

Polymorphism,  
Higher order functions

# Polymorphism (again)

- ▶ Types like `Integer`, `String`, `[Int]`, `(String, Int)` are simple to interpret
  - ▶ we know exactly what they are
- ▶ We have also seen *polymorphic* types
  - ▶ `[] :: [a]`
    - ▶ empty list with unknown content
  - ▶ `length :: [a] -> Int`
    - ▶ compute length of list without caring about the actual contents
  - ▶ `map :: (a->b) -> [a] -> [b]`
    - ▶ apply a function to each element in the list
    - ▶ the *type variables* connect the first and second argument; the elements of the list argument must be of the same type that the function takes
- ▶ A polymorphic type can be said to be *partially* known
  - ▶ type variables work the same as arguments to a function
  - ▶ replace (instantiate) type variables consistently with another type
- ▶ polymorph - many shapes

# Polymorphism (more)

- ▶ We have also seen types like
  - ▶ `index :: (Num t, Eq a) => a -> [a] -> t`
  - ▶ These can *not* be arbitrary types, but must fulfil additional *properties*
  - ▶ `t` must be a “number”
  - ▶ `a` must be testable for equality
- ▶ This is called *type classes* (or just classes) in Haskell
  - ▶ This is what makes *overloading* possible
  - ▶ Same operator used on different types, different implementations
- ▶ `(<) :: Ord a => a -> a -> Bool`
  - ▶ conceptually simple “less than”
  - ▶ different object compare in different ways
  - ▶ numbers
  - ▶ characters
  - ▶ strings
  - ▶ tuples
  - ▶ lists

# Type Classes

- ▶ A type class is defined by functions that need to be implemented by the types that belong to the type
  - ▶ `Eq a` consists of only `(==) :: a -> a -> Bool`
- ▶ We can make a type *belong* to a type class by defining the required function(s)
- ▶ Not really relevant for the types we define now, but will be later.

Type classes are a really beautiful concept!

# Abstraction (part 1)

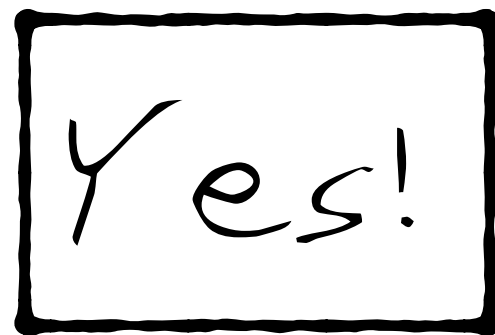
- ▶ Polymorphic types allow us to write functions without caring about the full details of the types involved.
- ▶ Overloading with type classes allow us to use the same operator (express the same concept) with different implementations.
- ▶ These are two instances of hiding details and leveraging concepts.
- ▶ Abstraction allows us to write programs that are
  - ▶ easy to understand
  - ▶ easy to maintain (change)
  - ▶ easy to test
- ▶ More to come..



Abstraction

# Higher Order Functions

- ▶ We have seen functions that can take another function as argument
  - ▶ `map :: (a->b) -> [a] -> [b]`
  - ▶ First argument is a function applied to each element in the list
- ▶ We have seen examples where we give the *name* of function
  - ▶ `map length ["lisp", "haskell", "erlang", "sml"]`
- ▶ ..the *value* of `length` is a *function*..
- ▶ ..but we don't have to name every value we use..
- ▶ ..can we write a function directly..?



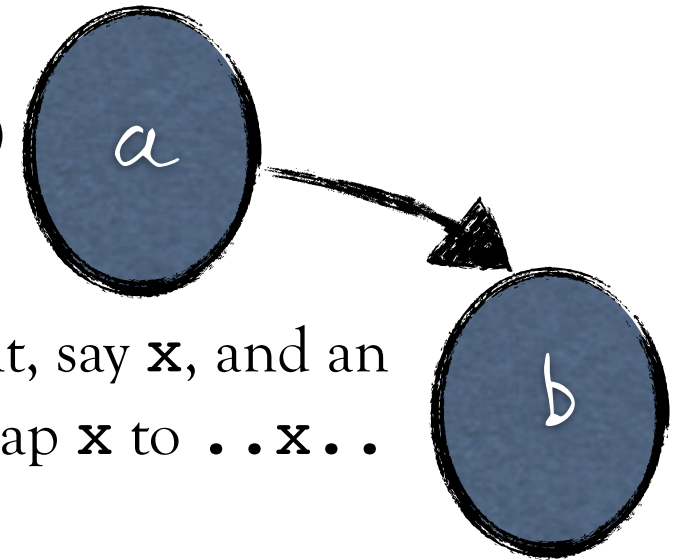
Yes!

# First Class Values

- ▶ In Haskell functions are said to be *first class* values
  - ▶ equivalent in treatment to *normal* values
- ▶ A name can be bound to a function
  - ▶ `add1 x = x + 1`
- ▶ A function can be passed as an argument to a function
  - ▶ As in `map` and `filter`
- ▶ A function can be returned as the result of a computation
  - ▶ We can *construct* functions at runtime
- ▶ We can write functions without having to give them a name
  - ▶ `..ok..`
- ▶ Why *first class*?
  - ▶ Unusual treatment of functions compared to other language
  - ▶ passing a function name as argument quite common
  - ▶ returning a function unusual
- ▶ Functional Programming Languages have it, and always have
  - ▶ They are first class!

# Constructing Functions (pt.1)

- ▶ When using **map** it is somewhat roundabout to define a named function and then passing the name, especially when the function is “small”
- ▶ What, then, is a function?
  - ▶ A mapping from one set (of values) to another set (of values)
  - ▶ Reflected in a function type  $a \rightarrow b$
- ▶ For *computing* a function (of one argument) we have an argument, say  $x$ , and an expression  $..x..$ , called a *body*, containing  $x$ , so we want to map  $x$  to  $..x..$ 
  - ▶  ~~$x \rightarrow ..x..$~~
  - ▶ Add something to distinguish from a function type
  - ▶  $\backslash x \rightarrow ..x..$ 
    - ▶  $\backslash x \rightarrow x+1$  is a function that adds 1
  - ▶ Original notation:  $\lambda x.x+1$ 
    - ▶ lambda calculus - the basis of functional programming (1931)
    - ▶ That's all you really need!



Anonymous function



# Using Anonymous Functions

- ▶ To add 1 to each element in a list
  - ▶ `map (\x->x+1) [19,31,43,59]`
- ▶ Filter numbers divisible by 7
  - ▶ `filter (\n->mod n 7==0) [1,19,47,117,35]`
- ▶ Write a function that filters numbers divisible by a given argument
  - ▶ `filterDiv :: Int -> [Int] -> [Int]`
  - ▶ Obvious to use `filter`

```
filterDiv :: Int -> [Int] -> [Int]
filterDiv n list =
    filter (\e->mod e n == 0) list
```

# Use of Function Arguments

- ▶ Nothing special when using a function argument
  - ▶ define and use as normal

```
map :: (a -> b) -> [a] -> [b]
map f [] = []
map f (x:xs) = (f x):map f xs
```

# List Function Patterns

- ▶ In the example for evaluating simple expressions, we had

- ▶ one function for *adding* all elements in a list
- ▶ one function for *multiplying* all elements in a list

- ▶ Very similar pattern

- ▶ Differences

- ▶ **Value** returned for empty list
- ▶ **Operation** in inductive case

- ▶ Can we write a function that can do both?

- ▶ Use common pattern
- ▶ *Inject* operation and value with arguments

```
sumlist [] = 0
sumlist (x:xs) = x + sumlist xs

multlist [] = 1
multlist (x:xs) = x * multlist xs
```

Value

Operation

```
fold :: (t1 -> t2 -> t2) -> t2 -> [t1] -> t2
fold f a [] = a
fold f a (x:xs) = f x (fold f a xs)
```

# Usage

```
fold :: (t1 -> t2 -> t2) -> t2 -> [t1] -> t2
fold f a [] = a
fold f a (x:xs) = f x (fold f a xs)
```

```
sumlist list = fold (+) 0 list

multlist list = fold (*) 1 list

-- factorial in another way
fac n = multlist [1..n]
```

- ▶ Note: + and \* *commutative*
  - ▶  $a+b=b+a$  and  $a*b=b*a$
- ▶ `fold (+) 0 [1,2,3]` is
  - ▶  $1 + (2 + (3 + 0)) = 6$
- ▶ - is *not* commutative
- ▶ `fold (-) 0 [1,2,3]` is
  - ▶  $1 - (2 - (3 - 0)) = 2$

- ▶ Alternatives:
  - ▶  $((1 - 2) - 3) - 0 = -5$
  - ▶  $((0 - 1) - 2) - 3 = -6$
- ▶ Second is a reasonable symmetric alternative
- ▶ Two definitions of **fold**

# Variants of fold

```
– fold from the right
foldr :: (t1 -> t2 -> t2) -> t2 -> [t1] -> t2
foldr f a [] = a
foldr f a (x:xs) = f x (foldr f a xs)

– fold from the left
foldl :: (t1 -> t2 -> t1) -> t1 -> [t2] -> t1
foldl f a [] = a
foldl f a (x:xs) = (foldl f (f a x) xs)
```

- ▶ Both exist in Haskell
- ▶ Note different types of function argument
  - ▶ **foldr**: first argument is the element in the list, second is the value
  - ▶ **foldl**: first argument is the value, second is the element in the list
- ▶ .. and note that the function argument takes *different* types
  - ▶ we can do more than just numeric operations
- ▶ **foldl** is an abstraction for list recursion with an accumulator

# Abstraction (part 2)

- ▶ With higher order list operations like `map`, `filter`, `foldl` and `foldr` we
  - ▶ hide the details of the list recursion
  - ▶ focus on the important part, i.e., the function being passed as argument
- ▶ A program using these will be
  - ▶ more compact
  - ▶ easier to read
  - ▶ easier to maintain
  - ▶ less cluttered with common patterns
- ▶ Side note:
  - ▶ `map` and `filter` can be defined in terms of a fold