

- d. Same as conditions a–c with the roles of the two countries reversed.
- e. The verification devices in both countries must be identical, and must be constructed jointly by scientists from both countries.
4. The purpose of this problem is to construct a long-distance coin flip using any two-to-one trapdoor function. For example, suppose that two chess players at distant parts of the world are playing chess by mail or telephone and want a fair way to determine who plays white. Or suppose that when making preparations for an international ice-hockey match, representatives of the two teams decide to flip a coin to see which country hosts the match, without having to arrange a meeting (or trust a third party) to “flip the coin.”

By a system of two-to-one trapdoor functions, we mean an algorithm which, given a key K_E of a suitable type, constructs a function $f: \mathcal{P} \rightarrow \mathcal{C}$ such that every element c in the image of f has exactly two preimages $p_1, p_2 \in \mathcal{P}$ such that $f(p_j) = c$; and an algorithm which, given a key K_D which “reverses K_E ,” can find both preimages of any c in the image of f . Here we assume that it is computationally infeasible to find K_D knowing only K_E . Given an element $p_1 \in \mathcal{P}$, notice that one can find the other element p_2 having the same image if one knows both K_E and K_D (namely, find both inverses of $f(p_1)$); but we assume that, knowing only K_E , one cannot feasibly compute the companion element p_2 for any p_1 at all.

Suppose that Player A (Aniuta) and Player B (Björn) want to use this set-up to flip a coin. Aniuta generates a pair of keys K_E and K_D and sends K_E (but *not* K_D) to Björn. Explain a procedure that has a 50%–50% chance of each player “winning” (give a suitable definition of “winning”), and that has adequate safeguards against cheating.

References for § IV.1

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