



Why Functional Programming Matters

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Overview

- ▶ Who is this guy?
- ▶ Misconceptions
- ▶ What is this strange language?
- ▶ Gluing functions together
- ▶ Gluing programs together
- ▶ Why now?

John Hughes

Trinity College, Cambridge

Oxford PhD

Chalmers University

Haskell

QuickCheck

QuviQ

Erlang



CHALMERS

Less is not more

Omitting features can't make a language more powerful.

Lack of side-effects cannot account for order of magnitude reduction in LOC.

goto structured programming analogy

Modular programs are easier to code, extend and test.

Good glue

“A chair can be made quite easily by making the parts - seat, legs, back etc. and sticking them together in the right way. But this depends on an ability to make joints and wood glue. Lacking that ability, the only way to make a chair is to carve it in one piece out of a solid block of wood, a much harder task.”

2 new kinds of glue: high-order functions & lazy evaluation

Miranda

$\text{listof } X ::= \text{nil} \mid \text{cons } X (\text{listof } X)$

$[]$ means nil

$[1]$ means $\text{cons } 1 \text{ nil}$

$[1,2,3]$ means $\text{cons } 1 (\text{cons } 2 (\text{cons } 3 \text{ nil}))$

$\text{sum nil} = 0$

$\text{sum (cons num list)} = \text{num} + \text{sum list}$

Gluing functions together

reduce f a

Replaces all occurrences of cons in a list by f, and all occurrences of nil by a.

copy = reduce cons nil

sum = reduce add 0

sum [1,2,3]

cons 1 (cons 2 (cons 3 nil))

add 1 (add 2 (add 3 0))

6

Append

append a b = reduce cons b a

append [1,2] [3,4]

reduce cons [3,4] [1,2]

(reduce cons [3,4]) (cons 1 (cons 2 nil))

cons 1 (cons 2 [3,4]))

[1,2,3,4]

Datatypes

$\text{treeof } X ::= \text{node } X \text{ (listof (treeof } X))$

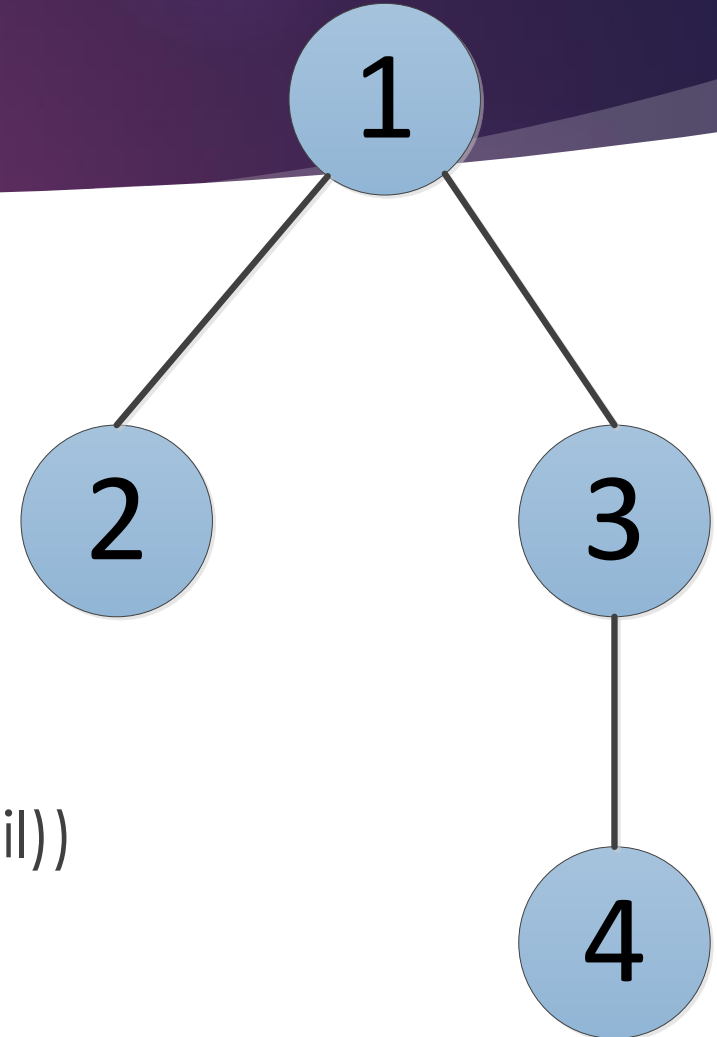
node 1

(cons (node 2 nil)

(cons (node 3

(cons (node 4 nil) nil))

nil))



Reduce Tree

$\text{redtree } f \ g \ a \ (\text{node label subtrees}) =$
 $f \ \text{label} \ (\text{redtree}' \ f \ g \ a \ \text{subtrees})$

$\text{redtree}' \ f \ g \ a \ (\text{cons subtree rest}) =$
 $g \ (\text{redtree } f \ g \ a \ \text{subtree}) \ (\text{redtree}' \ f \ g \ a \ \text{rest})$

$\text{redtree}' \ f \ g \ a \ \text{nil} = a$

Labels

labels = redtree cons append nil

cons 1

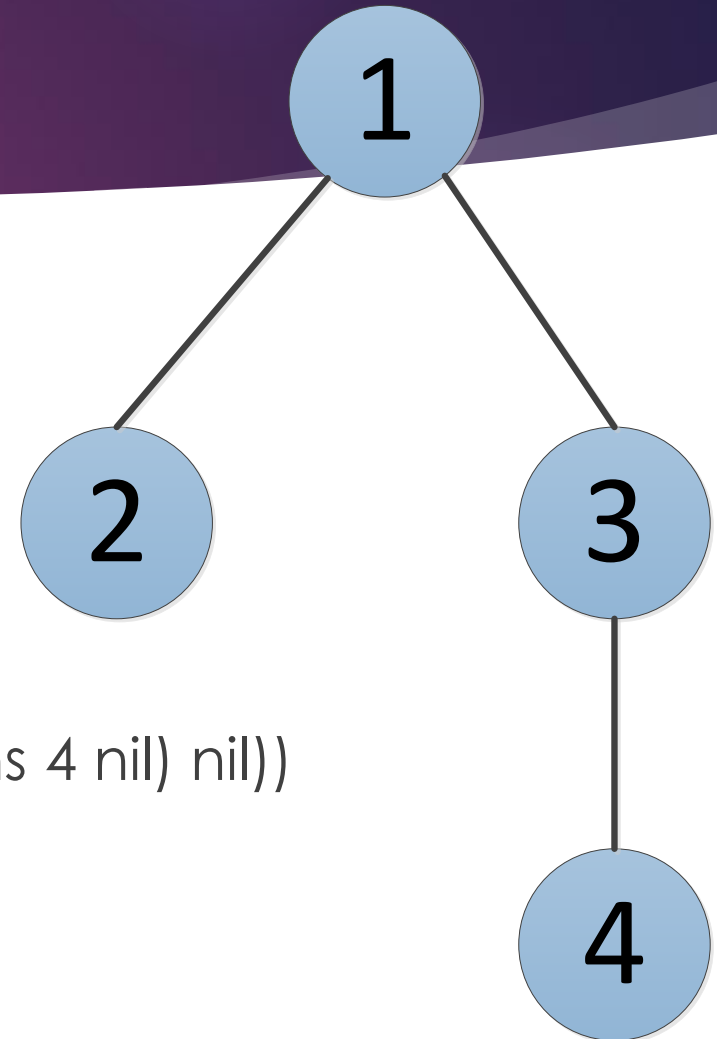
(append (cons 2 nil)

(append (cons 3

(append (cons 4 nil) nil))

nil))

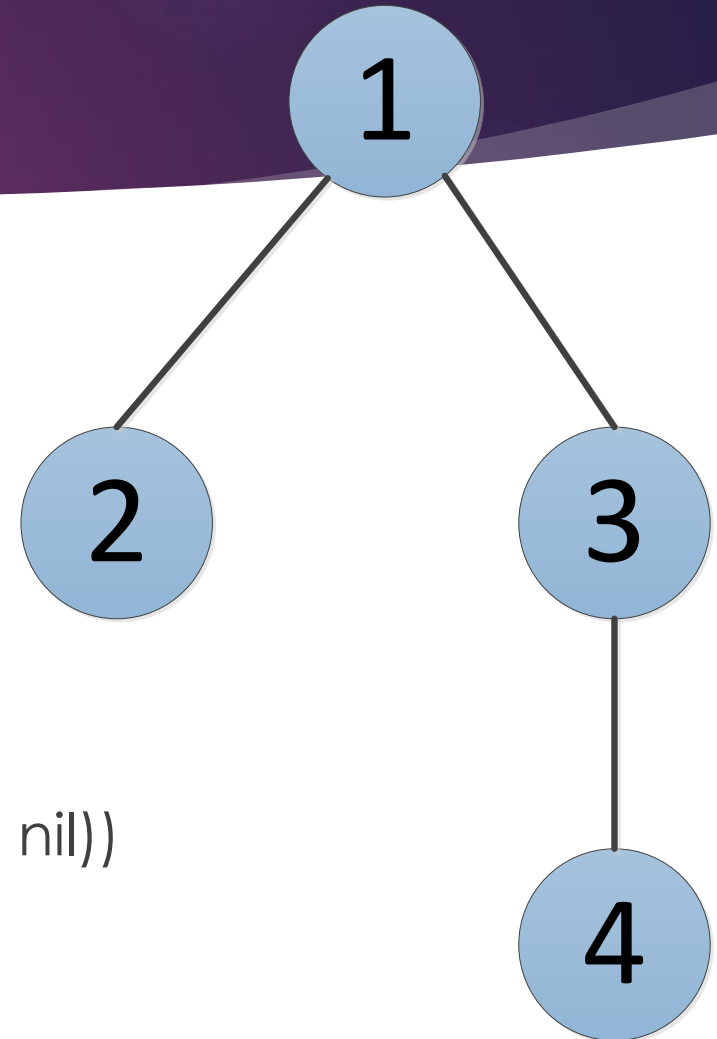
[1,2,3,4]



Map Tree

$(f . g) x = f (g x)$ $(\text{node} . f) x = \text{node} (f x)$
 $\text{maptree } f = \text{redtree } (\text{node} . f) \text{ cons nil}$

```
node (f 1)
  (cons (node (f 2) nil)
        (cons (node (f 3)
                     (cons (node (f 4) nil) nil))
              nil))
```



Simplicity & Performance

- ▶ Underscore
- ▶ React
- ▶ unidirectional dataflow
- ▶ Om
- ▶ Flux
- ▶ Mori
- ▶ Immutable



Gluing programs together

$(g . f)$

$g (f \text{ input})$

“ f is only started once g tries to read some input, and only runs for long enough to deliver the output g is trying to read. Then f is suspended and g is run until it tries to read another input. As an added bonus, if g terminates without reading all of f 's output then f is aborted.”

Square Roots

Newton Raphson algorithm:

Start with an initial approximation a_0 .

Compute better and better approximations based on previous.

Stop when approximations differ by less than a tolerance ϵ .

`sqrt a0 eps N = within eps (repeat (next N) a0)`

Generator

Generate a list of approximations.

```
repeat f a = cons a (repeat f (f a))  
[a, f a, f(f a), f(f(f a)), ...]
```

```
next N x = (x + N/x) / 2  
repeat (next N) a0
```


Selector

“[T]akes a tolerance and a list of approximations and looks down the list for two successive approximations that differ by no more than the given tolerance.”

within eps (cons a (cons b rest)) =
= b, if $\text{abs}(a-b) \leq \text{eps}$
= within eps (cons b rest), otherwise

Streams

- ▶ Lisp thunks
- ▶ ECMAScript Harmony
- ▶ yield
- ▶ generators
- ▶ coroutines
- ▶ callback hell
- ▶ Koa

Credits

It's Raining Haskell

an interview with Simon Peyton-Jones & John Hughes

<http://channel9.msdn.com/Blogs/Charles/YOW-2011-Simon-Peyton-Jones-and-John-Hughes-Its-Raining-Haskell>

Erlang Factory 2015 Panel Discussion

José Valim, Robert Virding, John Hughes & Guido van Rossum

<http://www.youtube.com/watch?v=oZwfi8JZ3kU>

More Credits

Immutability: Putting the Dream Machine to Work

by David Nolen

<https://www.youtube.com/watch?v=SiFwRtCnxv4>

Immutable Data in React & Flux

by Lee Byron

<http://www.youtube.com/watch?v=l7ldS-PbEgl>

Final Credits

Callbacks vs Coroutines

by TJ Holowaychuk

<http://medium.com/@tjholowaychuk/callbacks-vs-coroutines-174f1fe66127>