

Local Declarations, Overloading

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Program Design and Data Structures

Based on notes by Tjark Weber, Lars-Henrik Eriksson, Pierre Flener, Sven-Olof Nyström



Today

- Local Declarations, Scope
- Overloading
- Revision

Local Declarations, Scope

Local Declarations

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```

The identifiers declared between `let` and `in` are bound only until the end of the following expression.

Local Declarations: Example

```
Prelude> let x = 42 in x + 1  
43
```

```
Prelude> x
```

```
<interactive>:....: Not in scope: `x'
```

Scope

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Different programming languages have different scoping rules.

In Haskell, the scope of a top-level declaration is typically the entire program (i.e., a name that is declared at the top level may be used anywhere in the program code). The scope of a local declaration only extends to the end of the corresponding expression.

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Prelude> (x, y)  
(0,11)
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Shadowing: Some Remarks

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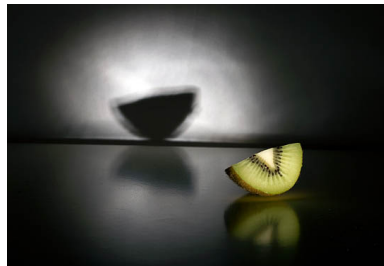
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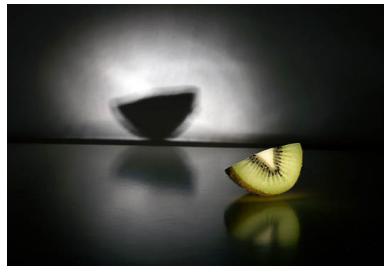
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- In principle, shadowing could be avoided by renaming so that all declarations use distinct names.
- However, in realistic programs shadowing typically does happen in some places (more or less by accident).



You need to know about shadowing because you should be able to read and understand realistic programs. Still, it is usually a good idea to avoid shadowing in your own code (by using distinct names).

Formal Arguments Bind Into the Function Body

The scope of a formal argument is the body of the function. Within the body, the argument shadows any earlier declaration of the same name.

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```
x = "shadowed"
```

```
add1 x = x + 1
```

```
Prelude> x
```

```
"shadowed"
```

```
Prelude> add1 42
```

```
43
```

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```
shadowedArgument x = let x=0 in x+1
```

```
Prelude> shadowedArgument 42
```

```
1
```

What's in a Name?

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By any other name would smell as sweet.*

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Names are only interpreted by human readers of the program. They serve to **explain** the purpose of values, functions, etc.

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Names are only interpreted by human readers of the program. They serve to **explain** the purpose of values, functions, etc.

Therefore, it is important to use **descriptive** names.

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(Of course, names that are declared outside the program, such as +, cannot be replaced in your code only.)

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We use software that can detect changed copies—even if you rename all identifiers!

Don't rename to cheat. (You won't cheat your way to a successful career anyway. Better try to pick up some effective study techniques.)

Overloading

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Prelude> 2.72 < 3.14
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We say that `+`, `-`, `*`, `/`, `<`, `>`, `==`, ... are **overloaded**.

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Even values from base types can be overloaded.

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Another example:

```
Prelude> maxBound :: Int
9223372036854775807
Prelude> maxBound :: Char
'\1114111'
```

Overloading: Type Inference

Sometimes, it is clear from the context (i.e., from the surrounding program code) what the concrete type of an overloaded function (or other value) is.

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Here, `<` has type `String->String->Bool`.

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```
Prelude> Data.Char.toUpper maxBound  
'\1114111'
```

Here, `maxBound` has type `Char`.

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Since `+` and `*` are available for different types (e.g., `Integer`, `Double`, ...), also `sumOfSquares` is available for these types:

```
Prelude> sumOfSquares 3
```

```
18
```

```
Prelude> sumOfSquares 3.0
```

```
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Some Standard Type Classes

Overloaded values live in **type classes**. A type class is a collection of types that support the same operation(s).

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Overloaded values live in **type classes**. A type class is a collection of types that support the same operation(s).

- **Eq** — types having equality/inequality operator
- **Ord** — types whose values can be ordered
- **Show** — types whose values can be converted **to** strings
- **Read** — types whose values can be converted **from** strings
- **Bounded** — types with both an upper and a lower bound.
- **Num** — numbers

The Eq typeclass

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class Eq a where  
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- Used in pattern matching (value patterns).
- Using == signals that this type class is to be used—that is, that the type of its arguments must belong to class Eq

The Ord typeclass

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class (Eq a) => Ord a where  
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- Allows to order values.
- Defined for base types, lists, tuples,...
- Used for sorting values.
- Assumes that values can be checked for equality: `(Eq a) => Ord a`.
 - `Ord` is a *subclass* of `Eq`.
 - `Eq` is a *superclass* of `Ord`.

The Show and Read typeclasses

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```
class Read a where  
  read :: String -> a
```

The Bounded typeclass

```
class Bounded a where  
  minBound, maxBound :: a
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- Examples: `Bool`, `Int`, `Char`, and tuples of bounded types.

The Num Typeclass

```
class (Eq a, Show a) => Num a where  
  (+), (-), (*) :: a -> a -> a  
  negate       :: a -> a  
  abs, signum  :: a -> a  
  fromInteger  :: Integer -> a
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- Examples: `Int`, `Integer`, `Double`, `Rational`, `Complex`, ...
- Can always be compared for equality (`Eq`) and printed (`Show`).
- Are not always ordered! (Why?)
- Has subclasses for standard division (`/`) and integral division (`mod`, `div`).

Reading Types with Class

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zero :: (Eq a, Num p) => a -> p  
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- then zero takes values of type a to values of type p .

Compare the logical formula

$$\forall a. \forall p. \text{Eq}(a) \wedge \text{Num}(p) \implies (a \implies p)$$

Revision

Basic Types in Haskell

- Integers (`Integer`): 0, 1, 2, -10, 42, ...
- Floating point numbers (`Double`): 0.0, -10.4, 3.14159, 1.234e6, ...
- Characters (`Char`): 'a', 'B', '1', ...
- Strings (`String`): "Hello", "PKD", "1", ...
- Booleans (`Bool`): `True`, `False`
- Unit (`()`): has only one value, written `()`
- ...

Note that 1, 1.0, '1' and "1" are all different.

Numeric Operations

For numeric types (`Num a`):

Functions	Type	Semantics
<code>+</code> , <code>-</code> , <code>*</code>	<code>a -> a -> a</code>	arithmetic operations
<code>==</code> , <code>/=</code>	<code>a -> a -> Bool</code>	equality, inequality
<code><</code> , <code><=</code> , <code>></code> , <code>>=</code>	<code>a -> a -> Bool</code>	comparison
<code>-</code> , <code>negate</code>	<code>a -> a</code>	negation (e.g., $-(2-3) \longrightarrow 1$)
<code>abs</code>	<code>a -> a</code>	absolute value (e.g., <code>abs (-1) \longrightarrow 1</code>)

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<code>==, /=</code>	<code>a->a->Bool</code>	equality, inequality
<code><, <=, >, >=</code>	<code>a->a->Bool</code>	comparison
<code>-, negate</code>	<code>a->a</code>	negation (e.g., $-(2-3) \rightarrow 1$)
<code>abs</code>	<code>a->a</code>	absolute value (e.g., <code>abs (-1) → 1</code>)

For integral types (`Integral a`):

Functions	Type	Semantics
<code>div, mod</code>	<code>a->a->a</code>	modulo (e.g., <code>mod 7 3 → 1</code>)

String Operations

Functions	Type ¹	Semantics
++	<code>String->String->String</code>	concatenation
length	<code>String->Int</code>	length
head	<code>String->Char</code>	first character
last	<code>String->Char</code>	last character
tail	<code>String->String</code>	the string without its first char
init	<code>String->String</code>	the string without its last char
take	<code>Int->String->String</code>	the string's first n characters
drop	<code>Int->String->String</code>	the string without its first n chars
!!	<code>String->Int->Char</code>	the string's n^{th} character

¹ These operations actually work for lists of any type, e.g. `[Int]`.

Conversions

Function	Type
<code>fromInteger</code>	<code>(Num a) => Integer -> a</code>
<code>toInteger</code>	<code>(Integral a) => a -> Integer</code>
<code>round</code>	<code>(RealFrac a, Integral b) => a -> b</code>
<code>truncate</code>	<code>(RealFrac a, Integral b) => a -> b</code>
<code>floor</code>	<code>(RealFrac a, Integral b) => a -> b</code>
<code>ceiling</code>	<code>(RealFrac a, Integral b) => a -> b</code>
<code>toEnum</code>	<code>(Enum a) => Int -> a</code>
<code>fromEnum</code>	<code>(Enum a) => a -> Int</code>
<code>show</code>	<code>(Show a) => a -> String</code>
<code>read</code>	<code>(Read a) => String -> a</code>

Operator Precedence and Fixity

How does Haskell know that $1 + 2 * 3$ should be evaluated as $1 + (2 * 3)$ rather than $(1 + 2) * 3$?

Every infix operator has a **precedence**. Higher precedences bind more tightly.

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Every infix operator has a **precedence**. Higher precedences bind more tightly.

How does Haskell know that $3 - 2 - 1$ should be evaluated as $(3 - 2) - 1$ rather than $3 - (2 - 1)$?

Infix operators may be **left-** or **right-associative**, meaning operations are grouped from the left (or right, respectively).

Comparison operators are not associative (what should $x == y >= z$ mean?).

Operator Precedence and Fixity

Prec	Left assoc	Non-assoc	Right assoc
9	!!		.
8			^, ^^, **
7	*, /, 'div', 'mod', 'rem', 'quot'		
6	+, -		
5			:, ++
4		==, /=, <, <=, >, >=	
3			&&
2			
1	>>, >>=		
0			\$, \$!, 'seq'

Pairs and Tuples

```
("PKD", 2019) :: (String, Integer)
```

```
(1, 2, 3) :: (Integer, Integer, Integer)
```

```
(3, "three", 3.0) :: (Integer, String, Double)
```

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The function `(,)` takes two arguments and returns the pair (2-tuple) that consists of these arguments (in the given order).

If-Then-Else, Declarations

```
if condition then trueValue else falseValue
```


If-Then-Else, Declarations

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```

Value declarations associate a value to an identifier.

```
x = 1  
myPi = 3.14159  
timesPi x = x * myPi
```

```
Prelude> x + x  
2  
Prelude> timesPi 2  
6.28318
```

Functions

Function declarations consist of the name, arguments, and body.

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takeLast n str = drop (length str - n) str
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Functions can also be anonymous.

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Prelude> (\n s -> drop (length s - n) s) 5 "computer"  
"puter"
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Prelude> (\n s -> drop (length s - n) s) 5 "computer"  
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```

Operators are just functions written between their arguments. We can convert between operators and functions using () and `f`.

```
Prelude> (+) 3 5  
8  
Prelude> 3 `takeLast` "PKD2019"  
"019"
```

Function Specifications

In this course, our function specifications will consist of

- 1 Function name and arguments
- 2 Function type — argument(s) and result
- 3 Description — what is the purpose of this function?
- 4 What the function assumes (precondition)
- 5 What the function computes (return value)
- 6 Side effects
- 7 Examples of function usage

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- 1 Function name and arguments
- 2 Function type — argument(s) and result
- 3 Description — what is the purpose of this function?
- 4 What the function assumes (precondition)
- 5 What the function computes (return value)
- 6 Side effects
- 7 Examples of function usage

Specifications may be given at different levels of detail.

Specifications may be given in natural language, or more formally.

Patterns

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Guarded patterns are of the form `pattern | expr`. Here, `expr` must have type `Bool` and will usually refer to names from `pattern`.

Uses of Patterns

Function declarations:

```
name pattern1 = expression1  
name pattern2 = expression2  
...  
name patternN = expressionN
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```
name pattern1 = expression1
name pattern2 = expression2
...
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```

Case expressions:

```
case expr of
  pattern1 -> expression1
  pattern2 -> expression2
  ...
  patternN -> expressionN
```

Haskell Sources

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- Top-level expressions are not allowed.
- Usually, each declaration is given on a separate line.
- The order of declarations doesn't matter.
- Multiple declarations of the same identifier are not allowed.