

# Introduction to Functional Programming and Haskell

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Program Design and Data Structures

Based on notes by Tjark Weber, Lars-Henrik Eriksson, Pierre Flener, Sven-Olof Nyström



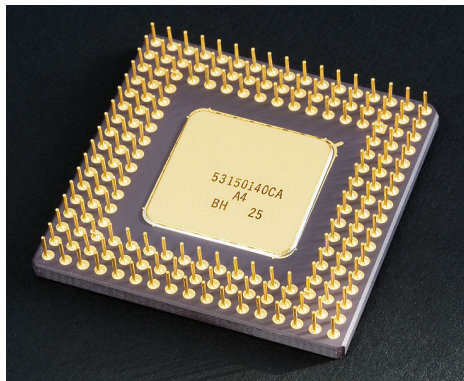
# Today

- Introduction to functional programming
- Introduction to Haskell

# Introduction to Functional Programming

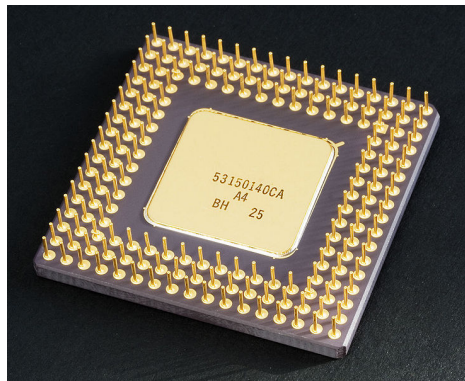
# Abstraction from Hardware

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However, high-level programming languages provide a **higher level of abstraction**.

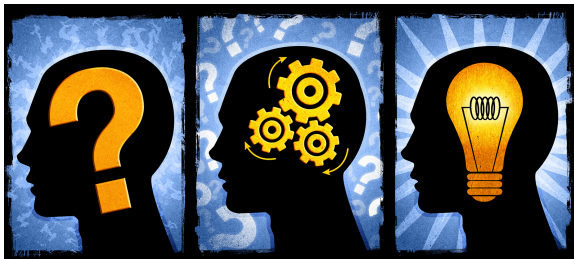
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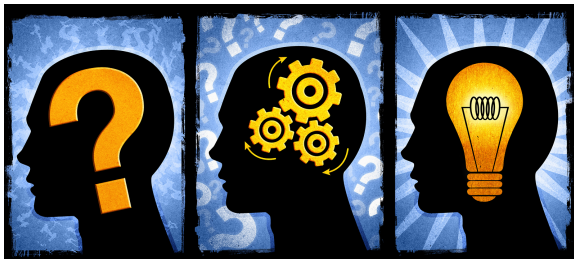
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Therefore, they can support advanced concepts that allow the programmer to focus on **solving problems** (rather than instructing a machine).



Programmers use many different high-level programming languages: e.g., C, C++, Java, PHP, Perl, Python, MATLAB, Lisp, Haskell.



# Programming Paradigms

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(e.g., C)

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# General-Purpose Languages

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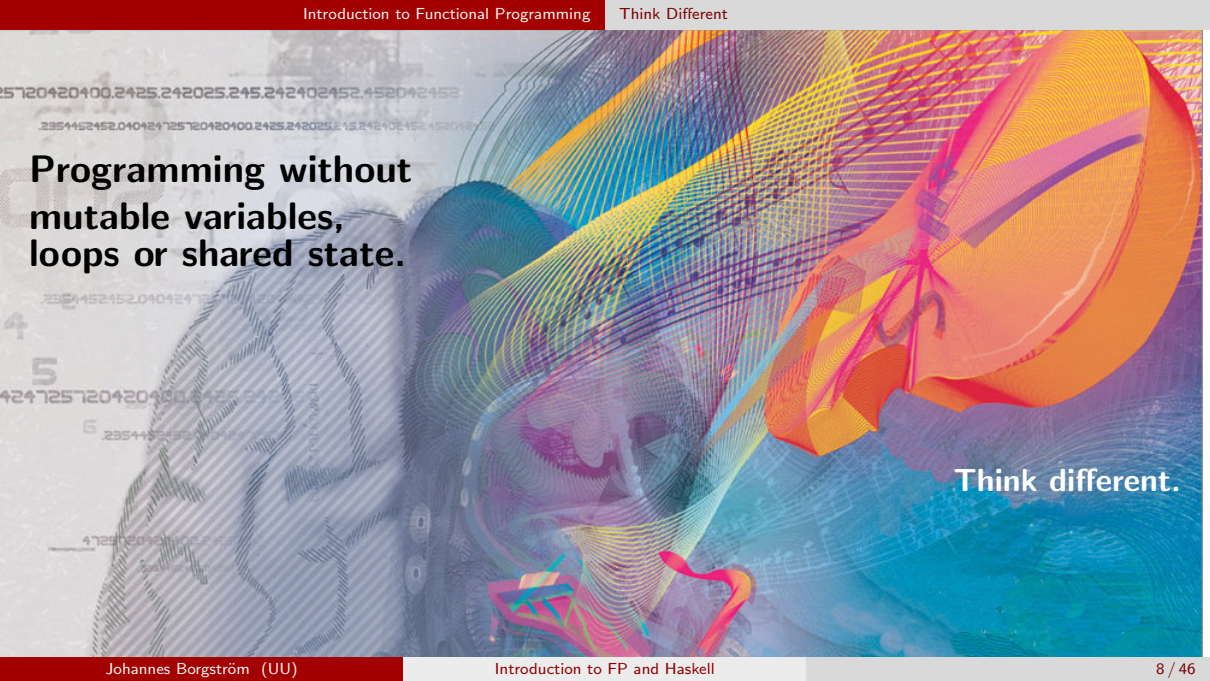
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In practice many “problems” are engineering challenges (rather than mathematical functions), and these can be much **easier** to solve in one language or another.



The background of the slide is a complex, abstract composition. It features a grayscale image of a person's face, which is partially obscured by a dense network of colorful, overlapping lines and curves in shades of blue, yellow, orange, and red. Scattered throughout the image are various numbers, including '25720420400.2425.242025.245.242402452.452042452', '2354452452.040424725720420400.2425.242025.215.242402452.452042452', '2354452452.040424725720420400.2425.242025.215.242402452.452042452', '424725720420400.2425.242025.215.242402452.452042452', '2354452452.040424725720420400.2425.242025.215.242402452.452042452', and '4725720420400.2425.242025.215.242402452.452042452'.

**Programming without  
mutable variables,  
loops or shared state.**

**Think different.**



CREDIT SUISSE



ERICSSON



facebook



Google



IBM



Microsoft

# Why Functional Programming?

- Shorter, more readable code

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# Functional Programming Languages

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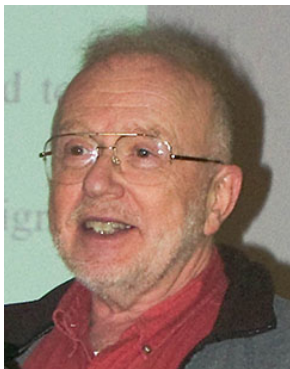
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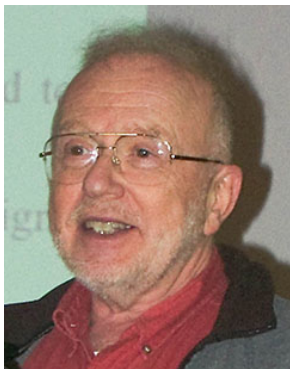
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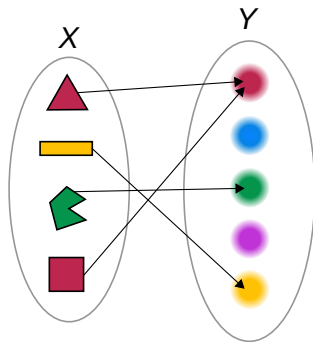
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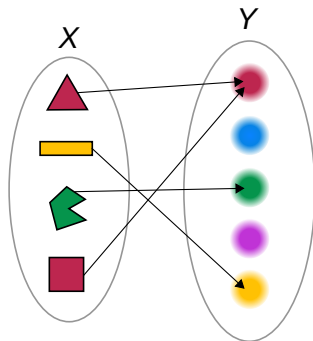
David Parnas (born 1941) is a Canadian early pioneer of software engineering. He developed the concept of information hiding in modular programming.

[http://en.wikipedia.org/wiki/David\\_Parnas](http://en.wikipedia.org/wiki/David_Parnas)

# Functions



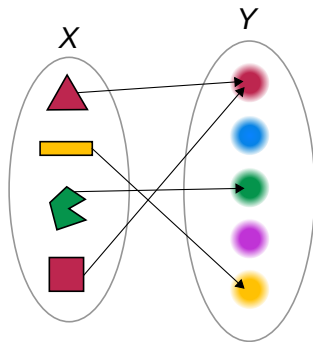
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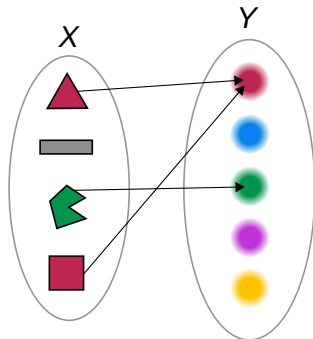


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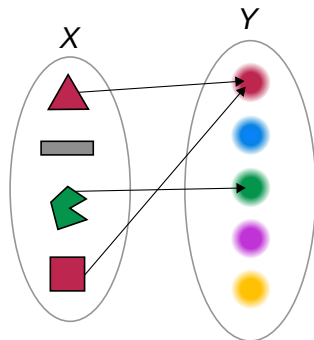
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# Partial Functions



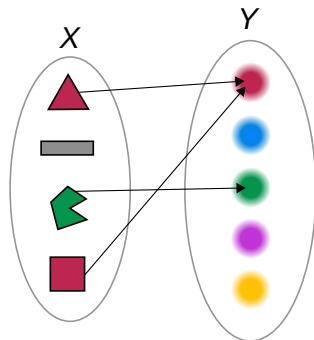


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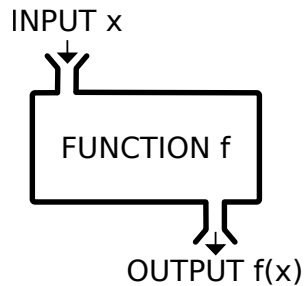
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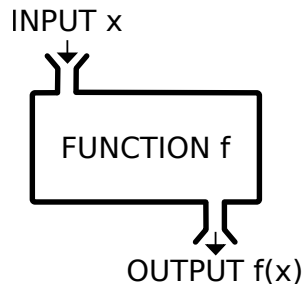
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A function may be given by an **expression** that determines the output value. For example,

```
double n = 2 * n
```

# Function Composition

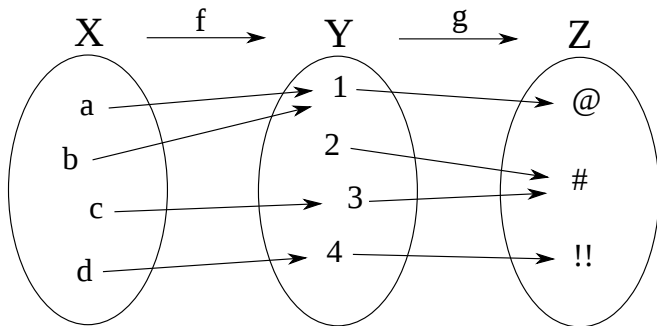
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Some aspects of the language are slightly intricate. We'll try to introduce these gently, but they may surprise you from time to time.



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(If you install from some other source, make sure that you get “Haskell Platform”, and not one of the cut-down distributions.)

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`Prelude>` is the GHCi prompt. You can now enter Haskell expressions, and GHCi will evaluate them for you!



# Haskell Expressions: Arithmetic

$2 + 15$

$49 * 100$

$1892 - 1472$

$5 / 2$

$50 * 100 - 4999$

$(50 * 100) - 4999$

$50 * (100 - 4999)$

$5 * (-3)$

# Haskell Expressions: Boolean Algebra

```
True && False
```

```
True && True
```

```
False || True
```

```
not False
```

```
not (True && True)
```

# GHCi: Interrupting and Leaving

To interrupt an ongoing computation, press **Ctrl+C**:

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To leave GHCi, type **:quit** at its prompt:

```
Prelude> :quit  
Leaving GHCi.
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Alternatively, press **Ctrl+D**.

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Well, right now it may look like that. But that's just because we haven't seen much of Haskell yet.

In fact, Haskell—like other modern languages—is Turing-complete. Any (mathematical) problem that can be solved by a computer program at all, can be solved by a Haskell program.

# Syntax

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(These are just examples. You can't really tell valid from invalid programs yet, because we haven't discussed the syntax of Haskell. We will introduce it bit by bit in the course, mostly through examples of valid programs.)

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However, GHC can't know what you had in mind ...

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- `+` means addition
- `2+15` means “the sum of 2 and 15”, i.e., 17
- `f x` means “the function `f` applied to the argument `x`”

# Syntax and Semantics: Common Descriptions

The syntax of programming languages is often defined very precisely, e.g., by formal grammars or syntax diagrams. Thus, it is (usually) possible to determine precisely—and by computer—whether a program is valid.



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It is possible—but a lot of work—to define the semantics of every program construct equally precisely. In practice, programmers mostly rely on informal descriptions (e.g., in English).

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<code>"Hello "++"world"</code>	<code>"Hello world"</code>
<code>not False</code>	<code>True</code>

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To evaluate functions applied to arguments, Haskell uses an evaluation strategy known as **non-strict evaluation**. Arguments are evaluated only if their value is needed in the evaluation of the function body. More on this later.

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Therefore, functional programming is well-suited for parallel computing (e.g., multi-core machines).

# Whitespace

You may freely insert spaces between expressions:

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Prelude> 1+2
```

```
3
```

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Later, we will see programs where spaces do matter.

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These comments may extend over several lines when used in Haskell programs.



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Not every expression can be evaluated to a (sensible) value: e.g.,

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Prelude> div 1 0
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*** Exception: divide by zero
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Later, we'll see how to recover from an exception, and how you can write your own functions that throw exceptions.

# Types

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A **(data) type** is a set of data with similar properties: e.g., possible values, operations that can be applied, and meaning.

At the hardware level, data is (mostly) just bits and bytes. Consequently, machine code is (mostly) untyped. However, many high-level languages offer data types.

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- Unit (`()`): has only one value, written `()`

Note that 1, 1.0, '1' and "1" are all different.

# Seeing Types in GHCi

You can use `:type expr` to have GHCi show the type of `expr`.

```
Prelude> :t 'a'
'a' :: Char
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You can also `:set +t` to see the type of all expressions that you enter.

```
Prelude> :set +t
Prelude> 'a'
'a'
it :: Char
Prelude> not True
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it :: Bool
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# Function Types

Also functions have types. For instance,

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Here, the first `Bool` is called the **argument type** and the second `Bool` is called the **result type** of the function.

`->` is a **type constructor**: when applied to two existing types, it constructs a new type (namely, a function type).

# Expression Language, Type Language

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We have seen ways to write Haskell types, such as

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Integer  
Bool -> Bool  
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```

When we write a program, it is clear from the context whether we are writing an **expression** or a **type**. For instance,

`expr :: Type`

# Type Checking

What should `32 + "1"` mean?

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In a **strongly typed** language, this is never allowed.

In a **weakly typed** language, the result depends (on the language). It could be 33, or **"321"**, or ...

# Static vs. Dynamic Type Checking

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In a language with **static** type checking, the compiler inspects the program and generates an error. Translation to machine code fails.

In a language with **dynamic** type checking, the compiler translates the program to machine code that performs type checking at run time, i.e., that generates an error when it is executed.

# Type Checking in Haskell

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# Type Inference

In many languages, the programmer needs to declare the type of variables, functions, etc.

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In Haskell, this is (mostly) not necessary. Instead, the types of expressions are automatically **inferred** by (a type inference algorithm that is part of) the Haskell implementation.



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This happens statically, i.e., without running the program.

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Again, **try to understand** these error messages! And be patient with GHC: it can't know what you really had in mind ...