CS 455: Principles of Database Systems

Review Guide 5: Transactions

- 1. Explain the distinction between a serial schedule and a serializable schedule.
- 2. What is a recoverable schedule? Why is recoverability of schedules desirable? Are there any circumstances under which it would be desirable to allow nonrecoverable schedules? Explain your answer.
- 3. Consider the following three transactions,

$$T_1: r_1(B); r_1(A); c_1;$$

 $T_2: r_2(A); w_2(B); c_2;$
 $T_3: r_3(A); w_2(A); c_3;$

and their schedule, $S: r_3(A); r_2(A); r_1(B); w_2(B); r_1(A); w_2(A); c_2; c_1; c_3;$

- (a) Give the precedence graph for *S*, and determine whether *S* is conflict serializable. If *S* is conflict serializable, give an equivalent serial ordering.
- (b) Is *S* possible using (basic) 2PL? If so, indicate in the schedule where locks can be acquired and released by T_1 , T_2 , T_3 . If not, explain why a 2PL schedule is not possible.
- 4. Consider a database with objects X and Y and assume that there are two transactions T_1 and T_2 :

$$T_1: r_1(X); r_1(Y); w_1(X); c_1;$$

 $T_2: r_2(X); r_2(Y); w_2(X); w_2(Y); c_2;$

- (a) Give an example nonserial schedule with actions of transactions T_1 and T_2 on objects X and Y that results in a RAW conflict.
- (b) Give an example nonserial schedule with actions of transactions T_1 and T_2 on objects X and Y that results in a WAR conflict.
- (c) Give an example nonserial schedule with actions of transactions T_1 and T_2 on objects X and Y that results in a WAW conflict.

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- (d) For each of the three schedules, show that Strict 2PL disallows the schedule.
- 5. What benefit does Strict 2PL provide? What disadvantages result?