

"Without the key, sir, excuse me if I believe the thing impossible."

"Do you wish me to name your key, madame?"

"If you please."

I then told her the key-word, which belonged to no language, and I saw her surprise. She told me that it was impossible, for she believed herself the only possessor of that word which she kept in her memory and which she had never written down.

I could have told her the truth — that the same calculation which had served me for deciphering the manuscript had enabled me to learn the word — but on a caprice it struck me to tell her that a genie had revealed it to me. This false disclosure fettered Madame d'Urfé to me. That day I became the master of her soul, and I abused my power. Every time I think of it, I am distressed and ashamed, and I do penance now in the obligation under which I place myself of telling the truth in writing my memoirs.

— Casanova, 1757, quoted in D. Kahn's *The Codebreakers*

The situation persisted for another 220 years after this encounter between Casanova and Madame d'Urfé: knowledge of how to encipher and knowledge of how to decipher were regarded as essentially equivalent in any cryptosystem. However, in 1976 W. Diffie and M. Hellman discovered an entirely different type of cryptosystem and invented "public key cryptography."

By definition, a public key cryptosystem has the property that someone who knows only how to encipher cannot use the enciphering key to find the deciphering key without a prohibitively lengthy computation. In other words the enciphering function  $f: \mathcal{P} \rightarrow \mathcal{C}$  is easy to compute once the enciphering key  $K_E$  is known, but it is very hard in practice to compute the inverse function  $f^{-1}: \mathcal{C} \rightarrow \mathcal{P}$ . That is, from the standpoint of realistic computability, the function  $f$  is not invertible (without some additional information — the deciphering key  $K_D$ ). Such a function  $f$  is called a *trapdoor function*. That is, a trapdoor function  $f$  is a function which is easy to compute but whose inverse  $f^{-1}$  is hard to compute without having some additional auxiliary information beyond what is necessary to compute  $f$ . The inverse  $f^{-1}$  is easy to compute, however, for someone who has this information  $K_D$  (the "deciphering key").

There is a closely related concept of a *one-way* function. This is a function  $f$  which is easy to compute but for which  $f^{-1}$  is hard to compute and cannot be made easy to compute even by acquiring some additional information. While the notion of a trapdoor function apparently appeared for the first time in 1978 along with the invention of the RSA public-key cryptosystem, the notion of a one-way function is somewhat older. What seems to have been the first use of one-way functions for cryptography was