

Suppose now that ϕ has length greater than 1 and that every combinator ψ shorter than ϕ is such that $\lambda_x\psi$ is equal to a combinator. We claim that $\lambda_x\phi$ is equal to a combinator. Indeed, $\phi = \psi^*\chi$, where ψ and χ are shorter than ϕ , hence

$$\lambda_x\phi = \lambda_x(\psi^*\chi) = (S^*(\lambda_x\psi))^*(\lambda_x\chi).$$

This is a combinator because $\lambda_x\psi$ and $\lambda_x\chi$ are combinators by our inductional assumption.

Using the methods of this proof, we can express 2 as a combinator:

$$\begin{aligned} 2 &= \lambda_f \lambda_x (f^*(f^*x)) \\ &= \lambda_f ((S^* \lambda_x f)^* \lambda_x (f^*x)) \\ &= \lambda_f ((S^* \lambda_x f)^* f) \quad \text{by R2} \\ &= (S^* (\lambda_f (S^* \lambda_x f)))^* \lambda_f f \\ &= (S^* (\lambda_f (S^* (K^* f))))^* I \\ &= (S^* ((S^* \lambda_f S)^* \lambda_f (K^* f)))^* I \\ &= (S^* ((S^* (K^* S))^* K))^* I \quad \text{by R2.} \end{aligned}$$

Exercises

1. Assuming that m, p and q are natural numbers expressed in the lambda calculus, show that $(m^p)^q = m^{(q \times p)}$.
2. Prove that $0^{\Sigma^* n} = 0$.
3. Prove that $I = (S^* K)^* K$.
4. Express m^n , $m \times n$, and $m + n$ in terms of I, K, S, m and n .

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Logic from Aristotle to Russell

Logic was not always regarded as a branch of mathematics, certainly not by Aristotle (384–322 BC), who was the first to write about logic in the West. Among the principles which he recognized are the following:

$$\begin{aligned}\neg\neg p &\iff p && \text{(double negation),} \\ p \vee \neg p & && \text{(excluded third),} \\ (p \Rightarrow q) &\iff (\neg q \Rightarrow \neg p) && \text{(contraposition).}\end{aligned}$$

He also looked at modal logic and showed how possibility can be defined in terms of necessity.

Aristotle's major concern was with a type of argument called the 'syllogism', which predominated in logical thinking for the next two thousand years. It dealt with four types of basic statements:

- SaP meaning *all S are P*,
- SeP meaning *no S are P*,
- SiP meaning *some S are P*,
- SoP meaning *some S are not P*.

He realized that PeS is equivalent to SeP and that PiS is equivalent to SiP and he adhered to a convention that SaP implies SiP. (Today we use

words differently: we assert that *all* unicorns have horns, but deny that *some* unicorns have horns. Evidently Aristotle did not believe in the empty set.)

A *syllogism* is an argument which infers one such basic statement from two others. Here are the first four ‘figures’ of the syllogism:

$$\begin{array}{c} \text{MaP} \\ \text{SaM} \\ \hline \text{SaP} \end{array} \qquad \begin{array}{c} \text{MeP} \\ \text{SaM} \\ \hline \text{SeP} \end{array} \qquad \begin{array}{c} \text{MaP} \\ \text{SiM} \\ \hline \text{SiP} \end{array} \qquad \begin{array}{c} \text{MeP} \\ \text{SiM} \\ \hline \text{SoP} \end{array}$$

William of Shyreswood (1250 AD) gave these syllogisms the names

barbara, celarent, darii, ferio,

— to make them easier to remember. There were more such figures, which we shall not discuss here. Here is a typical argument illustrating the ‘ferio’:

$$\begin{array}{c} \text{no minister is prudent} \\ \text{some socialists are ministers} \\ \hline \\ \text{some socialists are not prudent} \end{array}$$

The Stoics (200 BC), Philo of Megara in particular, essentially introduced truth tables into logic, thus anticipating Ludwig Wittgenstein (1889–1951). They discussed the problem of whether ‘*p or q*’ is true when both *p* and *q* are true and whether ‘if *p* then *q*’ is always true when *p* is false. They arrived at the modern conventions, expressed in the following *truth tables*:

$p \vee q$	$p \Rightarrow q$
—	—
T T T	T T T
T T F	T F F
F T T	F T T
F F F	F T F

The Stoics were committed to the view that there are only two truth values (T and F). In particular, they believed that a statement like ‘there will be a battle tomorrow’ is either true or false, although Aristotle seems to have had some second thoughts about this. This belief was associated in