



## Weekly Research Progress Report

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**# of hrs worked this week:** 30

Problems discussed in last week's report

### 1. Reading materials

- a. Summarize the new policies
- b. Summarize the test-cases produced by Redis
- c. Replicate the test-cases introduced by Redis

### 2. Integration of Jepsen I/O with RamCloud

- a. Configure the installation of five instances
- b. Implement the client-side of RamCloud

## Reading Materials

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### Generalized Isolation Level Definitions

ICDE, 2000

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#### Previous work

The isolation levels were named READ UNCOMMITTED, READ COMMITTED, REPEATABLE READ, and SERIALIZABLE.

#### Dirty Write

T2 cannot write an object  $x$  if an uncommitted transaction T1 has already modified  $x$ .

T1	W[X]		C or A							
T2		W[X]								

#### Dirty Read

T1 to acquire a long writelock and T2 to acquire (at least) a short-term read-lock, and proscribing P2 requires the use of long read and write locks

T1	W[X]		C or A							
T2		R[X]								

### Fuzzy Read

T1	R[X]		C or A							
T2		W[X]								

### Restrictiveness of Preventative Approach

T1	R(X,5)	W(X,1)				R(Y,5)	W(Y,9)	C		
T2			R(X,1)	R(Y,5)	C					

Optimistic and multi-version mechanisms disallow non-serializable histories

T1	R(X,5)	W(X,1)	R(Y,5)	W(Y,9)			C			
T2					R(X,1)	R(Y,9)		C		

P1 disallows it.

T1	W(z1)	W(x1)	W(y1)		C						
T2						R(x1)	W(y2)	C			
T3				W(x3)					R(y2)	W(z3)	C

- T1 ---(ww)---> T2
- T1 ---(wr)---> T2
- T1 ---(ww)---> T3

### New Generalized Isolation Specification

- **G0: Write Cycles**
  - $x1 << x2, y2 << y1$  (version depends on the actual operation)

T1	W(x1)				W(y1)	c				
T2		W(x2)	W(y2)	c						

- **G1a: Aborted Read**

- if history contains an aborted transaction T1 and a committed transaction T2 such that T2 has read some object (maybe via a predicate) modified by T1

T1	W(x1)		A							
T2		R(x1)	C							

- **G1b: Intermediate Reads**

- If it contains a committed transaction T2 that has read a version of object written by transaction T1 that was not T1's final modification of x.

T1	W(x1)		W(x1,1)		C					
T2		R(x1)		C						

- **G1c: Circular Information Flow**

- If DSG(H) contains a directed cycle consisting entirely of dependency edges

- **G2: Anti-dependency Cycles**

- if DSG(H) contains a directed cycle with one or more anti-dependency edges.

Level	Phenomena disallowed	Informal Description ( $T_i$ can commit only if:)
PL-1	G0	$T_i$ 's writes are completely isolated from the writes of other transactions
PL-2	G1	$T_i$ has only read the updates of transactions that have committed by the time $T_i$ commits (along with PL-1 guarantees)
PL-2.99	G1, G2-item	$T_i$ is completely isolated from other transactions with respect to data items and has PL-2 guarantees for predicate-based reads
PL-3	G1, G2	$T_i$ is completely isolated from other transactions, i.e., all operations of $T_i$ are before or after all operations of any other transaction

Bac