方知蓦然回首之时 那人却已不在灯火阑珊处

### 关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

# 洛谷3350 [ZJOI2016]旅行者

2019-02-24 题解

给定一个网格图, 边有边权, 给定M个查询, 求两点最短路。

考虑离线所有询问,然后分治,假设当前分治一条横穿中心的线,枚举每个线上节点更新当前所有询问的答案。

- 若某查询的两点在此线两侧,则最短路一定能在此次被更新
- 若某查询的两点在此线同侧,则最短路可能被当前更新,可以递归下去做。

可以证明每次选当前分治矩形的最长边,时间复杂度为 $O(n\sqrt{n}\log n)$ 。

#### 代码:

1

方知蓦然回首之时 那人却已不在灯火阑珊处

### 关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

```
template <class T> inline void read(T &x) {
    x = 0; register char c = getchar(); register bool f = 0;
    while (!isdigit(c)) f ^= c == '-', c = getchar();
    while (isdigit(c)) x = x * 10 + c - '0', c = getchar();
    if (f) x = -x;
template <class T> inline void print(T x) {
    if (x < 0) putchar('-'), x = -x;
    if (x > 9) print(x / 10);
    putchar('0' + x % 10);
template <class T> inline void print(T x, char c) { print(x), putchar(c); }
const int N = 1e^5 + 10, M = 210, mov[4][2] = \{\{0, 1\}, \{1, 0\}, \{0, -1\}, \{-1, 0\}\}
int n, m, p, ans[N];
std::vector <std::vector <int> > map[4], dis;
struct node {
    int x, y;
} u, v;
struct status {
    node u; int w;
    inline bool operator < (const status &other) const {</pre>
        return w > other.w;
} now; std::priority queue <status> q;
struct info {
    int xl, xr, yl, yr, id;
}; std::vector <info> vet;
template <class T> void resize(T &a, int n, int m) {
    for (int k = 0; k < 4; k++) {
        a.resize(n + 1);
        for (int i = 0; i <= n; i++) a[i].resize(m + 1);
    }
```

https://memset0.cn/luogu3350

方知蓦然回首之时 那人却已不在灯火阑珊处

### 关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

```
}
template <class T> void memset(T &a, int w, int xl, int xr, int yl, int yr) {
    for (int i = xl; i <= xr; i++)
        for (int j = yl; j <= yr; j++)
            a[i][i] = w;
}
void spfa(int xs, int ys, int xl, int xr, int yl, int yr) {
    memset(dis, 1e9, xl, xr, yl, yr);
    q.push((status)\{(node)\{xs, ys\}, 0\}), dis[xs][ys] = 0;
    while (q.size()) {
        now = q.top(), u = now.u, q.pop(); if (now.w != dis[now.u.x][now.u.y]
        for (int t = 0, w; t < 4; t++) {
            v.x = u.x + mov[t][0], v.y = u.y + mov[t][1], w = map[t][u.x][u.y]
            if (v.x < xl \mid | v.x > xr \mid | v.y < yl \mid | v.y > yr) continue;
            if (dis[u.x][u.y] + w < dis[v.x][v.y]) {
                dis[v.x][v.y] = dis[u.x][u.y] + w;
                q.push((status){v, dis[v.x][v.y]});
            }
}
void solve(int x1, int xr, int y1, int yr, std::vector <info> vet) {
    if (xl > xr || yl > yr || !vet.size()) return;
    std::vector <info> l_vet, r_vet;
    if ((xr - xl + 1) * (yr - yl + 1) < 100) {
        for (std::vector <info> ::iterator it = vet.begin(); it != vet.end();
            spfa(it->xl, it->yl, xl, xr, yl, yr);
            ans[it->id] = std::min(ans[it->id], dis[it->xr][it->yr]);
        } return;
    if ((xr - xl + 1) > (yr - yl + 1)) {
        int mid = (xl + xr) \gg 1;
        for (int i = yl; i <= yr; i++) {
            spfa(mid, i, xl, xr, yl, yr);
```

https://memset0.cn/luogu3350

方知蓦然回首之时 那人却已不在灯火阑珊处

#### 关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

```
for (std::vector <info> ::iterator it = vet.begin(); it != vet.en
                ans[it->id] = std::min(ans[it->id], dis[it->xl][it->yl] + dis
        for (std::vector <info> ::iterator it = vet.begin(); it != vet.end();
            if (it->xl < mid && it->xr < mid) l vet.push back(*it);</pre>
            if (it->xl > mid && it->xr > mid) r vet.push back(*it);
        solve(xl, mid - 1, yl, yr, l vet), l vet.clear();
        solve(mid + 1, xr, yl, yr, r vet), r vet.clear();
    } else {
        int mid = (vl + vr) >> 1;
        for (int i = xl; i <= xr; i++) {
            spfa(i, mid, xl, xr, yl, yr);
            for (std::vector <info> ::iterator it = vet.begin(); it != vet.en
                ans[it->id] = std::min(ans[it->id], dis[it->xl][it->yl] + dis
        for (std::vector <info> ::iterator it = vet.begin(); it != vet.end();
            if (it->yl < mid && it->yr < mid) l vet.push back(*it);</pre>
            if (it->vl > mid && it->vr > mid) r vet.push back(*it);
        }
        solve(xl, xr, yl, mid - 1, l vet), l vet.clear();
        solve(x1, xr, mid + 1, yr, r vet), r vet.clear();
   }
}
void main() {
    read(n), read(m), resize(dis, n, m);
    for (int i = 0; i < 4; i++) resize(map[i], n, m);
    for (int i = 1; i <= n; i++)
        for (int j = 1; j < m; j++)
            read(map[0][i][j]), map[2][i][j + 1] = map[0][i][j];
    for (int i = 1; i < n; i++)
        for (int j = 1; j <= m; j++)
            read(map[1][i][j]), map[3][i + 1][j] = map[1][i][j];
    read(p);
    for (int i = 1; i \le p; i++) ans[i] = 1e9;
    for (int i = 1, xl, xr, yl, yr; i <= p; i++) {
```

https://memset0.cn/luogu3350

#### 洛谷3350 [ZJOI2016]旅行者 - memset0 的博客

read(x1), read(y1), read(xr), read(yr);
vet.push\_back((info){x1, xr, y1, yr, i});

for (int i = 1; i <= p; i++) print(ans[i], '\n');</pre>

## memseto's Notebook

方知蓦然回首之时 那人却已不在灯火阑珊处 巧妙的思路 分治 最短路

} solve(1, n, 1, m, vet);

} signed main() { return ringo::main(), 0; }

关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

用户名 邮箱 网址 (选填) 可以在这里写评论哦 ~

T

方知蓦然回首之时 那人却已不在灯火阑珊处 MySQL 数据库基本操作学习笔记 上一篇 « UOJ 社区版安装笔记

在这里输入关键字哦~(回车搜索)

关于我 友情链接 文章聚合

Theme Ringo by memseto
Proudly powered by Typecho

ſ