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多项式复合逆学习笔记

2019-02-25 算法

定义:给定两个 n 次多项式 F(x) 和 G(x) ,若对于任意多项式 P(x) 都有 G(F(P))=P 则称 G(x) 为 F(x) 的复合逆在模 x^n 意义下的复合逆。可以证明,若两个多项式常数项为 0 且一次项不为 0 则复合逆唯一且满足 F(G(x))=G(F(x))=x。

遗憾的是,多项式复合逆没有 $o(n\log n)$ 的做法,但我们可以以 $O(n\log n)$ 的复杂度求出某一项,或者 $O(n^2)$ 的复杂度求出所有项。

拉格朗日反演即

$$[x^n]F(x) = rac{1}{n}[x^{-1}]rac{1}{G^n(x)}$$

可以证明

$$[x^n]F(x) = rac{1}{n}[x^{n-1}](rac{x}{G(x)})^n$$

后者可以直接快速幂(两只 \log),或者转换为 \ln 和 \exp ,可以参考 关于求 多项式 k 次幂的一些思考。

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如果需要证明可以参考 zjt 大爷的博客 (我是不会)。

NFLSOJ332 多项式复合逆 ▷ 2019-02-25

暴力做即可,甚至不需要多项式 \ln 和多项式 \exp 。复杂度 $O(n\log n)$ 。

代码:

```
author: memset0
    date: 2019.02.25 15:23:19
    website: https://memset0.cn/
// ============
#include <bits/stdc++.h>
#define 11 long long
#define poly std::vector <int>
#define for each(i, a) for (int i = 0, lim = a.size(); i < lim; ++i)</pre>
namespace ringo {
template <class T> inline void read(T &x) {
    x = 0; register char c = getchar(); register bool f = 0;
   while (!isdigit(c)) f ^= c == '-', c = getchar();
    while (isdigit(c)) x = x * 10 + c - '0', c = getchar();
    if (f) x = -x;
template <class T> inline void print(T x) {
   if (x < 0) putchar('-'), x = -x;
   if (x > 9) print(x / 10);
    putchar('0' + x % 10);
template <class T> inline void print(T x, char c) { print(x), putchar(c); }
inline void print(const poly &a) { for each(i, a) print(a[i], " \n"[i == li
inline void read(poly &a, int n) { for (int i = 0, x; i < n; i++) read(x), a.
const int N = 1e3 + 10, mod = 998244353;
```

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```
namespace poly namespace {
    const int M = N << 3, SIZE = sizeof(int);</pre>
    int w[M], rev[M];
    inline poly resize(poly f, int n) { return f.resize(n), f; }
    inline int dec(int a, int b) \{ a = b; return a < \emptyset \} a + mod : a; \}
    inline int sub(int a, int b) { a += b; return a >= mod ? a - mod : a; }
    inline int inv(int x) { return x < 2 ? 1 : (11)(mod - mod / x) * inv(mod
    inline int fpow(int a, int b) { int s = 1; for (; b; b >>= 1, a = (11)a *
    inline poly operator + (poly f, int a) { f[0] = sub(f[0], a); return f; }
    inline poly operator + (int a, poly f) { f[0] = sub(a, f[0]); return f; }
    inline poly operator - (poly f, int a) { f[0] = dec(f[0], a); return f; }
    inline poly operator - (int a, poly f) { for each(i, f) f[i] = dec(0, f[i])
    inline poly operator * (poly f, int a) { for each(i, f) f[i] = (ll)f[i] *
    inline poly operator * (int a, poly f) { for each(i, f) f[i] = (ll)f[i] *
    inline poly operator + (poly f, const poly &g) {
        f.resize(std::max(f.size(), g.size()));
        for each(i, f) f[i] = sub(i < f.size())? f[i] : 0, i < g.size()? g[i]
        return f;
    }
    inline poly operator - (poly f, const poly &g) {
        f.resize(std::max(f.size(), g.size()));
        for each(i, f) f[i] = dec(i < f.size() ? f[i] : 0, i < g.size() ? g[i
        return f;
    namespace cipolla namespace {
        int t, sqr w;
        typedef std::pair <int, int> pair;
        inline pair operator * (const pair &a, const pair &b) {
            return std::make pair(((11)a.first * b.first + (11)a.second * b.s
                ((11)a.first * b.second + (11)a.second * b.first) % mod);
        int cipolla(int x) {
            do t = rand() % mod; while (fpow(sqr w = dec((11)t * t % mod, x),
            pair s = std::make pair(1, 0), a = std::make pair(t, 1);
            for (int b = (mod + 1) >> 1; b; b >>= 1, a = a * a) if (b & 1) s
            return std::min(s.first, mod - s.first);
        }
```

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```
} using cipolla namespace::cipolla;
void ntt(int *a, int lim) {
    for (int i = 0; i < \lim; i++) if (i < rev[i]) std::swap(a[i], a[rev[i
    for (int len = 1; len < \lim; len <<= 1)
        for (int i = 0; i < \lim; i += (len << 1))
             for (int j = 0; j < len; j++) {
                 int x = a[i + j], y = (ll)w[j + len] * a[i + j + len] % m
                 a[i + j] = sub(x, y), a[i + j + len] = dec(x, y);
            }
int init(int len) {
    int \lim = 1, k = 0; while (\lim < \operatorname{len}) \lim <<= 1, ++k;
    for (int i = 0; i < \lim_{i \to +} |rev[i]| = (rev[i]) > 1 > 1 > 1 | ((i & 1))
    return lim;
}
void main init() {
    for (int len = 1, wn; (len \langle\langle 1\rangle\rangle \langle M\rangle; len \langle\langle =1\rangle\rangle {
        wn = fpow(3, (mod - 1) / (len << 1)), w[len] = 1;
        for (int i = 1; i < len; i++) w[i + len] = (ll)w[i + len - 1] * w
    }
inline poly operator * (const poly &f, const poly &g) {
    static int a[M], b[M];
    int lim = init(f.size() + g.size() - 1), inv_lim = inv(lim); poly h;
    memset(&a[f.size()], 0, (lim - f.size()) * SIZE); for each(i, f) a[i]
    memset(&b[g.size()], 0, (lim - g.size()) * SIZE); for each(i, g) b[i]
    ntt(a, lim), ntt(b, lim);
    for (int i = 0; i < \lim; i++) a[i] = (11)a[i] * b[i] % mod;
    std::reverse(a + 1, a + lim), ntt(a, lim);
    for (int i = 0, l = f.size() + g.size() - 1; <math>i < 1; i++) h.push back(
    return h;
inline poly inv(const poly &f) {
    static int a[M], b[M];
    poly g(1, inv(f[0]));
    for (int len = 2; (len >> 1) < f.size(); len <<= 1) {
        int lim = init(len << 1), inv lim = inv(lim);</pre>
```

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```
memset(&a[len], 0, len * SIZE); for (int i = 0; i < len; i++) a[i</pre>
        memset(&b[len], 0, len * SIZE); for (int i = 0; i < len; i++) b[i</pre>
        ntt(a, lim), ntt(b, lim);
        for (int i = 0; i < \lim; i++) a[i] = (11)a[i] * b[i] % mod * b[i]
        std::reverse(a + 1, a + lim), ntt(a, lim), g.resize(len);
        for each(i, g) g[i] = dec(sub(g[i], g[i]), (ll)a[i] * inv lim % m
    } return g.resize(f.size()), g;
inline poly sqrt(const poly &f) {
    poly g(1, cipolla(f[0]));
    for (int len = 2; (len \Rightarrow 1) < f.size(); len \iff 1)
        g = resize(resize(resize(g * g, len) + f, len) * inv(resize(2 * g
    return g.resize(f.size()), g;
inline poly deri(const poly &f) {
    poly g;
    for (int i = 0; i < f.size() - 1; i++) g.push back((l1)(i + 1) * f[i
    return g.push back(∅), g;
}
inline poly inte(poly f) {
    poly g(1, 0);
    for (int i = 0; i < f.size() - 1; i++) g.push back((11)inv(i + 1) * f
    return g;
inline poly ln(const poly &f) { return inte(resize(deri(f) * inv(f), f.si
inline poly exp(const poly &f) {
   poly g(1, 1);
    for (int len = 2; (len >> 1) < f.size(); len <<= 1)
        g = resize(g * (1 - ln(resize(g, len)) + resize(f, len)), len);
   return g.resize(f.size()), g;
inline poly fpow(poly a, int b) {
    int n = a.size(); poly s(1, 1);
    for (; b; b >>= 1, a = resize(a * a, n))
        if (b \& 1) s = resize(s * a, n);
    return s;
}
```

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```
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} using namespace poly namespace;
int n; poly f, g, h, t;
void main() {
   read(n), read(f, n), g = poly(1), t = poly(1, 1);
   h = f, h.erase(h.begin()), h.push_back(₀), h = inv(h);
   for (int i = 1; i < n; i++) {
       t = t * h, t.resize(n);
       g.push_back((ll)inv(i) * t[i - 1] % mod);
   } print(g);
} signed main() { return ringo::main init(), ringo::main(), 0; }
多项式对数函数
                  多项式快速幂
                                  多项式指数函数
                                                     多项式复合逆
```

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邮箱

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可以在这里写评论哦~

1

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LOJ6287 诗歌 上一篇«

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