

Written Questions

1. a) 2^{100} , 2^{2013} , $4\log(n)$, $\log(n^{10})$, $70n$, $n\log(n)$, $2n\log(n^2)$, $n^{1.03}$, $5n^6$, $8n^6+5n^2$, $2n^{6.5}$, 2.2^n , 2.5^n

b) $n\log(n)$ is the big-theta of $2n\log(n^2)$

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2. $O() = C + C + Cn*(C + Cn*(C) + C) + C$
> removing constant terms

$O() = n*(n)$
> simplifying the expression
Big-O = $O(n^2)$

 $\Omega() = C + C + C*(C + C*(C) + C) + C$
> simplifying the expression
Big-Omega = $\Omega(1)$

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3. a) i) $f(n)$ is in $\theta(g(n))$
ii) $f(n)$ is in $O(g(n))$
iii) $f(n)$ is in $O(g(n))$
iv) $f(n)$ is in $O(g(n))$
v) $f(n)$ is in $O(g(n))$
vi) $f(n)$ is in $\Omega(g(n))$
vii) $f(n)$ is in $\Omega(g(n))$
viii) $f(n)$ is in $O(g(n))$
ix) $f(n)$ is in $O(g(n))$
x) $f(n)$ is in $\Omega(g(n))$
xi) $f(n)$ is in $\Omega(g(n))$

b) i) time quadruples when n doubles, which hints towards a Big-O that is $O(n^2)$
ii) the formula $\text{executionTime}^{(1/n)}$ gives a value very close to 1 for all n that are shown,
I therefore assume that the Big-O is $O(C^n)$ where C is an undefined constant.

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Programming Questions

My output/times :

```
1
    algone    : 1 0ms
    algtwo    : 1 0ms
    algthree  : 1 0ms
```

10

	algone	:	55	0ms
	algtwo	:	55	0ms
	algthree	:	55	0ms
100	algone	:	5050	0ms
	algtwo	:	5050	0ms
	algthree	:	5050	0ms
1000	algone	:	500500	0ms
	algtwo	:	500500	4ms
	algthree	:	500500	0ms
10000	algone	:	50005000	1ms
	algtwo	:	50005000	130ms
	algthree	:	50005000	0ms
100000	algone	:	5000050000	0ms
	algtwo	:	5000050000	11432ms
	algthree	:	5000050000	0ms
1000000	algone	:	500000500000	0ms
	algtwo	:	500000500000	1273096ms
	algthree	:	500000500000	0ms

Other computer :

1	algone	:	1	0ms
	algtwo	:	1	0ms
	algthree	:	1	0ms
10	algone	:	55	0ms
	algtwo	:	55	0ms
	algthree	:	55	0ms
100	algone	:	5050	0ms
	algtwo	:	5050	0ms
	algthree	:	5050	0ms
1000	algone	:	500500	0ms
	algtwo	:	500500	4ms
	algthree	:	500500	0ms
10000	algone	:	50005000	1ms
	algtwo	:	50005000	22ms
	algthree	:	50005000	0ms
100000	algone	:	5000050000	1ms
	algtwo	:	5000050000	2153ms
	algthree	:	5000050000	0ms

```
1000000
  algone      : 500000500000 3ms
  algtwo      : 500000500000 213577ms
  algthree    : 500000500000 0ms
```

Order of magnitude :

```
  algone      :  $O() = C + n \cdot C = O(n)$ 
  algtwo      :  $O() = C + n \cdot (n \cdot C) = O(n^2)$ 
  algthree    :  $O() = C = O(1)$ 
```

Comment on findings :

Although the numbers are much bigger on my computer, there is a clear and very large increase in execution time when using algorithm 2, and a smaller one on algorithm 1.

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Observations :

The execution time generally increased at a pace similar to the jumps in number of multiplications, but since the execution time is so small there were a lot of very large jumps (in micro seconds) in execution time.

Recurrence equations :

```
M(n) = M(n-3)+3
M(3) = 0
M(2) = 0
M(1) = 0
M(0) = 0
```

Asymptotic characterization :

```
M(n)      = M(n-3)+3
M(n-3)    = M(n-6)+3+3
M(n-6)    = M(n-9)+3+3+3
> M(k)    =  $3 \cdot k / 3$ 
Big-O     =  $O(n)$ 
```

Tail recursion :

No, the method `Multiply(A,n)` is not tail recursive since the return value is then multiplied by a number before being returned again.

Here is an example of the same method in tail-recursive form :

```
public static double multiply(int[] A, int n, int result) {  
    if (n < 4) {  
        return 0;  
    } else {  
        return multiply(A, n-3, A[n-1]*result);  
    }  
}
```