

**Programação e Algoritmia**

**--X--**

**Programming and Algorithms**

1 – Object-Oriented Programming

# Non-linear data structures

With less implementation details 😊

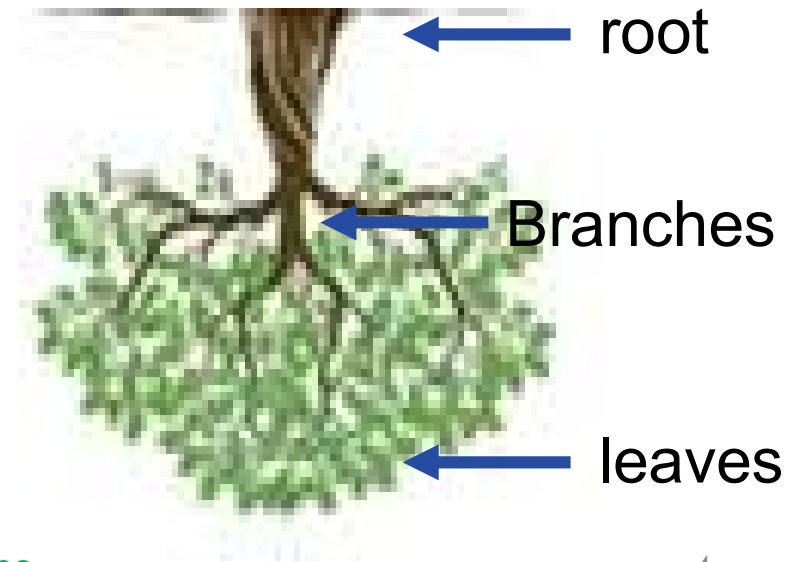
# Trees

A brief introduction

# What are trees?

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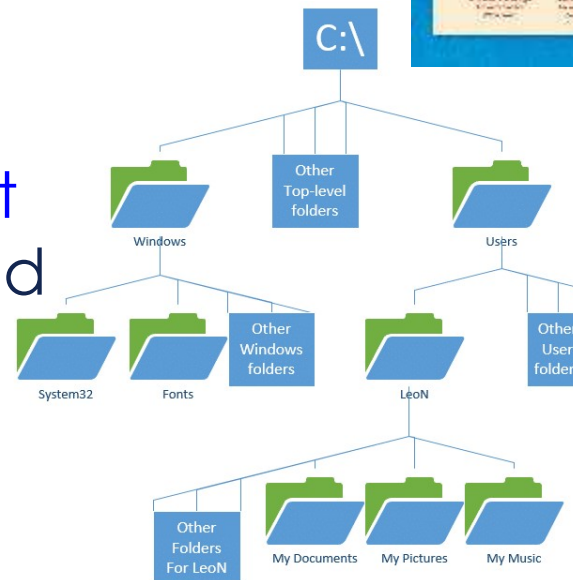
- ❖ A tree is a hierarchical form of data structure.
- ❖ In a tree\* there is a **parent-child relationship** between items.
  - Instead of items followed each other as in lists, queues and stacks
- ❖ To visualize what trees look like, imagine a tree growing up from the ground.
- ❖ Trees in programming are normally drawn downward, so you would be better off imagining the root structure of the tree growing downward.



\* More precisely in trees of interest to us, with oriented branches.

# Applications

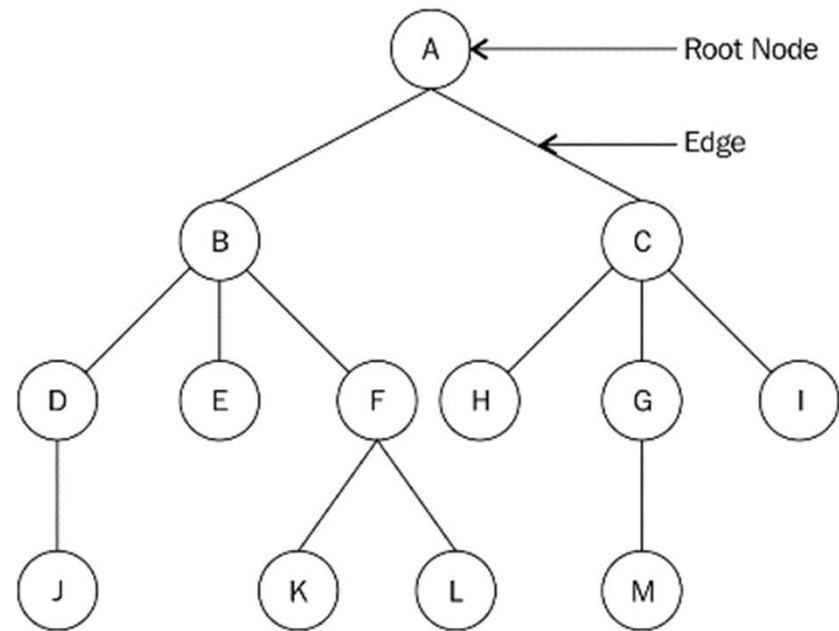
- ❖ Organograms
- ❖ Taxonomies
- ❖ Representation of hierarchical information
  - Example: filesystems
- ❖ The HTML/XML Document Object Model is organized in the form of a tree.
- ❖ Fast search
- ❖ ...



# Terminology

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- ❖ To understand trees, we need to first understand the basic ideas on which they rest.
- ❖ Terms associated with a Tree:
- ❖ **Node:**
  - Each circled alphabet represents a node.
  - A node is any structure that holds data.
- ❖ **Root node:**
  - It is the node **where the tree begins**
  - It is the only **node without a parent node**.
  - Is the only node from which all other nodes come.
  - The root node in our tree is the node A.
- ❖ **Edge:**
  - The connection between two nodes.



# Terminology (cont.)

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## ❖ Parent:

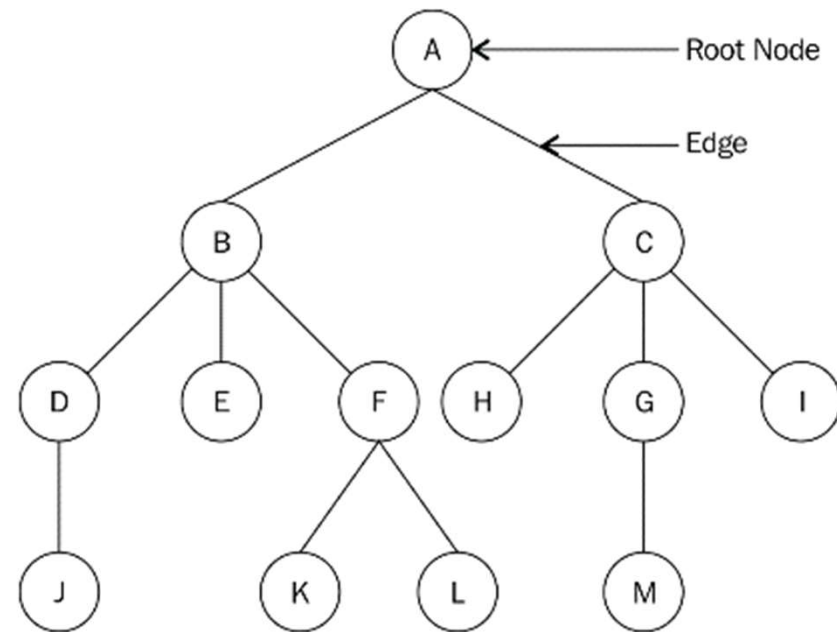
- A node in the tree with other connecting nodes is the parent of those nodes
- Node B is the parent of nodes D, E, and F.

## ❖ Child:

- This is a node connected to its parent.
- Nodes B and C are children of node A, the parent and root node.

## ❖ Sibling:

- All nodes with the same parent are siblings.
- This makes the nodes B and C siblings.



# Terminology (cont.)

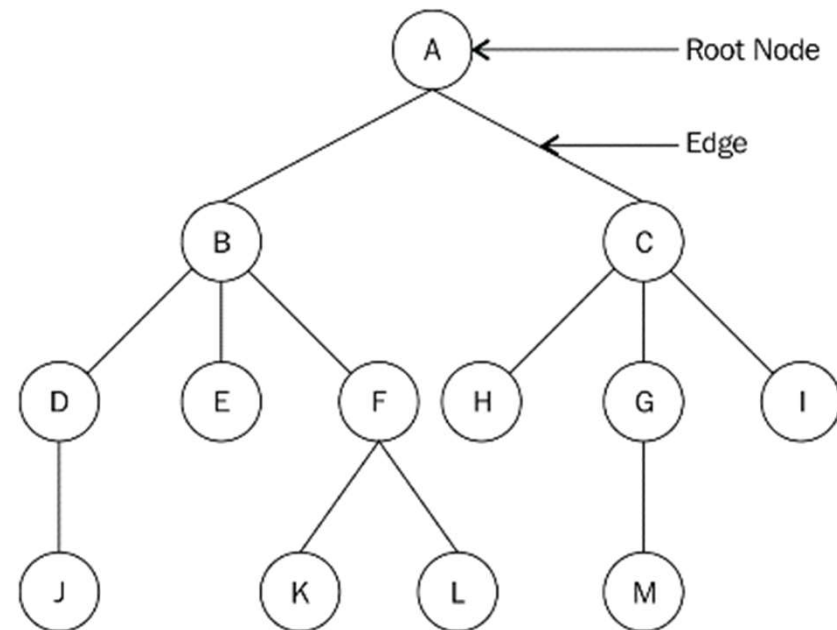
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## ❖ Degree:

- Number of children of the node.
- The degree of node A is 2.

## ❖ Leaf node:

- This is a node with a degree of 0.
- Nodes J, E, K, L, H, M, and I are all leaf nodes.





# Binary trees

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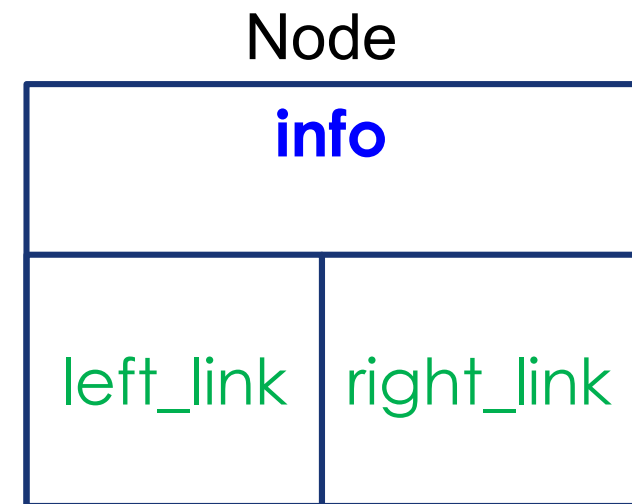
- ❖ Binary trees are a specific type of trees
- ❖ In a binary tree each node has a maximum of 2 children
  - They are designated as left child and right child

# Tree nodes

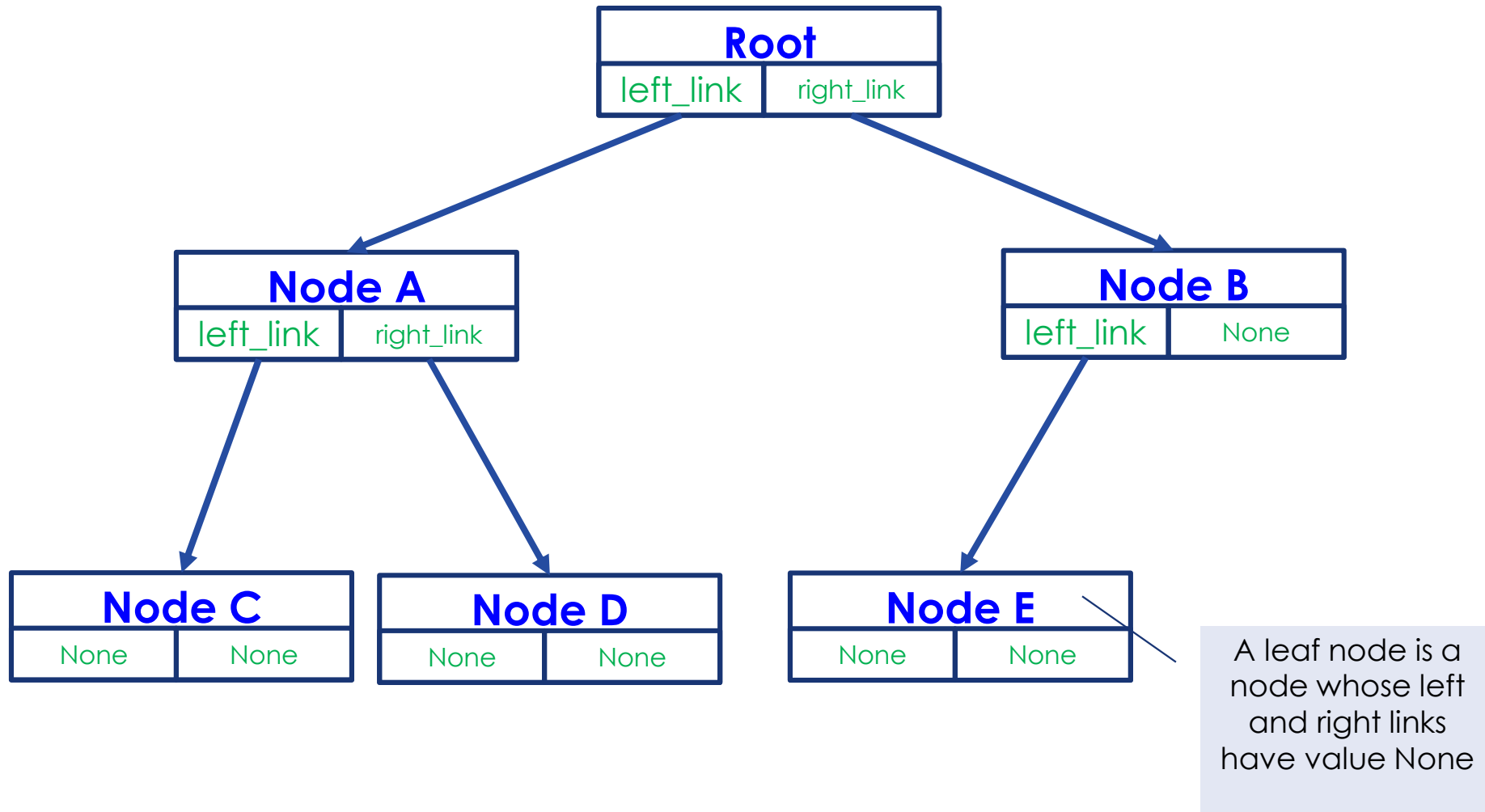
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- ❖ Like Linked Lists, **trees are built up of nodes**.
  - But now the nodes that make up a tree need to contain data about the parent-child relationship
- ❖ A simple binary tree node class in Python:

```
class Node:  
    def __init__(self, value):  
        self.info = value  
        self.right_link = None  
        self.left_link = None
```



# Example (binary tree)



# binarytree library

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- ❖ In Python, a binary tree can be represented in different ways with different data structures(dictionary, list) and class representation for a node.
- ❖ For example, [binarytree](#) library helps to directly implement a binary tree.
  - [Binarytree Module in Python - GeeksforGeeks](#)
- ❖ This module does not come pre-installed with Python's standard utility module.
  - To install it type the below command in the terminal.
  - `pip install binarytree`

# Creating a simple tree

```
from binarytree import Node
root = Node(3)
root.left = Node(6)
root.right = Node(8)
root.left.left = Node(10)

# Getting binary tree
print('Binary tree :', root)

# Getting list of nodes
print('List of nodes :', list(root))

# Getting inorder of nodes
print('Inorder of nodes :', root.inorder)

# Checking tree properties
print('Size of tree :', root.size)
print('Height of tree :', root.height)
```

Binary tree :

```
  3
 / \
6   8
/
10
```

List of nodes : [Node(3), Node(6), Node(8), Node(10)]

Inorder of nodes : [Node(10), Node(6), Node(3), Node(8)]

Size of tree : 4

Height of tree : 2

# (some) Tree operations

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## ❖ Insertion

## ❖ Traversals

- Since a binary tree is a non-linear data structure, there is more than one way to traverse through the tree data, including:

**inorder traversal**

**preorder traversal**

**postorder traversal.**

# Inorder Traversal

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- ❖ In an inorder traversal, the **left** child is visited first, followed by the **parent** node, then followed by the **right** child.

```
def inorder(node):  
    if node:  
        # call inorder on the left subtree until it reaches a leaf  
        inorder(node.left)  
  
        # Once we reach a leaf, we print the data  
        print(node.val) # binarytree node uses val  
  
        # Now, since the left subtree and the root has been printed,  
        #call inorder on right subtree until we reach a leaf node.  
        inorder(node.right)
```

# Example (inorder)

```
from binarytree import Node
```

```
root = Node(3)
```

```
root.left = Node(6)
```

```
root.right = Node(8)
```

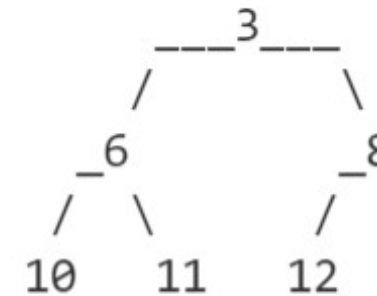
```
root.left.left = Node(10)
```

```
root.left.right = Node(11)
```

```
root.right.left = Node(12)
```

```
print(root)
```

```
inorder(root)
```



10  
6  
11  
3  
12  
8



# Preorder

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- ❖ In a preorder traversal, the **root node is visited first**, followed by the left child, then the right child.

```
def preorder(node):  
    if node:  
        # Print the value of the root node first  
        print(node.val)  
  
        # Recursively call preorder on the left subtree  
        preorder(node.left)  
  
        # Recursively call preorder on the right subtree  
        preorder(node.right)
```

# Example (preorder)

---

# For the tree,

#        10

#       /   \

#     34     89

#   /   \   /   \

# 20   45 56   54

# Preorder traversal: 10 34 20 45 89 56 54

# Postorder

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- ❖ In a postorder traversal, the left child is visited first, followed by the right child, then the root node.

```
def postorder(node):  
    if node:  
        # call postorder on the left subtree  
        postorder(node.left)  
  
        # call postorder on the right subtree  
        postorder(node.right)  
  
        # Print the value of the root node  
        print(node.val)
```

# Example (postorder)

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# For the tree,

#        10

#       /   \

#     34     89

#   /   \   /   \

# 20    45 56   54

# Postorder traversal: 20 45 34 56 54 89 10

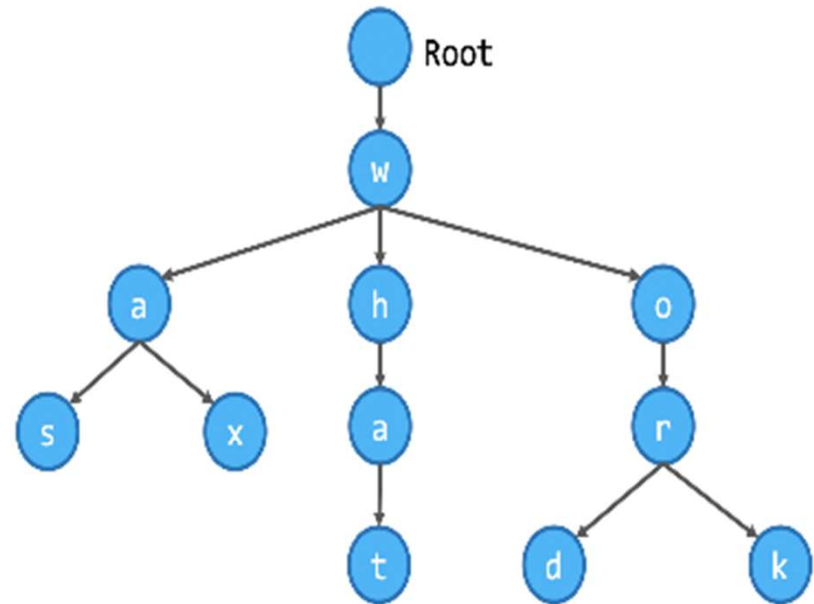
# Tries

It is not a typo 😊

# Trie

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- ❖ It is a tree-like data structure made up of nodes.
- ❖ Each node may have none, one or more children.
- ❖ When used to store a vocabulary, each node is used to store a character
  - consequently each "branch" of the trie represents a unique word.
- ❖ Example at right:
  - a trie with five words (was, wax, what, word, work) stored in it.



# Applications

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- ❖ Trie is a very useful data structure.
- ❖ It is commonly used to represent a dictionary for looking up words in a vocabulary.
- ❖ For example, consider the task of implementing a search bar with **auto-completion** or **query suggestion**.
  - When the user enters a query, the search bar will automatically suggest common queries **starting with** the characters input by the user.

# How does a Trie Work?

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- ❖ There are two major operations that can be performed on a trie, namely:
  - **Inserting** a word into the trie
  - **Searching** for words using a prefix
- ❖ Both operations involves **traversing the trie by starting from the root** node.



# Advantages

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- ❖ Fast access to the start of all words having specified initial characters
- ❖ Lower memory requirements
  - No need to repeat the characters

# Implementation example

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- ❖ You can find an implementation in Python at
  - [Implementing Trie in Python | Albert Au Yeung](#)
- ❖ It will be used in Practical Classes

# Example of use

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```
def main():  
  
    trie = trie_from_file("words.txt")  
  
    word = input('Beginning of the Word ?')  
    while len (word) > 0:  
  
        res = trie.query(word)  
  
        for w in res:      # res is a list  
            print(w[0])  
  
        word = input('Beginning of the Word ?')
```

# Graphs

# Graphs

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- ❖ Graphs have become a powerful means of modelling and capturing data in real-world scenarios such as social media networks, web pages and links, and locations and routes in GPS.
- ❖ If you have **a set of objects that are related to each other**, then you can represent them using a graph.

# Graphs

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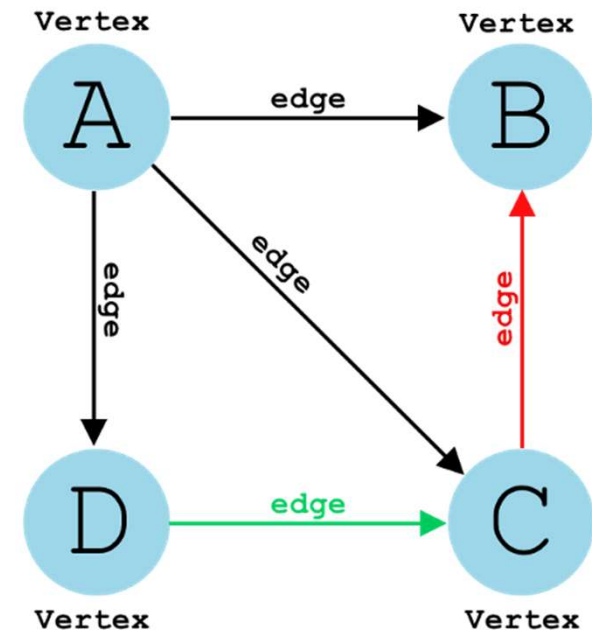
- ❖ Are complex, non-linear data structures

- ❖ Are characterized by a group of vertices...

- ❖ ... connected by edges.

- ❖ PT:

- Vertices: Vértices (ou nós)
- Edges: Ligações (ou arestas ou arcos)



# Graphs

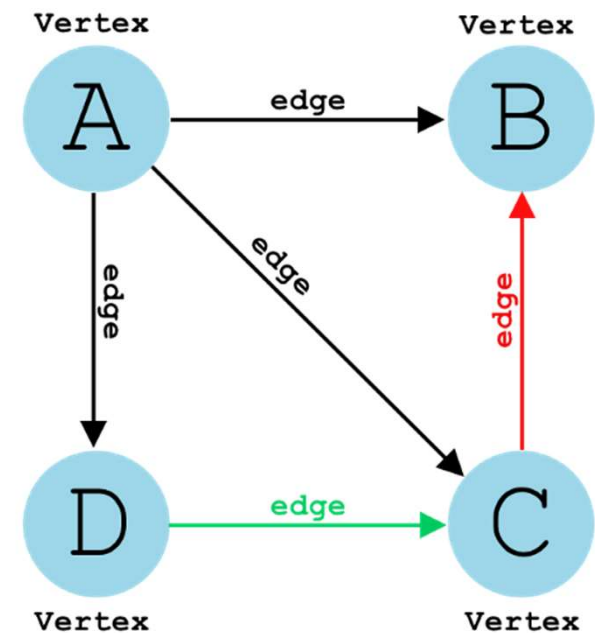
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## ❖ Vertices

- Represent entities in a graph.
- Every vertex has a value associated with it.
- Example: if we represent a list of cities using a graph, the vertices will represent the cities.

## ❖ Edges

- Represent the relationship between the vertices in the graph.
- Edges may or may not have a value associated with them.
- Example: if we represent a list of cities using a graph, the edges will represent the path between the cities.



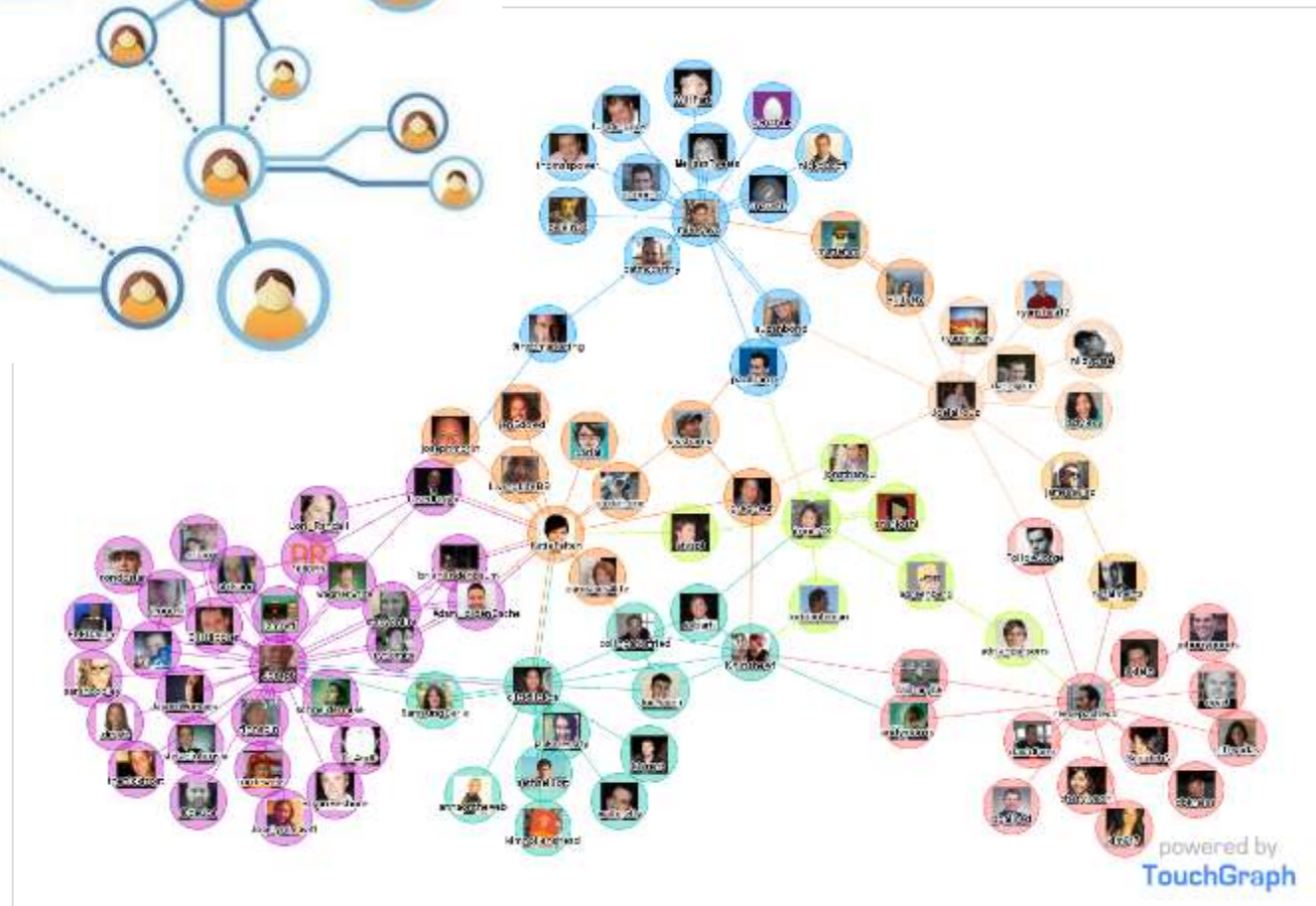
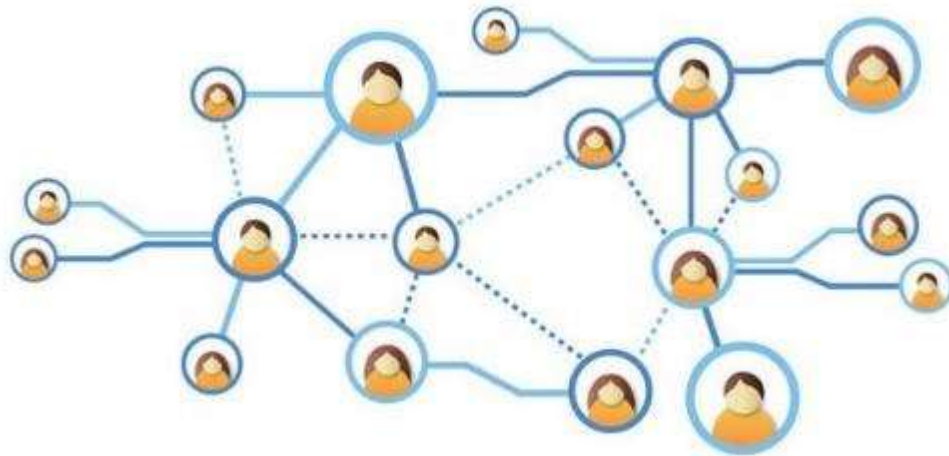
# Applications of Graphs

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- ❖ Graphs are used everywhere, from schooling to business.
  - Especially in the fields of computer science, physics, and chemistry.
- ❖ A few other applications of graphs are:
  - To visualize organized data.
  - Directed Graphs are used in Google's [Page Ranking Algorithm](#).
  - [Social Networks](#) use graphs to represent different users as vertices and edges to represent the connections between them.
  - In a mapping application, graphs are used to represent places and the path (distance) between them.

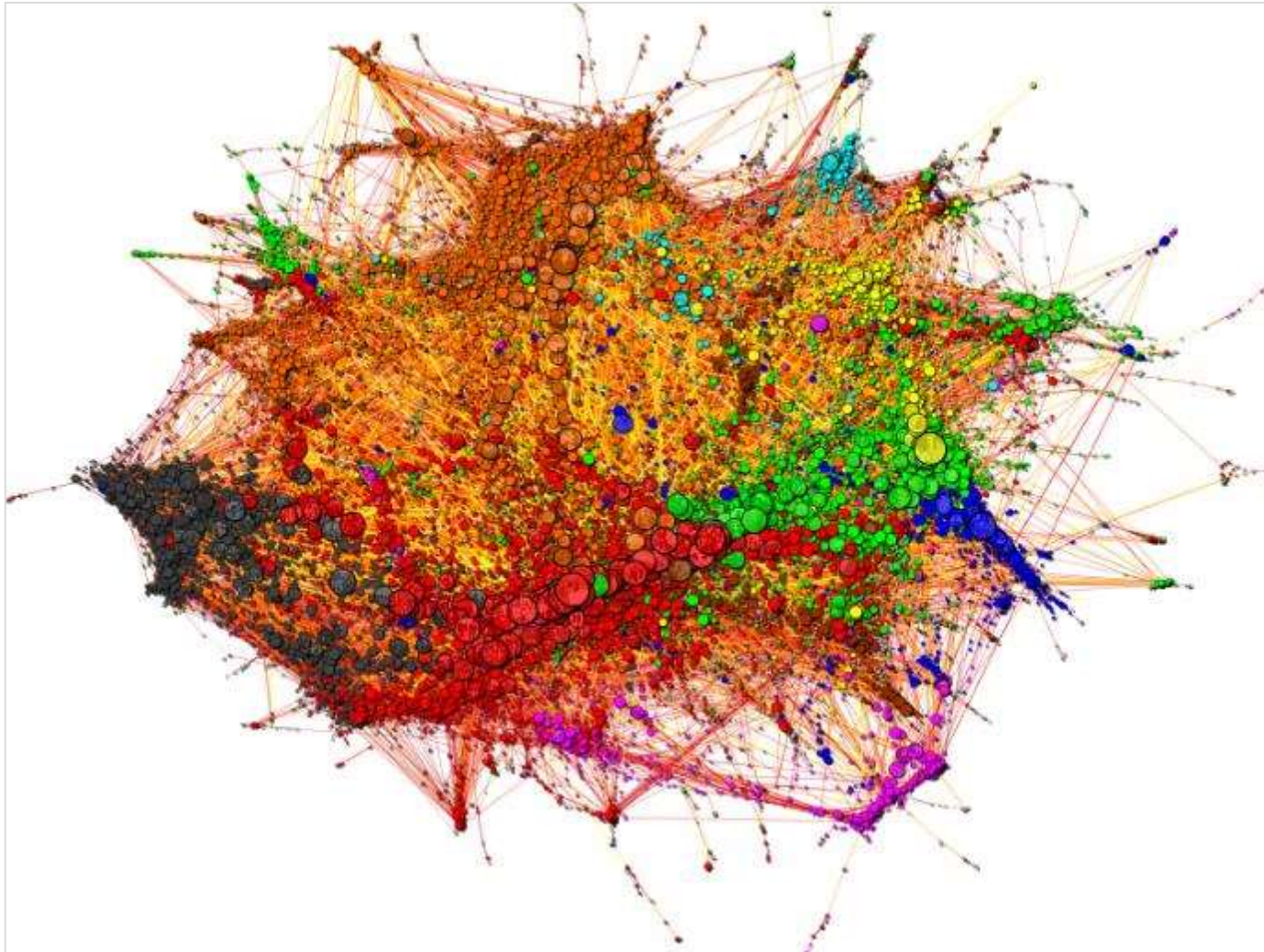


# Applications of Graphs – Social Network



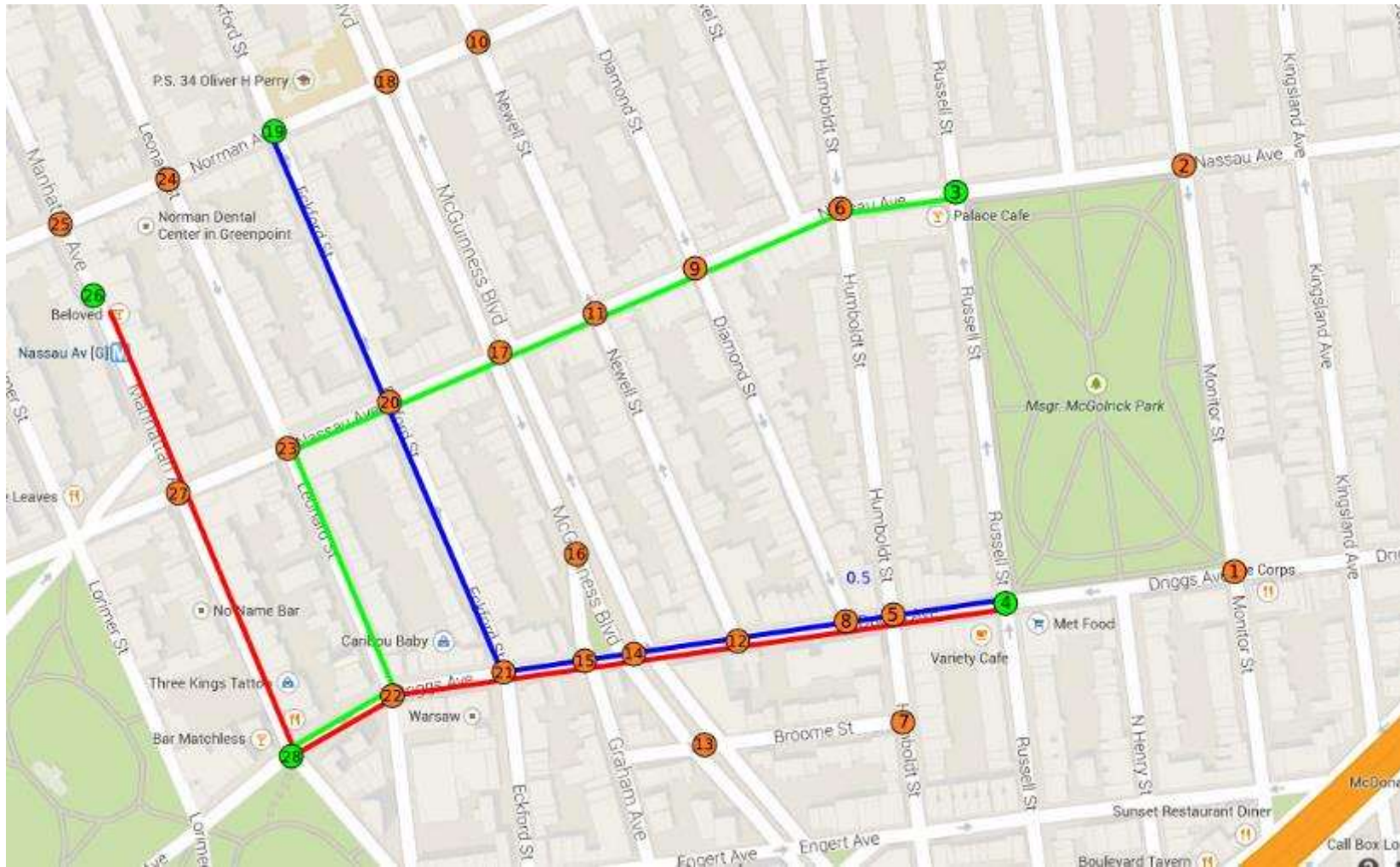
# Applications of Graphs – Web

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# Applications of Graphs



# Types of graphs

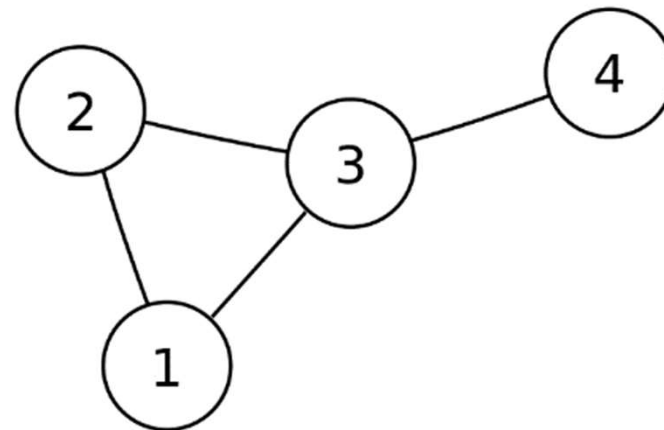
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- ❖ There are many types of graphs, based on weights, direction, interconnectivity, and special properties:

- ❖ **Undirected Graphs**

- In an undirected graph, the edges have **no direction**.
- If there is a path from vertex X to vertex Y, then there is a path from vertex Y to vertex X. Edge (X, Y) represents the edge connecting vertex X to vertex Y.

**That is,  $\text{edge}(X, Y) == \text{edge}(Y, X)$**



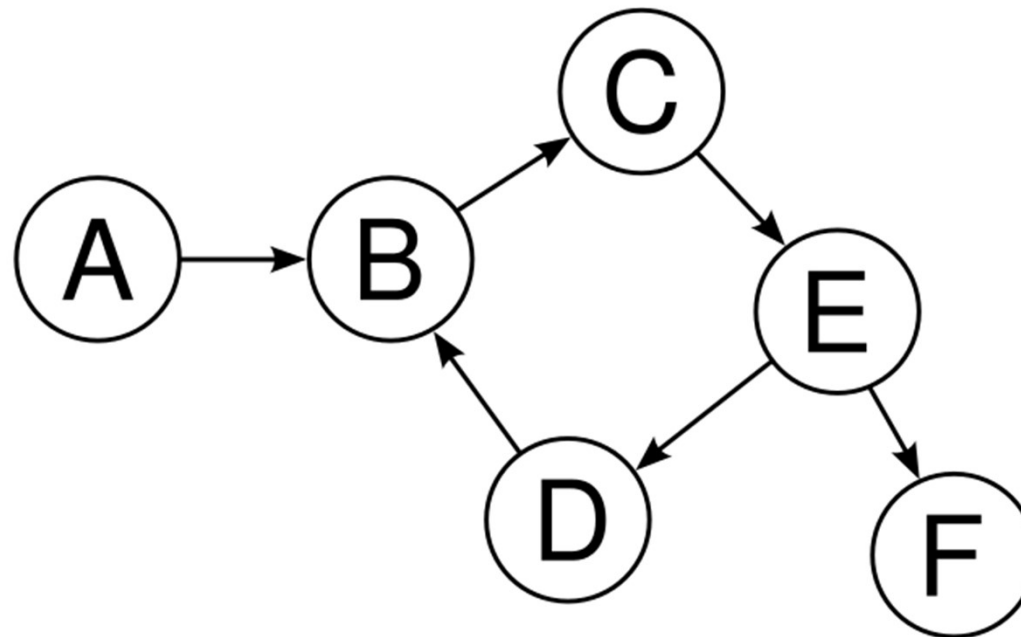
# Types of graphs

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## ❖ Directed Graphs

- In a directed graph or digraph, the edges have an orientation.
- If there is a path from vertex X to vertex Y, then there isn't necessarily a path from vertex Y to vertex X.

That is,  $\text{edge}(X, Y) \neq \text{edge}(Y, X)$



# Types of graphs

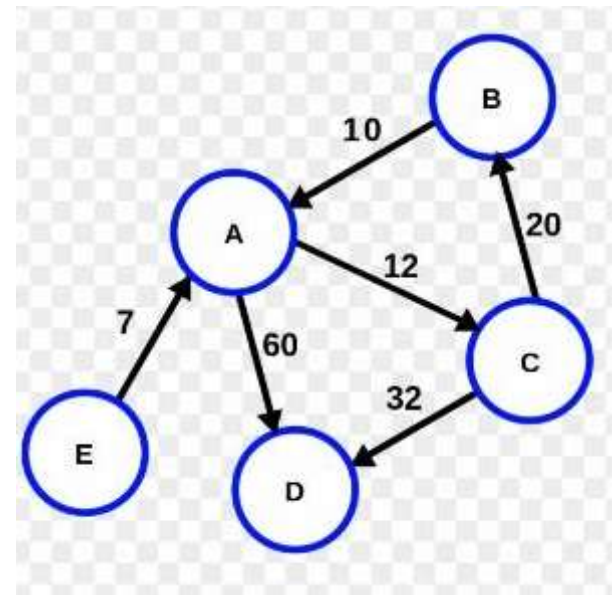
## ❖ Weighted Graphs

- A weighted graph has a value associated with every edge.
- The value may represent quantities like cost, distance, time, etc., depending on the graph.
- An edge of a weighted graph is represented as,  $(u, v, w)$ .

**u -> Source vertex**

**v -> Destination vertex**

**w -> Weight associated to go from u to v.**



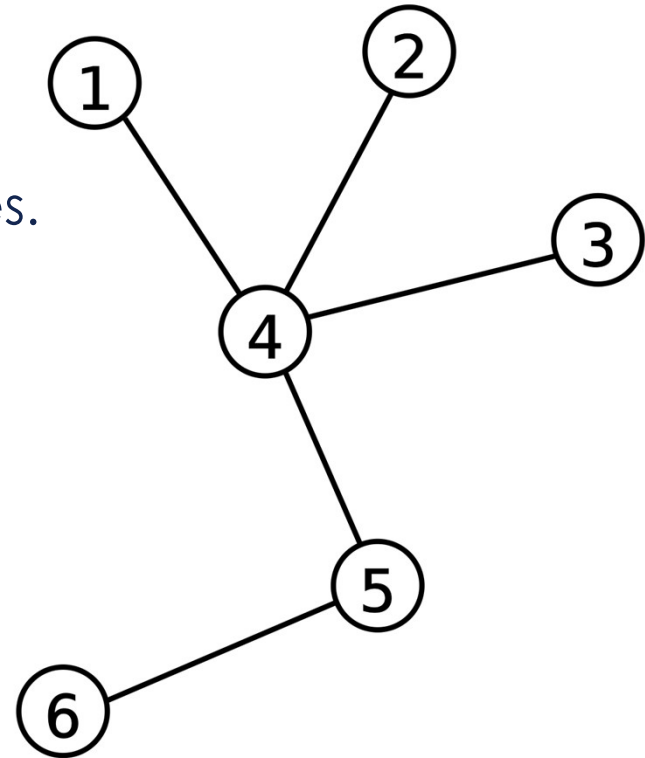
# Special Graphs

## ❖ Trees

- A Tree is a connected graph without cycles.

**A cycle in a graph is a sequence with the first and last vertices in the repeating sequence.**

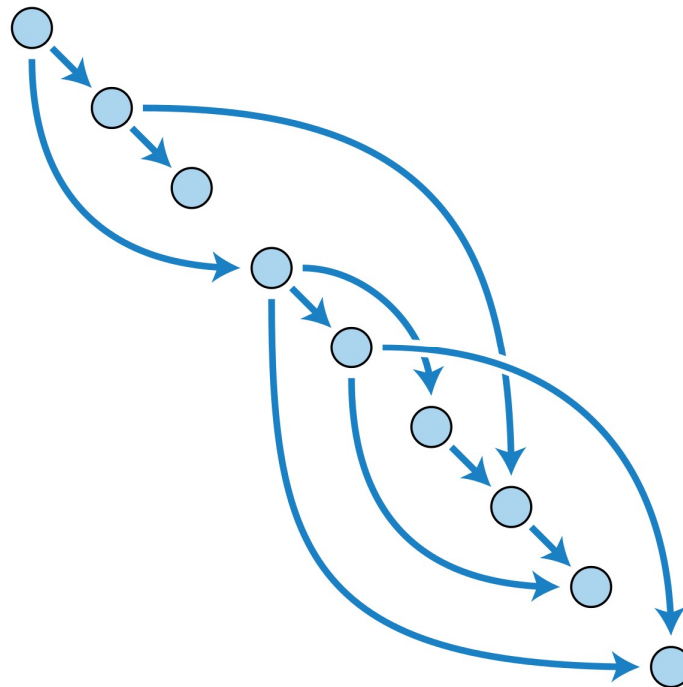
- They have  $X$  vertices and  $X-1$  edges.



# Special Graphs

## ❖ Directed Acyclic Graphs

- Directed Acyclic Graphs or DAGs are graphs with no directed cycles.
- They represent structures with dependencies.

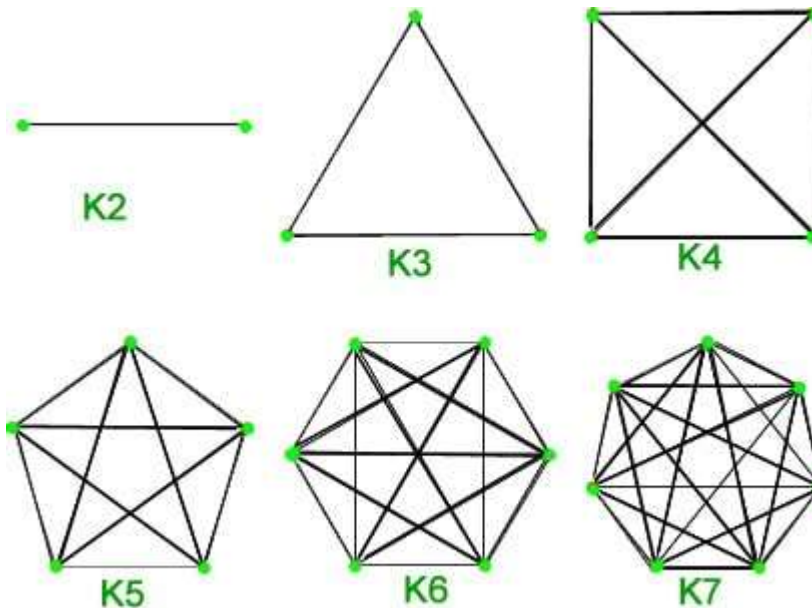




# Special Graphs

## ❖ Complete Graphs

- Complete graphs have a **unique edge between every pair of vertices**.
- A complete graph  $n$  vertices have  $(n*(n-1)) / 2$  edges and are represented by  $K_n$ .



# Representing Graphs

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- ❖ There are multiple ways of using data structures to represent a graph.
- ❖ The three most common ways are:
  - Adjacency Matrix
  - Adjacency List
  - Edge List

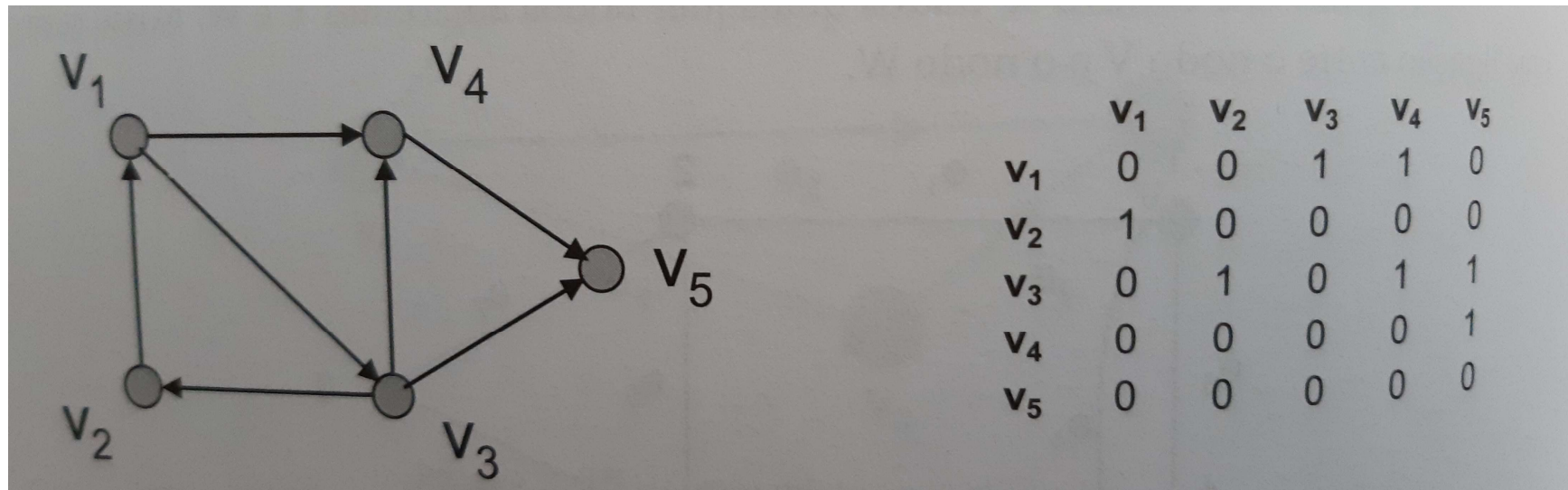
# Adjacency Matrix

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- ❖ It is a very simple way to represent a graph using a 2D array
- ❖ In a weighted graph, the element  $A[i][j]$  represents the cost of moving from vertex  $i$  to vertex  $j$ .
- ❖ In an unweighted graph, the element  $A[i][j]$  represents a Boolean value that determines if a path exists from vertex  $i$  to vertex  $j$ .
  - If  $A[i][j] == 0$ , then no path from vertex  $i$  to vertex  $j$  exists.
  - If  $A[i][j] == 1$ , there is a path from vertex  $i$  to vertex  $j$ .
- ❖ For example, a snake game can be represented by using an adjacency matrix.
  - This enables us to use various algorithms to find the shortest path to finish the game.
- ❖ Similarly, many shortest path algorithms use an adjacency matrix.

# Example

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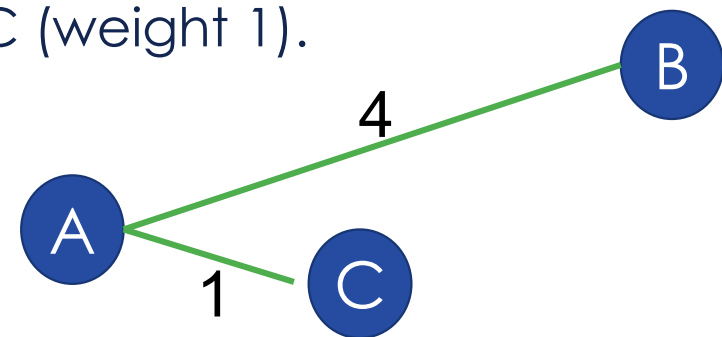


# Adjacency List

- ❖ An adjacency list represents a graph as a list that has vertex-edge mappings.

- ❖ Example:

- $A \rightarrow [(B, 4), (C, 1)]$
- represents an adjacency list where the vertex A is connected to B (weight 4) and C (weight 1).



- ❖ This works really well for sparse graphs.

# Edge list

---

❖ An edge list represents the graph as an unstructured list of edges.

❖ Example:

graph = [(C, A, 4), (A, C, 1), (B, C, 6),  
(A, B, 4), (C, B, 1), (C, D, 2)]

# Algorithms for Graphs

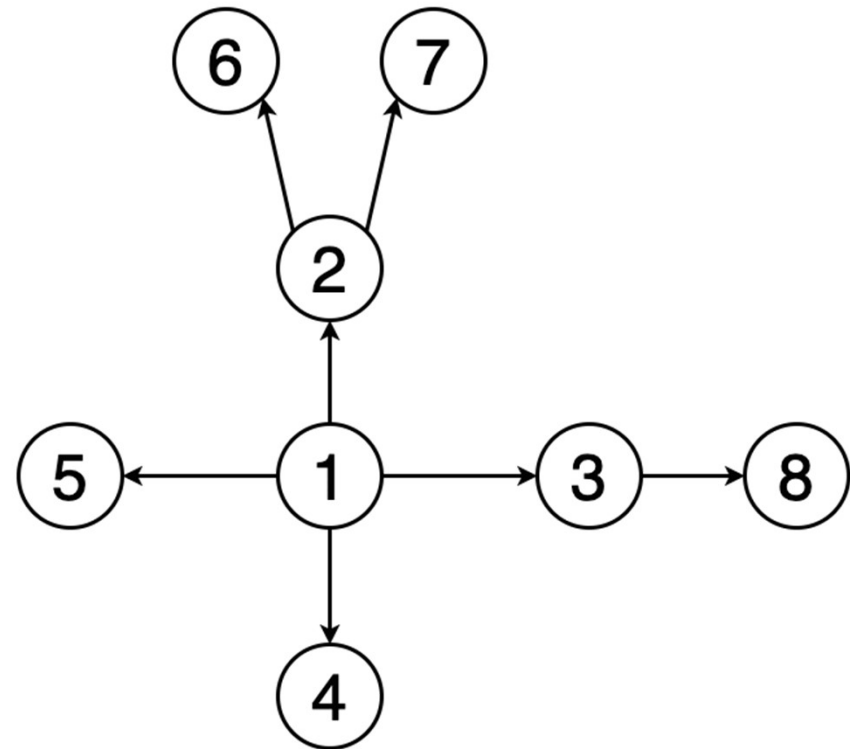
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- ❖ Traversing is one of the fundamental operations that can be performed on graphs.
  - Needed, for example, to display the graph and search
- ❖ In next slides the most relevant are Visually explained

# Breadth-first search (BFS)

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- ❖ In breadth-first search (BFS), we start at a particular vertex and **explore all its neighbors at the present depth** before moving on to the vertices in the next level.
- ❖ Unlike trees, graphs can contain cycles  
(a path where the first and last vertices are the same)
- ❖ Hence, we must **keep track of the visited vertices**.
  - a queue is used.



- ❖ Note how vertices are discovered (yellow) and get visited (red).



# Applications of BFS

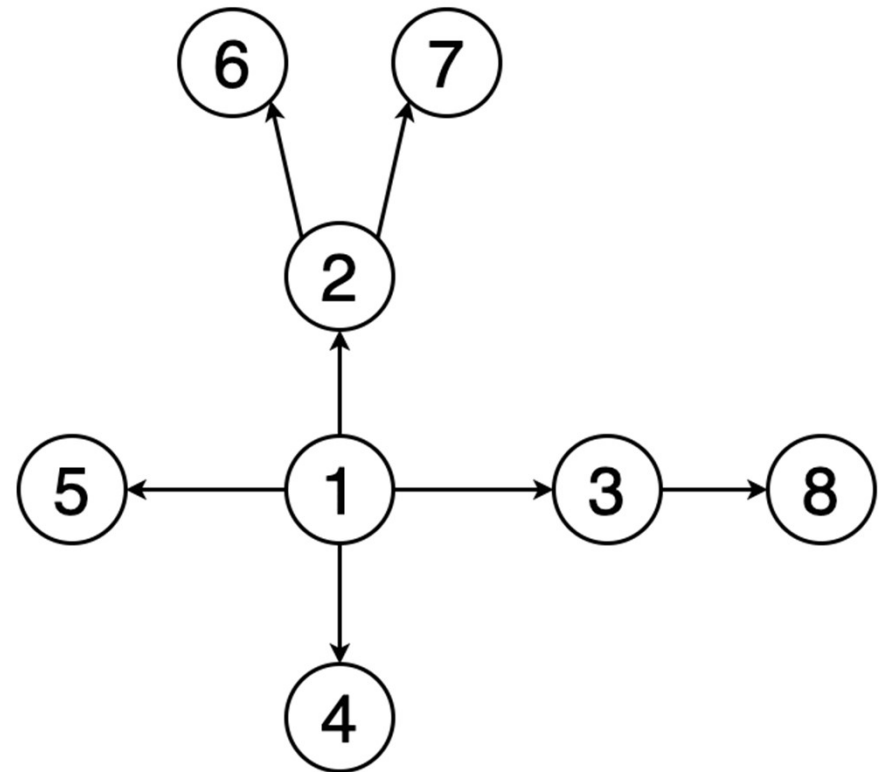
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- ❖ Used to determine the shortest paths and minimum spanning trees.
- ❖ Used by search engine crawlers to build indexes of web pages.
- ❖ Used to search on social networks.
- ❖ Used to find available neighbor nodes in peer-to-peer networks such as BitTorrent.

# Depth-first search (DFS)

---

- ❖ In depth-first search (DFS) we start from a particular vertex and **explore as far as possible along each branch** before retracing back (backtracking)
- ❖ In DFS also we must keep track of the visited vertices.
- ❖ When implementing DFS, we use a stack to support backtracking.



- ❖ Note how it traverses to the depths and backtracks.

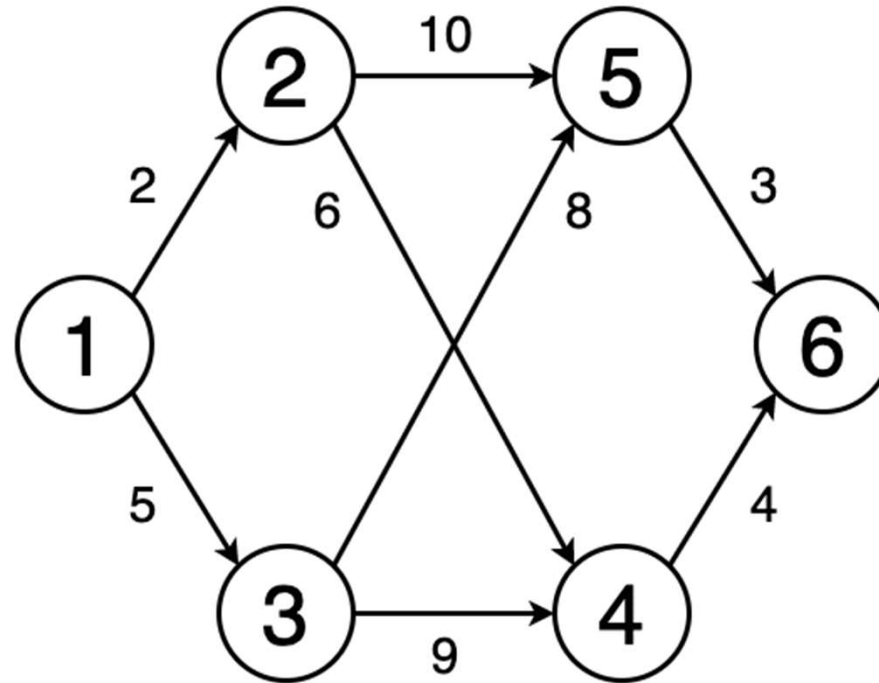
# Applications of DFS

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- ❖ Used to find a path between two vertices.
- ❖ Used to detect cycles in a graph.
- ❖ Used in topological sorting.
- ❖ Used to solve puzzles having only one solution
  - e.g., mazes

# Shortest path

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- ❖ The **shortest path** is a path in the graph such that the sum of the weights of the edges that should be travelled is minimum.
- ❖ In the animation the shortest path from vertex 1 to vertex 6 is determined

# Shortest path (cont.)

---

## ❖ Algorithms

- Dijkstra's shortest path algorithm
- Bellman–Ford algorithm

## ❖ Applications

- Used to find directions to travel from one location to another in mapping software like Google maps or Apple maps.
- Used in networking to solve the min-delay path problem.
- Used in abstract machines to determine the choices to reach a certain goal state via transitioning among different states

**e.g., can be used to determine the minimum possible number of moves to win a game**

# (Example of) Modules for graphs

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## ❖ NetworkX

- Python package for the creation, manipulation, and study of the structure, dynamics, and functions of complex networks.
- <https://networkx.org/>
- Data structures for graphs, digraphs, and multigraphs
- Many standard graph algorithms

## ❖ Python-igraph

- It is a library for creating, manipulating and analyzing graphs.
- It is intended to be as powerful (i.e. fast) as possible to enable working with large graphs
- <https://igraph.org/python/>

# Examples using NetworkX

- ❖ use of **Dijkstra's algorithm** to find the shortest weighted path:

```
>>> G = nx.Graph()
>>> e = [('a', 'b', 0.3), ('b', 'c', 0.9), ('a', 'c', 0.5), ('c', 'd', 1.2)]
>>> G.add_weighted_edges_from(e)
>>> print(nx.dijkstra_path(G, 'a', 'd'))
['a', 'c', 'd']
```

- Note the use of class Graph

- ❖ Drawing graphs

```
>>> import matplotlib.pyplot as plt
>>> G = nx.cubical_graph()
>>> subax1 = plt.subplot(121)
>>> nx.draw(G) # default spring_layout
>>> subax2 = plt.subplot(122)
>>> nx.draw(G, pos=nx.circular_layout(G), node_color='r', edge_color='b')
```

## Example 2 - Shortest path

```
import networkx as nx
```

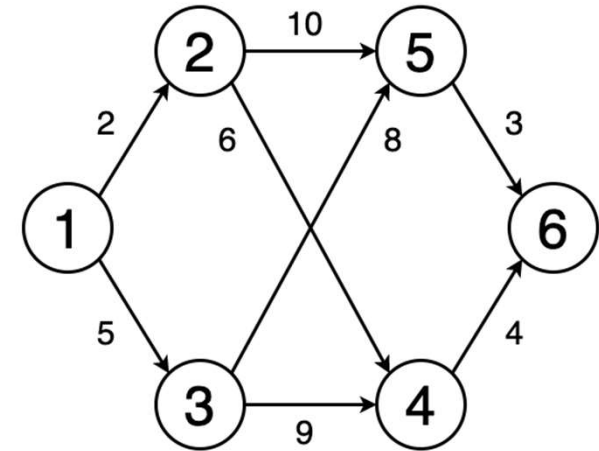
```
g = nx.Graph()
e = [ ('1', '2', 2),
      ('1', '3', 5),
      ('2', '4', 6),
      ('3', '5', 8),
      ('2', '5', 10),
      ('3', '4', 9),
      ('5', '6', 3),
      ('4', '6', 4)]
```

```
g.add_weighted_edges_from(e)
```

```
print(nx.dijkstra_path(g, '1', '6'))
```

Result:

```
['1', '2', '4', '6']
```





# More information

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## ❖ 10 Graph Algorithms Visually Explained

- A quick introduction to 10 basic graph algorithms with examples and visualisations
- <https://towardsdatascience.com/10-graph-algorithms-visually-explained-e57faa1336f3>

## ❖ You can check out the implementations of graph algorithms found in the [networkx](#) and [igraph](#) python modules.

- Python- igraph manual :  
<https://igraph.org/python/tutorial/latest/tutorial.html>

## ❖ You can read about python-igraph in article [Newbies Guide to Python-igraph](#).

# Object Orient Analysis, Design and Programming

Basics