ISMT for SMT COMP 2023

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1 Overview

We propose a solver *ismt* for the model validation and a derived solver *yices-ismt* for the single query on QF_NIA that combines *yices2* [1] and *ismt*, where *yices* is the winner solver of QF_NIA of UNSAT performance in SMT-COMP 2021. The solver *ismt* is our experimental SMT solver, composed of a parser, a preprocessor, and a theory solver. The preprocessor performs a sequence of calls to formula rewriters. The theory solver consists of four parts: collector, decider, searcher, and resolver. In addition, *ismt* implements optional algorithm plugins invoked in the theory solver, such as bit-blasting, interval arithmetic, etc.

2 Dependencies

For completeness, *yices-ismt* invokes the SMT solver *yices2* 2.6.2 as a sub-solver. Rewriters and algorithms related to polynomial and interval are implemented on the libpoly library v0.1.11 [2]. The back-end SAT solver of the bit-blasting algorithm is from CaDiCal 1.5.2 [3].

3 Implementation

Ismt. Currently, *ismt* only supports QF_NIA. We propose a series of bit-width decision heuristics as a decider. After intervals for all variables are (partially or) completely decided, the searcher attempts to find a model in the space via an improved bit-blasting method from [4].

yices-ismt. Assume that the original SMT formula is ϕ , ismt searches in a finite space $F = \bigotimes_{x \in V} I(x)$, where V is the set of variables and $I(x) = l_x \le x \le u_x$ is the interval constraint for variable x. If ismt has not found a model in F, the combined solver yices-ismt will generate a lemma ψ to rule out the space where $\psi = \bigvee_{x \in V} \neg I(x)$. Otherwise, if ismt runs out of resources during searching, $\psi = \top$. We implement the solver yices-ismt with a portion of time-bound allocation

 $\{0.1, 0.8, 0.1\}$ for sequential execution of $\{yices2(\phi), ismt(\phi), yices2(\phi \wedge \psi)\}$. If the first two end earlier, $yices2(\phi \wedge \psi)$ runs within the remaining time. yices-ismt returns $\{sat, unsat\}$ once yices or ismt results $\{sat, unsat\}$, and unknown when the time budget is exhausted.

References

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