

## Discrete Math 2: Experience One

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**Rules:** Complete the following problems by working in groups of at most three, and submit solutions (one document per group) typeset using L<sup>A</sup>T<sub>E</sub>X by blah blah blah.

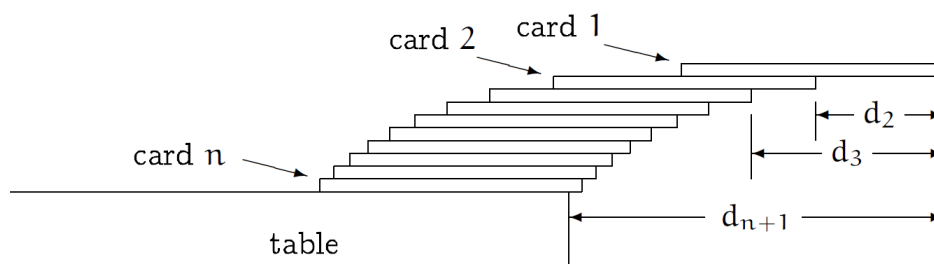
**Harmonic Numbers.** Recall that the **harmonic series** is the series whose terms are the reciprocals of positive integers, and that we defined the  $n^{\text{th}}$  **harmonic number**  $H_n$  to be

$$H_n = \sum_{i=1}^n \frac{1}{i}.$$

Their name supposedly comes from the fact that a tone with wavelength  $1/n$  is the  $n^{\text{th}}$  *harmonic* of a tone whose wavelength is 1.

**Card Trick.** The numbers  $H_n$  arise naturally in simple situations.

Suppose you have  $n$  cards and a table. Create the largest possible overhang by stacking the cards up over the table's edge (see the figure below for clarification).



Assume the edges of the cards are parallel to the edge of the table (we cannot rotate the cards to make corners stick out further), and assume each card is 2 units long (this will make calculations a bit easier).

**Problem 1:** Determine with proof the maximum overhang you can achieve with 52 cards, measuring in terms of card lengths.

Set-up: Referring to the figure for clarification, put  $d_k$  equal to the distance from the extreme edge of the top card to the corresponding edge of the  $k^{\text{th}}$  card from the top. So,  $d_1 = 0$ , and  $d_{k+1}$  should be the center of gravity of the first  $k$  cards.

Note that the center of gravity of  $k$  objects with respective weights  $w_1, \dots, w_k$  and with respective centers of gravity at positions  $p_1, \dots, p_k$ , is at position

$$\frac{w_1 p_1 + \dots + w_k p_k}{w_1 + \dots + w_k}.$$

Assume the cards have homogeneous density and equal weight.

Derive the recurrence

$$\begin{aligned} k d_{k+1} &= k + d_1 + \dots + d_{k-1} + d_k, \quad k \geq 0; \\ (k-1) d_k &= k-1 + d_1 + \dots + d_{k-1}, \quad k \geq 1. \end{aligned}$$

Deduce

$$d_{k+1} = d_k + \frac{1}{k}, \quad \text{and hence} \quad d_{k+1} = H_k.$$

Subtract the second recurrence from the first:

$$kd_{k+1} = k + d_1 + d_2 + \dots + d_{k-1} + d_k$$

$$(k-1)d_k = k-1 + d_1 + d_2 + \dots + d_{k-1}$$

Most of the right-hand terms will cancel out, leaving us with:

$$kd_{k+1} - kd_k + d_k = 1 + d_k$$

$$kd_{k+1} - kd_k = 1$$

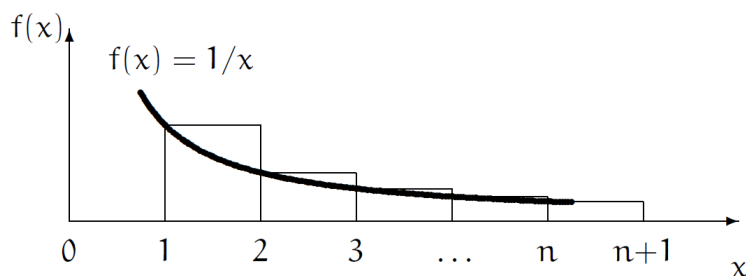
$$d_{k+1} - d_k = \frac{1}{k}$$

$$d_{k+1} = d_k + \frac{1}{k}$$

In class we proved (or will prove) that the harmonic series ( $H_n$  with  $n \rightarrow \infty$ ) *diverges* (meaning that if you specify some, even huge-a\$\$, number  $M$ , you can find an  $n$  so that  $H_n > M$ ). The proof that the harmonic series diverges showed (will show) that (with base of logarithm being 2)

$$\frac{\lfloor \log n \rfloor + 1}{2} < H_n \leq \lfloor \log n \rfloor + 1;$$

so, the proof method allows to determine  $H_n$  within a factor of 2. But a better bound can be had.

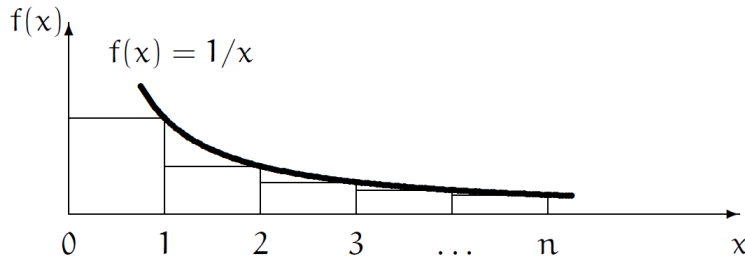


The natural logarithm is defined as the area under the curve  $y = 1/x$ ; the area between 1 and  $n$  is less than the area of the  $n$  rectangles in the figure above. So

$$\int_1^n \frac{dx}{x} = \ln n < H_n.$$

But placing the rectangles differently, as in the figure below, gives

$$\ln n < H_n < \ln n + 1, \quad \text{for } n > 1.$$



Now, we now know  $H_n$  within an error of at most 1. But we can do better.

we can use the plug and chug method by finding the 53<sup>rd</sup> harmonic number and that should give us the value of the max hang over length of 52 cards.

We've already established that a low approximation for  $H_n$  is  $\ln(n)$ . In class we talked about how  $\gamma$  is a constant for the area of  $H_n$  that isn't under the curve  $1/x$ , and thus is left out when we take  $\ln(n)$ . So all that remains to do is find  $\ln(n) + \gamma$ , which will give us an approximation for  $H_n$ . The precision of this approximation is only limited by our approximation of  $\gamma$ .

$$\begin{aligned}
 n &= 1000000 \\
 H_n &= 14.392726722864989 \\
 \gamma + \ln(n) &= 14.392726272863225 \\
 n &= 51 \\
 H_n &= 4.518813181466678 \\
 \gamma + \ln(n) &= 4.509041347623278 \\
 n &= 52 \\
 H_n &= 4.538043950697447 \\
 \gamma + \ln(n) &= 4.5284594334803785 \\
 n &= 53 \\
 H_n &= 4.556911875225749 \\
 \gamma + \ln(n) &= 4.547507628451074 \\
 \text{This is our answer} \\
 H_n &= \gamma + \ln(n)
 \end{aligned}$$

**Problem 2:** Derive the power series

$$\ln\left(\frac{x}{x-1}\right) = \frac{1}{x} + \frac{1}{2x^2} + \frac{1}{3x^3} + \frac{1}{4x^4} + \dots, \quad (1)$$

and deduce from (1) the following relationship:

$$\begin{aligned}
 \ln n - \ln 1 &= \sum_{k=2}^n \left( \frac{1}{k} + \frac{1}{2k^2} + \frac{1}{3k^3} + \frac{1}{4k^4} + \dots \right) \\
 &= (H_n - 1) + \frac{1}{2} (H_n^{(2)} - 1) + \frac{1}{3} (H_n^{(3)} - 1) + \frac{1}{4} (H_n^{(4)} - 1) + \dots,
 \end{aligned}$$

where the notation  $H_n^{(r)}$  means  $\sum_{k=1}^n \frac{1}{k^r}$ , the  $r^{\text{th}}$  order harmonic number.

**Note:** Taking the limit  $n \rightarrow \infty$  of the  $r^{\text{th}}$  order harmonic number and we have what is now called the “*Riemann zeta function*”

$$\zeta(r) = \lim_{n \rightarrow \infty} H_n^{(r)} = \sum_{k \geq 1} \frac{1}{k^r}.$$

It perhaps ought to be called the “*Euler zeta function*” when  $r$  is a real number, but that’s another story. Riemann studied  $\zeta(r)$  and its *analytic continuation* into the complex plane, allowing  $r$  to be complex. and discovered many astonishing properties of the function including some that have to do with the distribution of prime numbers. The “*Riemann hypothesis*” (should be the “*Riemann conjecture*”) is that if  $\zeta(r) = 0$ , then the real part of  $r$  is equal to  $1/2$ . There is a \$1M bounty on the conjecture.

Anyway, from what was derived from (1), we have

$$H_n - \ln n = 1 - \frac{1}{2} (H_n^{(2)} - 1) - \frac{1}{3} (H_n^{(3)} - 1) - \frac{1}{4} (H_n^{(4)} - 1) - \dots.$$

As  $n \rightarrow \infty$ , the right-hand side becomes the limiting value

$$1 - \frac{1}{2} (\zeta(2) - 1) - \frac{1}{3} (\zeta(3) - 1) - \frac{1}{4} (\zeta(4) - 1) - \dots,$$

which is now known as *Euler’s constant* and gets denoted by the Greek letter  $\gamma$ . What you’ve just worked through is what Euler did, and it establishes the relationship

$$\lim_{n \rightarrow \infty} (H_n - \ln n) = \gamma.$$

First, to derive the power series, begin with the FGF,  $\frac{1}{1-x}$ . The FGF gives us the sequence of all 1’s. From there, take  $-\ln 1 - x$ , which gives us the sequence  $x + \frac{x^2}{2} + \frac{x^3}{3} + \dots$ . Now all we need to do is move every  $x$  to the denominator and we’ll have something that looks just like the power series above. To do this, replace  $x$  with  $x^{-1}$ , and then simplify:

$$-\ln(1 - \frac{1}{x}) = \frac{1}{x} + \frac{1}{2x^2} + \frac{1}{3x^3} + \dots$$

$$-\ln \frac{x-1}{x} = \frac{1}{x} + \frac{1}{2x^2} + \frac{1}{3x^3} + \dots$$

$$\ln \frac{x}{x-1} = \frac{1}{x} + \frac{1}{2x^2} + \frac{1}{3x^3} + \dots$$

This is the generating function we’ve been looking for.

Next, notice that  $\ln \frac{x}{x-1}$  is the same thing as  $\ln x - \ln x - 1$ . Consider adding this to itself an egregious number of times (starting at  $n = 2$ ):

$$\begin{aligned} \ln 2 - \ln 1 &= \frac{1}{2} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{4^2} + \dots \\ + \ln 3 - \ln 2 &= \frac{1}{3} + \frac{1}{2 \cdot 3^2} + \frac{1}{3^3} + \frac{1}{4 \cdot 3^4} + \dots \\ + \ln 4 - \ln 3 &= \frac{1}{4} + \frac{1}{2 \cdot 4^2} + \frac{1}{3 \cdot 4^3} + \frac{1}{4 \cdot 4^4} + \dots \\ &\dots \\ + \ln x - \ln(x-1) &= \frac{1}{x} + \frac{1}{2x^2} + \frac{1}{3x^3} + \frac{1}{4x^4} + \dots \\ = \ln(x) - \ln(1) &= \sum_{k=2}^n \left( \frac{1}{k} + \frac{1}{2k^2} + \frac{1}{3k^3} + \frac{1}{4k^4} + \dots \right) \end{aligned}$$

And thus, the relationship is deduced. Tadah!

**Problem 3:** Use a computer algebra system (which will have very good approximations for  $\zeta(r)$ ) to compute  $\gamma$  to at least 10 digits of accuracy.

We as a group of nerds... I mean computer Science majors, decided to write our own program for computing gamma. We imported (used somebody else’s version) zeta which has been checked for extended approximations to

run our calculations. We then provide a value into the gamma function which gives a closer approximation with the higher the input value. The  $f(x)$  function then calculates the value of the Euler-Mascheroni constant (gamma) using the Riemann zeta function and prints the result. - Python Code -

```
from scipy.special import zeta
def main():
    print(gamma(1000))
def gamma(x):
    return 1 + f(x)
def f(x):
    sum = 0
    for n in range(2, x):
        sum -= (1 / n) * (zeta(n) - 1)
    return sum
if __name__ == "__main__":
    main()
```

- Output -

0.577215664901533

**Frustrated Caterpillar.** You and a friend have a meter-long rubber\* band and decide to mess with a caterpillar named Ting. Ting is at one end of the rubber band and you and your friend are at either end of the rubber band. Your friend remains stationary, and releases Ting who can crawl at 1 centimeter per minute. Ting heads toward the other end of the rubber band, where you are (because you have a caterpillar treat) and you stretch the rubber band 1 meter at the end of the each minute. Ting keeps their relative position on the rubber band as you stretch it (so, at the end of the first minute, after Ting has advanced 1 cm and is 1% from the start and 99% from you, and you stretch the rubber band to 2 meters, Ting will be 2 cm from the starting point and 198 cm from you).

**Problem 4:** Determine with proof how long it will take Ting to reach you. (Use  $\gamma$  from the previous part.)

At the beginning when the rubber band is only 1 meter long and Ting crawls 1 cm for the first time, Ting is 1 percent of the way across the rubber band. After the rubber band is stretched, Ting is now 2 cm along a 2 meter long rubber band, so Ting is still 1 percent of the way across the rubber band. When Ting crawls another cm and the rubber band is stretched another meter, Ting is 1.5 percent of the way across the rubber band. This is a harmonic series, as shown by the math below:

$$nth = \frac{1}{n}$$

^(this is the percentage moved at the nth minute)

$$\sum_{n=1} \frac{1}{n} \geq 100$$

$$\gamma \approx 0.577215664901533$$

$$\gamma \approx H_n - \ln(n)$$

$$H_n \approx \gamma + \ln(n)$$

Since this is a harmonic series, all we need to do is find  $H_n$  where  $H_n \geq 100$ , because when that happens, Ting will be 100 percent of the way across the rubber band. This can be expressed like so:

$$H_n \geq 100 - \gamma$$

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\*... or some sort of material that has the mystical properties required for this problem

$$\begin{aligned}
100 &\leq \gamma + \ln(n) \\
e^{100-\gamma} &\leq e^{\ln(n)} \\
e^{100-\gamma} &\leq n \\
e^{100-0.577215664901533} &\leq n \\
1.5092688622 \times 10^{43} &\leq n
\end{aligned}$$

Therefore, without further ado, the final answer is that it will take Ting  $1.5092688622 \times 10^{43}$  minutes to reach the end of the rubber band.

**Generating Functions.** A **partition** of  $n \in \mathbb{Z}^+$  **into**  $k$  **parts** is an equation of the form  $n = \lambda_1 + \lambda_2 + \cdots + \lambda_k$ , for some  $k$ , with  $\lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_k$ , and  $\lambda_i > 0$ ; each  $\lambda$  is a **part**. The number  $k$  of parts may not be specified, in which case the terminology is “*a partition of  $n$* ”.

Define  $p(n)$  to be the number of partitions of  $n$ . For convenience, we take  $p(0)$  to be 1. It might behoove you to verify the sequence  $(p(n))_{n \geq 0}$  begins

$$(p(n))_{n \geq 0} = (1, 1, 2, 3, 5, 7, 11, 15, 22, \dots);$$

that is,  $p(2) = 2, p(5) = 7, p(6) = 11$ , and so on.

Define  $\mathcal{E}(x)$  to be the generating function for the sequence  $(p(n))_{n \geq 0}$ ; that is  $\mathcal{E}(x) = \sum_{n \geq 0} p(n)x^n$ . In class it was observed that

$$\mathcal{E}(x) = \prod_{k \geq 0} \frac{1}{1 - x^k}.$$

**Problem 5:** Define  $p_{\text{no } 1}(n)$  to be the number of partitions of  $n$  with no part equal to 1. Prove that  $p_{\text{no } 1}(n) = p(n) - p(n-1)$ .

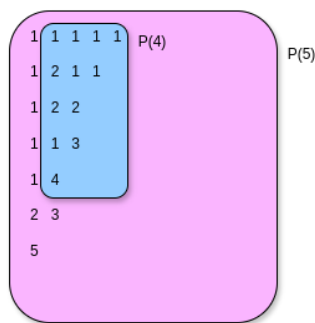
To start off, let's get an idea of what the sequence looks like for  $P_{\text{no } 1}$

$$\begin{aligned}
p_{\text{no } 1}(0) &= 1, \quad p(0) = 1 \\
p_{\text{no } 1}(1) &= 0, \quad p(1) = 1 \\
p_{\text{no } 1}(2) &= 1, \quad p(2) = 2 \\
p_{\text{no } 1}(3) &= 1, \quad p(3) = 3 \\
p_{\text{no } 1}(4) &= 2, \quad p(4) = 5 \\
p_{\text{no } 1}(5) &= 2, \quad p(5) = 7 \\
p_{\text{no } 1}(6) &= 4, \quad p(6) = 11 \\
p_{\text{no } 1}(7) &= 4, \quad p(7) = 15
\end{aligned}$$

The values of the unique partitions above (the right side equation) can be referenced in the “Partition function” section of the following wikipedia page: [wikipedia page](#).

We can set a base case for 0 and 1, and then work with the remaining sequence. We can create a combinatorial argument for this. Assuming that  $p_{\text{no } 1}(n) = p(n) - p(n-1)$ , then we can move terms around and get that the combinations with no 1 plus the combinations that do contain ones is equal to the total possible combinations of

partitions. We can show this looking at the figure below with  $n = 5$  as our example:



For this case,  $7 - 5 = 2$  which follows with the sequence we got above, and further shows that the provided relationship is in fact correct.

**Problem 6:** Define  $p_k(n)$  to be the number of partitions of  $n$  with at most  $k$  parts. Prove that the generating function for  $(p_k(n))_{n \geq 0}$  is

$$\sum_{n \geq 0} p_k(n)x^n = \frac{1}{(1-x)(1-x^2)(1-x^3)(1-x^4) \cdots (1-x^k)},$$

and [calculate]  $p(27)$  and  $p_5(27)$  using the generating function relationships just proved.

For the generating function above, the  $n$ th coefficient is the number of ways to partition a set of size  $n$  where the largest part of the partition is size  $k$ . This is because the generating function is multiplying all the ways to partition a set of size  $n$  with no parts larger than 1, all the ways to partition with no parts larger than 2, etc. etc. up to every way to partition with no parts larger than  $k$ .

This counts the same thing as the number of partitions of  $n$  with at most  $k$  parts. If you drew out some partition of dots where each row was a part and the most dots in any row is  $k$ , this is the number of ways to partition a set of size  $n$  with no parts being larger than  $k$ .

However, if you let the columns represent the partitions instead of the rows, you would have  $k$  columns, all of various sizes (since the most dots in any row is  $k$ ). *This* is counting the number of partitions of  $n$  with at most  $k$  parts, which is what this generating function is supposed to be counting. Therefore, the generating function for  $(p_k(n))_{n \geq 0}$  is

$$\sum_{n \geq 0} p_k(n)x^n = \frac{1}{(1-x)(1-x^2)(1-x^3)(1-x^4) \cdots (1-x^k)}$$

To find  $p_5(27)$ , find the 27th coefficient on the generating function, where  $(1-x^i)$  has  $1 \leq i \leq 5$ . This can be done using the following Maxima code:

```
g(x) := product(1-x^i, i, 1, 5);
f(x) := 1 / g(x);

taylor(f(x), x, 0, 27);
```

This returns a series, the 27th term of which has the coefficient 480. Therefore,  $p_5(27) = 480$ .

$p(27)$  is asking us to find the number of partitions of  $n$  where  $n$  is 27, with at most  $n$  parts (because you can't have more than  $n$  parts with only  $n$  items you're partitioning). Therefore, to find  $p(27)$ , do the same thing as above but for  $1 \leq i \leq 27$ . Doing this, the coefficient on the 27th term of the series is 3010. And so,  $p(27) = 3010$ .

**Catalan Numbers.** Define  $C_n$  to be the number of valid strings of  $n$  pairs of parentheses; that is, the empty string is a valid string, and if  $A$  and  $B$  are valid strings of parentheses, then so are  $A(B)$  and  $(A)B$ . The sequence  $(C_n)_{n \geq 0}$  begins

$$(C_n)_{n \geq 0} = (1, 1, 2, 5, 14, \dots).$$

**Problem 7:** Prove that the numbers  $C_n$  can be created via the recurrence relation

$$C_{n+1} = C_n C_0 + C_{n-1} C_1 + C_{n-2} C_2 + \dots + C_1 C_{n-1} + C_0 C_n \text{ for } n \geq 0$$

$$C_0 = 1$$

Alrighty then, let's go ahead and get crack-a-lackin! First let's set up the recurrence relation as a sum for visual appeal.

$$C_{n+1} = C_n C_0 + C_{n-1} C_1 + C_{n-2} C_2 + \dots + C_1 C_{n-1} + C_0 C_n$$

$$C_{n+1} = \sum_{i=0}^n C_{n-i} C_i$$

$$C_n = \sum_{i=0}^{n-1} C_{n-i-1} C_i$$

below we will plug in numbers to test our base cases. this is more to prove to ourselves that the recurrence for plain ol'  $C_n$  we deduced above does in fact work.

$$C_0 = 1$$

$$C_1 = \sum_{i=0}^{1-1} C_{1-0-1} C_0 = \sum_{i=0}^0 C_0 C_0 = C_0 C_0 = 1$$

$$C_2 = \sum_{i=0}^{2-1} C_{2-i-1} C_i = \sum_{i=0}^1 C_{1-i} C_i = C_1 C_0 + C_0 C_1 = 2$$

$$C_3 = \sum_{i=0}^{3-1} C_{3-i-1} C_i$$

$$C_3 = C_2(C_0) + C_1(C_1) + C_0(C_2)$$

$$C_3 = 2(1) + 1(1) + 1(2) = 5$$

$$C_4 = \sum_{i=0}^3 C_{3-i} C_i$$

$$C_4 = C_3(C_0) + C_2(C_1) + C_1(C_2) + C_0(C_3)$$

$$C_4 = 5(1) + 2(1) + 1(2) + 1(5) = 14$$

$$1, 1, 2, 5, 14$$

It works! So we shall use  $C_n = \sum_{i=0}^{n-1} C_{n-i-1} C_i$  in Problem 8



**Problem 8:** Define  $C(x)$  to be the generating function for the Catalan numbers; that is,  $C(x) = \sum_{n \geq 0} C_n x^n$ . Show

that  $C(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$ .

Define  $[n = 0]$  to give 1 if true and 0 if false

$$C_n = \sum_{i=0}^{n-1} C_{n-i-1} C_i + [n = 0]$$

Define  $C(x) = \sum_{n \geq 0} C_n x^n$

$$C(x) = \sum_{n \geq 0} \left( \sum_{i=0}^{n-1} C_{n-i-1} C_i \right) x^n + \sum_{n \geq 0} [n = 0] x^n$$

$$C(x) = \sum_{n \geq 0} \left( \sum_{i=0}^{n-1} C_{n-i-1} C_i \right) x^n + 1$$

$$C_n = \sum_{i=0}^{n-1} C_{n-i-1} C_i = C_{n-1} C_0 + C_{n-2} C_1 + C_{n-3} C_2 + \dots + C_2 C_{n-3} + C_1 C_{n-2} + C_0 C_{n-1}$$

$$\begin{aligned} C_n^2 &= \left( \sum_{i=0}^{n-1} C_{n-i-1} C_i \right) \left( \sum_{i=0}^{n-1} C_{n-i-1} C_i \right) \\ &= (C_{n-1} C_0 + C_{n-2} C_1 + \dots + C_1 C_{n-2} + C_0 C_{n-1}) (C_{n-1} C_0 + C_{n-2} C_1 + \dots + C_1 C_{n-2} + C_0 C_{n-1}) \\ &= \sum_{i=0}^n C_i C_{n-i} = C_n C_0 + C_{n-1} C_1 + C_{n-2} C_2 + \dots + C_2 C_{n-2} + C_1 C_{n-1} + C_0 C_n \end{aligned}$$

Well hey,  $C_n^2 = C_{n+1}$ . That means that  $C(x)^2 = \sum_{n \geq 0} C_{n+1} x^n$ .

$$C(x)^2 = \sum_{n \geq 0} C_{n+1} x^n$$

$$xC(x)^2 = \sum_{n \geq 0} C_{n+1} x^{n+1}$$

We have to add 1 to account for  $C_0$

$$xC(x)^2 + 1 = C(x)$$

$$xC(x)^2 - C(x) + 1 = 0$$

Using the quadratic equation to solve for  $C(x)$  with  $a = x$ ,  $b = -1$ , and  $c = 1$

$$C(x) = \frac{1 \pm \sqrt{(-1)^2 - 4x}}{2x}$$

$$C(x) = \frac{1 \pm \sqrt{1 - 4x}}{2x}$$

Because  $C(x) = \frac{1 + \sqrt{1 - 4x}}{2x}$  goes to infinity when approaching 0 we can throw that solution out. The other solution goes to 1 which would satisfy  $C_0 = 1$ .

$$C(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$$

**Problem 9:** Show that  $C_n = \frac{1 \cdot 3 \cdot 5 \cdots (2n-1)}{(n+1)!} \cdot 2^n$ , and hence, finally, that  $C_n = \frac{1}{n+1} \binom{2n}{n}$ . (This is the most common expression for the  $n^{\text{th}}$  Catalan number.)

$$C(x) = \frac{1 - \sqrt{1-4x}}{2x}$$

Using Newtons Binomial Theorem!

$$(1+x)^a = \sum_{n \geq 0} \binom{a}{n} x^n$$

$$\sqrt{1-4x} = (1+(-4x))^{\frac{1}{2}} = \sum_{n \geq 0} \binom{\frac{1}{2}}{n} (-4x)^n$$

So

$$C(x) = \frac{1 - \sqrt{1-4x}}{2x} = \frac{1 - \sum_{n \geq 0} \binom{\frac{1}{2}}{n} (-4x)^n}{2x}$$

When  $n = 0$ ,  $\binom{\frac{1}{2}}{0} (-4x)^0 = 1 * 1 = 1$ . This means that we can start the summation at 1 and get rid of the 1 at the start.

$$C(x) = -\frac{1}{2} \left( \sum_{n \geq 1} \binom{\frac{1}{2}}{n} (-4)^n x^{n-1} \right)$$

$$C(x) = -\frac{1}{2} \left( \sum_{n \geq 1} \binom{\frac{1}{2}}{n} (-4)^n x^{n-1} \right)$$

Index shift

$$C(x) = -\frac{1}{2} \left( \sum_{n \geq 0} \binom{\frac{1}{2}}{n+1} (-4)^{n+1} x^n \right)$$

$$C_n = -\frac{1}{2} (-4)^{n+1} \binom{\frac{1}{2}}{n+1}$$

Expanding the binomial coefficient

$$\binom{n}{k} = \frac{n!}{k!}$$

$$C_n = -\frac{1}{2} (-4)^{n+1} \frac{\left(\frac{1}{2}\right)^{n+1}}{(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1} \left(\frac{1}{2}\right)^{n+1}}{2(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1} \left(\left(\frac{1}{2}\right)\left(\frac{1}{2}-1\right)\left(\frac{1}{2}-2\right)\left(\frac{1}{2}-3\right)\dots\left(\frac{1}{2}-n+1\right)\left(\frac{1}{2}-n\right)\right)}{2(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1} \left(\left(\frac{1}{2}\right)\left(\frac{1}{2}-1\right)\left(\frac{1}{2}-2\right)\left(\frac{1}{2}-3\right)\dots\left(\frac{3}{2}-n\right)\left(\frac{1-2n}{2}\right)\right)}{2(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1} \left(\left(\frac{1}{2}\right)\left(\frac{-1}{2}\right)\left(\frac{-3}{2}\right)\left(\frac{-5}{2}\right)\dots\left(\frac{3-2n}{2}\right)\left(\frac{1-2n}{2}\right)\right)}{2(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1} \left(\left(\frac{1}{2^{n+1}}\right)(-1)(-3)(-5)\dots(3-2n)(1-2n)\right)}{2(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1}((1)(-1)(-3)(-5)...(3-2n)(1-2n))}{(2^{n+2})(n+1)!}$$

$$C_n = \frac{(-1)(-4)^{n+1}((-1)(-3)(-5)...(3-2n)(1-2n))}{(2^{n+2})(n+1)!}$$

$$C_n = \frac{(-1)^{n+1}(4)^{n+1}(-1)^{n+1}((1)(3)(5)...(2n-3)(2n-1))}{(2^{n+2})(n+1)(n)!}$$

$$C_n = \frac{(2)^{n+1}(2)^{n+1}}{(2^{n+2})(n+1)} * \frac{(1)(3)(5)...(2n-3)(2n-1)}{(n)!}$$

$$C_n = \frac{(2)^{n+1}}{(2)(n+1)} * \frac{(1)(3)(5)...(2n-3)(2n-1)}{(n)!}$$

$$C_n = \frac{((1)(3)(5)...(2n-3)(2n-1))}{(n+1)!} \cdot 2^n$$

There is the first form that is asked for in the question prompt

Now time for the next form

$$C_n = \frac{(1)(3)(5)...(2n-1)}{(n+1)!} \cdot 2^n$$

$$C_n = \frac{1}{n+1} * \frac{((1)(3)(5)...(2n-1))2^n}{n!}$$

Working backwards

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

$$C_n = \frac{1}{n+1} \frac{(2n)!}{n!((2n)-n)!}$$

$$C_n = \frac{1}{n+1} \frac{(2n)(2n-1)(2n-2)(2n-3)(2n-4)...(4)(3)(2)(1)}{(n!)^2}$$

$$C_n = \frac{1}{n+1} \frac{((2n)(2n-2)(2n-4)...(4)(2))((2n-1)(2n-3)...(3)(1))}{(n!)^2}$$

$$C_n = \frac{1}{n+1} \frac{2^n((n)(n-1)(n-2)...(2)(1))((2n-1)(2n-3)...(3)(1))}{(n!)^2}$$

$$C_n = \frac{1}{n+1} \frac{2^n(n!)((2n-1)(2n-3)...(3)(1))}{(n!)^2}$$

$$C_n = \frac{1}{n+1} \frac{2^n((2n-1)(2n-3)...(3)(1))}{n!}$$

Well boom there you go they are worked out to be the same thing

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

**Problem 10:** In how many ways can a  $2 \times 2 \times n$  pillar be built out of  $2 \times 1 \times 1$  bricks?

Here is the recurrence relation.

$$S_0 = 1$$

$$S_1 = 2$$

$$S_2 = 9$$

$$S_n = 2S_{n-1} + S_{n-2} + 4 \sum_{i=2}^n S_{n-i}$$

We got this recurrence by looking at the layers of blocks and all the combinations of blocks when you consider a number of layers starting from the top.

If the top layer is all horizontal blocks, then we can shave it off the pillar and look at layer  $(n-1)$ . And since there are two ways for the top layer of blocks to be made up of only horizontal blocks (both blocks facing North/South, or both blocks facing East/West), we have  $2S_{n-1}$  in the recurrence. If the top layer is all vertical blocks, then we can shave it off the pillar (removing a height of 2 from the pillar) and look at the next flat layer, which is layer  $(n-2)$ . There is only one way for the top layer to be all vertical blocks, and so we have  $S_{n-2}$  in the recurrence. And finally, if the top layer is made of a combination of vertical and horizontal blocks, we don't know how many blocks we'll have to remove before hitting the next flat layer in the pillar. The closest it could be is just below the topmost vertical blocks (hence starting the summation at  $i=2$ ). It could also be that there are no all-horizontal or all-vertical layers in the pillar. This is where the  $\sum_{i=2}^n S_{n-i}$  in the recurrence comes from, because we may end up having to check every layer. And since there are four ways to have a mixed vertical/horizontal top layer (vertical blocks are on the North side, East side, South side, or West side), we multiply this part of the recurrence by 4 to get  $4 \sum_{i=2}^n S_{n-i}$ .

Make it work for any  $n$ . Define  $S_n = 0$  if  $n < 0$ . Define  $[n = 0]$  to give 1 if true and 0 otherwise.

$$S_n = 2S_{n-1} + S_{n-2} + 4 \sum_{i=2}^n S_{n-i} + [n = 0]$$

$$S_0 = 2S_{-1} + S_{-2} + 4 \sum_{i=2}^0 S_{0-i} + 1 = 2(0) + 0 + 4(0) + 1 = 1$$

$$S_1 = 2S_0 + S_{-1} + 4 \sum_{i=2}^1 S_{1-i} + 0 = 2(1) + 0 + 4(0) = 2$$

$$S_2 = 2S_1 + S_0 + 4 \sum_{i=2}^2 S_{2-i} + 0 = 2(2) + 1 + 4(1) = 4 + 5 = 9$$

It works!

Define  $S(x) = \sum_{n \geq 0} S_n x^n$ .

Multiply everything by the summation and  $x^n$

$$S(x) = 2 \sum_{n \geq 0} S_{n-1} x^n + \sum_{n \geq 0} S_{n-2} x^n + 4 \sum_{n \geq 0} \left( \sum_{i=2}^n S_{n-i} \right) x^n + \sum_{n \geq 0} [n = 0] x^n$$

$$S(x) = 2xS(x) + x^2S(x) + 4 \sum_{n \geq 0} \left( \sum_{i=2}^n S_{n-i} \right) x^n + 1$$

$$S(x) = 2xS(x) + x^2S(x) + 4 \sum_{n \geq 0} \left( \left( \sum_{i=0}^n S_i \right) - S_n - S_{n-1} \right) x^n + 1$$

$$S(x) = 2xS(x) + x^2S(x) + 4\left(\sum_{n \geq 0} \left(\sum_{i=0}^n S_i\right)x^n - \sum_{n \geq 0} S_n x^n - \sum_{n \geq 0} S_{n-1} x^n\right) + 1$$

$$S(x) = 2xS(x) + x^2S(x) + 4\left(\frac{1}{1-x}S(x) - S(x) - xS(x)\right) + 1$$

$$S(x) = 2xS(x) + x^2S(x) + 4\frac{1}{1-x}S(x) - 4S(x) - 4xS(x) + 1$$

$$S(x) - 2xS(x) - x^2S(x) - 4\frac{1}{1-x}S(x) + 4S(x) + 4xS(x) = 1$$

$$S(x)(1 - 2x - x^2 - 4\frac{1}{1-x} + 4 + 4x) = 1$$

$$S(x)(5 + 2x - x^2 - 4\frac{1}{1-x}) = 1$$

$$S(x) = \frac{1}{5 + 2x - x^2 - 4\frac{1}{1-x}}$$

$$S(x) = \frac{1-x}{(x+1)(x^2-4x+1)}$$

$$S(x) = \frac{1-x}{-(x+1)(-x+\sqrt{3}+2)(x+\sqrt{3}-2)}$$

$$S(x) = \frac{x-1}{(x+1)(-x+\sqrt{3}+2)(x+\sqrt{3}-2)}$$

$$S(x) = \frac{A}{x+1} + \frac{B}{-x+\sqrt{3}+2} + \frac{C}{x+\sqrt{3}-2}$$

$$x-1 = A(-x+\sqrt{3}+2)(x+\sqrt{3}-2) + B(x+1)(x+\sqrt{3}-2) + C(x+1)(-x+\sqrt{3}+2)$$

Cutting to the point after using smart online math tools,  $A = \frac{1}{3}$ ,  $B = -\frac{1}{6}$ ,  $C = -\frac{1}{6}$

$$S(x) = \frac{\frac{1}{3}}{x+1} - \frac{\frac{1}{6}}{-x+\sqrt{3}+2} - \frac{\frac{1}{6}}{x+\sqrt{3}-2}$$

$$S(x) = \frac{1}{3} * \frac{1}{x+1} - \frac{1}{6} * \frac{1}{-x+\sqrt{3}+2} - \frac{1}{6} * \frac{1}{x+\sqrt{3}-2}$$

Using the reverse polynomial thing you get

$$S(x) = \sum_{n \geq 0} \frac{1}{6} (2(-1)^n + (2-\sqrt{3})^{n+1} + (2+\sqrt{3})^{n+1}) x^n$$

$$S_n = \frac{1}{6} (2(-1)^n + (2-\sqrt{3})^{n+1} + (2+\sqrt{3})^{n+1})$$

Checked for  $n = 0, 1, 2, 3$

**Problem 11:** A form of DNA found stuck to a space probe has been analyzed and found to consist of five different molecules, call them a, b, c, d, e. Research shows that the pairs cd, ce, ed, and ee cannot occur consecutively in a string of this mysterious DNA, but any string without these *forbidden pairs* can occur. So, for example bbdca is not possible, but bbdca is. How many different mysterious DNA strings of length  $n$  are possible? Note that when  $n = 2$ , there are 21 different strings because the left and right ends of the string are distinguishable.

I shall call the function  $P$  for probe.  $P$  is the number of different mysterious DNA strings of length  $n$  that are possible.

These numbers come from a python program.

$$P_0 = 1$$

$$P_1 = 5$$

$$P_2 = 21$$

$$P_3 = 89$$

$$P_4 = 377$$

$$P_5 = 1597$$

$$P_n = 4P_{n-1} + P_{n-2}$$

Here is our reasoning for the ratio above:

Condition on the last letter in the sequence. The following is a table of possible letters:

Second-to-last Letter	Last Letter
a, b, c, d, e	a
a, b, c, d, e	b
a, b, c, d, e	c
a, b, d	d
a, b, d	e

a, b, and c have no constraints on the letters that can come before them, so a, b, and c each have 5 possible preceding letters.

d and e can only have a, b, and d come before them. a and b are unrestricted with five possible letters before them each, so d and e collectively have  $5 * 4 = 20$  possible letters for the  $(n - 1)$ th letter. But we had to go one layer deeper to get that number, so divide 20 by 4 (two for d as last letter and two for e as last letter), which gives us another 5 for the  $(n - 1)$ th letter.

d is restricted, so if the last letter is d or e and the second-to-last letter is d, then the third-to-last letter can be a, b, or d. Again, a and b are unrestricted with 5 possible preceding letters. So  $5 * 4 = 20$ . But divide that by 4 (two for d two for e) and you get 5. This counts toward  $(n - 2)$  because now we're looking at third-to-last letters.

Thus, the ratio of  $4P_{n-1} + P_{n-2}$ .

Make the recursion work for any n

Define  $P_n = 0$  if  $n < 0$ .

$$P_n = 4P_{n-1} + P_{n-2} + [n = 1] + [n = 0]$$

$$P_0 = 4P_{-1} + P_{-2} = 4(0) + 0 + 0 + 1 = 1$$

$$P_1 = 4P_0 + P_{-1} + [n = 0] = 4(1) + 0 + 1 + 0 = 5$$

$$P_2 = 4P_1 + P_0 + [n = 1] + [n = 0] = 4(5) + 1 + 0 + 0 = 21$$

$$P_3 = 4P_2 + P_1 + [n = 1] + [n = 0] = 4(21) + 5 + 0 + 0 = 89$$

$$P_4 = 4P_3 + P_2 + [n = 1] + [n = 0] = 4(89) + 21 + 0 + 0 = 377$$

$$P_5 = 4P_4 + P_3 + [n = 1] + [n = 0] = 4(377) + 89 + 0 + 0 = 1597$$

It works!

Define  $P(x) = \sum_{n \geq 0} P_n x^n$

$$P(x) = 4 \sum_{n \geq 0} P_{n-1} x^n + \sum_{n \geq 0} P_{n-2} x^n + \sum_{n \geq 0} [n = 1] x^n + \sum_{n \geq 0} [n = 0] x^n$$

$$P(x) = 4x \sum_{n \geq 0} P_{n-1} x^{n-1} + x^2 \sum_{n \geq 0} P_{n-2} x^{n-2} + x + 1$$

Prove that  $\sum_{n \geq 0} P_{n-1} x^{n-1} = \sum_{n \geq 0} P_n x^n$

$$\sum_{n \geq 0} P_{n-1} x^{n-1} = P_{-1} x^{-1} + P_0 x^0 + P_1 x^1 + P_2 x^2 + \dots$$

$$= 0 + P_0 x^0 + P_1 x^1 + P_2 x^2 + \dots$$

So they are equal!

$$P(x) = 4xP(x) + x^2P(x) + x + 1$$

$$P(x) - 4xP(x) - x^2P(x) = x + 1$$

$$P(x)(1 - 4x - x^2) = x + 1$$

$$P(x) = \frac{x + 1}{1 - 4x - x^2}$$

Checked in maxima its the right power series

Need to use the reverse polynomial on  $1 - 4x - x^2$

$$P(x) = \frac{x + 1}{(1 - (2 - \sqrt{5})x)(1 - (2 + \sqrt{5})x)}$$

$$P(x) = \frac{A}{1 - (2 - \sqrt{5})x} + \frac{B}{1 - (2 + \sqrt{5})x}$$

$$P(x) = \frac{\sqrt{5} - 3}{2\sqrt{5}(-(2 - \sqrt{5})x + 1)} - \frac{\sqrt{5} + 3}{2\sqrt{5}((2 + \sqrt{5})x - 1)}$$

$$P(x) = \left(\frac{\sqrt{5} - 3}{2\sqrt{5}}\right) \frac{1}{(\sqrt{5} - 2)x + 1} - \left(\frac{\sqrt{5} + 3}{2\sqrt{5}}\right) \frac{1}{(\sqrt{5} + 2)x - 1}$$

$$P(x) = \left(\frac{\sqrt{5} - 3}{2\sqrt{5}}\right) \frac{1}{1 + (\sqrt{5} - 2)x} - \left(\frac{\sqrt{5} + 3}{2\sqrt{5}}\right) \frac{1}{-1 + (\sqrt{5} + 2)x}$$

$$P(x) = \left(\frac{\sqrt{5} - 3}{2\sqrt{5}}\right) \frac{1}{1 - (2 - \sqrt{5})x} + \left(\frac{\sqrt{5} + 3}{2\sqrt{5}}\right) \frac{1}{1 - (\sqrt{5} + 2)x}$$

$$P(x) = \left(\frac{\sqrt{5} - 3}{2\sqrt{5}}\right) \sum_{n \geq 0} (2 - \sqrt{5})^n x^n + \left(\frac{\sqrt{5} + 3}{2\sqrt{5}}\right) \sum_{n \geq 0} (\sqrt{5} + 2)^n x^n$$

$$P_n = \left(\frac{\sqrt{5} - 3}{2\sqrt{5}}\right) (2 - \sqrt{5})^n + \left(\frac{\sqrt{5} + 3}{2\sqrt{5}}\right) (2 + \sqrt{5})^n$$

Checked for  $n = 0, 1, 2, 3, 4, 5$

**Derangements.** Recall that a *derangement* of  $[n]$  is a function  $f : [n] \rightarrow [n]$  that is 1-to-1 and onto with the property that  $f(i) \neq i$  for every  $i \in [n]$ . Define the number of derangements to be  $D_n$ .

**Problem 12:** Prove that  $D_n$  satisfies the recurrence

$$D_n = (n-1)D_{n-1} + (n-1)D_{n-2} \text{ for } n \geq 1;$$

$$D_0 = 1,$$

and therefore also the recurrence

$$D_n = nD_{n-1} + (-1)^n \text{ for } n \geq 1$$

$$D_0 = 1.$$

Combinatorial Proof of  $D_n = (n-1)D_{n-1} + (n-1)D_{n-2}$ :

We will refer to the problem of a hat checker at a party manages to improperly return every hat to every hat wear-er. For any derangement, some hat has to take the first hat's spot. Call this hat  $i$ . (And note that there are  $n-1$  possibilities for  $i$ , since  $i$  can be any hat that isn't the first hat.) From here, there are two possible scenarios. Either the 1st hat is in the  $i$ th spot, or the first hat is not in the  $i$ th spot.

If the first hat is in the  $i$ th spot, then we know the places of two hats and only have to find the derangement for the remaining  $n-2$  hats. This is represented by  $(n-1)D_{n-2}$ .

If the first hat is *NOT* in the  $i$ th spot, then the rest of the derangement can be found by *temporarily* putting the 1st hat in the  $i$ th spot, and then finding the derangement of all hats *except* the first one. This is represented by  $(n-1)D_{n-1}$ .

Therefore,  $D_n = (n-1)D_{n-1} + (n-1)D_{n-2}$ .

Now using the recurrence just previously found we will prove the second recurrence

$$D_n = (n-1)D_{n-1} + (n-1)D_{n-2}$$

$$D_n = nD_{n-1} - D_{n-1} + (n-1)D_{n-2}$$

$$D_n - nD_{n-1} = -(D_{n-1} + (n-1)D_{n-2})$$

$$D_n - nD_{n-1} = -(-(D_{n-2} + (n-2)D_{n-3}))$$

$$D_n - nD_{n-1} = -(-(-(D_{n-3} + (n-3)D_{n-4})))$$

$$\vdots$$

$$D_n - nD_{n-1} = -(-(-(\dots - (0-1))) \dots)$$

$$D_n - nD_{n-1} = (-1)^n$$

$$D_n = nD_{n-1} + (-1)^n$$

**Problem 13:** Use an exponential generating function to find a closed formula for  $D_n$  (the formula will involve a sum, initially, but it can be significantly reduced).

$$D_0 = 1$$

$$D_n = nD_{n-1} + (-1)^n$$

Need to make this recurrence work for any  $n$ . Define  $D_n = 0$  if  $n < 0$ .

$$D_0 = 0D_{-1} + (-1)^0 = 1$$

$$D_1 = 1D_0 + (-1)^1 = 1 - 1 = 0$$



$$D_2 = 2D_1 + (-1)^2 = 2(0) + 1 = 1$$

$$D_3 = 3D_2 + (-1)^3 = 3(1) + -1 = 2$$

Alright the recurrence works. Define  $D(x) = \sum_{n \geq 0} D_n \frac{x^n}{n!}$

$$D_n = nD_{n-1} + (-1)^n$$

$$D(x) = n \sum_{n \geq 0} D_{n-1} \frac{x^n}{n!} + \sum_{n \geq 0} (-1)^n \frac{x^n}{n!}$$

$$D(x) = \sum_{n \geq 0} D_{n-1} \frac{x^n}{(n-1)!} + (1 \frac{x^0}{0!} - 1 \frac{x^1}{1!} + 1 \frac{x^2}{2!} - 1 \frac{x^3}{3!} + 1 \frac{x^4}{4!} - \dots)$$

$$D(x) = x \sum_{n \geq 0} D_{n-1} \frac{x^{n-1}}{(n-1)!} + e^{-x}$$

$$D(x) = xD(x) + e^{-x}$$

$$D(x) - xD(x) = e^{-x}$$

$$D(x)(1 - x) = e^{-x}$$

$$D(x) = \frac{e^{-x}}{1 - x}$$

Checked with Maxima and Verified the first 4 terms.

$$D(x) = (e^{-x}) \left( \frac{1}{1-x} \right)$$

$$D(x) = \sum_{n \geq 0} \left( \sum_{i=0}^n \binom{n}{i} (-1)^i (n-i)! \right) \frac{x^n}{n!}$$

$$D_n = \sum_{i=0}^n \frac{n!}{i!(n-i)!} (-1)^i (n-i)!$$

$$D_n = \sum_{i=0}^n \frac{n!}{i!} (-1)^i$$

$$D_n = n! \sum_{i=0}^n \frac{1}{i!} (-1)^i$$

$$D_n = \lfloor \frac{n!}{e} + \frac{1}{2} \rfloor$$

Checked n for 0,1,2,3

**Problem 14:** Use the tactics from the unit on the *Discrete Taylor Series* to find a formula for  $D_n$ .

1. Prove the identity  $n! = \sum_{k \geq 0} D_k \binom{n}{k}$ .

Well here is the difference table for  $n!$

n	0	1	2	3	4	5
$f(n)$	1	1	2	6	24	120
$\Delta^1(1)f(n)$	0	1	4	18	96	600
$\Delta^2(2)f(n)$	1	3	14	78	504	3720
$\Delta^3(3)f(n)$	2	11	64	426	3216	27240
$\Delta^4(4)f(n)$	9	53	362	2790	24024	229080
$\Delta^5(5)f(n)$	44	309	2428	21234	205056	2170680
$\Delta^6(6)f(n)$	265	2119	18806	183822	1965624	22852200
$\Delta^7(7)f(n)$	1854	16687	165016	1781802	20886576	264398280
$\Delta^8(8)f(n)$	14833	148329	1616786	19104774	243511704	3332744760
$\Delta^9(9)f(n)$	133496	1468457	17487988	224406930	3089233056	45440868120
$\Delta^{10}(10)f(n)$	1334961	16019531	206918942	2864826126	42351635064	666166856520
$\Delta^{11}(11)f(n)$	14684570	190899411	2657907184	39486808938	623815221456	10446911529000
$\Delta^{12}(12)f(n)$	176214841	2467007773	36828901754	584328412518	9823096307544	174478419885720
$\Delta^{13}(13)f(n)$	2290792932	34361893981	547499510764	9238767895026	164655323578176	3091496076405240
$\Delta^{14}(14)f(n)$	32071101049	513137616783	8691268384262	155416555683150	2926840752827064	57915148833808680
$\Delta^{15}(15)f(n)$	481066515734	8178130767479	146725287298888	2771424197143914	54988308080981616	1143668772912038280
$\Delta^{16}(16)f(n)$	7697064251745	138547156531409	2624698909845026	52216883883837702	1088680464831056664	
$\Delta^{17}(17)f(n)$	130850092279664	2486151753313617	49592184973992676	1036463580947218962		
$\Delta^{18}(18)f(n)$	2355301661033953	47106033220679059	986871395973226286			
$\Delta^{19}(19)f(n)$	44750731559645106	939765362752547227				
$\Delta^{20}(20)f(n)$	895014631192902121					

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You can clearly see that  $D_k$  is the 0th term in all the derivatives.

$$n! = \sum_{k \geq 0} D_k \binom{n}{k} = \sum_{k \geq 0} \Delta^{(k)}(0)! \binom{n}{k}$$

Because of linearity we can focus on these two

$$D_k = \Delta^{(k)}(0)!$$

$$k! \sum_{i=0}^k \frac{1}{i!} (-1)^i = \Delta^{(k)}(0)!$$

$$\sum_{i=0}^k \frac{1}{i!} (-1)^i = \frac{\Delta^{(k)}(0)!}{k!}$$

They both converge to  $\frac{1}{e}$  therefor they are the same.

$$\frac{1}{e} = \frac{1}{e}$$

- Examine the difference table for  $n!$ , and create its discrete Taylor series.

Look at the difference table in 14 part 1.

$$n! = \sum_{k \geq 0} D_k \binom{n}{k}$$

- Examine the difference table for  $d_n = (-1)^n D_n$ , and derive a closed formula for  $D_n$  (which you already know from Problem 13).

The difference table for  $d_n = (-1)^n D_n$

n	0	1	2	3	4	5	6
f(n)	1	0	1	-2	9	-44	265
$\Delta^1 f(n)$	-1	1	-3	11	-53	309	-2119
$\Delta^2 f(n)$	2	-4	14	-64	362	-2428	18806
$\Delta^3 f(n)$	-6	18	-78	426	-2798	21234	-183822
$\Delta^4 f(n)$	24	-96	504	-3216	24024	-205056	1965624
$\Delta^5 f(n)$	-120	600	-3720	27240	-229080	2170680	-22852200
$\Delta^6 f(n)$	720	-4320	30960	-256320	2399760	-25022880	287250480
$\Delta^7 f(n)$	-5040	35280	-287280	2656080	-27422640	312273360	-3884393520
$\Delta^8 f(n)$	40320	-322560	2943360	-30878720	339696000	-4196666880	56255149440
$\Delta^9 f(n)$	-362880	3265920	-33022080	369774720	-4536362880	60451816320	-869007242880
$\Delta^{10} f(n)$	3628800	-36288000	402796800	-4906137600	64988179200	-929459059200	14266826784000
$\Delta^{11} f(n)$	-39916800	439084800	-5308934400	69894316800	-994447238400	15196285843200	-248112809683200
$\Delta^{12} f(n)$	479001600	-5748019200	75203251200	-1064341555200	16190733081600	-263309095526400	4557208289356800
$\Delta^{13} f(n)$	-6227020800	80951270400	-1139544806400	17255074636800	-279499828608000	4820517384883200	-88166812070937600
$\Delta^{14} f(n)$	87178291200	-1220496076800	18394619443200	-296754903244800	5100017213491200	-92987329455820800	1792259345728051307
$\Delta^{15} f(n)$	-1307674368000	19615115520000	-315149522688000	5396772116736000	-98087346669312000	1885246675183872113	
$\Delta^{16} f(n)$	20922789888000	-3347646382080000	57119216394240000	-1034841187860480061983334021853184119			
$\Delta^{17} f(n)$	-3556874280960000	60466862776320000	-1091960404254720862086818140639232125				
$\Delta^{18} f(n)$	64023737057280000	-1152427267031040062196014181064704131					
$\Delta^{19} f(n)$	-1216451004088320062311256907767808137						
$\Delta^{20} f(n)$	2432902008176640143						

Well if you take the all the derivatives evaluated at 0 and raise them to the  $-1$  and sum them all up they will clearly converge to  $\frac{1}{e}$ .

$$n! = (-1)^n \Delta^{(n)} d_0$$

$$\frac{1}{e} = \sum_{n \geq 0} (\Delta^{(n)} d_0)^{-1}$$

$$(-1)^n n! = \Delta^{(n)} d_0$$

$$\frac{1}{e} = \sum_{n \geq 0} ((-1)^n n!)^{-1}$$

$$\frac{n!}{e} = n! \sum_{n \geq 0} ((-1)^n n!)^{-1}$$

$$\lfloor \frac{n!}{e} + \frac{1}{2} \rfloor = \lfloor \frac{1}{2} + n! \sum_{n \geq 0} ((-1)^n n!)^{-1} \rfloor$$

Well hey that is  $D_n$

$$D_n = \lfloor \frac{n!}{e} + \frac{1}{2} \rfloor = \lfloor \frac{1}{2} + n! \sum_{n \geq 0} ((-1)^n n!)^{-1} \rfloor$$