

and hence  $g(X)|h(X)$ . Thus  $g(X)$  is the polynomial of least possible degree (up to a constant factor) with roots  $\alpha, \alpha^2, \dots, \alpha^{d-1}$ .

If  $c(X)$  is a code polynomial in the polynomial code generated by  $g(X)$ , then  $c(X) = a(X)g(X)$  for some  $a(X) \in F[X]$  and, therefore,  $\alpha, \alpha^2, \dots, \alpha^{d-1}$  are roots of  $c(X)$ .

We know that in a group code the minimum distance of the code equals the minimum of the weights of non-zero code words. Since polynomial codes are group codes, it follows that the code generated by  $g(X)$  has minimum distance at least  $d$  if there is no code word  $c_0c_1 \dots c_{n-1}$  with less than  $d$  non-zero entries. Suppose, to the contrary, that a code word has less than  $d$  non-zero entries. Then the corresponding code polynomial is of the form

$$c(X) = b_1X^{n_1} + b_2X^{n_2} + \dots + b_{d-1}X^{n_{d-1}}$$

where  $b_1, b_2, \dots, b_{d-1} \in F$  and also, we may assume that

$$n_1 > n_2 > \dots > n_{d-1} \geq 0$$

Since the code is of length  $n$ , every code polynomial is of degree at most  $n-1$  and, therefore,  $n_1 \leq n-1$  ( $\leq q^r-2$ ). As already pointed out,  $\alpha, \alpha^2, \dots, \alpha^{d-1}$  are roots of  $c(X)$  and we have

$$\begin{aligned} b_1\alpha^{n_1} + b_2\alpha^{n_2} + \dots + b_{d-1}\alpha^{n_{d-1}} &= 0 \\ b_1\alpha^{2n_1} + b_2\alpha^{2n_2} + \dots + b_{d-1}\alpha^{2n_{d-1}} &= 0 \\ \dots & \dots \dots \\ b_1\alpha^{(d-1)n_1} + b_2\alpha^{(d-1)n_2} + \dots + b_{d-1}\alpha^{(d-1)n_{d-1}} &= 0 \end{aligned}$$

or

$$\mathbf{A} \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_{d-1} \end{pmatrix} = 0 \quad (4.2)$$

where

$$\mathbf{A} = \begin{pmatrix} \alpha^{n_1} & \alpha^{n_2} & \dots & \alpha^{n_{d-1}} \\ \alpha^{2n_1} & \alpha^{2n_2} & \dots & \alpha^{2n_{d-1}} \\ \dots & \dots & \ddots & \dots \\ \alpha^{(d-1)n_1} & \alpha^{(d-1)n_2} & \dots & \alpha^{(d-1)n_{d-1}} \end{pmatrix}$$

The determinant of  $\mathbf{A}$  is a **Vandermonde determinant** and we know that

$$\det \mathbf{A} = \prod_{i>j} (\alpha^{n_i} - \alpha^{n_j})$$

The element  $\alpha$  of  $K$  being primitive and

$$q^r - 1 > n_1 > n_2 > \dots > n_{d-1} \geq 0$$

we have  $\alpha^{n_i} - \alpha^{n_j} \neq 0$  for  $i \neq j$  and, therefore,  $\det \mathbf{A} \neq 0$ . Now (4.2) is a system of  $d - 1$  homogeneous linear equations in  $d - 1$  variables  $b_1, \dots, b_{d-1}$  and  $\det \mathbf{A} \neq 0$ . Therefore the system of equations admits only the zero solution and  $c(X) = 0$ . Hence there is no non-zero code word with less than  $d$  non-zero entries and the code has minimum distance at least  $d$ .

### Examples 4.8

#### Case (i)

Construct a binary BCH code of length 7 and minimum distance 3.

Here  $n = 7$  and so we need to construct an extension of  $\mathbb{B}$  of degree  $r$  where  $2^r \geq 7 + 1 = 8$ . Thus we take  $r = 3$ . We know from Example 4.7 (Case (iii)) that  $X^3 + X + 1$  is a primitive polynomial of degree 3 over  $\mathbb{B}$ . Therefore

$$K = \mathbb{B}[X]/\langle X^3 + X + 1 \rangle$$

is a field of order 8 and

$$\alpha = X + \langle X^3 + X + 1 \rangle$$

is a primitive element of  $K$ . Then  $\alpha$  satisfies the relations  $\alpha^3 + \alpha + 1 = 0$ , and  $\alpha^7 = 1$  and  $X^3 + X + 1$  is the minimal polynomial of  $\alpha$ . Since  $\alpha$  and  $\alpha^2$  have the same minimal polynomial (Proposition 4.2), the generator polynomial of the required BCH code is  $X^3 + X + 1$ .

The message polynomials are of degree at most 3. If

$$a(X) = a_0 + a_1X + a_2X^2 + a_3X^3$$

is an arbitrary message polynomial, the corresponding code polynomial is  $a(X)(X^3 + X + 1)$  and so the corresponding code word is

$$(a_0, a_1 + a_0, a_2 + a_1, a_3 + a_2 + a_0, a_3 + a_1, a_2, a_3)$$

The encoding polynomial has 3 non-zero terms and, therefore, the code has minimum distance 3.

If we had started with the primitive polynomial  $X^3 + X^2 + 1$ , the corresponding BCH code with code word length 7 and minimum distance at least 3 is the polynomial code with encoding polynomial  $X^3 + X^2 + 1$ .

#### Case (ii)

Next we construct a binary BCH code of length 15 and minimum distance 5.

Here  $n = 15 \geq 2^4 - 1$  and so we need to construct an extension  $K$  of  $\mathbb{B}$  of degree 4. We have seen earlier (Example 4.3 Case (ii)) that  $X^4 + X + 1$  is a primitive polynomial and so

$$\alpha = X + \langle X^4 + X + 1 \rangle$$

is a primitive element of

$$K = \mathbb{B}[X]/\langle X^4 + X + 1 \rangle$$

The minimal polynomial of  $\alpha$  is  $m_1(X) = X^4 + X + 1$ . Also

$$m_2(X) = m_4(X) = m_1(X)$$

(Proposition 4.2). We next have to find the minimal polynomial of  $\alpha^3$ . The elements  $\alpha^3, \alpha^6, \alpha^{12}, \alpha^9$  have the same minimal polynomial and so

$$\begin{aligned} m_3(X) &= (X - \alpha^3)(X - \alpha^6)(X - \alpha^9)(X - \alpha^{12}) \\ &= X^4 + X^3(\alpha^3 + \alpha^6 + \alpha^9 + \alpha^{12}) + X^2(\alpha^9 + \alpha^{12} + \alpha^{15} + \alpha^{15} + \alpha^{18} + \alpha^{21}) \\ &\quad + X(\alpha^{18} + \alpha^{21} + \alpha^{24} + \alpha^{27}) + \alpha^{30} \\ &= X^4 + X^3(\alpha^3 + \alpha^3 + \alpha^2 + \alpha^3 + \alpha + \alpha^3 + \alpha^2 + \alpha + 1) \\ &\quad + X^2(\alpha^9 + \alpha^{12} + \alpha^3 + \alpha^6) + X(\alpha^3 + \alpha^6 + \alpha^9 + \alpha^{12}) + 1 \\ &= X^4 + X^3 + X^2 + X + 1 \end{aligned}$$

Therefore, the encoding polynomial of the BCH code with minimum distance at least 5 is

$$\begin{aligned} g(X) &= (X^4 + X + 1)(X^4 + X^3 + X^2 + X + 1) \\ &= X^8 + X^7 + X^6 + X^4 + 1 \end{aligned}$$

Since the encoding polynomial has 5 non-zero terms, the minimum distance of the code is exactly 5. The code being of length 15, a message polynomial is of degree at most 6. Let

$$a(X) = a_0 + a_1X + a_2X^2 + a_3X^3 + a_4X^4 + a_5X^5 + a_6X^6$$

be an arbitrary message polynomial. The code word corresponding to the code polynomial  $a(X)g(X)$  is

$$(a_0, a_1, a_2, a_3, a_4 + a_0, a_5 + a_1, a_6 + a_2 + a_0, a_0 + a_1 + a_3, a_0 + a_1 + a_2 + a_4, \\ a_1 + a_2 + a_3 + a_5, a_2 + a_3 + a_4 + a_6, a_3 + a_4 + a_5, a_4 + a_5 + a_6, a_5 + a_6, a_6)$$

### Case (iii)

Find a generator polynomial of the binary BCH code of length 31 and minimum distance 5, it being given that  $X^5 + X^2 + 1$  is an irreducible polynomial over  $\mathbb{B}$ .

### Solution

The polynomial  $X^5 + X^2 + 1$  being irreducible,

$$F = \mathbb{B}[X]/\langle X^5 + X^2 + 1 \rangle$$

is a field of order 32. Since  $F^* = F \setminus \{0\}$  is a cyclic group of order 31 – a prime – every non-identity element of  $F^*$  is a primitive element of  $F$ . In particular

$$\alpha = X + \langle X^5 + X^2 + 1 \rangle$$

is a primitive element of  $F$  and  $X^5 + X^2 + 1$  is the minimal polynomial of  $\alpha$ .

Let  $m_i(X)$  be the minimal polynomial of  $\alpha^i$ ,  $1 \leq i \leq 4$ . Observe that  $\alpha, \alpha^2, \alpha^4$  have the same minimal polynomial. So,

$$m_1(X) = m_2(X) = m_4(X) = X^5 + X^2 + 1$$

We now have to find  $m_3(X)$ . As  $\alpha^3, \alpha^6, \alpha^{12}, \alpha^{24}, \alpha^{17}$  have the same minimal polynomial and  $\{3, 6, 12, 24, 17\}$  is the complete cyclotomic class modulo 31 relative to 2,

$$m_3(X) = (X - \alpha^3)(X - \alpha^6)(X - \alpha^{12})(X - \alpha^{17})(X - \alpha^{24})$$

In  $m_3(X)$ , the coefficient of  $X^4$  is

$$\begin{aligned} \alpha^3 + \alpha^6 + \alpha^{12} + \alpha^{17} + \alpha^{24} &= \alpha^3 + \alpha^3 + \alpha + \alpha^{12} + (\alpha^3 + \alpha)^4 + (\alpha^2 + 1)(\alpha^6 + \alpha^2) \\ &= \alpha + \alpha^{12} + \alpha^{12} + \alpha^4 + \alpha^8 + \alpha^6 + \alpha^4 + \alpha^2 \\ &= \alpha + \alpha^2 + \alpha^6 + \alpha^8 \\ &= \alpha + \alpha^2 + \alpha^3 + \alpha + \alpha^3(\alpha^2 + 1) \\ &= \alpha^5 + \alpha^2 \\ &= 1 \end{aligned}$$

The coefficient of  $X^3$  is

$$\begin{aligned} \alpha^9 + \alpha^{15} + \alpha^{20} + \alpha^{27} + \alpha^{18} + \alpha^{23} + \alpha^{30} + \alpha^{29} + \alpha^5 + \alpha^{10} \\ &= \alpha^9 + \alpha^{17} + \alpha^7(\alpha^2 + 1)^4 + \alpha^{20} + (\alpha^4 + 1)(\alpha^8 + 1) + \alpha^9(\alpha^8 + 1) + \alpha^7 \\ &= \alpha^9 + \alpha^{17} + \alpha^{15} + \alpha^7 + \alpha^8 + 1 + \alpha^{12} + \alpha^8 + \alpha^4 + 1 + \alpha^{17} + \alpha^9 + \alpha^7 \\ &= \alpha^{15} + \alpha^{12} + \alpha^4 \\ &= (\alpha^2 + 1)^3 + \alpha^2(\alpha^4 + 1) + \alpha^4 \\ &= \alpha^6 + \alpha^4 + \alpha^2 + 1 + \alpha^6 + \alpha^2 + \alpha^4 \\ &= 1 \end{aligned}$$

The coefficient of  $X^2$  is

$$\begin{aligned} \alpha^{21} + \alpha^{26} + \alpha^2 + \alpha + \alpha^8 + \alpha^{13} + \alpha^4 + \alpha^{11} + \alpha^{16} + \alpha^{22} \\ &= \alpha^{23} + \alpha^2 + \alpha + \alpha^{10} + \alpha^4 + \alpha^{13} + \alpha^2(\alpha^8 + 1) \\ &= \alpha^3(\alpha^8 + 1) + \alpha + \alpha^4 + \alpha^3(\alpha^4 + 1) \\ &= \alpha^{11} + \alpha + \alpha^4 + \alpha^7 \\ &= \alpha(\alpha^4 + 1) + \alpha + \alpha^4 + \alpha^2(\alpha^2 + 1) \\ &= \alpha^5 + \alpha^2 \\ &= 1 \end{aligned}$$

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The coefficient of  $X$  is

$$\begin{aligned}
 \alpha^7 + \alpha^{14} + \alpha^{19} + \alpha^{25} + \alpha^{28} &= \alpha^7 + \alpha^{16} + \alpha^5(\alpha^8 + 1) + \alpha^8(\alpha^8 + 1) \\
 &= \alpha^7 + \alpha^{13} + \alpha^5 + \alpha^8 \\
 &= \alpha^2(\alpha^2 + 1) + \alpha^{10} + \alpha^5 \\
 &= \alpha^4 + \alpha^2 + \alpha^4 + 1 + \alpha^2 + 1 \\
 &= 0
 \end{aligned}$$

The constant term is

$$\alpha^{3+6+12+17+24} = \alpha^{62} = 1$$

Hence

$$m_3(X) = X^5 + X^4 + X^3 + X^2 + 1$$

Therefore the generator polynomial of the BCH code of length 31 over  $\mathbb{B}$  is

$$\begin{aligned}
 g(X) &= \text{LCM}\{m_1(X), m_2(X), m_3(X), m_4(X)\} \\
 &= \text{LCM}\{m_1(X), m_3(X)\} \\
 &= m_1(X)m_3(X) \\
 &= (X^5 + X^2 + 1)(X^5 + X^4 + X^3 + X^2 + 1) \\
 &= X^{10} + X^9 + X^8 + X^6 + X^5 + X^3 + 1
 \end{aligned}$$

### Case (iv)

Find a generator polynomial for a 5-error-correcting binary BCH code of length 63, it being given that  $X^6 + X + 1$  is a primitive polynomial over  $\mathbb{B}$ .

### Solution

The length of the code is  $63 = 2^6 - 1$  and we need to find an extension of  $\mathbb{B}$  of degree 6. It being given that  $X^6 + X + 1$  is a primitive polynomial over  $\mathbb{B}$ ,

$$K = \mathbb{B}[X]/\langle X^6 + X + 1 \rangle$$

is a field of order  $2^6$ ,

$$\alpha = X + \langle X^6 + X + 1 \rangle$$

is a primitive element of  $K$  and  $X^6 + X + 1$  is the minimal polynomial of  $\alpha$ .

Since the code we are looking for is to be 5-error-correcting, the minimum distance of the code is at least  $2 \times 5 + 1 = 11$  (Theorem 1.2). We thus have to find the minimal polynomials  $m_i(X)$  of  $\alpha^i$  for  $1 \leq i \leq 10$ . It follows from

Proposition 4.2 that

$$\begin{aligned}
 m_1(X) &= m_2(X) = m_4(X) = m_8(X) = m_{16}(X) = m_{32}(X) \\
 m_3(X) &= m_6(X) = m_{12}(X) = m_{24}(X) = m_{48}(X) = m_{33}(X) \\
 m_5(X) &= m_{10}(X) = m_{20}(X) = m_{40}(X) = m_{17}(X) = m_{34}(X) \\
 m_7(X) &= m_{14}(X) = m_{28}(X) = m_{56}(X) = m_{49}(X) = m_{35}(X) \\
 m_9(X) &= m_{18}(X) = m_{36}(X)
 \end{aligned}$$

Thus  $m_1(X)$ ,  $m_3(X)$ ,  $m_5(X)$ ,  $m_7(X)$  are of degree 6 each while  $m_9(X)$  is of degree 3. Also then the encoding polynomial of the BCH code we are looking for is

$$g(X) = m_1(X)m_3(X)m_5(X)m_7(X)m_9(X)$$

which is of degree 27. Now

$$\begin{aligned}
 \alpha^6 &= \alpha + 1 \Rightarrow \alpha^{12} = \alpha^2 + 1 \\
 \alpha^{24} &= \alpha^4 + 1 \\
 \alpha^{48} &= \alpha^8 + 1 = \alpha^3 + \alpha^2 + 1 \\
 \alpha^{33} &= (\alpha^4 + 1)(\alpha^4 + \alpha^3) = \alpha^8 + \alpha^7 + \alpha^4 + \alpha^3 = \alpha^3 + \alpha^2 + \alpha^2 + \alpha + \alpha^4 + \alpha^3 \\
 &= \alpha^4 + \alpha
 \end{aligned}$$

Therefore

$$\begin{aligned}
 m_3(X) &= (X + \alpha^3)(X + \alpha + 1)(X + \alpha^2 + 1)(X + \alpha^4 + 1)(X + \alpha^3 + \alpha^2 + 1) \\
 &\quad \times (X + \alpha^4 + \alpha) \\
 &= [X^2 + X(\alpha^3 + \alpha + 1) + \alpha^4 + \alpha^3][X^2 + X(\alpha^4 + \alpha^2) + \alpha^4 + \alpha^2 + \alpha] \\
 &\quad \times [X^2 + X(\alpha^4 + \alpha^3 + \alpha^2 + \alpha + 1) + \alpha^3 + \alpha^2 + \alpha + 1] \\
 &= [X^4 + X^3(\alpha^4 + \alpha^3 + \alpha^2 + \alpha + 1) + X^2(\alpha^4 + \alpha^2) + X(\alpha^5 + \alpha^2 + 1) \\
 &\quad + \alpha^4 + \alpha^3 + 1][X^2 + X(\alpha^4 + \alpha^3 + \alpha^2 + \alpha + 1) + \alpha^3 + \alpha^2 + \alpha + 1] \\
 &= X^6 + X^4(\alpha^8 + \alpha^6 + \alpha^4 + \alpha^2 + 1 + \alpha^4 + \alpha^2 + \alpha^3 + \alpha^2 + \alpha + 1) \\
 &\quad + X^3(\alpha^8 + \alpha^2 + \alpha^3) + X^2(\alpha^9 + \alpha^8 + \alpha^6 + \alpha^4 + \alpha^2 + \alpha) \\
 &\quad + X(\alpha^2 + \alpha^6 + \alpha^7) + (\alpha^7 + \alpha^2 + \alpha + 1) \\
 &= X^6 + X^4 + X^3 + X^2 + X + 1
 \end{aligned}$$

Again

$$\begin{aligned}
 \alpha^{10} &= \alpha^5 + \alpha^4 & \alpha^{20} &= \alpha^5 + \alpha^4 + \alpha^3 + \alpha^2 \\
 \alpha^{40} &= \alpha^5 + \alpha^3 + \alpha^2 + \alpha + 1 \\
 \alpha^{17} &= (\alpha^5 + \alpha + 1)(\alpha + 1) = \alpha^5 + \alpha^2 + \alpha \\
 \alpha^{34} &= \alpha^{10} + \alpha^4 + \alpha^2 = \alpha^5 + \alpha^2
 \end{aligned}$$

Therefore

$$\begin{aligned}
m_5(X) &= (X + \alpha^5 + \alpha^4)(X + \alpha^5 + \alpha^4 + \alpha^3 + \alpha^2)(X + \alpha^5 + \alpha^3 + \alpha^2 + \alpha + 1) \\
&\quad \times (X + \alpha^5 + \alpha^2 + \alpha)(X + \alpha^5 + \alpha^2)(X + \alpha^5) \\
&= [X^2 + X(\alpha^3 + \alpha^2) + \alpha^5 + \alpha^4 + \alpha + 1] \\
&\quad \times [X^2 + X(\alpha^3 + 1) + \alpha^5 + \alpha^4 + \alpha^3 + \alpha^2 + \alpha] \\
&\quad \times [X^2 + \alpha^2 X + \alpha^5 + \alpha^4 + \alpha^2 + \alpha] \\
&= [X^4 + X^3(\alpha^2 + 1) + X^2(\alpha^3 + \alpha^2 + 1 + \alpha^5 + \alpha^3 + \alpha^2 + \alpha + 1) \\
&\quad + X(\alpha^5 + \alpha^2 + 1) + \alpha^4 + 1][X^2 + \alpha^2 X + \alpha^5 + \alpha^4 + \alpha^2 + \alpha] \\
&= X^6 + X^5(\alpha^2 + 1 + \alpha^2) + X^4(\alpha^5 + \alpha^4 + \alpha^2 + \alpha + \alpha^4 + \alpha^2 + \alpha^5 + \alpha) \\
&\quad + X^3[(\alpha^2 + 1)(\alpha^5 + \alpha^4 + \alpha^2 + \alpha) + \alpha^2(\alpha^5 + \alpha) + \alpha^5 + \alpha^2 + 1] \\
&\quad + X^2[\alpha^4 + 1 + \alpha^2(\alpha^5 + \alpha^2 + 1) + (\alpha^5 + \alpha)(\alpha^5 + \alpha^4 + \alpha^2 + \alpha)] \\
&\quad + X[\alpha^2(\alpha^4 + 1) + (\alpha^5 + \alpha^2 + 1)(\alpha^5 + \alpha^4 + \alpha^2 + \alpha)] \\
&\quad + (\alpha^4 + 1)(\alpha^5 + \alpha^4 + \alpha^2 + \alpha) \\
&= X^6 + X^5 + X^2 + X + 1
\end{aligned}$$

For finding  $m_7(X)$ :

$$\begin{aligned}
\alpha^7 &= \alpha^2 + \alpha \\
\alpha^{14} &= \alpha^4 + \alpha^2 \\
\alpha^{28} &= \alpha^4 + \alpha^3 + \alpha^2 \\
\alpha^{56} &= \alpha^4 + \alpha^3 + \alpha^2 + \alpha + 1 \\
\alpha^{49} &= \alpha^4 + \alpha^3 + \alpha \\
\alpha^{35} &= \alpha^3 + \alpha + 1
\end{aligned}$$

Then

$$\begin{aligned}
\alpha^7 + \alpha^{14} + \alpha^{28} + \alpha^{56} + \alpha^{49} + \alpha^{35} &= 0 \\
\alpha^7 \times \alpha^{14} \times \alpha^{28} \times \alpha^{56} \times \alpha^{49} \times \alpha^{35} &= \alpha^{189} = 1
\end{aligned}$$

Sum of the products of these 6 powers of  $\alpha$  taken 5 at a time

$$\begin{aligned}
&= \alpha^{63-7} + \alpha^{63-14} + \alpha^{63-28} + \alpha^{63-56} + \alpha^{63-49} + \alpha^{63-35} \\
&= \alpha^{56} + \alpha^{49} + \alpha^{35} + \alpha^7 + \alpha^{14} + \alpha^{28} \\
&= 0
\end{aligned}$$

Since  $\alpha$  is an element of a field  $K$  and

$$\alpha^7 \times \alpha^{56} = \alpha^{14} \times \alpha^{49} = \alpha^{28} \times \alpha^{35} = 1$$

sum of products of these elements taken 2 at a time

$$\begin{aligned}
 &= \text{sum of products of these elements taken 4 at a time} \\
 &= \alpha^7(\alpha^{14} + \alpha^{28} + \alpha^{35} + \alpha^{49} + \alpha^{56}) + \alpha^{14}(\alpha^{28} + \alpha^{35} + \alpha^{49} + \alpha^{56}) \\
 &\quad + \alpha^{28}(\alpha^{35} + \alpha^{49} + \alpha^{56}) + \alpha^{35}(\alpha^{49} + \alpha^{56}) + \alpha^{49} \times \alpha^{56} \\
 &= \alpha^7 \times \alpha^7 + \alpha^{14}(\alpha^7 + \alpha^{14}) + (1 + \alpha^{14} + \alpha^{21}) + (\alpha^{21} + \alpha^{28}) + \alpha^{42} \\
 &= \alpha^{14} + \alpha^{21} + \alpha^{28} + 1 + \alpha^{14} + \alpha^{21} + \alpha^{21} + \alpha^{28} + \alpha^{42} \\
 &= 1 + \alpha^{21} + \alpha^{42} \\
 &= 1 + (\alpha^2 + \alpha)^3 + (\alpha^2 + \alpha)^6 \\
 &= 1 + (\alpha^5 + \alpha^4 + \alpha^3 + \alpha + 1) + (\alpha^5 + \alpha^4 + \alpha^3 + \alpha + 1)^2 \\
 &= 1 + (\alpha^5 + \alpha^4 + \alpha^3 + \alpha + 1) + (\alpha^5 + \alpha^4 + \alpha^3 + \alpha^2 + \alpha + 1 + \alpha^2 + 1) \\
 &= 0
 \end{aligned}$$

Sum of the products of these powers of  $\alpha$  taken 3 at a time

$$\begin{aligned}
 &= (\alpha^{49} + \alpha^{56} + \alpha^7 + \alpha^{14}) + (\alpha^7 + \alpha^{21} + \alpha^{28}) + (\alpha^{28} + \alpha^{35}) + \alpha^{49} \\
 &\quad + (\alpha^{14} + \alpha^{28} + \alpha^{35}) + (\alpha^{35} + \alpha^{42}) + \alpha^{56} + (\alpha^{49} + \alpha^{56}) + \alpha^7 + \alpha^{14} \\
 &= \alpha^7 + \alpha^{14} + \alpha^{21} + \alpha^{28} + \alpha^{35} + \alpha^{42} + \alpha^{49} + \alpha^{56} \\
 &= \alpha^{21} + \alpha^{42} \\
 &= 1
 \end{aligned}$$

Therefore

$$m_7(X) = X^6 + X^3 + 1$$

We are now left with computing  $m_9(X)$ .

$$\begin{aligned}
 m_9(X) &= (X + \alpha^9)(X + \alpha^{18})(X + \alpha^{36}) \\
 &= X^3 + X^2(\alpha^9 + \alpha^{18} + \alpha^{36}) + X(\alpha^{27} + \alpha^{45} + \alpha^{54}) + 1 \\
 &= X^3 + X^2[(\alpha^4 + \alpha^3) + (\alpha^4 + \alpha^3)^2 + (\alpha^4 + \alpha^3)^4] \\
 &\quad + X[\alpha^{27} + \alpha^{45} + \alpha^{54}] + 1
 \end{aligned}$$

Now

$$\begin{aligned}
 \alpha^{18} &= (\alpha^4 + \alpha^3)^2 = \alpha^8 + \alpha^6 = \alpha^3 + \alpha^2 + \alpha + 1 \\
 \alpha^{36} &= (\alpha^3 + \alpha^2 + \alpha + 1)^2 = \alpha^4 + \alpha^2 + \alpha \\
 \alpha^{27} &= (\alpha^4 + \alpha^3)(\alpha^3 + \alpha^2 + \alpha + 1) = \alpha^7 + \alpha^3 = \alpha^3 + \alpha^2 + \alpha \\
 \alpha^{54} &= (\alpha^3 + \alpha^2 + \alpha)^2 = \alpha^4 + \alpha^2 + \alpha + 1 \\
 \alpha^{45} &= (\alpha^3 + \alpha^2 + \alpha + 1)(\alpha^3 + \alpha^2 + \alpha) = \alpha^4 + \alpha^3 + 1
 \end{aligned}$$