

# Outline

Given a resource semiring  $\mathbf{R}$  with elements  $s, t, u, \dots \in R$ , define the following equivalence relation on the set  $R^T := \{0_{\mathbf{R}}\} \cup \mu X.(R \setminus \{0_{\mathbf{R}}\}) + X \times X$ :

$$\tau \sim \rho \iff \begin{cases} \tau = s, \rho = t \text{ and } s =_R t \\ \tau = s, \rho = (\rho', \rho'') \text{ and } \tau \sim \rho' \text{ and } \tau \sim \rho'' \\ \tau = (\tau', \tau''), \rho = t \text{ and } \tau' \sim \rho \text{ and } \tau'' \sim \rho \\ \tau = (\tau', \tau''), \rho = (\rho', \rho'') \text{ and } \tau' \sim \rho' \text{ and } \tau'' \sim \rho'' \end{cases}$$

**Claim.**  $R^T / \sim$  can be endowed with usage semiring structure.

Adjust the  $R$ -LCA to account for the tree structure by augmenting the operation  $!_\tau : \mathcal{A} \rightarrow \mathcal{A}$  and the elements  $W_{\pi\rho}, F_\rho, \delta_{\pi\rho}$ :

$$!_\tau^T x := \begin{cases} [\ulcorner true \urcorner, !_s x] & \text{if } s = \text{contract}(\tau) \\ [\ulcorner false \urcorner, E(!_{\tau'}^T x, !_{\tau''}^T x)], & \text{if } \langle \tau', \tau'' \rangle = \text{contract}(\tau) \end{cases}$$

The outermost pairing with Bool allows us to inspect/ the shape of the resource annotation internally within the BCI algebra.

Futhermore, it is also possible to define  $W_{\pi\rho}, F_\rho, \delta_{\pi\rho}$ , but that involves quite a few subtleties.

As an example usage of tree resource annotation, let's define the semantics of  $\oplus$ -types and check the rules are sound wrt. to it.

$\llbracket (A \oplus B)(\gamma) \rrbracket := (|A(y)| \sqcup |B(y)|, \models_{A \oplus B(\gamma)})$ , where

$$a \models_{A \oplus B(\gamma)} (i, x) \text{ iff } \begin{aligned} & (\exists a_x \in \mathcal{A}. a = [\ulcorner true \urcorner, a_x] \wedge a_x \models_{A(\gamma)} x \wedge \ulcorner true \urcorner \models_{Bool(\gamma)} i) \vee \\ & (\exists b_x \in \mathcal{A}. a = [\ulcorner false \urcorner, b_x] \wedge b_x \models_{B(\gamma)} x \wedge \ulcorner false \urcorner \models_{Bool(\gamma)} i) \end{aligned}$$

$$\begin{array}{c} \frac{0\Gamma \vdash A \quad 0\Gamma \vdash B}{0\Gamma \vdash A \oplus B} \oplus\text{-type} \qquad \frac{\Gamma \vdash M \overset{\sigma}{:} A}{\Gamma \vdash \mathbf{inl} M \overset{\sigma}{:} A \oplus B} \mathbf{inl} \qquad \frac{\Gamma \vdash N \overset{\sigma}{:} B}{\Gamma \vdash \mathbf{inr} N \overset{\sigma}{:} A \oplus B} \mathbf{inr} \\[10pt] \frac{\Gamma \vdash M \overset{\sigma}{:} A \oplus B \quad \Gamma', a \overset{\rho}{:} A \vdash T_a \overset{\sigma}{:} C[\mathbf{inl} a/x] \quad \Gamma', b \overset{\pi}{:} B \vdash T_b \overset{\sigma}{:} C[\mathbf{inr} b/x] \quad 0\Gamma = 0\Gamma'}{\Gamma' + \langle \rho, \pi \rangle \Gamma \vdash \mathbf{case}(M, T_a, T_b) \overset{\sigma}{:} C[M/x]} \oplus\text{-elim} \end{array}$$

Figure 1: Rules for  $\oplus$ -types

WLOG, let's verify the soundness of **inl** and **case**.

Given a realiser  $a_M$  for  $\forall \gamma \in |\Gamma|. M(\gamma)$ , let  $a_{\mathbf{inl}} := \lambda^* a_\gamma. [\ulcorner true \urcorner, a_M \cdot a_\gamma]$

$$\begin{aligned}
\forall \gamma \in |\Gamma|, a_\gamma \in \mathcal{A}. a_\gamma \Vdash_\Gamma y &\implies a_M \cdot a_y \Vdash_{A \oplus B(\gamma)} M(\gamma), \\
\forall \gamma' \in |\Gamma'|, s \in |A(\gamma')|, a'_\gamma \in \mathcal{A}, a_s \in \mathcal{A}. a'_\gamma \Vdash_\Gamma \gamma' \wedge a_s \Vdash_{A(\gamma)} s &\implies a_T \cdot [a'_\gamma, !_\rho a_s] \Vdash_{C(\gamma', \text{inl}(s))} \\
T_a(\gamma', s), \\
\forall \gamma' \in |\Gamma'|, t \in |B(\gamma')|, a'_\gamma \in \mathcal{A}, a_t \in \mathcal{A}. a'_\gamma \Vdash_\Gamma \gamma' \wedge a_t \Vdash_{B(\gamma)} t &\implies b_T \cdot [a'_\gamma, !_\pi a_t] \Vdash_{C(\gamma', \text{inr}(t))} \\
T_b(\gamma', t),
\end{aligned}$$

Problem : We end up with something of the shape  $!_{\langle \pi, \rho \rangle}([\ulcorner \text{boolean} \urcorner, x])$  and want to select the projection depending on the inner boolean value.