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Quadratic residue codes

8.1 INTRODUCTION

Let p be an odd prime. Let Q denote the set of all quadratic residues mod p and N the set of all quadratic non-residues mod p . Let s be another prime which is a quadratic residue mod p . Then $s \in Q$ and it follows (from Proposition 5.6) that Q is closed with respect to multiplication by s . Therefore, Q is partitioned as a disjoint union of cyclotomic cosets modulo p under multiplication by s . Similarly, N is partitioned as a union of cyclotomic cosets modulo p under multiplication by s . Let α be a primitive p th root of unity in some extension of the field $\text{GF}(s)$. By Euler's theorem, there exists a positive integer m such that $s^m \equiv 1 \pmod{p}$. Let ρ be a primitive element of an extension $\text{GF}(s^m)$ of $\text{GF}(s)$ of degree m . We may then take

$$\alpha = \rho^{(s^m - 1)/p}$$

It follows from Theorem 7.3 that

$$q(x) = \prod_{i \in Q} (x - \alpha^i) \quad n(x) = \prod_{j \in N} (x - \alpha^j) \quad (8.1)$$

are polynomials with coefficients in $\text{GF}(s)$.

Lemma 8.1

$$x^p - 1 = (x - 1)q(x)n(x)$$

Proof

As every α^i is a p th root of unity, every root of $q(x)$ and every root of $n(x)$ is a root of $x^p - 1$. Therefore, $q(x)$ and $n(x)$ divide $x^p - 1$. Also $Q \cap N = \emptyset$ and $q(x)n(x)|x^p - 1$. Clearly $x - 1|x^p - 1$ and 1 is neither a root of $q(x)$ nor of $n(x)$. Therefore

$$(x - 1)q(x)n(x)|(x^p - 1)$$

But both polynomials are monic and of the same degree p . Therefore

$$x^p - 1 = (x - 1)q(x)n(x)$$

Set

$$\mathcal{R} = \text{GF}(s)[x]/\langle x^p - 1 \rangle$$

where $\langle x^p - 1 \rangle$ denotes the ideal of $\text{GF}(s)[x]$ generated by $x^p - 1$.

Definition 8.1

Quadratic residue codes \mathcal{F} , \mathcal{N} , $\bar{\mathcal{F}}$ and $\bar{\mathcal{N}}$ are the cyclic codes of length p over $\text{GF}(s)$ generated by the polynomials $q(x)$, $n(x)$, $(x - 1)q(x)$ and $(x - 1)n(x)$ respectively, i.e. these are the ideals in \mathcal{R} generated by the respective polynomials.

It is clear that

$$\bar{\mathcal{F}} \subseteq \mathcal{F} \quad \text{and} \quad \bar{\mathcal{N}} \subseteq \mathcal{N}$$

As

$$\text{degree of } q(x) = \text{degree of } n(x) = \frac{p-1}{2}$$

both \mathcal{F} and \mathcal{N} are linear codes over $\text{GF}(s)$ of dimension

$$p - \frac{p-1}{2} = \frac{p+1}{2}$$

each. Similarly, $\bar{\mathcal{F}}$ and $\bar{\mathcal{N}}$ are linear codes over $\text{GF}(s)$ of dimension $(p-1)/2$ each.

Let $s = 2$ and $p = 7$. Then, 1, 2, 4 are quadratic residues mod 7 and 3, 5, 6 are quadratic non-residues mod 7. Consider the field

$$\mathbb{B}[x]/\langle x^3 + x + 1 \rangle$$

of order 8. Then

$$\alpha = x + \langle x^3 + x + 1 \rangle$$

is a primitive 7th root of unity having $x^3 + x + 1$ as its minimal polynomial over \mathbb{B} . Now

$$\begin{aligned} q(x) &= (x + \alpha)(x + \alpha^2)(x + \alpha^4) \\ &= x^3 + x^2(\alpha + \alpha^2 + \alpha^4) + x(\alpha^3 + \alpha^5 + \alpha^6) + 1 \\ &= x^3 + x[\alpha^3 + \alpha^2(\alpha + 1) + (\alpha + 1)^2] + 1 \\ &= x^3 + x(\alpha^2 + \alpha^2 + 1) + 1 \\ &= x^3 + x + 1 \end{aligned}$$

So, the quadratic residue code \mathcal{F} of length 7 over \mathbb{B} is generated by

$$x^3 + x + 1 + \langle x^7 + 1 \rangle = q(x) + \langle x^7 + 1 \rangle$$

Observe that

$$x^7 + 1 = (x + 1)(x^3 + x + 1)(x^3 + x^2 + 1)$$

and, therefore, the other quadratic residue code \mathcal{N} is generated by

$$x^3 + x^2 + 1 = n(x)$$

Now

$$\begin{aligned} q(x^3) + \langle x^7 + 1 \rangle &= (x^3)^3 + x^3 + 1 + \langle x^7 + 1 \rangle \\ &= x^9 + x^3 + 1 + \langle x^7 + 1 \rangle \\ &= n(x) + \langle x^7 + 1 \rangle \in \mathcal{N} \\ q(x^5) + \langle x^7 + 1 \rangle &= (x^5)^3 + x^5 + 1 + \langle x^7 + 1 \rangle \\ &= x^{15} + x^5 + 1 + \langle x^7 + 1 \rangle \\ &= (x^{10} + x^4 + x^2) + (x^4 + x^3 + x) \\ &\quad + x^3 + x^2 + 1 + \langle x^7 + 1 \rangle \\ &= (x^3 + x^2 + 1)(x^2 + x + 1) + \langle x^7 + 1 \rangle \\ q(x^6) + \langle x^7 + 1 \rangle &= x^{12} + x^6 + 1 + \langle x^7 + 1 \rangle \\ &= (x^3 + x^2 + 1)^2 + \langle x^7 + 1 \rangle \end{aligned}$$

Thus

$$q(x^n) + \langle x^7 + 1 \rangle \in \mathcal{N}$$

whenever n is any quadratic non-residue mod 7.

Observe that the map $x \rightarrow x^5$ maps

$$x^6 + x^2 + 1 + \langle x^7 + 1 \rangle = (x^3 + x + 1)^2 + \langle x^7 + 1 \rangle \in \mathcal{F}$$

onto

$$\begin{aligned} (x^5)^6 + (x^5)^2 + 1 + \langle x^7 + 1 \rangle &= x^{30} + x^{10} + 1 + \langle x^7 + 1 \rangle \\ &= x^3 + x^2 + 1 + \langle x^7 + 1 \rangle \end{aligned}$$

which generates \mathcal{N} and the map $x \rightarrow x^6$ maps

$$x^5 + x^4 + 1 + \langle x^7 + 1 \rangle = (x^3 + x + 1)(x^2 + x + 1) + \langle x^7 + 1 \rangle \in \mathcal{F}$$

onto the generator

$$x^3 + x^2 + 1 + \langle x^7 + 1 \rangle$$

of \mathcal{N} . Thus for every non-residue n modulo 7, there is an element

$$a(x) + \langle x^7 + 1 \rangle \in \mathcal{F}$$

which maps onto the generator

$$x^3 + x^2 + 1 + \langle x^7 + 1 \rangle$$

of \mathcal{N} . Also the map $x \rightarrow x^n$ determines a permutation

$$\sigma: \{0, 1, 2, \dots, 6\} \rightarrow \{0, 1, 2, \dots, 6\}$$

as $\text{g.c.d.}(n, 7) = 1$. For example, for $n = 3$, the permutation determined is

$$\sigma = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 \\ 0 & 3 & 6 & 2 & 5 & 1 & 4 \end{pmatrix} = (1 \ 3 \ 2 \ 6 \ 4 \ 5)$$

In general

$$\sigma = \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 \\ 0 & \bar{1} & \bar{2} & \bar{3} & \bar{4} & \bar{5} & \bar{6} \end{pmatrix}$$

where for $1 \leq i \leq 6$, \bar{i} denotes the least non-negative remainder on dividing ni by 7. The degrees of the generators of \mathcal{F} and \mathcal{N} being 3 each, the number of elements in the two codes are the same. Hence, it follows that the quadratic residue codes \mathcal{F} and \mathcal{N} are equivalent. Similarly, the expurgated quadratic residue codes $\bar{\mathcal{F}}$ and $\bar{\mathcal{N}}$ generated by

$$(x + 1)(x^3 + x + 1) \quad \text{and} \quad (x + 1)(x^3 + x^2 + 1)$$

respectively are also equivalent.

Theorem 8.1

The quadratic residue codes \mathcal{F} and \mathcal{N} of length p over $\text{GF}(s)$ generated by $q(x)$ and $n(x)$ are equivalent. Also the expurgated quadratic residue codes $\bar{\mathcal{F}}$ and $\bar{\mathcal{N}}$ of length p over $\text{GF}(s)$ are equivalent.

Proof

Let n be a fixed quadratic non-residue mod p . Then there exists a positive integer r such that

$$nr \equiv 1 \pmod{p}$$

As 1 is always a quadratic residue, nr is a quadratic residue. But, then n being a non-residue it follows that r is a quadratic non-residue mod p . For any $i \in Q$, it follows that ir is a non-residue mod p . Now

$$q(x^n) = \prod_{i \in Q} (x^n - \alpha^i)$$

where α is a primitive p th root of unity in some extension of $\text{GF}(s)$. Also

$$nr \equiv 1 \pmod{p} \Rightarrow \alpha^i = \alpha^{nr^i} = (\alpha^{ri})^n$$

so that α^{ir} is a root of $q(x^n)$. This is so for every $i \in Q$ and so

$$n(x) | q(x^n)$$

Hence the map induced by $x \rightarrow x^n$ maps code words from \mathcal{F} onto code words in \mathcal{N} . Again

$$1 = nr + pt$$

for some integer t and, therefore,

$$x + \langle x^p - 1 \rangle = x^{nr+pt} + \langle x^p - 1 \rangle$$

If $t < 0$, then $-t > 0$ and

$$\begin{aligned} x^{-pt} + \langle x^p - 1 \rangle &= (x^p)^{-t} + \langle x^p - 1 \rangle \\ &= 1 + \langle x^p - 1 \rangle \end{aligned}$$

and so

$$\begin{aligned} x + \langle x^p - 1 \rangle &= (x^{nr+pt} + \langle x^p - 1 \rangle)(x^{-pt} + \langle x^p - 1 \rangle) \\ &= x^{nr} + \langle x^p - 1 \rangle \\ &= (x^n)^r + \langle x^p - 1 \rangle \end{aligned}$$

It follows that the map

$$F[x]/\langle x^p - 1 \rangle \rightarrow F[x]/\langle x^n - 1 \rangle$$

induced by $x \rightarrow x^n$ is onto and hence one-one as well. Therefore, the restriction of this map to

$$\mathcal{F} \rightarrow \mathcal{N}$$

is also one-one and the two spaces being of the same dimension, the map is onto as well. Furthermore, the map $x \rightarrow x^n$ determines the permutation σ of the set $\{0, 1, 2, \dots, p-1\}$ given by

$$\sigma = \begin{pmatrix} 0 & 1 & 2 & \cdots & p-1 \\ 0 & \bar{1} & \bar{2} & \cdots & \frac{p-1}{p-1} \end{pmatrix}$$

where for $1 \leq i \leq p-1$, \bar{i} denotes the least non-negative remainder where ni is divided by p . Thus \mathcal{F} and \mathcal{N} are equivalent.

Equivalence of $\bar{\mathcal{F}}$ and $\bar{\mathcal{N}}$ follows similarly.

8.2 SOME EXAMPLES OF QUADRATIC RESIDUE CODES

Consider the $[7, 4, 3]$ binary Hamming code \mathcal{C} . A generator matrix of this code is (see p. 115)

$$\mathbf{G} = \begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 1 \end{pmatrix}$$

Consider the field

$$\text{GF}(2^3) = \mathbb{B}[X]/\langle x^3 + x + 1 \rangle$$

Then,

$$\alpha = x + \langle x^3 + x + 1 \rangle$$

is a primitive 7th root of unity. Now 1, 2, 4 are quadratic residues modulo 7 and so $q(x)$ is a cubic polynomial having α as a root. Also, the minimal polynomial