Haskell Basics

Johannes Borgström johannes.borgstrom@it.uu.se

Program Design and Data Structures

Based on notes by Tjark Weber, Lars-Henrik Eriksson, Pierre Flener, Sven-Olof Nyström

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Examples: 1, -25

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Integer is an **arbitrary precision** type: for each number, it uses as many bits as needed. Thus, it will hold any number no matter how big, up to the limit of your machine's memory.

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Note the incorrect result! Integer overflows are a frequent cause of programming errors.

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Examples: 1.0,
$$3e4 (= 30000.0)$$
, $3.4e-2 (= 0.034)$

Value syntax: one or more digits with decimal point and/or in scientific notation (letter e followed by an integer). Negative numbers are preceded by –

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Haskell also offers a type Float that has even less precision than Double (but uses less memory). Don't use Float unless you Really Know What You Are Doing!

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```
Prelude > abs ((1.0 + 1e16 - 1e16) - 1.0) < 10.0 True
```

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String is actually just an abbreviation for [Char], i.e., lists of characters.

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See http://unicode-table.com/en/ for the Unicode table.

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[] ²	Char->String
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${\tt fromEnum}$	Char->Int
show	Integer->String
show	Double->String

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² E.g., ['a']. This isn't actually a function, but special syntax.

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Prelude> read "42"

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Prelude> read "42" :: Integer

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Even for functions like read, it is *not* necessary to indicate the desired type when this can be determined from the program context:

```
Prelude> read "42" + 1
43

Prelude> read "42" + 1.0
43.0
```

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Explanation: True and False are (the only) values of type Bool, just like 1 and 42 are values of type Integer.

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Functions with result type Bool are sometimes called **predicates**.

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Example:

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(With strict evaluation, this expression would throw an exception.)

The Haskell Prelude

Many other useful functions are provided by the Haskell Prelude: see

http://hackage.haskell.org/package/base/docs/Prelude.html

These are available by default in all Haskell programs.

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However, we have already seen many examples of **infix operators**, i.e., functions that take two (or more) arguments and are written *between* their arguments. For instance,

```
1 + 2
2.72 == 3.14
"foo" ++ "bar"
```

Infix vs. Prefix Notation

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To use a prefix function in infix position, simply enclose the function in backticks. For instance,

```
Prelude > 7 `mod` 3
1
```

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Infix Operators: Precedence

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*, /, `div`, `mod`	7	
+, -	6	
++	5	
==, /=, <, <=, >, >=	4	
&&	3	
11	2	

Source: http://www.haskell.org/onlinereport/haskell2010/haskellch4.html#x10-820004.4.2

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(It is possible to change these precedence values.)

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(It is possible to change the associativity.)

GHCi's : info

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Tuples

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Examples:

```
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Fine points:

- There are no 1-tuples in Haskell: (42) is just 42 in parentheses.
- It is possible to have tuples of tuples: e.g., ((1,2),(3,4)).
- The value () is the only "0-tuple". Its type is ().

Product Types

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```
(1, 2, 3) :: (Integer, Integer, Integer)
("PKD", 2019) :: (String, Integer)
(3, "three", 3.0) :: (Integer, String, Double)
```

Pairs: (,)

The function (,) takes two arguments and returns the pair (2-tuple) that consists of these arguments (in the given order).

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```
Prelude> (,) "PKD" 2019 ("PKD",2019)

Prelude> (,) 1 2 (1,2)

Prelude> (,) 2 1 (2,1)
```

Pairs: fst and snd

Functions fst and snd extract the first and second, respectively, component of a pair.

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2
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Later, we will see how to extract components from n-tuples for n > 2.

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```
if 2+2==5 then "hel"++"lo" else "good"++"bye"
```

Conditional Expressions: Type

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All parts are mandatory. if condition then trueValue is not a valid expression. (What should be its value when condition is false?!)

Exercises

• Express each of the following expressions as if-then-else expressions. In other words, find condition, trueValue and falseValue such that the given expression is equivalent to if condition then trueValue else falseValue

- 1 E | | I
- E && F

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Express each of the following expressions as if-then-else expressions. In other words, find condition, trueValue and falseValue such that the given expression is equivalent to

- E || F
- 2 E && F
- Evaluate (step-by-step) the following expression:

if condition then trueValue else falseValue

```
if 1 + 2 < 4 then length ("hel"++"lo!") else 4 'div' 2
```

Value Declarations

Value declarations associate a value to an identifier.

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```
Prelude> x = 1
Prelude> myPi = 3.14159
Prelude> twoPi = 2.0 * myPi
Prelude> (@@) = "use descriptive identifiers!"

Prelude> x + x
2
Prelude> twoPi
6.28318
```

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Prelude> x + x
2
Prelude> twoPi
6.28318
```

Note: value declarations are not expressions!

Identifiers: Syntax

Note that 3+(-2) is different from 3+-2! Haskell thinks that +- is an (undeclared) identifier.

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```
Prelude > 3+(-2)
1

Prelude > 3+-2

<interactive >: 3:2:
    Not in scope: `+-'
    ...
```

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Bindings and Environments

The execution of a declaration, say $x = \exp r$, creates a **binding**: the identifier x is *bound* to the value of the expression expr.

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In GHCi, you can use :show bindings to show the current bindings and their type:

```
Prelude > : show bindings
x :: Integer = 1
myPi :: Double = 3.14159
twoPi :: Double = _
(@@) :: [Char] = _
```

Changing Environments

Later declarations of the same identifier change the environment.

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```
Prelude > x = 1
Prelude > : show bindings
x :: Integer = 1
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Prelude > : show bindings
x :: Integer = 2
Prelude > x = "x"
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```

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```

Declarations in Haskell are similar to defining equations in mathematics.

GHCi's it

The identifier it is always bound to the value of the last expression that was evaluated.

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```
Prelude > 1 + 2
3
Prelude > : show bindings
it :: Integer = 3
Prelude> it
3
Prelude > "another " ++ "expression"
"another expression"
Prelude> it
"another expression"
```

Execution of Declarations

Remember that Haskell is non-strict? When a declaration, say $x = \exp x$, is executed, the value of $\exp x$ is not computed right away. Instead, it is computed later, when (and if) x is evaluated.

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```
Prelude> x = 1 + 1 :: Int
Prelude> :show bindings
x :: Int = _

Prelude> x -- force evaluation of x
2
Prelude> :show bindings
x :: Int = 2
it :: Int = 2
```

Execution of Declarations (cont.)

Therefore, it is possible to bind x to an expression whose evaluation would result in a runtime error. The runtime error will only be generated when (and if) x is evaluated.

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```
Prelude> x = 1 `div` 0
Prelude> x
*** Exception: divide by zero
```

Identifiers Are Not Variables

.. and declarations are not variable updates.

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```
Prelude> x = 10
Prelude> addX y = x+y
Prelude> addX 42
52

Prelude> x = 20
Prelude> addX 42
52
```

Practical Matters

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Remember: attendance (at 7 out of 9 labs this Fall) is mandatory!