x) \rightarrow f_2(x)$ Since all claments in an cavivalence class have the same image, it is well defined .. To show injectives let x, x, EX ather x If factory) = factory) I then FCx) = fCxx) Bolx, ~xxx Innerling = Exxy = Exxy and lufa is injective To show succeetive let g: X > X/N From earlier, f(x) = f2 (g(x)).

By assumption, f(x) is surjective, meaning that famust be surjective. as shown back as Hw Is. Sinker fais I surjective and the injectives and the injectives and the injectives and because secrespondance exists. 4 Comparability. X = y then X = y = y = X = WLOG X < Y So X = xy , but: f x = y = y then X < y as x +4. Anticeflexivity : YXGIR, XXX, X2=X2 and X=X 5) Let L= Elllis a lover bound of A3 Donote sup = band YlEL, 156 Let U= Eu lu is an upper bound of L3 ACU as YaGA and YlEL, Lea as lisalauer bound on A so Aris an upper bound on L. So, YuEU, be u and bea, so b=maxl meaning b=infA. 6a) Let bEB=Cbyb2. bn)
Since B, EA, Ba EA2. then b, EB, EA, b, EB2 SA2... and b, GA, b2 GA2....
So b GA, XA2X. = A and BSA

b) Let b=Cbuba - JEB and each B, is nenempty.

Since BEA, then bEA, so b, EA, b2EAa

But BEB; and so b; EB; is

R.CA: otherwise, the stranger restriction and B; EA: otherwise, if A; EB; then JbEB; s.t. b#A; which doesn't work as BEA. Therefore, the converse of If A snonempty then no A; can be compty. If JA; s.t. A; = \$\phi\$ then A = \$\phi\$ as far any AxB, if B=\$\phi\$ then Ax\$\phi=\$\phi\$. Therefore no A; can be empty if A is nenempty. The Converse, if each A; is notempty then A is nonempty holds. The contrapositive, if A is empty than JA; that is empty holds, as discussed previously; so the anginal statement, the converse, holds. d) (A, UBi) × (A2 UB2) X.... PAUB Note this as TTCA: UB; ) = AUB Notex The B Gartesian xproductionis calantable, A UB = CA, XA2 X. ) UCB, X B2 X. But TTCA;UB; ) = CA, WBI, ) & CA, WB2) W. \_\_ Arnez. By properties of certesian product

Arnez. CA, XA2) UCA, XB2) UCB, XB2) UCB, XB2) Then this becomes all permutations of A; and B; as addres toms so = CA, xA2x - A)VCA, xB2X -and so on, but A=A, xA2X... and B=B,xB2x are in this set so AUBE TTCA: UB; )

Note this holds if A or B is empty as wlog A is empty AUB=B but JA; s.t. A; UB; = B; you slill have the rest of the product. Note if Bi=b this still holds as then B= \$\phi\$ For TTCA; NB; ) = ANB
Using a similar argument as earlier, for an i TTCA: nB:)=(A, xA2X. xA1) nCA, xB2X. ) CB, xBax. xBi) But the intersect on restricts B So unticainsi) = ANB - - ANB here This holds regardless if A an Bis empty
as if an A is amply then AinBi = \$\phi\$ so Tt CA inBi) = \$\phi\$
and A = \$\phi\$ so AnB = \$\phi\$.