

The first step in the rho method is to choose an easily evaluated map from  $\mathbf{Z}/n\mathbf{Z}$  to itself, namely, a fairly simple polynomial with integer coefficients, such as  $f(x) = x^2 + 1$ . Next, one chooses some particular value  $x = x_0$  (perhaps  $x_0 = 1$  or  $2$ , or perhaps it is a randomly generated integer) and computes the successive iterates of  $f$ :  $x_1 = f(x_0)$ ,  $x_2 = f(f(x_0))$ ,  $x_3 = f(f(f(x_0)))$ , etc. That is, we define

$$x_{j+1} = f(x_j), \quad j = 0, 1, 2, \dots$$

Then we make comparisons between different  $x_j$ 's, hoping to find two which are in different residue classes modulo  $n$  but in the same residue class modulo some divisor of  $n$ . Once we find such  $x_j, x_k$ , we have  $\text{g.c.d.}(x_j - x_k, n)$  equal to a proper divisor of  $n$ , and we are done.

**Example 1.** Let us factor 91 by choosing  $f(x) = x^2 + 1$ ,  $x_0 = 1$ . Then we have  $x_1 = 2$ ,  $x_2 = 5$ ,  $x_3 = 26$ , etc. We find that  $\text{g.c.d.}(x_3 - x_2, n) = \text{g.c.d.}(21, 91) = 7$ , so 7 is a factor. Of course, this is a trivial example: we could have found the factor 7 faster by trial division.

In the rho method it is important to choose a polynomial  $f(x)$  which maps  $\mathbf{Z}/n\mathbf{Z}$  to itself in a rather disjointed, "random" way. For example, we shall later see that  $f(x)$  must not be a linear polynomial, and in fact, should not give a 1-to-1 map.

Let us suppose that  $f(x)$  is a "random" map from  $\mathbf{Z}/n\mathbf{Z}$  to itself, and compute how long we expect to have to wait before we have two iterations  $x_j$  and  $x_k$  such that  $x_j - x_k$  has a nontrivial common factor with  $n$ . We do this by finding for a fixed divisor  $r$  of  $n$  (which, in practice, is not yet known to us) the average (taken over all maps from  $\mathbf{Z}/n\mathbf{Z}$  to itself and over all values  $x_0$ ) of the first index  $k$  such that there exists  $j < k$  with  $x_j \equiv x_k \pmod{r}$ . In other words, we regard  $f(x)$  as a map from  $\mathbf{Z}/r\mathbf{Z}$  to itself and ask how many iterations are required before we encounter the first repetition of values  $x_k = x_j$  in  $\mathbf{Z}/r\mathbf{Z}$ .

**Proposition V.2.1.** *Let  $S$  be a set of  $r$  elements. Given a map  $f$  from  $S$  to  $S$  and an element  $x_0 \in S$ , let  $x_{j+1} = f(x_j)$  for  $j = 0, 1, 2, \dots$ . Let  $\lambda$  be a positive real number, and let  $\ell = 1 + \lceil \sqrt{2\lambda r} \rceil$ . Then the proportion of pairs  $(f, x_0)$  for which  $x_0, x_1, \dots, x_\ell$  are distinct, where  $f$  runs over all maps from  $S$  to  $S$  and  $x_0$  runs over all elements of  $S$ , is less than  $e^{-\lambda}$ .*

**Proof.** The total number of pairs is  $r^{r+1}$ , because there are  $r$  choices of  $x_0$ , and for each of the  $r$  different  $x \in S$  there are  $r$  choices of  $f(x)$ . How many pairs  $(f, x_0)$  are there for which  $x_0, x_1, \dots, x_\ell$  are distinct? There are  $r$  choices for  $x_0$ , there are  $r - 1$  choices for  $f(x_0) = x_1$  (since this cannot equal  $x_0$ ), there are  $r - 2$  choices for  $f(x_1) = x_2$ , and so on, until  $f(x)$  has been defined for  $x = x_0, x_1, \dots, x_{\ell-1}$ . Then the value of  $f(x)$  for each of the  $r - \ell$  remaining  $x$  is arbitrary, i.e., there are  $r^{r-\ell}$  possibilities for those values. Hence, the total number of possible ways of choosing  $x_0$  and assigning the values  $f(x)$  so that  $x_0, \dots, x_\ell$  are distinct is: