

How many bit operations does this take? In each step you have either 1 or 2 multiplications of numbers which are less than  $m^2$ . And there are  $k - 1$  steps. Since each step takes  $O(\log^2(m^2)) = O(\log^2 m)$  bit operations, we end up with the following estimate:

**Proposition I.3.6.** Time( $b^n \bmod m$ ) =  $O((\log n)(\log^2 m))$ .

**Remark.** If  $n$  is very large in Proposition I.3.6, you might want to use the corollary of Proposition I.3.5, replacing  $n$  by its least nonnegative residue modulo  $\varphi(m)$ . But this requires that you know  $\varphi(m)$ . If you do know  $\varphi(m)$ , and if  $\text{g.c.d.}(b, m) = 1$ , so that you can replace  $n$  by its least nonnegative residue modulo  $\varphi(m)$ , then the estimate on the right in Proposition I.3.6 can be replaced by  $O(\log^3 m)$ .

As a final application of the multiplicativity of the Euler  $\varphi$ -function, we prove a formula that will be used at the beginning of Chapter II.

**Proposition I.3.7.**  $\sum_{d|n} \varphi(d) = n$ .

**Proof.** Let  $f(n)$  denote the left side of the equality in the proposition, i.e.,  $f(n)$  is the sum of  $\varphi(d)$  taken over all divisors  $d$  of  $n$  (including 1 and  $n$ ). We must show that  $f(n) = n$ . We first claim that  $f(n)$  is multiplicative, i.e., that  $f(mn) = f(m)f(n)$  whenever  $\text{g.c.d.}(m, n) = 1$ . To see this, we note that any divisor  $d|mn$  can be written (in one and only one way) in the form  $d_1 \cdot d_2$ , where  $d_1|m$ ,  $d_2|n$ . Since  $\text{g.c.d.}(d_1, d_2) = 1$ , we have  $\varphi(d) = \varphi(d_1)\varphi(d_2)$ , because of the multiplicativity of  $\varphi$ . We get all possible divisors  $d$  of  $mn$  by taking all possible pairs  $d_1, d_2$  where  $d_1$  is a divisor of  $m$  and  $d_2$  is a divisor of  $n$ . Thus,  $f(mn) = \sum_{d_1|m} \sum_{d_2|n} \varphi(d_1)\varphi(d_2) = (\sum_{d_1|m} \varphi(d_1))(\sum_{d_2|n} \varphi(d_2)) = f(m)f(n)$ , as claimed. Now to prove the proposition suppose that  $n = p_1^{\alpha_1} \cdots p_r^{\alpha_r}$  is the prime factorization of  $n$ . By the multiplicativity of  $f$ , we find that  $f(n)$  is a product of terms of the form  $f(p^\alpha)$ . So it suffices to prove the proposition for  $p^\alpha$ , i.e., to prove that  $f(p^\alpha) = p^\alpha$ . But the divisors of  $p^\alpha$  are  $p^j$  for  $0 \leq j \leq \alpha$ , and so  $f(p^\alpha) = \sum_{j=0}^{\alpha} \varphi(p^j) = 1 + \sum_{j=1}^{\alpha} (p^j - p^{j-1}) = p^\alpha$ . This proves the proposition for  $p^\alpha$ , and hence for all  $n$ .

### Exercises

1. Describe all of the solutions of the following congruences:
  - (a)  $3x \equiv 4 \pmod{7}$ ; (d)  $27x \equiv 25 \pmod{256}$ ;
  - (b)  $3x \equiv 4 \pmod{12}$ ; (e)  $27x \equiv 72 \pmod{900}$ ;
  - (c)  $9x \equiv 12 \pmod{21}$ ; (f)  $103x \equiv 612 \pmod{676}$ .
2. What are the possibilities for the last hexadecimal digit of a perfect square? (See Exercise 7 of § I.1.)
3. What are the possibilities for the last base-12 digit of a product of two consecutive positive odd numbers?