

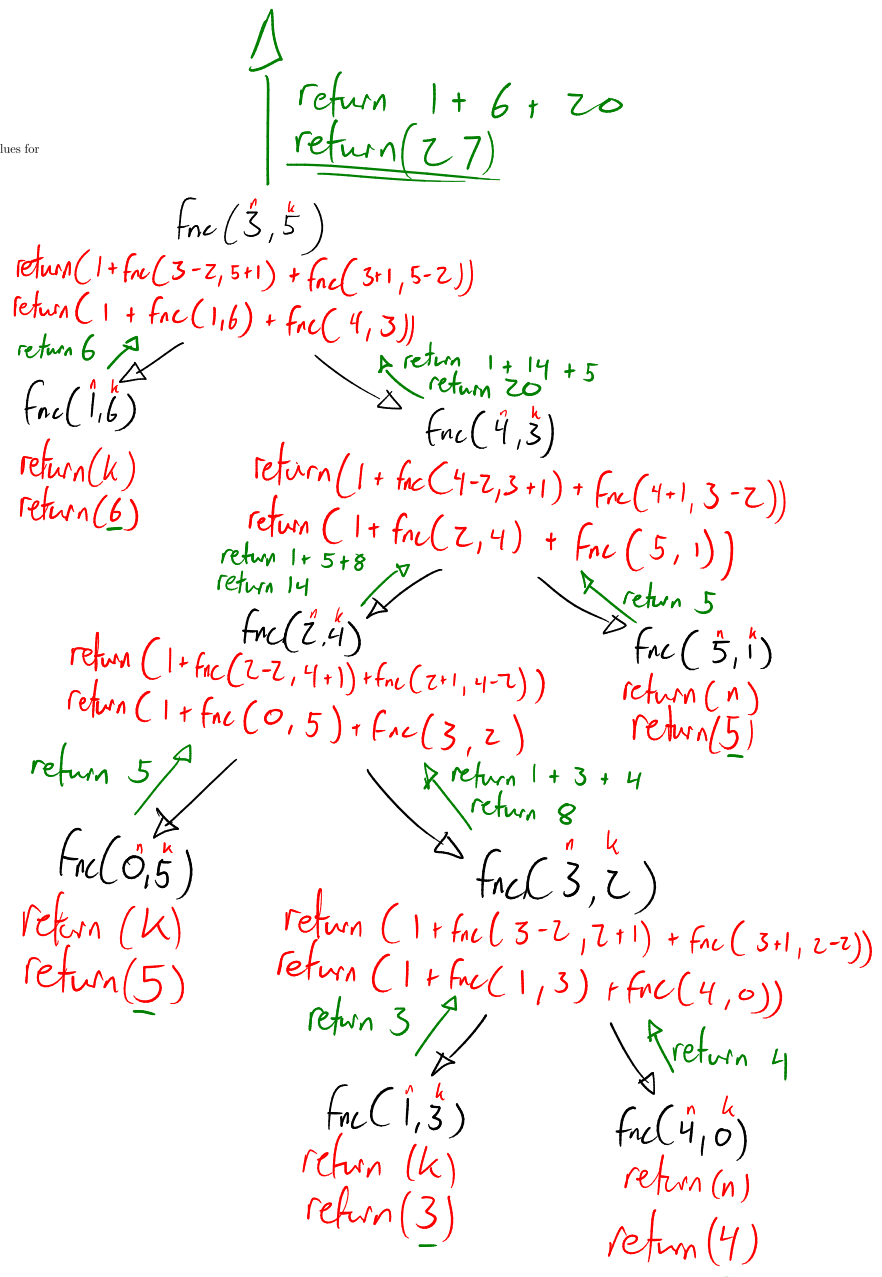
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signed: Ar. 15/2/2021

1. (10 marks)

```
public static int fnc(int n, int k) {
    if ( n <= 1 ) { return k; }
    else if ( k <= 1 ) { return n; }
    else { return (1 + fnc(n-2,k+1) + fnc(n+1,k-2)); }
}
```

assuming there is a closing bracket:



2. (10 marks)

Apply the parentheses matching algorithm discussed in Lecture 4 (slide 15) to the following input sequence of tokens:

$$[(c * a) + \{(\{c - b\} * d) / (d + \{((f * g) + h) * b\})\}]$$

There are three types of parentheses in this sequence: (), [], and { }. This algorithm checks, using a stack, if all parentheses in the input sequence properly match. Show the content of the stack when the algorithm terminates, indicating clearly which end of the stack is the top.

You do not have to show intermediate steps – show only the final stack.

top \longrightarrow

return false

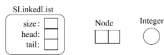
(
{
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$$\cancel{[(c * a) + \{(\{c - b\} * d) / (d + \{((f * g) + h) * b\})\}]}$$

3. (10 marks)

This question refers to the `SLinkedList<E>` class from Lecture 2 (discussed in lectures and L&T). We are assuming that this class has methods `size`, `elementAtHead`, `elementAtTail`, `insertAtHead`, `insertAtTail`, and `removeAtHead`.

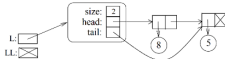
Using the following graphical representation of objects of classes `SLinkedList`, `Node` and `Integer`:



and indicating references to objects using arrows and null references using crosses, the state of the variables and objects after the following code is executed:

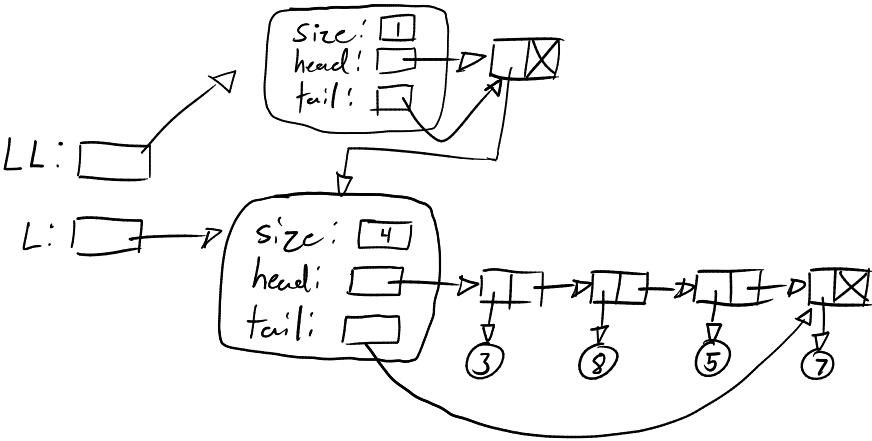
```
SLinkedList<Integer> L = new SLinkedList<Integer>();
L.insertAtHead(8);
L.insertAtHead(0);
SLinkedList<SLinkedList<Integer>> LL;
```

can be represented by the following diagram:



Draw the full diagram representing the state of all variables and objects after the following additional code is executed:

```
LL = new SLinkedList<SLinkedList<Integer>>();
LL.insertAtHead(L);
LL.insertAtHead(new SLinkedList<Integer>());
LL.elementAtTail().insertAtTail(7);
LL.elementAtHead().insertAtHead(3);
```



```
public static <E> void compress(Stack<E> s1, Stack<E> s2) {
```

```
    // record initial size to know where to stop
    int initialSize = s2.size();
```

```
    // remove nulls and push not nulls into s2
    while(!s1.isEmpty()) {
```

```
        if(s1.top() != null) {
            s2.push(s1.pop());
        } else {
            s1.pop();
        }
    }
```

```
    // push s1 values that are currently in s2 back into s1
    while(s2.size() != initialSize){
        s1.push(s2.pop());
    }
}
```

4. (10 marks) Give code for the method compress in the class Stack shown below. This method has two arguments, the "main" stack s1 and an auxiliary stack s2, both of the generic type Stack<E>. This generic type Stack<E> is the interface discussed in Lecture 4 (and is included in the net.datastructures package) for the Goodrich-Tamassia-Goldwasser textbook). The objective of the method compress is to remove all null elements from the stack s1. The remaining (non-null) elements should be kept on s1 in their initial order. The stack s2 should be used as a temporary storage for the elements from s1. At the end of the computation of this method, stack s2 should have the same content as at the beginning of the computation. The method should not use any arrays.

Now the method main below for an example of the expected behaviour of the method compress. Remark: The method compress only knows that s1 and s2 are implementations of the interface Stack<E>, but does not know anything about the details of those implementations. This means that method compress can access s1 and s2 only by using the interface methods.

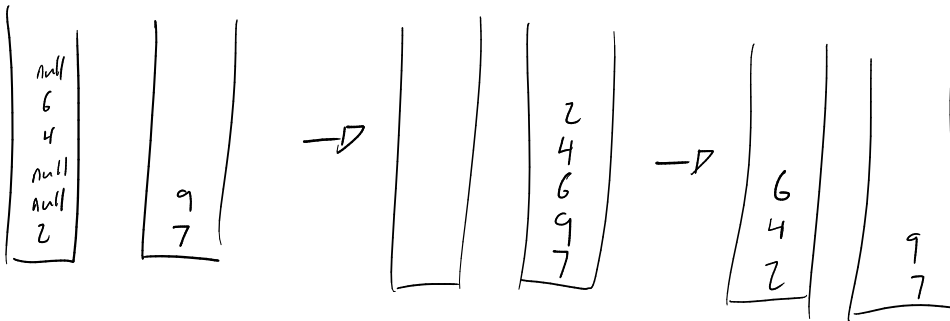
```
package lab04TheCoursework.cw1;

import net.datastructures.Stack;
import net.datastructures.ArrayStack;

public class Stack {

    public static <E> void compress(Stack<E> s1, Stack<E> s2) {
        ... // YOUR CODE REPLACES THIS HERE
    }

    public static void main(String[] args) {
        // test method compress
        Stack<Integer> S = new ArrayStack<Integer>(10);
        S.push(2); S.push(null); S.push(null); S.push(4); S.push(6); S.push(null);
        Stack<Integer> X = new ArrayStack<Integer>(10);
        X.push(7); X.push(9);
        System.out.println("stack S: " + S);
        // print: "stack S: [2, null, null, 4, 6, null]"
        System.out.println("stack X: " + X);
        // print: "stack X: [7, 9]"
        compress(S, X);
        System.out.println("stack S: " + S);
        // should print: "stack S: [2, 4, 6]"
        System.out.println("stack X: " + X);
        // should print: "stack X: [7, 9]"
    }
}
```



5. (10 marks)

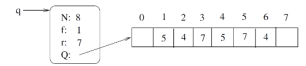
This question refers to the array-based "Circular Queue" implementation of the Queue data structure included in Lecture 4. Assume that class `CircularQueue<E>` is this implementation in Java. The statement:

```
Queue<Integer> q = new CircularQueue<Integer>(8);
```

creates the empty queue `q`, which has the internal representation shown in the following diagram:



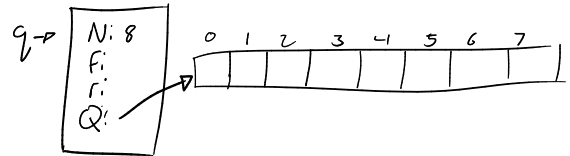
After seven **enqueue** operations and one **dequeue** operation the internal representation of the queue `q` is:



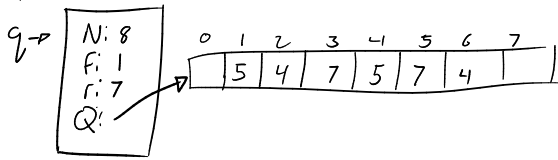
Trace now the computation of the following code:

```
while ( (q.size() > 1) && (q.front() < 10) ) {
    q.enqueue(q.dequeue() + q.dequeue());
}
```

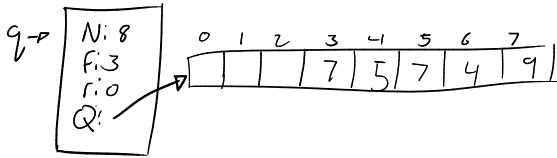
and show the **internal** representation of the queue `q` (in the form as in the above diagrams) after this computation.
Remark: you do not need to implement the class `CircularQueue<E>`.



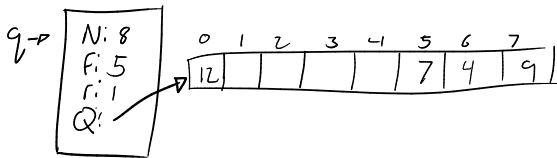
initial state:



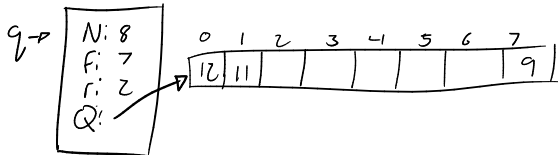
round 1:



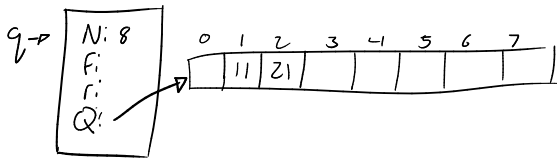
round 2:



round 3:



round 4:



STOP