

# Chapter Six

## Planar Graphs

### 6.1 Planar Graphs :-

**Def :-** A graph or multi graph which can be drawn in the plane so that its edges do not cross is said to be planar .

**EX (1) :-**



Fig (1)

**\*Maps :-**

**Def :-** A particular planar representation of a finite planar multi graph is called a map . A given map divides the plane into various regions .

**EX (2) :-**

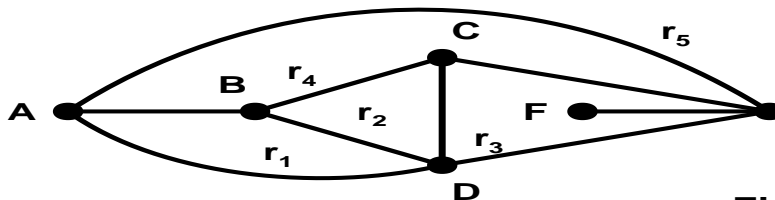


Fig (2)

The map with six vertices and nine edges divides the plane into five regions .

**Theorem 6.1 :-** The sum of the degrees of the regions of a map is equal to twice the number of edges .

The degree of the regions of fig (2) :-

$\deg(r_1) = 3$  ,  $\deg(r_2) = 3$  ,  $\deg(r_3) = 5$  ,  $\deg(r_4) = 4$  ,  $\deg(r_5) = 3$

The sum of the degree is 18 , which , as expected , is twice the number of edges .

## 6.2 Euler's Formula :-

Theorem :- ( Euler formula )

$$V - E + R = 2$$

Where  $V$  is the number of vertices ,  $E$  is the number of edges , and  $R$  is the number of

regions For Fig (2) .  $6 - 9 + 5 = 2$

Theorem 6.3 :- let  $G$  be a connected planar graph with  $p$  vertices and  $q$  edges , where

$$P \geq 3 . \text{ Then } q \leq 3p - 6$$

Proof :- let  $r$  be the number of regions in a planar representation of  $G$  by Euler formula .

$$P - q + r = 2$$

Now the sum of the degree of the regions equals  $2q$  by .

Theorem (6.1) :- [ The sum of the degree of the regions of a map is equal to twice the number of edges )

But each regions has degree 3 or more : hence .

$$2q \geq 3r$$

Thus  $r \leq 2q / 3$  . substituting this in Euler s formula gives .

$$2 = p - q + r \leq p - q + 2q / 3$$

$$\text{Or } 2 \leq p - q / 3$$

Multiplying the inequality by 3 gives  $6 \leq 3P - q$

Which gives us our result .  $q \leq 3P - 6$

EX(3) :- For Fig (3) :  $k_5$

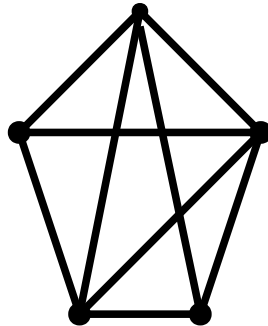


Fig (3)

$$P = 5, q = 10$$

$$q \leq 3p - 6 \rightarrow 10 \leq 3(5) - 6 = 9$$

which is impossible . Thus is non planar .

Theorem : ( kuratowski ) :-

A graph is non planar if and only if it contains a sub graph homeomorphic to  $K_{3,3}$  or  $K_5$ .

### 6.3 Colored Graphs :-

A vertex coloring , or simply coloring , of a graph  $G$  is an assignment of colors to the vertices of  $G$  such that adjacent vertices have different colors.

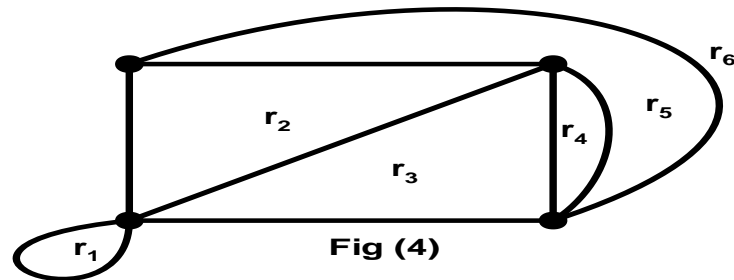
Def :- we say that  $G$  is  $n$  – colorable if there exists a coloring of  $G$  which uses  $n$  colors ( since the word “ color “ is used a noun ).

Theorem :- A planar graph  $G$  is 5 – colorable .

### 6.4 Four color theorem :-

Consider a map  $M$  ( i.e a planar representation of a finite planar multi graph ) . Two regions of  $M$  are said to be adjacent if they have an edge in common .

**EX(4):-** In the Fig (4)



The regions  $r_2$  and  $r_3$  are adjacent but the regions  $r_3$  and  $r_5$  are not . By a coloring of  $M$  we mean an assignment of a color to each region of  $M$  such that adjacent regions have different colors . A map  $M$  is  $n$  – colorable if there exists a coloring of  $M$  which uses  $n$  colors . The map in Fig (4) is 3 – colorable since the region could be painted as follows :-

$r_1$  red ,  $r_2$  white ,  $r_3$  red ,  $r_4$  white ,  $r_5$  red ,  $r_6$  blue .

### **6.5 Trees :-**

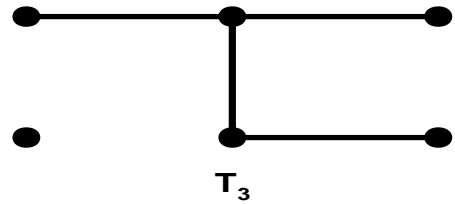
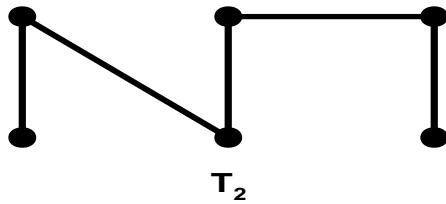
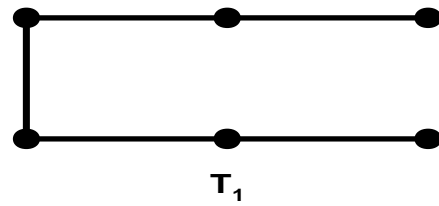
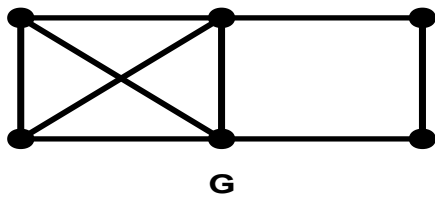
**Def :-** A tree is a connected graph with no cycles .

**Theorem :-** Trees are 2 – colorable

**Def :-** A sub graph  $T$  of a graph  $G$  is called a spanning tree of  $G$  if  $T$  is a tree and  $T$  include all the vertices of  $G$

**Def :-** The minimal spanning tree of  $G$  is a spanning tree of  $G$  such that the sum of the length of the edges is minimal among all spanning trees of  $G$  .

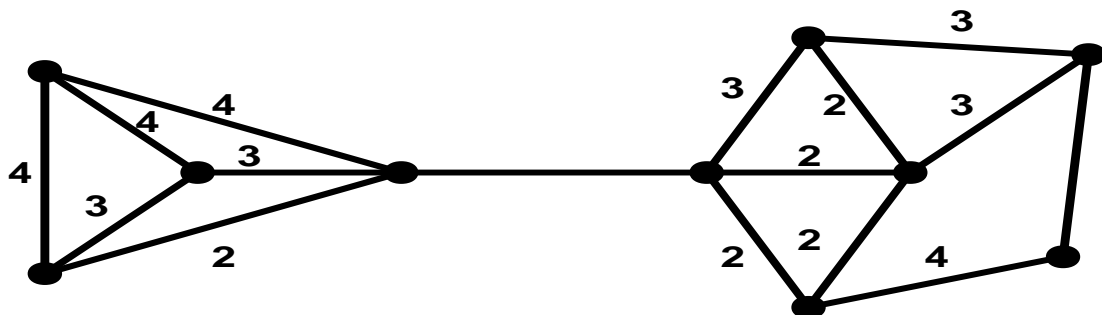
EX (5) :-



$G$  : is a graph

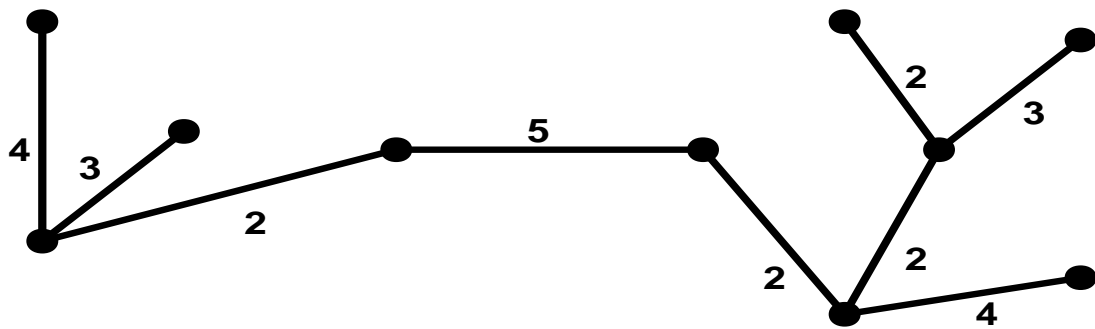
$T_1, T_2, T_3$  : are spanning tree of  $G$

EX (6) :- Find the minimal spanning tree of  $G$



$G$  ( graph )

Sol :-



T ( minimal spanning tree )

## 6.6 Rooted Trees :-

**Def :-** A rooted tree  $R$  consists of a tree graph together with a designated vertex  $r$  called the root of the tree .

**\*Level :-** The length of the path from the root  $r$  to  $v$  is called the level or depth of  $v$

**\*Leaf :-** Those vertices with degree one , other than  $r$  , are called the leaves of the rooted tree .

**\*Branch :-** A directed path from a vertex to a leaf is called a branch .

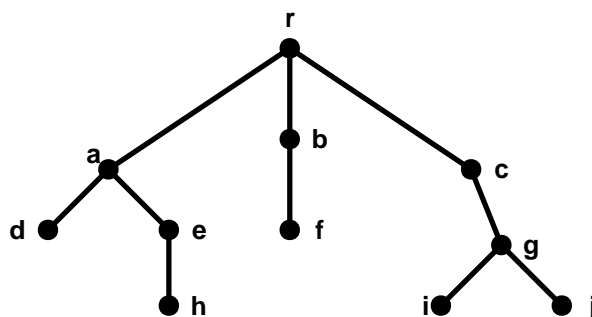


Fig (5)

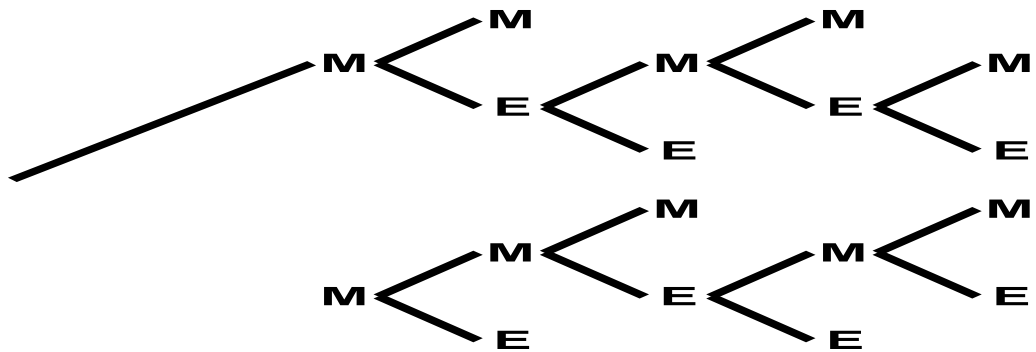
- Fig (5) shows a rooted tree :-

The root  $r$  is at the top of the tree .

The tree has five leaves , d , f , h , i , and j .

- The level of a is 1 , the level of f is 2 and the level of j is 3 .

EX (7) :- Fig (6) shows the various ways , the tournament can proceed find the leaves which correspond to the ways .



Sol :-

Observe that there ten leaves which correspond to the ten ways

MM , MEMM , MEMEM , MEME E , MEE , EMM , EMEMM , EMEME ,  
EMEE , EE

