

### 4.3. Lagrange Interpolation

Throughout this section we shall assume  $F$  is a fixed field and that  $t_0, t_1, \dots, t_n$  are  $n + 1$  distinct elements of  $F$ . Let  $V$  be the subspace of  $F[x]$  consisting of all polynomials of degree less than or equal to  $n$  (together with the 0-polynomial), and let  $L_i$  be the function from  $V$  into  $F$  defined for  $f$  in  $V$  by

$$L_i(f) = f(t_i), \quad 0 \leq i \leq n.$$

By part (i) of Theorem 2, each  $L_i$  is a linear functional on  $V$ , and one of the things we intend to show is that the set consisting of  $L_0, L_1, \dots, L_n$  is a basis for  $V^*$ , the dual space of  $V$ .

Of course in order that this be so, it is sufficient (cf. Theorem 15 of Chapter 3) that  $\{L_0, L_1, \dots, L_n\}$  be the dual of a basis  $\{P_0, P_1, \dots, P_n\}$  of  $V$ . There is at most one such basis, and if it exists it is characterized by

$$(4-11) \quad L_j(P_i) = P_i(t_j) = \delta_{ij}.$$

The polynomials

$$(4-12) \quad \begin{aligned} P_i &= \frac{(x - t_0) \cdots (x - t_{i-1})(x - t_{i+1}) \cdots (x - t_n)}{(t_i - t_0) \cdots (t_i - t_{i-1})(t_i - t_{i+1}) \cdots (t_i - t_n)} \\ &= \prod_{j \neq i} \left( \frac{x - t_j}{t_i - t_j} \right) \end{aligned}$$

are of degree  $n$ , hence belong to  $V$ , and by Theorem 2, they satisfy (4-11).

If  $f = \sum_i c_i P_i$ , then for each  $j$

$$(4-13) \quad f(t_j) = \sum_i c_i P_i(t_j) = c_j.$$

Since the 0-polynomial has the property that  $0(t) = 0$  for each  $t$  in  $F$ , it follows from (4-13) that the polynomials  $P_0, P_1, \dots, P_n$  are linearly independent. The polynomials  $1, x, \dots, x^n$  form a basis of  $V$  and hence the dimension of  $V$  is  $(n + 1)$ . So, the independent set  $\{P_0, P_1, \dots, P_n\}$  must also be a basis for  $V$ . Thus for each  $f$  in  $V$

$$(4-14) \quad f = \sum_{i=0}^n f(t_i) P_i.$$

The expression (4-14) is called **Lagrange's interpolation formula**. Setting  $f = x^i$  in (4-14) we obtain

$$x^i = \sum_{i=0}^n (t_i)^i P_i.$$

Now from Theorem 7 of Chapter 2 it follows that the matrix

$$(4-15) \quad \begin{bmatrix} 1 & t_0 & t_0^2 & \cdots & t_0^n \\ 1 & t_1 & t_1^2 & \cdots & t_1^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & t_n & t_n^2 & \cdots & t_n^n \end{bmatrix}$$

is invertible. The matrix in (4-15) is called a **Vandermonde matrix**; it is an interesting exercise to show directly that such a matrix is invertible, when  $t_0, t_1, \dots, t_n$  are  $n + 1$  distinct elements of  $F$ .

If  $f$  is any polynomial over  $F$  we shall, in our present discussion, denote by  $f\sim$  the polynomial function from  $F$  into  $F$  taking each  $t$  in  $F$  into  $f(t)$ . By definition (cf. Example 4, Chapter 2) every polynomial function arises in this way; however, it may happen that  $f\sim = g\sim$  for two polynomials  $f$  and  $g$  such that  $f \neq g$ . Fortunately, as we shall see, this unpleasant situation only occurs in the case where  $F$  is a field having only a finite number of distinct elements. In order to describe in a precise way the relation between polynomials and polynomial functions, we need to define the product of two polynomial functions. If  $f, g$  are polynomials over  $F$ , the product of  $f\sim$  and  $g\sim$  is the function  $f\sim g\sim$  from  $F$  into  $F$  given by

$$(4-16) \quad (f\sim g\sim)(t) = f\sim(t)g\sim(t), \quad t \text{ in } F.$$

By part (ii) of Theorem 2,  $(fg)(t) = f(t)g(t)$ , and hence

$$(fg)\sim(t) = f\sim(t)g\sim(t)$$

for each  $t$  in  $F$ . Thus  $f\sim g\sim = (fg)\sim$ , and is a polynomial function. At this point it is a straightforward matter, which we leave to the reader, to verify that the vector space of polynomial functions over  $F$  becomes a linear algebra with identity over  $F$  if multiplication is defined by (4-16).

**Definition.** Let  $F$  be a field and let  $\mathfrak{A}$  and  $\mathfrak{A}\sim$  be linear algebras over  $F$ . The algebras  $\mathfrak{A}$  and  $\mathfrak{A}\sim$  are said to be **isomorphic** if there is a one-to-one mapping  $\alpha \rightarrow \alpha\sim$  of  $\mathfrak{A}$  onto  $\mathfrak{A}\sim$  such that

$$(a) \quad (c\alpha + d\beta)\sim = c\alpha\sim + d\beta\sim$$

$$(b) \quad (\alpha\beta)\sim = \alpha\sim\beta\sim$$

for all  $\alpha, \beta$  in  $\mathfrak{A}$  and all scalars  $c, d$  in  $F$ . The mapping  $\alpha \rightarrow \alpha\sim$  is called an **isomorphism** of  $\mathfrak{A}$  onto  $\mathfrak{A}\sim$ . An isomorphism of  $\mathfrak{A}$  onto  $\mathfrak{A}\sim$  is thus a vector-space isomorphism of  $\mathfrak{A}$  onto  $\mathfrak{A}\sim$  which has the additional property (b) of 'preserving' products.

**EXAMPLE 4.** Let  $V$  be an  $n$ -dimensional vector space over the field  $F$ . By Theorem 13 of Chapter 3 and subsequent remarks, each ordered basis  $\mathfrak{B}$  of  $V$  determines an isomorphism  $T \rightarrow [T]_{\mathfrak{B}}$  of the algebra of linear operators on  $V$  onto the algebra of  $n \times n$  matrices over  $F$ . Suppose now that  $U$  is a fixed linear operator on  $V$  and that we are given a polynomial

$$f = \sum_{i=0}^n c_i x^i$$

with coefficients  $c_i$  in  $F$ . Then

$$f(U) = \sum_{i=0}^n c_i U^i$$

and since  $T \rightarrow [T]_{\mathfrak{A}}$  is a linear mapping

$$[f(U)]_{\mathfrak{A}} = \sum_{i=0}^n c_i [U^i]_{\mathfrak{A}}.$$

Now from the additional fact that

$$[T_1 T_2]_{\mathfrak{A}} = [T_1]_{\mathfrak{A}} [T_2]_{\mathfrak{A}}$$

for all  $T_1, T_2$  in  $L(V, V)$  it follows that

$$[U^i]_{\mathfrak{A}} = ([U]_{\mathfrak{A}})^i, \quad 2 \leq i \leq n.$$

As this relation is also valid for  $i = 0, 1$  we obtain the result that

$$(4-17) \quad [f(U)]_{\mathfrak{A}} = f([U]_{\mathfrak{A}}).$$

In words, if  $U$  is a linear operator on  $V$ , the matrix of a polynomial in  $U$ , in a given basis, is the same polynomial in the matrix of  $U$ .

**Theorem 3.** *If  $F$  is a field containing an infinite number of distinct elements, the mapping  $f \rightarrow f^{\sim}$  is an isomorphism of the algebra of polynomials over  $F$  onto the algebra of polynomial functions over  $F$ .*

*Proof.* By definition, the mapping is onto, and if  $f, g$  belong to  $F[x]$  it is evident that

$$(cf + dg)^{\sim} = df^{\sim} + dg^{\sim}$$

for all scalars  $c$  and  $d$ . Since we have already shown that  $(fg)^{\sim} = f^{\sim}g^{\sim}$ , we need only show that the mapping is one-to-one. To do this it suffices by linearity to show that  $f^{\sim} = 0$  implies  $f = 0$ . Suppose then that  $f$  is a polynomial of degree  $n$  or less such that  $f' = 0$ . Let  $t_0, t_1, \dots, t_n$  be any  $n + 1$  distinct elements of  $F$ . Since  $f^{\sim} = 0$ ,  $f(t_i) = 0$  for  $i = 0, 1, \dots, n$ , and it is an immediate consequence of (4-14) that  $f = 0$ . ■

From the results of the next section we shall obtain an altogether different proof of this theorem.

## Exercises

1. Use the Lagrange interpolation formula to find a polynomial  $f$  with real coefficients such that  $f$  has degree  $\leq 3$  and  $f(-1) = -6$ ,  $f(0) = 2$ ,  $f(1) = -2$ ,  $f(2) = 6$ .
2. Let  $\alpha, \beta, \gamma, \delta$  be real numbers. We ask when it is possible to find a polynomial  $f$  over  $R$ , of degree not more than 2, such that  $f(-1) = \alpha$ ,  $f(1) = \beta$ ,  $f(3) = \gamma$  and  $f(0) = \delta$ . Prove that this is possible if and only if

$$3\alpha + 6\beta - \gamma - 8\delta = 0.$$

3. Let  $F$  be the field of real numbers,

$$A = \begin{bmatrix} 2 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$p = (x - 2)(x - 3)(x - 1).$$

(a) Show that  $p(A) = 0$ .

(b) Let  $P_1, P_2, P_3$  be the Lagrange polynomials for  $t_1 = 2, t_2 = 3, t_3 = 1$ .

Compute  $E_i = P_i(A), i = 1, 2, 3$ .

(c) Show that  $E_1 + E_2 + E_3 = I, E_i E_j = 0$  if  $i \neq j, E_i^2 = E_i$ .

(d) Show that  $A = 2E_1 + 3E_2 + E_3$ .

4. Let  $p = (x - 2)(x - 3)(x - 1)$  and let  $T$  be any linear operator on  $R^4$  such that  $p(T) = 0$ . Let  $P_1, P_2, P_3$  be the Lagrange polynomials of Exercise 3, and let  $E_i = P_i(T), i = 1, 2, 3$ . Prove that

$$E_1 + E_2 + E_3 = I, \quad E_i E_j = 0 \quad \text{if } i \neq j,$$

$$E_i^2 = E_i, \quad \text{and} \quad T = 2E_1 + 3E_2 + E_3.$$

5. Let  $n$  be a positive integer and  $F$  a field. Suppose  $A$  is an  $n \times n$  matrix over  $F$  and  $P$  is an invertible  $n \times n$  matrix over  $F$ . If  $f$  is any polynomial over  $F$ , prove that

$$f(P^{-1}AP) = P^{-1}f(A)P.$$

6. Let  $F$  be a field. We have considered certain special linear functionals on  $F[x]$  obtained via ‘evaluation at  $t$ ’:

$$L(f) = f(t).$$

Such functionals are not only linear but also have the property that  $L(fg) = L(f)L(g)$ . Prove that if  $L$  is any linear functional on  $F[x]$  such that

$$L(fg) = L(f)L(g)$$

for all  $f$  and  $g$ , then either  $L = 0$  or there is a  $t$  in  $F$  such that  $L(f) = f(t)$  for all  $f$ .

## 4.4. Polynomial Ideals

In this section we are concerned with results which depend primarily on the multiplicative structure of the algebra of polynomials over a field.

**Lemma.** Suppose  $f$  and  $d$  are non-zero polynomials over a field  $F$  such that  $\deg d \leq \deg f$ . Then there exists a polynomial  $g$  in  $F[x]$  such that either

$$f - dg = 0 \quad \text{or} \quad \deg(f - dg) < \deg f.$$

*Proof.* Suppose

$$f = a_m x^m + \sum_{i=0}^{m-1} a_i x^i, \quad a_m \neq 0$$

and that

$$d = b_n x^n + \sum_{i=0}^{n-1} b_i x^i, \quad b_n \neq 0.$$

Then  $m \geq n$ , and

$$f - \left(\frac{a_m}{b_n}\right)x^{m-n}d = 0 \quad \text{or} \quad \deg \left[ f - \left(\frac{a_m}{b_n}\right)x^{m-n}d \right] < \deg f.$$

Thus we may take  $g = \left(\frac{a_m}{b_n}\right)x^{m-n}$ . ■

Using this lemma we can show that the familiar process of 'long division' of polynomials with real or complex coefficients is possible over any field.

**Theorem 4.** *If  $f, d$  are polynomials over a field  $F$  and  $d$  is different from 0 then there exist polynomials  $q, r$  in  $F[x]$  such that*

- (i)  $f = dq + r$ .
- (ii) either  $r = 0$  or  $\deg r < \deg d$ .

*The polynomials  $q, r$  satisfying (i) and (ii) are unique.*

*Proof.* If  $f$  is 0 or  $\deg f < \deg d$  we may take  $q = 0$  and  $r = f$ . In case  $f \neq 0$  and  $\deg f \geq \deg d$ , the preceding lemma shows we may choose a polynomial  $g$  such that  $f - dg = 0$  or  $\deg(f - dg) < \deg f$ . If  $f - dg \neq 0$  and  $\deg(f - dg) \geq \deg d$  we choose a polynomial  $h$  such that  $(f - dg) - dh = 0$  or

$$\deg [f - d(g + h)] < \deg (f - dg).$$

Continuing this process as long as necessary, we ultimately obtain polynomials  $q, r$  such that  $r = 0$  or  $\deg r < \deg d$ , and  $f = dq + r$ . Now suppose we also have  $f = dq_1 + r_1$  where  $r_1 = 0$  or  $\deg r_1 < \deg d$ . Then  $dq + r = dq_1 + r_1$ , and  $d(q - q_1) = r_1 - r$ . If  $q - q_1 \neq 0$  then  $d(q - q_1) \neq 0$  and

$$\deg d + \deg (q - q_1) = \deg (r_1 - r).$$

But as the degree of  $r_1 - r$  is less than the degree of  $d$ , this is impossible and  $q - q_1 = 0$ . Hence also  $r_1 - r = 0$ . ■

**Definition.** *Let  $d$  be a non-zero polynomial over the field  $F$ . If  $f$  is in  $F[x]$ , the preceding theorem shows there is at most one polynomial  $q$  in  $F[x]$  such that  $f = dq$ . If such a  $q$  exists we say that  $d$  divides  $f$ , that  $f$  is divisible by  $d$ , that  $f$  is a multiple of  $d$ , and call  $q$  the quotient of  $f$  and  $d$ . We also write  $q = f/d$ .*

**Corollary 1.** *Let  $f$  be a polynomial over the field  $F$ , and let  $c$  be an element of  $F$ . Then  $f$  is divisible by  $x - c$  if and only if  $f(c) = 0$ .*

*Proof.* By the theorem,  $f = (x - c)q + r$  where  $r$  is a scalar polynomial. By Theorem 2,

$$f(c) = 0q(c) + r(c) = r(c).$$

Hence  $r = 0$  if and only if  $f(c) = 0$ . ■

**Definition.** Let  $F$  be a field. An element  $c$  in  $F$  is said to be a **root** or a **zero** of a given polynomial  $f$  over  $F$  if  $f(c) = 0$ .

**Corollary 2.** A polynomial  $f$  of degree  $n$  over a field  $F$  has at most  $n$  roots in  $F$ .

*Proof.* The result is obviously true for polynomials of degree 0 and degree 1. We assume it to be true for polynomials of degree  $n - 1$ . If  $a$  is a root of  $f$ ,  $f = (x - a)q$  where  $q$  has degree  $n - 1$ . Since  $f(b) = 0$  if and only if  $a = b$  or  $q(b) = 0$ , it follows by our inductive assumption that  $f$  has at most  $n$  roots. ■

The reader should observe that the main step in the proof of Theorem 3 follows immediately from this corollary.

The formal derivatives of a polynomial are useful in discussing multiple roots. The **derivative** of the polynomial

$$f = c_0 + c_1x + \cdots + c_nx^n$$

is the polynomial

$$f' = c_1 + 2c_2x + \cdots + nc_nx^{n-1}.$$

We also use the notation  $Df = f'$ . Differentiation is linear, that is,  $D$  is a linear operator on  $F[x]$ . We have the higher order formal derivatives  $f'' = D^2f$ ,  $f^{(3)} = D^3f$ , and so on.

**Theorem 5 (Taylor's Formula).** Let  $F$  be a field of characteristic zero,  $c$  an element of  $F$ , and  $n$  a positive integer. If  $f$  is a polynomial over  $F$  with  $\deg f \leq n$ , then

$$f = \sum_{k=0}^n \frac{(D^k f)}{k!} (c)(x - c)^k.$$

*Proof.* Taylor's formula is a consequence of the binomial theorem and the linearity of the operators  $D$ ,  $D^2$ ,  $\dots$ ,  $D^n$ . The binomial theorem is easily proved by induction and asserts that

$$(a + b)^m = \sum_{k=0}^m \binom{m}{k} a^{m-k} b^k$$

where

$$\binom{m}{k} = \frac{m!}{k!(m - k)!} = \frac{m(m - 1) \cdots (m - k + 1)}{1 \cdot 2 \cdots k}$$

is the familiar binomial coefficient giving the number of combinations of  $m$  objects taken  $k$  at a time. By the binomial theorem

$$\begin{aligned} x^m &= [c + (x - c)]^m \\ &= \sum_{k=0}^m \binom{m}{k} c^{m-k}(x - c)^k \\ &= c^m + mc^{m-1}(x - c) + \cdots + (x - c)^m \end{aligned}$$