

# ISMT for SMT COMP 2022

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## 1 Overview

We propose a solver *ismt* for the model validation and a wrapper solver *yices-ismt* for the single query on QF\_NIA that combines *yices2* [1] and *ismt*, where *yices* is the winner solver of QF\_NIA of UNSAT performance in SMT-COMP 2021. The solver *ismt* is our experimental SMT solver, composed of a parser, a preprocessor, and a theory solver. The preprocessor performs a sequence of calls to formula rewriters. The theory solver consists of four parts: collector, decider, searcher, and resolver. In addition, *ismt* implements optional algorithm plugins invoked in the theory solver, such as bit-blasting, interval arithmetic, etc.

## 2 Dependencies

For completeness, *yices-ismt* invokes the SMT solver *yices2* 2.6.2 as a sub-solver. Rewriters and algorithms related to polynomial and interval are implemented on the libpoly library v0.1.11 [2]. The back-end SAT solver of the bit-blasting algorithm is from CaDiCal 1.5.2 [3].

## 3 Implementation

**Ismt.** Currently, *ismt* only supports QF\_NIA. We propose a series of bit-width decision heuristics as a decider. After intervals for all variables are (partially or) completely decided, the searcher attempts to find a model in the space via an improved bit-blasting method from [4].

**yices-ismt.** Assume that the original SMT formula is  $\phi$ , *ismt* searches in a finite space  $F = \bigotimes_{x \in V} I(x)$ , where  $V$  is the set of variables and  $I(x) = l_x \leq x \leq u_x$  is the interval constraint for variable  $x$ . If *ismt* has not found a model in  $F$ , the combined solver *yices-ismt* will generate a lemma  $\psi$  to rule out the space where  $\psi = \bigvee_{x \in V} \neg I(x)$ . Otherwise, if *ismt* runs out of resources during searching,  $\psi = \top$ . We implement the solver *yices-ismt* with a portion of time-bound allocation

$\{0.1, 0.8, 0.1\}$  for sequential execution of  $\{yices2(\phi), ismt(\phi), yices2(\phi \wedge \psi)\}$ . If the first two end earlier,  $yices2(\phi \wedge \psi)$  runs within the remaining time.  $yices-ismt$  returns  $\{sat, unsat\}$  once  $yices$  or  $ismt$  results  $\{sat, unsat\}$ , and *unknown* when the time budget is exhausted.

## References

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