

How many bit operations does this take? In each step you have either 1 or 2 multiplications of numbers which are less than m^2 . And there are $k - 1$ steps. Since each step takes $O(\log^2(m^2)) = O(\log^2 m)$ bit operations, we end up with the following estimate:

Proposition I.3.6. $\text{Time}(b^n \bmod m) = O((\log n)(\log^2 m))$.

Remark. If n is very large in Proposition I.3.6, you might want to use the corollary of Proposition I.3.5, replacing n by its least nonnegative residue modulo $\varphi(m)$. But this requires that you know $\varphi(m)$. If you do know $\varphi(m)$, and if $\text{g.c.d.}(b, m) = 1$, so that you can replace n by its least nonnegative residue modulo $\varphi(m)$, then the estimate on the right in Proposition I.3.6 can be replaced by $O(\log^3 m)$.

As a final application of the multiplicativity of the Euler φ -function, we prove a formula that will be used at the beginning of Chapter II.

Proposition I.3.7. $\sum_{d|n} \varphi(d) = n$.

Proof. Let $f(n)$ denote the left side of the equality in the proposition, i.e., $f(n)$ is the sum of $\varphi(d)$ taken over all divisors d of n (including 1 and n). We must show that $f(n) = n$. We first claim that $f(n)$ is multiplicative, i.e., that $f(mn) = f(m)f(n)$ whenever $\text{g.c.d.}(m, n) = 1$. To see this, we note that any divisor $d|mn$ can be written (in one and only one way) in the form $d_1 \cdot d_2$, where $d_1|m$, $d_2|n$. Since $\text{g.c.d.}(d_1, d_2) = 1$, we have $\varphi(d) = \varphi(d_1)\varphi(d_2)$, because of the multiplicativity of φ . We get all possible divisors d of mn by taking all possible pairs d_1, d_2 where d_1 is a divisor of m and d_2 is a divisor of n . Thus, $f(mn) = \sum_{d_1|m} \sum_{d_2|n} \varphi(d_1)\varphi(d_2) = \left(\sum_{d_1|m} \varphi(d_1)\right) \left(\sum_{d_2|n} \varphi(d_2)\right) = f(m)f(n)$, as claimed. Now to prove the proposition suppose that $n = p_1^{\alpha_1} \cdots p_r^{\alpha_r}$ is the prime factorization of n . By the multiplicativity of f , we find that $f(n)$ is a product of terms of the form $f(p^\alpha)$. So it suffices to prove the proposition for p^α ; i.e., to prove that $f(p^\alpha) = p^\alpha$. But the divisors of p^α are p^j for $0 \leq j \leq \alpha$, and so $f(p^\alpha) = \sum_{j=0}^{\alpha} \varphi(p^j) = 1 + \sum_{j=1}^{\alpha} (p^j - p^{j-1}) = p^\alpha$. This proves the proposition for p^α ; and hence for all n .

Exercises

1. Describe all of the solutions of the following congruences:

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|--------------------------------|------------------------------------|
| (a) $3x \equiv 4 \pmod{7}$; | (d) $27x \equiv 25 \pmod{256}$; |
| (b) $3x \equiv 4 \pmod{12}$; | (e) $27x \equiv 72 \pmod{900}$; |
| (c) $9x \equiv 12 \pmod{21}$; | (f) $103x \equiv 612 \pmod{676}$. |

- What are the possibilities for the last hexadecimal digit of a perfect square? (See Exercise 7 of §I.1.)
- What are the possibilities for the last base-12 digit of a product of two consecutive positive odd numbers?