

# Confluence of Graph Rewriting with Interfaces

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University College London

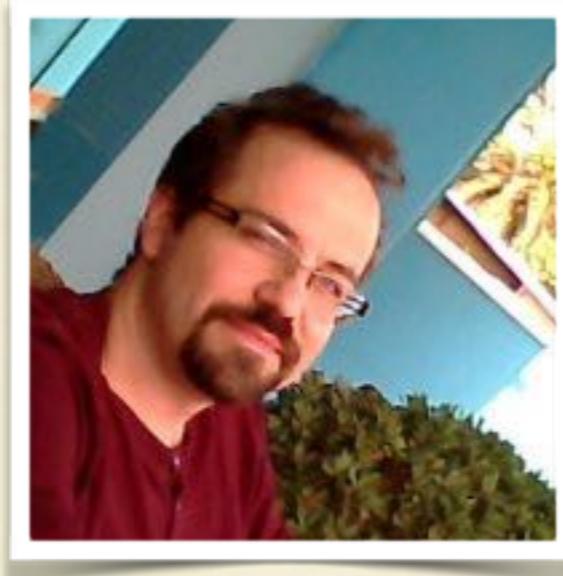
April 26, 2017



# Joint work with



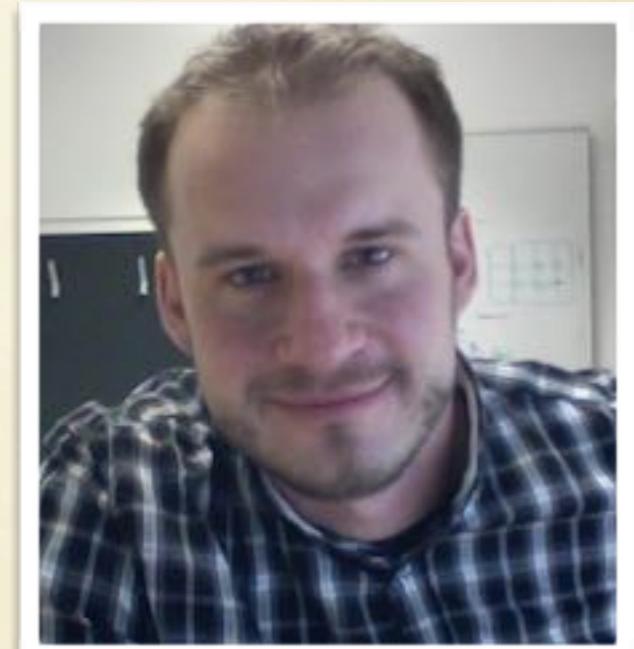
Filippo  
Bonchi



Fabio  
Gadducci



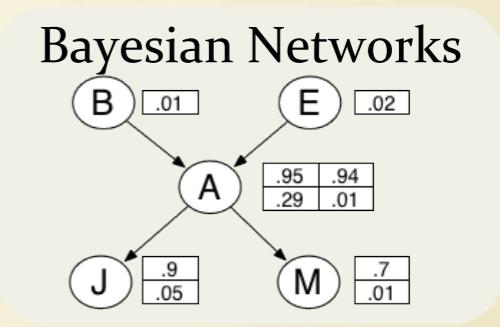
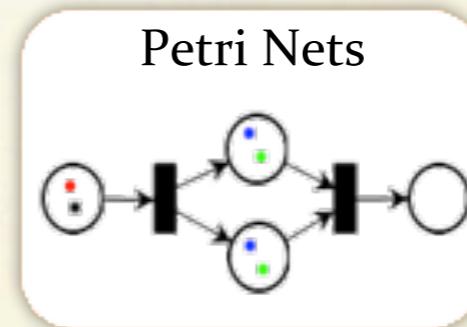
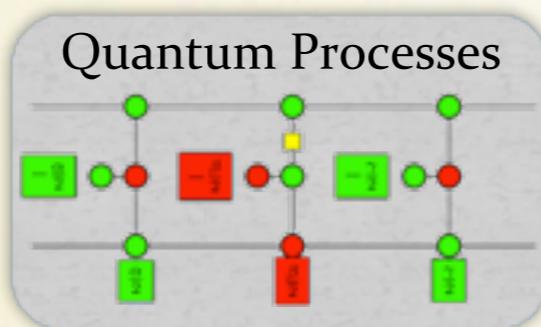
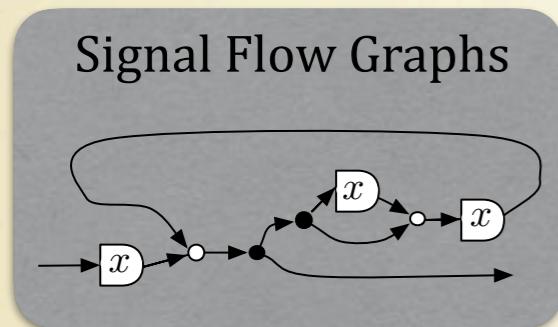
Aleks  
Kissinger



Paweł  
Sobociński

# Context and Motivation

# The algebraic approach to network diagrams



Diagrammatic languages studied using the compositional methods of programming language theory.

Syntax

+

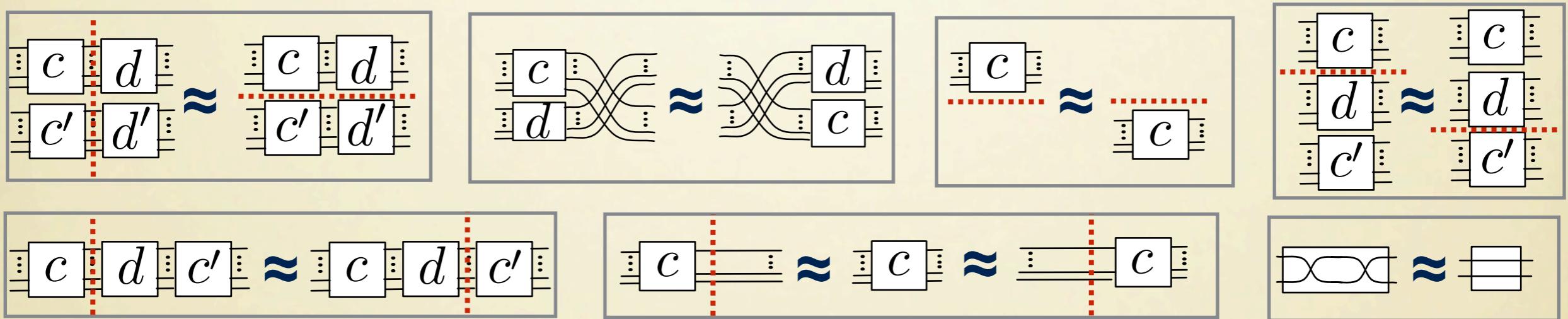
Equations

# Graphical syntax

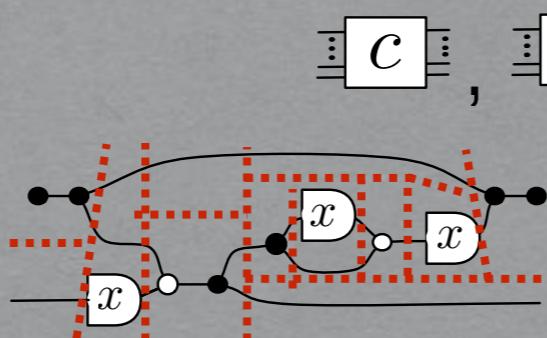
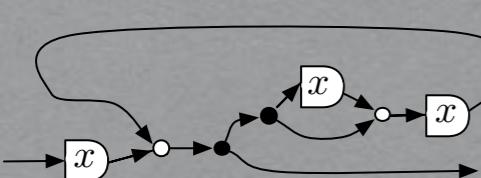
$$o \in \Sigma$$

$\boxed{c}, \boxed{d} ::= \square \mid \square \mid \square \mid \boxed{o} \mid \begin{array}{c} c \\ \hline d \end{array} \mid \boxed{c} \boxed{d}$

Quotiented by the *laws of symmetric monoidal categories*:



Signal Flow Graphs



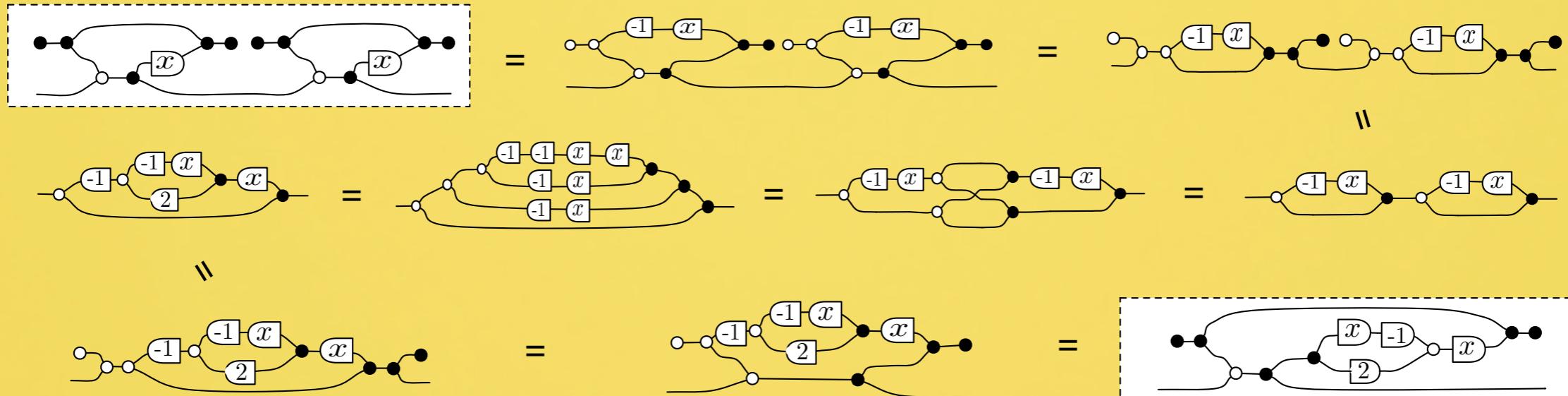
$$\boxed{c}, \boxed{d} ::=$$

$$\begin{array}{c} \bullet \\ \square \\ \bullet \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} \bullet \\ \square \\ \bullet \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} k \\ \square \\ k \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} x \\ \square \\ x \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} D \\ \square \\ D \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} o \\ \square \\ o \\ \square \\ \square \\ \square \end{array}$$

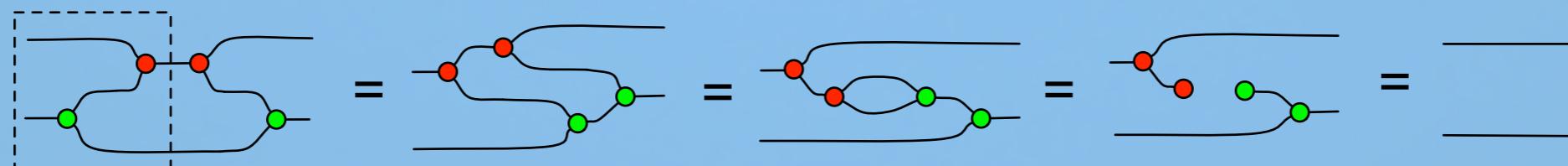
$$\begin{array}{c} c \\ \hline d \\ \square \\ \square \\ \square \\ \square \end{array} \mid \begin{array}{c} c \\ \hline d \\ \square \\ \square \\ \square \\ \square \end{array}$$

# Equational theories of diagrams

## Signal flow graphs: circuit equivalence



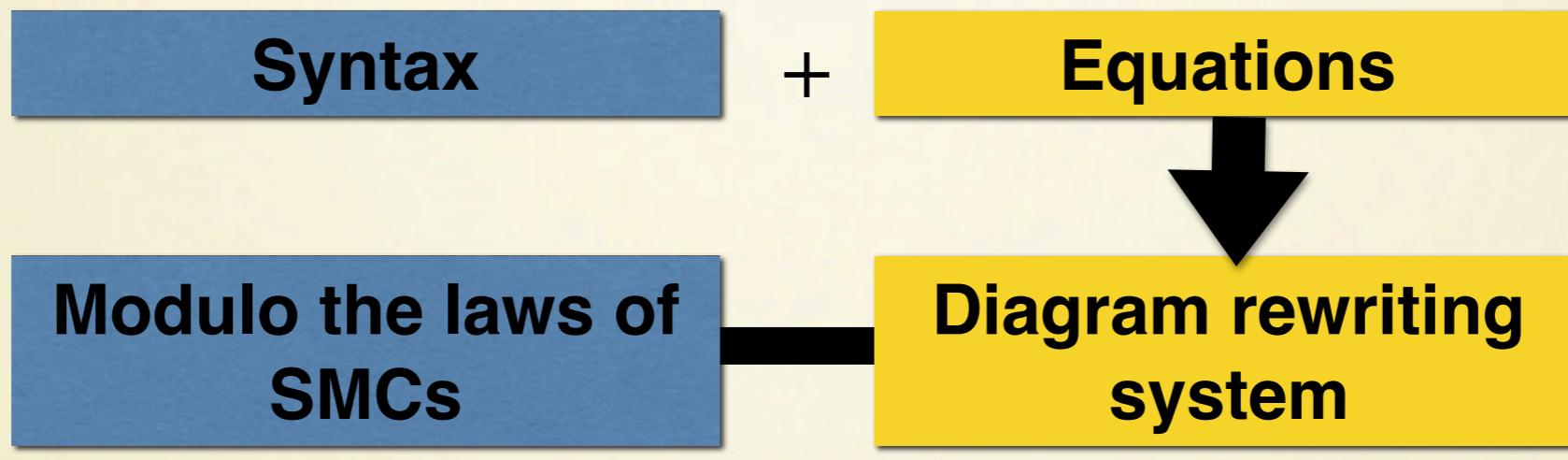
## Quantum processes: the CNOT gate is unitary



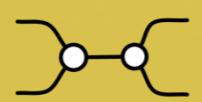
## Bayesian reasoning: disintegration



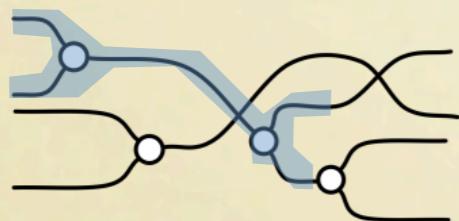
# Perspective of this work



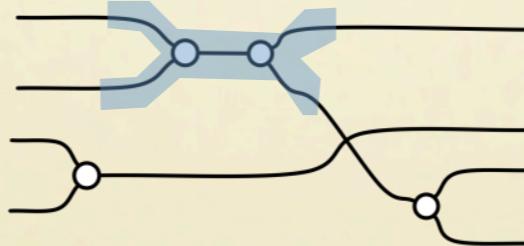
(R)



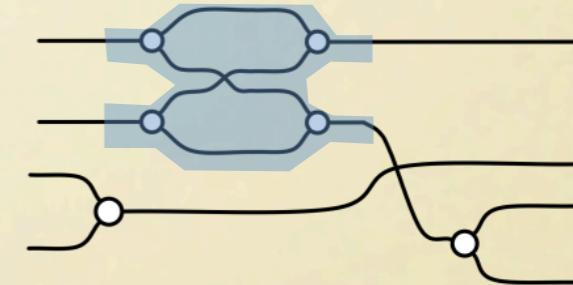
$\Rightarrow$



$\approx_{SMC}$



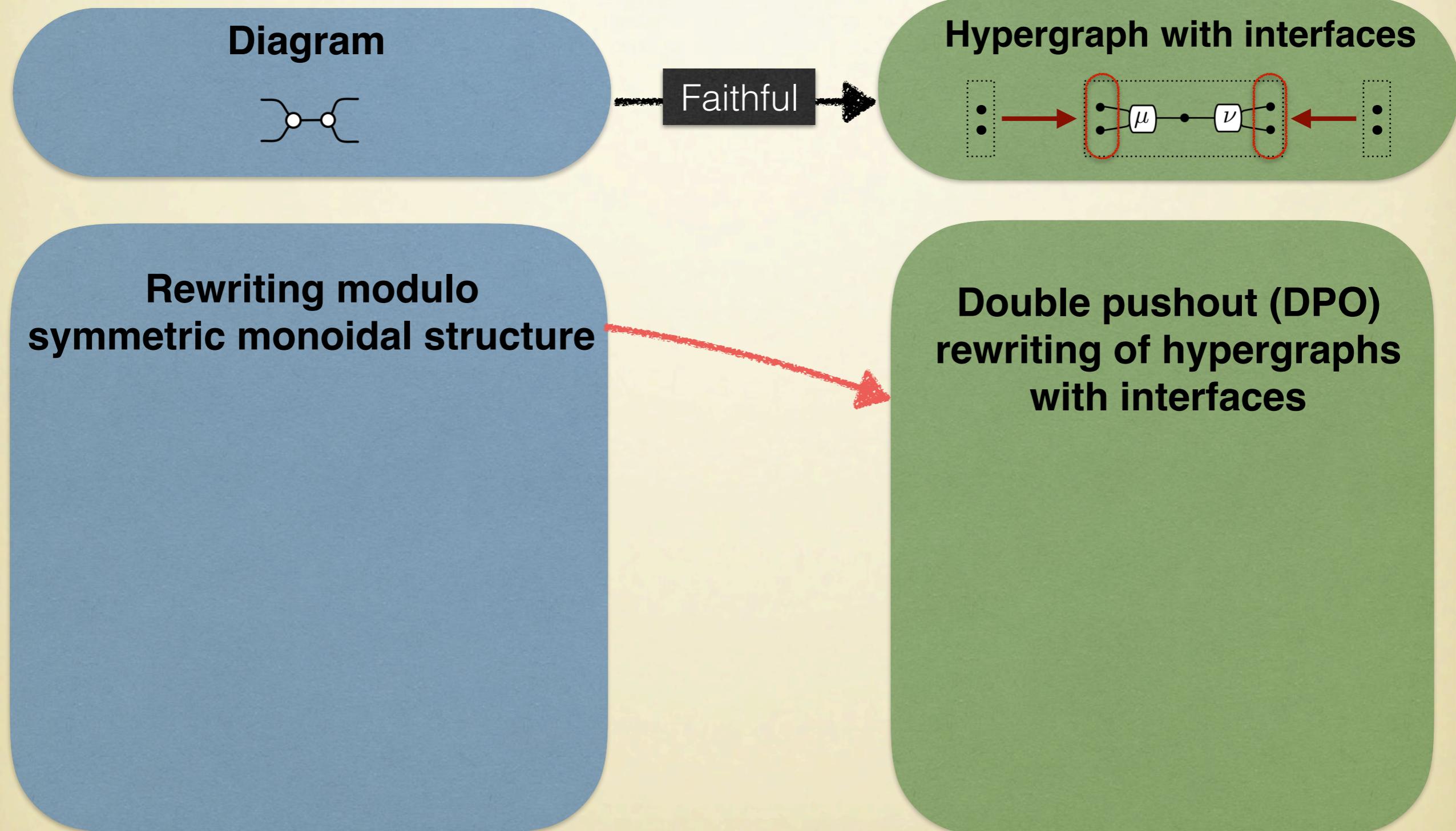
$\Rightarrow_R$



**Solution: diagram rewriting is interpreted as graph rewriting with interfaces**

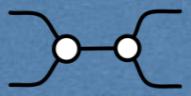
# Implementing Diagram Rewriting

# The graph interpretation (LICS'16)

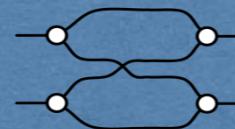


# DPO rewriting with interfaces

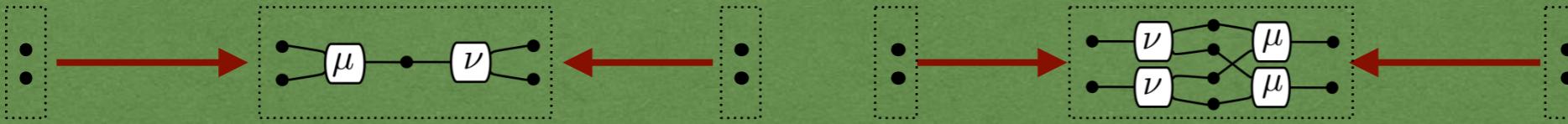
(R)



⇒



Syntax

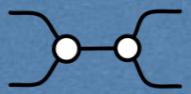


Hypergraphs  
with interfaces

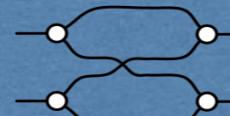


# DPO rewriting with interfaces

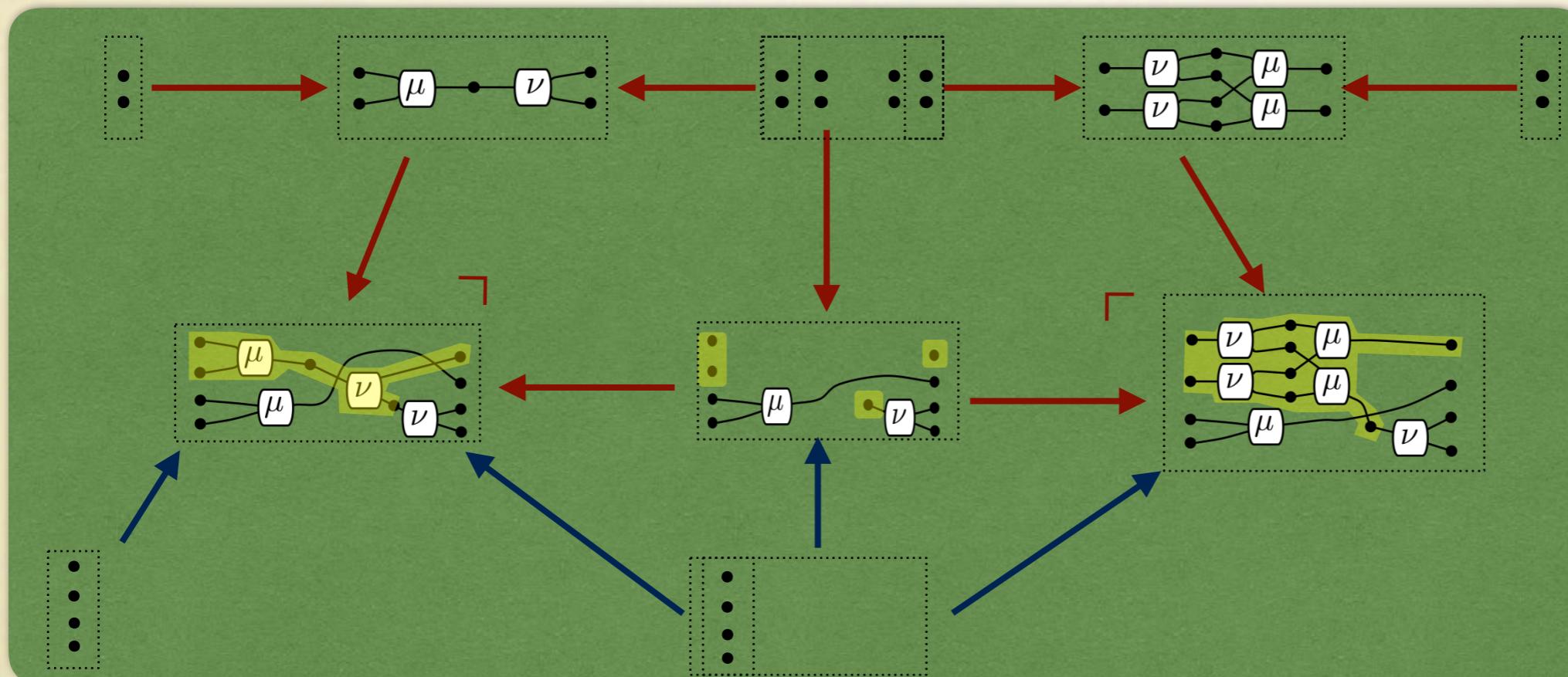
(R)



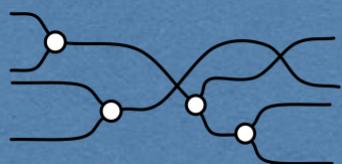
$\Rightarrow$



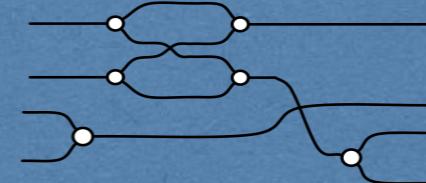
Syntax



Syntax

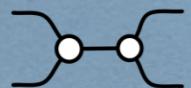


$\Rightarrow_R$



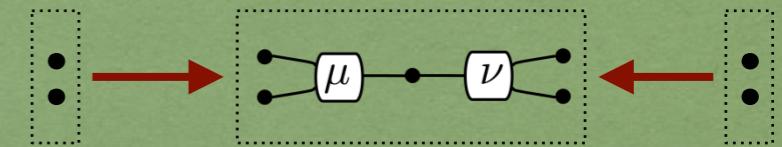
# The graph interpretation (LICS'16)

Diagram



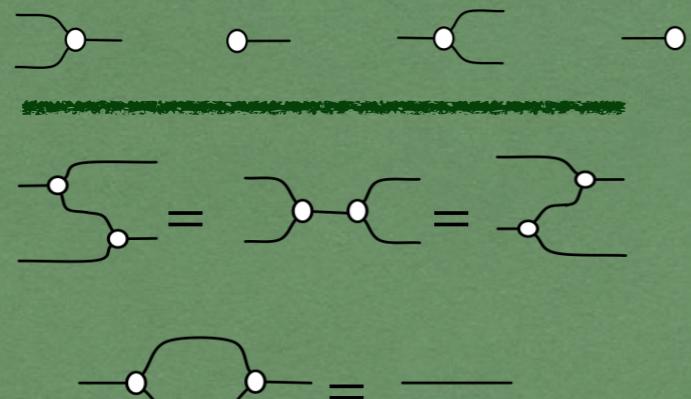
Faithful

Hypergraph with interfaces



Rewriting modulo  
symmetric monoidal structure

Rewriting modulo  
SM + Frobenius structure



Complete  
but unsound

Sound &  
Complete

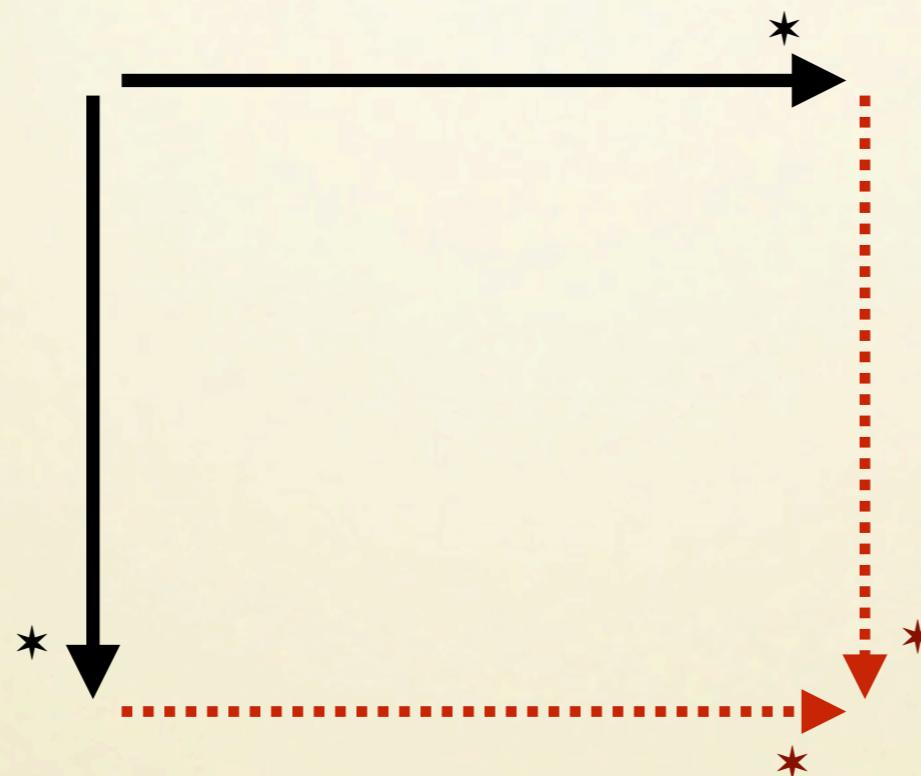
Sound &  
Complete

Double pushout (DPO)  
rewriting of hypergraphs  
with interfaces

Convex DPO rewriting

# Confluence

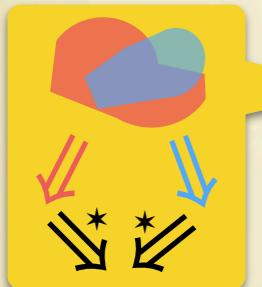
# Confluence, abstractly



If  $E$  is confluent & terminating  
then  $x \stackrel{E}{=} y$  becomes decidable.

# Decidability of Confluence

In term rewriting, confluence is **decidable** for terminating systems



All the critical pairs  
are joinable



The system is  
confluent

(Knuth-Bendix)

In DPO (hyper)graph rewriting, confluence is **undecidable** (Plump)

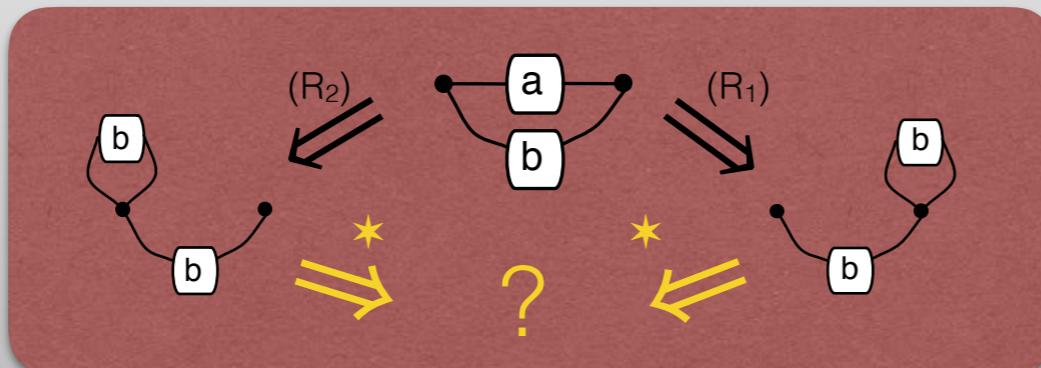
$$\bullet \xrightarrow[0]{a} \bullet \xrightarrow[1]{a} \bullet \Rightarrow \bullet \xrightarrow[0]{b} \bullet \quad (R_1)$$

$$\bullet \xrightarrow[0]{a} \bullet \xrightarrow[1]{a} \bullet \Rightarrow \bullet \xrightarrow[0]{b} \bullet \quad (R_2)$$

$$\bullet \xleftarrow{(R_2)} \bullet \xrightarrow{(R_1)} \bullet$$

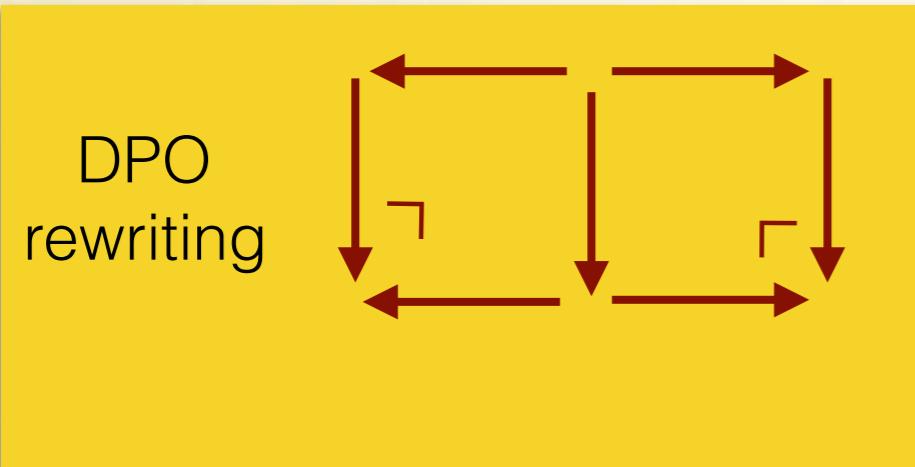
$$\bullet \xleftarrow{(R_1)} \bullet \xrightarrow{(R_2)} \bullet$$

All the critical pairs  
are joinable...



... but the system  
is not confluent.

# Interfaces to the Rescue

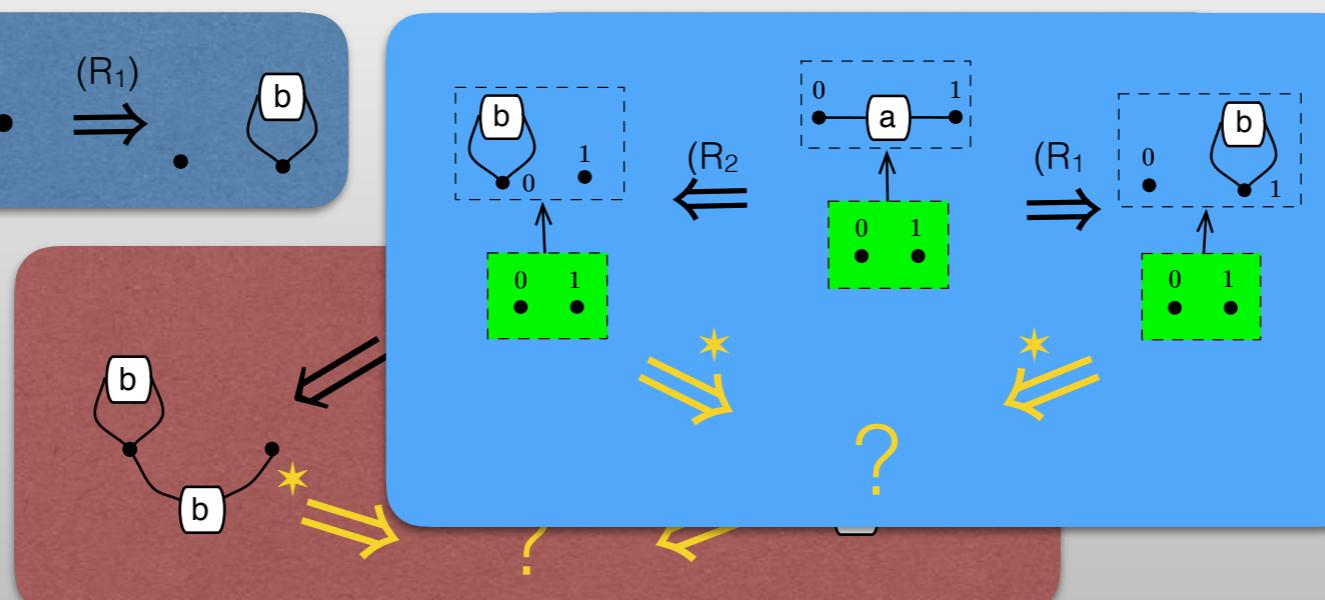
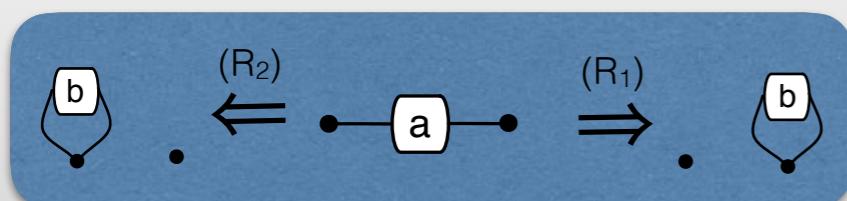


VS



$$\begin{array}{ccc} \overset{0}{\bullet} & \boxed{a} & \overset{1}{\bullet} \\ \Rightarrow & & \end{array} \quad \begin{array}{c} \bullet \\ \text{---} \\ \text{---} \\ \text{---} \\ \bullet \end{array} \quad (R_1)$$

$$\begin{array}{ccc} \overset{0}{\bullet} & \boxed{a} & \overset{1}{\bullet} \\ \Rightarrow & & \end{array} \quad \begin{array}{c} \bullet \\ \text{---} \\ \text{---} \\ \text{---} \\ \bullet \end{array} \quad (R_2)$$



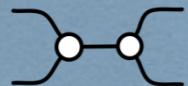
All the transition pairs  
are enabled ...

... but the system  
is not confluent.

**Theorem** In DPO rewriting with *interfaces*, confluence is decidable.

# The graph interpretation

## Diagram

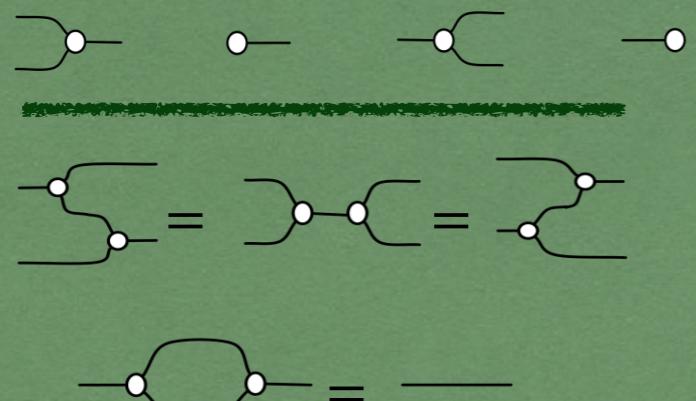


## Corollary II

Confluence is decidable for connected systems

## Rewriting modulo symmetric monoidal structure

## Rewriting modulo SM + Frobenius structure



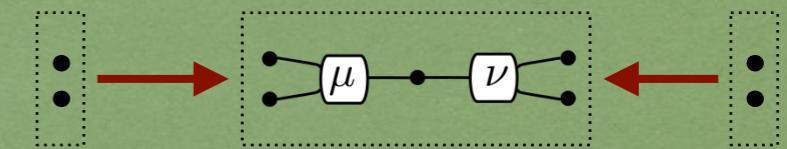
## Corollary I

Confluence is decidable

Sound & Complete

Sound & Complete

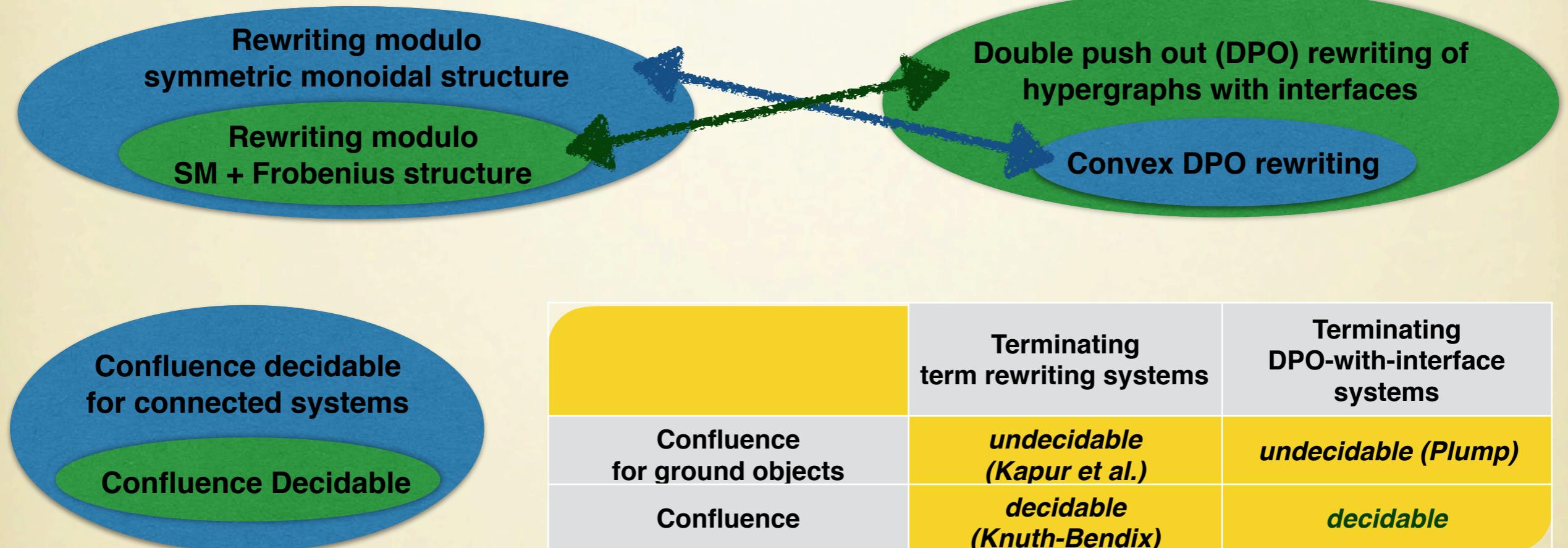
## Hypergraph with interfaces



## Double pushout (DPO) rewriting of hypergraphs with interfaces

## Convex DPO rewriting

# Conclusions



## Applications:

Algebras of network diagrams in quantum theory, control theory, concurrency, ....

Any adhesive category

What's next: termination

