Outline

$$\frac{0\Gamma \vdash A \quad 0\Gamma \vdash B}{0\Gamma \vdash \rho A \oplus \pi B}\ \oplus\text{-type} \qquad \frac{\Gamma \vdash S_1 \overset{\sigma}{:} A}{\rho\Gamma \ \vdash \mathbf{inl}\, S_1 \overset{\sigma}{:} \rho A \oplus \pi B}\ \text{inl} \qquad \frac{\Gamma \vdash S_2 \overset{\sigma}{:} B}{\pi\Gamma \vdash \mathbf{inr}\, S_2 \overset{\sigma}{:} \rho A \oplus \pi B}\ \text{inr}$$

$$\frac{\begin{array}{c} 0\Gamma, x \overset{0}{:} \rho A \oplus \pi B \vdash C \\ \Gamma \vdash M \overset{\sigma}{:} \rho A \oplus \pi B \qquad \Gamma', a \overset{\rho}{:} A \vdash T_1 \overset{\sigma}{:} C[\mathbf{inl}\, a/x] \qquad \Gamma', b \overset{\pi}{:} B \vdash T_2 \overset{\sigma}{:} C[\mathbf{inr}\, b/x] \qquad 0\Gamma = 0\Gamma' \end{array}}{\Gamma' + \Gamma \vdash \mathbf{case}(M, T_1, T_2) \overset{\sigma}{:} C[M/x]}\ \oplus\text{-elim}$$

$$\frac{\Gamma \vdash S_1 \overset{\sigma}{:} A \qquad \Gamma \vdash M \overset{\sigma}{:} \rho A \oplus \pi B \qquad \Gamma', a \overset{\rho}{:} A \vdash T_1 \overset{\sigma}{:} C[\mathbf{inl}\, a/x] \qquad 0\Gamma = 0\Gamma'}{\Gamma' + \rho\Gamma \vdash \mathbf{case(inl}(S_1), T_1, T_2) \equiv T_1[S_1/a]}\ \oplus\text{-comp}$$

Figure 1: Rules for $\oplus$-type

We give the following semantics for the $\oplus$-type:
$|\rho A \oplus \pi B\ (\gamma)| := |A(\gamma)| \sqcup |B(\gamma)|$
$a \vDash_{\rho A \oplus \pi B\ (\gamma)} (i, x)$ iff $(\exists b. a = [!_\rho b, \ulcorner true \urcorner] \wedge b \vDash_{A(\gamma)} x \wedge i = 0) \vee$
$(\exists c. a = [!_\pi c, \ulcorner false \urcorner] \wedge c \vDash_{B(\gamma)} x \wedge i = 1)$

Claim. *The rules are sound when interpreted wrt to the given semantics for $\oplus$-types and realisability model.*

Proof. The underlying set-theoretic functions are immediate. For the realisers, let

- $a_{\mathbf{inl}} := \lambda^* x.[F_\rho \cdot\ !_\rho a_{s_1} \cdot x, \ulcorner true \urcorner]$

  if $a_{s_1} \cdot a_\gamma \vDash s_1$, *then* $a_{\mathbf{inl}} \cdot !_\rho a_\gamma \vDash \mathbf{inl} s_1$; $a_{\mathbf{inl}} \cdot !_\rho a_\gamma \rightsquigarrow [!_\rho(a_{s_1} \cdot a_\gamma), \ulcorner true \urcorner]$

- $a_{\mathbf{case}} := \begin{array}{l} \lambda^* x.\ \text{let } [a'_\gamma, a_\gamma] = x, \\ \quad [a, b] = a_m \cdot a_\gamma \text{ in} \\ \quad E(b, a_{T_1}, a_{T_2}) \cdot [a'_\gamma, a] \end{array}$

  assuming $a_m \cdot a_\gamma \vDash M$, $a_{T_1} \cdot [a'_\gamma, !_\rho a_a] \vDash T_1$, $a_{T_2} \cdot [a'_\gamma, !_\pi a_b] \vDash T_2$, *then we want to find* $a_{\mathbf{case}}$, *s.t.* $a_{\mathbf{case}} \cdot [a'_\gamma, a_\gamma] \vDash \mathbf{case}(M, T_1, T_2)$.

  if $a_m \cdot a_\gamma = [!_\rho a_a, \ulcorner true \urcorner]$, *then* $a_{\mathbf{case}} \cdot [a_\gamma, a'_\gamma] \rightsquigarrow E(\ulcorner true \urcorner, a_{T_1}, a_{T_2}) \cdot [a'_\gamma, !_\rho a_a] \rightsquigarrow a_{T_1} \cdot [a'_\gamma, !_\rho a_a]$
  if $a_m \cdot a_\gamma = [!_\pi a_b, \ulcorner false \urcorner]$, *then* ...

$\square$

Claim. *There is a bijection:*

$RTm(\Gamma, \Pi(x \overset{\tau}{:} \rho A \oplus \pi B)\, C) \cong RTm(\Gamma, \Pi(y \overset{\tau\rho}{:} A)\, C[\mathbf{inl} y/x]) \times RTm(\Gamma, \Pi(z \overset{\tau\pi}{:} B)\, C[\mathbf{inl} z/x]))$

(natural in Γ).

Proof. Given a term $\Gamma \vdash f \overset{1}{:} (x \overset{\tau}{:} \rho A \oplus \pi B) \to C$, we can derive another term $\Gamma \vdash f^l \overset{1}{:} (y \overset{\tau\rho}{:} A) \to C[\mathbf{inl} y/x]$:

$$\dfrac{\dfrac{\Gamma \vdash f : (x \overset{\tau}{:} \rho A \oplus \pi B) \to C}{\Gamma, y \overset{0}{:} A \vdash f \overset{1}{:} (x \overset{\tau}{:} \rho A \oplus \pi B) \to C}\,Weak \quad \dfrac{\dfrac{\vdash 0\Gamma, y \overset{1}{:} A}{0\Gamma, y \overset{1}{:} \vdash y \overset{1}{:} A}\,var}{0\Gamma, y \overset{\rho}{:} A \vdash \mathbf{inl} y \overset{1}{:} \rho A \oplus \pi B}\,inl}{\dfrac{\Gamma, y \overset{\tau\rho}{:} A \vdash f(\mathbf{inl} y) \overset{1}{:} C[\mathbf{inl} y/x]}{\Gamma \vdash \lambda y \overset{\tau\rho}{:} A.\, f(\mathbf{inl} y) : (y \overset{\tau\rho}{:} A) \to C[\mathbf{inl} y/x]}\,Lam}\,App$$

Analogously, we can obtain $\Gamma \vdash f^r \overset{1}{:} (z \overset{\tau\pi}{:} B) \to C[\mathbf{inr} y/x]$.
Now suppose we have terms $\Gamma \vdash l \overset{1}{:} (y \overset{\tau\rho}{:} A) \to C[\mathbf{inl} y/x]$ and $\Gamma \vdash r \overset{1}{:} (z \overset{\tau\pi}{:} B) \to C[\mathbf{inr} z/x]$.
Using the isomorphism $\Lambda^{\mathcal{L}}$, we get judgements $\Gamma, y \overset{\tau\rho}{:} A \vdash l^* : C[\mathbf{inl} y/x]$ and $\Gamma, z \overset{\tau\pi}{:} B \vdash r^* : C[\mathbf{inr} z/x]$. $\square$

Now we can focus on the relational part

$$\begin{aligned}
(\rho A \oplus \pi B)(\gamma)_R &:= A(\gamma)_R \sqcup B(\gamma)_R \\
(\rho A \oplus \pi B)(\gamma)_{refl} &:= A(\gamma)_{refl} \sqcup B(\gamma)_{refl} \\
(\rho A \oplus \pi B)(\gamma)_\sigma &:= A(\gamma)_\sigma \sqcup B(\gamma)_\sigma, \quad \sigma \in \{\text{src}, \text{tgt}\}
\end{aligned}$$

Suppose A, B are discrete. Then $\rho A \oplus \pi B$ is also discrete. (Since coproducts preserve isos).

1 Spellchecking

2 Polynomial functors in a linear setting

2.1 Category of closed types and linear functions

2.2 Internal representation

2.3 External representation (using adjoints)

2.4 Properties of quantative polynomial functors

3 Algebras for quantative polynomial functors

3.1 $\mathbb{N}$

3.2 Lists

3.3 Trees

3.4 Induction principle

4 Rules for W-types in QTT

5 Generalising to non-empty contexts

6 Parametricity and W-types