· GCD - greatest common divisor - polynomial, of the highest possible degree, that is or torctor of both the two original polynomials

1) tind roots of polyhomials:

$$\alpha: d_{1,2} = \frac{-5 \pm 75^2 - 4\alpha c}{2\alpha} = \frac{-7 \pm 749 - 4 \cdot 1.6}{2 \cdot 1} = \frac{-7 \pm 725}{2} = \frac{-1}{2}$$

Is this approach effective? Tes, but only for opvaduatic equations (see the tormula for calculating roots of cubic equation)

· GCD - how to calculate effectively?

=) Extended Euclidean algorithm

Theory: greatest common divisor g=gcd(a,5) can be expressed by using the Bézout's lemma as in:

where

g=gcd(a,s) ... greatest common divisor

a, b known (input) polynomials, We are looking ton the gcd of these polynomials

Piar... pair of unknown coprime polynomials. Two polynomials are coprime it and only it they share no roots. Hence, gcd(P,or)=1.

Similarly, the least common multiple l= landa, b) can be expressed by using coprine polynomials pair r,s as in:

$$l = lcm(a,b) = ar = -b \cdot s$$

Hence

$$o(r+b)s=0$$

again gcd (r,s)= 1 => polynomials h,s are coprime.

Now we howe two equations:

In matrix torm:

However, we also want to obtain polynomials paying.

=) Multiplication by identity matrix:

$$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} P & r \\ q & s \end{bmatrix} = \begin{bmatrix} P & r \\ q & s \end{bmatrix}$$

$$=) \begin{bmatrix} a & b \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} p & r \\ q & s \end{bmatrix} = \begin{bmatrix} 9 & 0 \\ p & r \\ q & s \end{bmatrix}$$

Textended Euclidean algorithm computes "output" matrix

orplying a sequence of column operations on the input matrix.

Extended Evolidean algorithm's overview:

a, b are input polynomials; g is ged of input polynomials; and play and his are coprine polynomials pairs.

Allowed columns operations:

-> collumns coin be exchanged

$$\begin{bmatrix} C_1 & C_2 \end{bmatrix} \sim \begin{bmatrix} C_2 & C_1 \end{bmatrix}$$

Syntax of this operation:

C1 (k+1) = C2 (k) ... coflumn C, in algorithm's step k+7 is equal to collumn Cz in algorithm's step k $C_2(k+1) = C_1(k)$ collumn C_2 in algorithm's stop k+1 is equal to collumn on in algorithm's step K

$$\begin{bmatrix} d+1 & d+2 \\ 0 & 0 \\ 0 & 1 \end{bmatrix} = \frac{C_0(k+1)}{C_0(k+1)} = \frac{C_0(k)}{C_0(k+1)} \begin{bmatrix} d+2 & d+1 \\ 0 & 1 \\ 0 & 1 \end{bmatrix}$$

-) arbitrary column can be multiplied with nonzero constant

Syntax of this operation:

 $C_1(k+1) = d \cdot C_1(k)$... Column C_1 in algorithm's step k+1 is equal to d-times column C_1 in algorithm's step k

> any column can be multiplied with ensitting polynomial and result con be added to second column

 $C_2(k+1) = C_2(k) + x \cdot C_1(k)$. Solumn C_2 in algorithm's step k+1 is equal to sum of C_2 in k and C_1 in k that is multiplied by polynomial x

Ex:

Ex: Find gcd
$$(a, b)$$
, where $a = d^2 + 7d + 6$
 $b = d^2 - 5d - 6$

$$\begin{bmatrix} 0 & 5 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} d^2 + 72 + 6 \\ 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} d^{2} + 72 + 6 & d^{2} - 52 - 6 \end{bmatrix} c_{2}(1) = c_{2}(0) - c_{1}(0) \begin{bmatrix} d^{2} + 72 + 6 & -122 - 12 \\ 1 & 0 \end{bmatrix} c_{2}(2) = \frac{7}{2}(2)$$

$$\begin{bmatrix} d^{2} + 72 + 6 & d + 1 \\ 7 & \frac{1}{12} \\ 0 & -\frac{7}{12} \end{bmatrix} = G(2) - d \cdot G(2) \begin{bmatrix} 6d + 6 & d + 1 \\ 1 - \frac{7}{12}d & \frac{7}{12} \\ \frac{1}{12}d & -\frac{7}{12} \end{bmatrix}$$

$$\begin{bmatrix} d+1 & d+1 \\ \frac{1}{6} - \frac{1}{72} d & \frac{1}{72} \\ \frac{1}{2} d & -\frac{1}{72} \end{bmatrix} \sim \begin{bmatrix} c_2(5) = c_2(4) - c_1(4) \\ \frac{1}{6} - \frac{1}{72} d & \frac{1}{72} d - \frac{1}{72} d \end{bmatrix} = \begin{bmatrix} g & 0 \\ p & h \end{bmatrix}$$

$$\begin{bmatrix} \frac{1}{6} - \frac{1}{72} d & \frac{1}{72} d - \frac{1}{12} \\ \frac{1}{72} d & -\frac{1}{12} d - \frac{1}{12} \end{bmatrix} = \begin{bmatrix} g & 0 \\ p & h \end{bmatrix}$$

$$a = d + 1$$

 $b = d^2 + 2 - 2$

$$\begin{bmatrix} d+1 & d^{2}+d-2 \\ 0 & 1 \end{bmatrix} = C_{2}(0) - d \cdot C_{1}(0) \begin{bmatrix} d+1 & -2 \\ 1 & -d \\ 0 & 1 \end{bmatrix} = C_{2}(2) = C_{2}(1)$$

$$\begin{bmatrix} 2+7 & 1 \\ 7 & \frac{d}{2} \\ 0 & -\frac{7}{2} \end{bmatrix} = C_{1}(2) - 2 \cdot C_{2}(2)$$

$$\begin{bmatrix} 1 & 1 \\ 1 - \frac{d^{2}}{2} & \frac{d}{2} \\ \frac{d}{2} & -\frac{7}{2} \end{bmatrix}$$

$$E_X$$
: What is the least common multiple of a, b

$$a = d + 7$$

$$b = d^2 + d - 2$$

(see previous example)

$$\begin{aligned}
r &= \frac{d^2}{2} + \frac{d}{2} - 1 \\
S &= -\frac{d}{2} - \frac{1}{2}
\end{aligned}$$

$$\begin{aligned}
l &= |cm(a, 5)| = O \cdot r = (d+1) \cdot (\frac{d^2}{2} + \frac{d}{2} - 1) = \\
&= \frac{d^3}{2} + \frac{d^2}{2} - d + \frac{d^2}{2} + \frac{d}{2} - 1 = \frac{d^3}{2} + d^2 - \frac{d}{2} - 1
\end{aligned}$$

$$\begin{aligned}
l &= |cm(a, 5)| = -b \cdot s = (d^2 + d - 2) \cdot (\frac{d}{2} + \frac{1}{2}) = \\
&= \frac{d^3}{2} + \frac{d^2}{2} - d + \frac{d^2}{2} + \frac{d}{2} - 1 = \frac{d^3}{2} + d^2 - 1
\end{aligned}$$